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Lab Title:	Implementing Red Black Tree & Treemap interface

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1 Time analysis of Functions

n : the number of nodes in the tree

1.1 Red Black Tree

1. `getRoot()`: $\mathcal{O}(1)$
2. `isEmpty()`: $\mathcal{O}(1)$
3. `clear()`: $\mathcal{O}(1)$
4. `search(T)`: $\mathcal{O}(\log n)$ as it uses method `findNode(T)` that has also run time $\mathcal{O}(\log n)$
5. `contains(T)`: $\mathcal{O}(\log n)$ as it uses method `findNode(T)`
6. `insert(T, V)`: consists of 3 steps
 - (a) ordinary BST insert for balanced tree which its worst run time is $\mathcal{O}(\log n)$
 - (b) the new node is already colored Red so the coloring run time is $\mathcal{O}(1)$
 - (c) fixing the properties of red black tree using method `redRed(node)` which its worst run time is $\mathcal{O}(\log n)$so the total worst run time $\mathcal{O}(\log n)$
7. `delete(T)`: consists of 4 steps
 - (a) using `findNode(T)` to get the node to be deleted in $\mathcal{O}(\log n)$
 - (b) in case it's an internal node find its replacement using `findMin(node)` $\mathcal{O}(\log n)$
 - (c) remove the node and replace it with a new node by method `removeAndReplace(node, node)` in $\mathcal{O}(1)$
 - (d) fixing the properties of red black tree using method `doubleBlack(node)` which its worst run time is $\mathcal{O}(\log n)$so the total worst run time $\mathcal{O}(\log n)$
8. `leftRotate(node)` & `rightRotate(node)`: $\mathcal{O}(1)$

2 Tree Map

1. `ceilingEntry(T)`: $\mathcal{O}(\log n)$ as it uses method `findNode(T)`
2. `ceilingKey(T)`: $\mathcal{O}(\log n)$ as it uses method `ceilingEntry(T)`
3. `clear()`: $\mathcal{O}(1)$ as it uses the red black tree `clear()` method
4. `containsKey(T)`: $\mathcal{O}(\log n)$ as it uses the red black tree `contains(T)` method
5. `containsValue(V)`: $\mathcal{O}(n)$ as it uses `values()` method

6. `entrySet()`: $\mathcal{O}(n)$ as it uses method `InorderTraversal(node)` which its run time is $\mathcal{O}(n)$
7. `firstEntry()`: $\mathcal{O}(\log n)$
8. `firstKey()`: $\mathcal{O}(\log n)$ as it uses `firstEntry()`
9. `floorEntry(T)`: $\mathcal{O}(\log n)$ as it uses `findNode(T)`
10. `floorKey(T)`: $\mathcal{O}(\log n)$ as it uses `floorEntry(T)`
11. `get(T)`: $\mathcal{O}(\log n)$ as it uses the red black tree `search()` method
12. `headMap(T)` & `headMap(T, boolean)`: $\mathcal{O}(n)$ as it uses method `InorderTraversal(node)`
13. `keySet()`: $\mathcal{O}(n)$ as it uses method `InorderTraversal(node)`
14. `lastEntry()`: $\mathcal{O}(\log n)$
15. `lastKey()`: $\mathcal{O}(\log n)$ as it uses `lastEntry()`
16. `pollFirstEntry()`: $\mathcal{O}(\log n)$ as it uses `firstEntry()` and `remove(T)`
17. `pollLastEntry()`: $\mathcal{O}(\log n)$ as it uses `lastEntry()` and `remove(T)`
18. `put(T, V)`: $\mathcal{O}(\log n)$ as it uses the red black tree `insert(T, V)` method
19. `putAll(map)`: $\mathcal{O}(n \log n)$ as it uses the red black tree `insert(T, V)` method n times
20. `remove(T)`: $\mathcal{O}(\log n)$ as it uses the red black tree `delete(T)` method
21. `size()`: $\mathcal{O}(n)$ as it uses method `InorderTraversal(node)`
22. `values()`: $\mathcal{O}(n)$ as it uses method `InorderTraversal(node)`