Final Exam

Program: IS Course: Time Series Database Course Code: IS

Level: 3 Lecturer: Dr. Hosam Refaat Date: 18/6/2023

Total Pages: 5 Total Marks: 70 Time Allowed: 120 min

Form

	followi	200	stions errect answe	C3	, c 2, c	\	2.3.1	(70 ma
- R1(A).w2(A).	-3(A).w1	(A),w3(A), Co	mmit3	. Commit2.	Commit1		170 1111
A) Rec	overable	B)	Cascadeless	(C)	Strict	Di	Not recoverable	
-	_		Commit2, A		outer		1.ot ice vetacie	
	overable .		Cascadeless	(C)	Strict	D	Not recoverable	
			(A), (3(A))	-	outer		110t recoverable	
	overable	(B)	Cascadeless	(C)	Strict	100	Not recoverable	
_			W1(A) W2 (A	_		1	17	
	overable	(B)		(C)	Strict	(D)	Not recoverable	
			D C2, C1	10)				
	overable	B		(C)	Strict	.Dr	Not recoverable	
	X), (W2(X	() (W1(X		mmit1				
	coverable			(C)	Strict	D)	Not recoverable	
7- R10	X), W2()	O WIO	Commit2, C	Commit	1	All the second of	Spirite Street Street Street	
All the second s	ecoverable			(C)	Strict	(D)	Not recoverable	
18- R1	X), R2(X), W1(X	Commit1, W	/2(X), C				
	coverable			2	Strict	D)	Not recoverable	
		The same of the sa	iit1, W2(X), C	ommit2		1		
	coverable	B)		L)	Strict	(D)	None of these	
	-), R1(Y)	, R2(X), Com	mit1, C	commit2			
	coverable	B)		(C)	Strict	(D)	None of these	
11	system h	as at mo	ost one user at	t a time				
	gle-User			ultiuser			E) None of these	
12- If the	DDMS o	loes not	have concurren	t contro	d, what cloud	be the prob	lem for the transact	ion show belo
Time	T1	T2						
1	Read (X);		1					
2	X := X - 10;							
3	Write (X);							
4		Read (X);						
5		X := X +20;						
- 6		Write (X):						
7		COMMIT;]					
8	Read (Y);							
9	ROLLBACK							
A) n	o probler	m			B) Dirty Rea	ad and inco	rrect summary	
	ost upda		ty read		D) Lost upd	ate and inc	orrect summary	
			ead and incorr	ect sum			•	
			the can be oc			a cobodulo	ic	

Examination Committee

Dr.Hosam Refaat



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						CONTRACTOR OF THE PERSON
		7,	12			
		med hard VI:				
		read_iterm(X); X:=X-N;				
			read_item(X);			
			X≥X+M;			
		write_item(X); read_item(Y);				
			write_item(X);			
		Y:=Y+N;				
		write_iterrs(Y);				
Lost Update	B) Dirty Read	C) Incor	rect Summary	D)	None of these	
14. The type of prob				e is		
		τ,	Τ ₂			
		read_item(X); X:=X-N;				
		write_Bern(X);				
			read_item(X);			
			X:nX+M;			
		mm and 100	write_serre(x);			
		Abort				
A) Lost Update	Dirty Read	-	ect Summary	D)	None of these] .
15- Basic unit of data	a transfer from th	e disk to the co	mputer main i	memory is	one record. (tru	ie/false)×
16- The deadlock av	oidance algorithn	n discovers that				
17- Every Cascadele	ansaction. (true/fa	ilse)\/	Malan)		10/10	1, eable
18- Precedence Gra				heduler ar		
(true/false) X	pii Aigoritiiii is u.	sed vandate the	concurrent sc	nedurer ar	e commer equivi	icit of not.
19- All transaction t	hat in the "partia	lly committed"	state will be co	ommitted s	uccessfully. (tru	e/false) X
20-Starvation occur	[The state of the s				
	er. (true/false) U		estampoentrajo so est∎ so aprin			
21- The concurrent						
22- the is ope		ollback except	that it applies t	to a single o	peration rather	than to a
whole transaction						_
A Undo	B) Redo	C) Both A and			None of these	_
	g: concurrent exe					1
24- process	B) Parallel sing: processes are					J
A) Interleaved	B) Parallel	C) Both	A and B D) None of	these	1
25- is logical	unit of database			A STATE OF THE PARTY OF THE PAR	The state of the s	read -
	insert or update					
A) Interleaved proce		nsaction) application [program	D) None of t	hese
-						

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1	Lost Update	B)	Temporary	Update	1) I	ocorrect Sun	mary	D) N	one of these	
\$ 1-	This schedule car	be c	assified as	RI	(x), W	L(x), R1(y)	. W1()	y), Ç1 , R	12(x),(W2(x)	, C2
A)	Recoverable	B)	Cascadeless	The second second	Strict	Married Street, Street	The same of the sa	ot Recove	AND DESCRIPTION OF THE PARTY OF]
42.	This schedule car	n be c	assified as	R2	(x), R1(x)	, W1(x), R1()), W1(y)_C1, W2(x), C2	
A)	Recoverable	B)		-	Strict			one of the		1
	The oper				hat all ti	e operatio	ns of a	committe	ed transaction	have been
-	applied successfu									
A)		(5)	Redo		h A and B		D)		None of these	
	Two schedules a								database.	
_	Result equivalent			_			ializable	e D)	Serializable	
45	The following so			-	alent. (tr	ue/false)				
		-	, RI(y), W2(y	500				_	T19002	
			, R2(x), R2(y		lant (tr	no falca)	WIT	() - B	2(X)	
46	- The following so		f(x), $WI(y)$, $F(x)$			ueriaise	0,0			
			(x). Rhy), W				Wic	رالا الم الرالا	15(3)	
4	- Real-Time Data					se computi				to operate in
	a timely manner									•
4	8-Being serializabl	e is th	e same as b	eing seria	al. (true	(false) ×				
4	9- The concurrency	cont	rol method i	nay deci	de to ab	ort the tra	nsactio	n, to be i	restarted late	r, because it
	violates serializa	bility.	(true/false)		1 .					
	0- A schedule is vie								ue/false)	
	1-Every Recoveral					A STATE OF THE PARTY OF THE PAR				
5	2- The transaction			ed auton	atically	by the sys	tem and	d is used	to identify on	e or more
	transaction. (tru			able (ton	· o/folos	/				
5	3- Every Strict Sch	edule	s is Kecover	abie. (tri	d no vel	×			C 41- 4	. that mice
)	 4 is not executed their deadlines. 		atter its dea	anne an	uno vai	ue is gaine	by the	e system	from the task	s that miss
	soft-deadline		Firm-Dea	dline	(C)	Hard-Dead	line	(D)	None of these	\neg
	5- the lazy video s	vstem			()	nato-i)cat	inic	10)	Notic of these	
	soft-deadline	V	B) Firm-De		C	Hard-Dead	lline	(D)	None of these	
-	6- Missing a		can result i						rione er mese	
_	soft-deadline		B) Firm-De		e	Hard-Dead		D)	None of these	
	7-The operator s				e is an e					
A	soft-deadline		B) Firm-De			Hard-Dea	dline	D)	None of these	
5	8- The _scheduler	r can a	rbitrarily su	spend and				e task wi	thout affecting	its behavior.
) Recoverable	B		ASSESSMENT OF THE OWNER, WHEN THE PARTY NAMED IN	Preen			Non-preer		
5	9- After a transac	tion r	eaches its c	ommit pe	oint, any	portion of	the log	that has	not been writ	ten to the disk
	yet must now b	e wri	ten to the d	isk. (tru	e/false)	(
(60-Transaction fa	ilure i	nay also oc	cur beca	use of er	roneous pa	ramete	r values	or because of a	logical
	programming									



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- The Journal keeps track of all transaction operations that affect the values of database items. (true/false)
- The transaction operations are stored in Journal after affecting the value of database file. (true/false
3- In read operation, it's not necessary to bring the disk block into the main memory buffer. (true/false
4- The System Log contain the this line [write_item,T,X] to record that transaction T has changed the
value of database item X. (true/false)
5-Each transaction should be ended explicitly after implicit Begin. (true/false)×
6- A schedule S is serializable if it is equivalent to some serial schedule of the same n transactions.
(true/false)
7- Local errors occur when certain conditions necessitate cancellation of the transaction is violated.
(true/false)
8- Basic operations in the database are read and write. (true/false)
59- The basic unit for database read operation from the computer storage to the main memory is one record. (true/false)
70- Some disk blocks may lose their data because of a read or write malfunction or because of a disk read/write head crash. (true/false)
71- Strict Schedules can be requiring cascaded rollback. (true/false)
72-Temporal-consistency is maintaining consistency between the actual state of the environment and the
state as reflected or perceived by the system. (true/false)
73- A schedule S is serial if, for every transaction T participating in the schedule, all the operations of T are
executed consecutively in the schedule. (true/false) The following schedules are conflict equivalent (true/false) 5 3 3 3 , R (8), W (9)
74- The following schedules are conflict equivalent. (Hue) alse)
$S: RI(x), W2(x), RI(y), W2(y)$ $R_1(X) = \omega_2(X)$
S: $R1(x)$, $R1(y)$, $W2(x)$, $W2(y)$ V $R_1(y) - W_2(y)$
75 is a computer monitoring the state of the environment supplying the environment with the
appropriate driving signals. A) Controlled System By Controlling System C) Both A and B D) None of these
76- A schedule S is said to be if it is conflict equivalent to some serial schedule S'.
A) Result equivalent B) Conflict equivalent Conflict serializable D) Serializable
77 is the underlying application.
A) Controlled System B) Controlling System C) Both A and B D) None of these
8- Two schedules are said to be if the order of any two conflicting operations is the same in both schedules.
A) Result equivalent B) Conflict equivalent C) Conflict serializable D) Serializable
79- Deadline is timely and logically correct execution is considered to be critical.
A) Soft B) Firm Pland D) None of these

does not cause a system failure or compromises the system's integrity

B) Firm-Deadline

C) Hard-Deadline

80-Missing a

soft-deadline

D) None of these



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A schedule S is _					
tem that T reads	have committed.				
Recoverable	B) Cascadeless		Strict	D)	
					same database items
	aved in a way tha	t makes t	the value o	of some da	abase item incorrect.
	B) Dirty Read				D) None of thes
					itabase, may cause a t
as a fund withdra	wal from that acc	ount, to l	be canceled	d. this pro	blem is called
) Local errors	B) exception cor	nditions	er	Both A an	d B D) None of the
9-This schedule can	be classified as _				
R2(X), $R1($	X), $WI(X)$, $RI(Y)$, Y	WI(Y), CI	I, W2(X), C	2	
Recoverable	B) Cascadeless	25	Strict	(D)	Not Recoverable
30-This schedule can					
	(X), $R1(Y)$, $W1(Y)$.				
A) Recoverable	B) Cascadeless		Strict		Not Recoverable
	ne where every tra	ansaction	reads only	the items	that are written by c
transactions.	D Committee	10	Christ	ID	None of these
A) Recoverable	Cascadeless		Strict	(D)	None of these
22 The following se	chedula can be clas	cified ac			
32-The following so		The second secon	C1, W3(v), C	3, RZ(v), WZ	(z), W2(y), C2
R1(x), R2	chedule can be class 2(x), R1(z), R3(x), R3(B) Cascadeless	y), W1(x),	C1, W3(y), C	3, R2(y), W2	(z), W2(y), C2 Not Recoverable
A) Recoverable	2(x), R1(z), R3(x), R3(y), W1(x),	Strict		
A) Recoverable 33- The following s	2(x), R1(z), R3(x), R3(B) Cascadeless	ssified as	Strict	D)	Not Recoverable
A) Recoverable 33- The following s	B) Cascadeless schedule can be class	ssified as	Strict	D)	Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable	B) Cascadeless schedule can be class (z), R3(x), R3(y), W (z), R3(x), R3(y), W (B) Cascadeless	ssified as 1(x), W3(y	Strict /), R2(y), W: Strict	D) 2(z), W2(y),	Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1	B) Cascadeless schedule can be class (z), R3(x), R3(y), W (z), R3(x), R3(y), W (B) Cascadeless	ssified as 1(x), W3(y C)	Strict /), R2(y), W: Strict	D) 2(z), W2(y),	Not Recoverable C1, C2, C3 Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c	B) Cascadeless (z), R3(x), R3(y), W (z), R3(x), R3(y), W (a) B) Cascadeless (a) Cascadeless (a) Cascadeless (a) Cascadeless	ssified as 1(x), W3(y C) R1(Strict (), R2(y), W: Strict (x), R2(x), W2(x)	D) 2(z), W2(y), x), W1(x), C2, I	Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s	B) Cascadeless (z), R3(x), R3(y), W (z), R3(x), R3(y), W (a) B) Cascadeless (a) Cascadeless (a) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless	ssified as 1(x), W3(y C) R1(ssified as	Strict /), R2(y), W: Strict x), R2(x), W2(x)	D) 2(z), W2(y), x), W1(x), C2, I	Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s	B) Cascadeless (z), R3(x), R3(y), W (z), R3(x), R3(y), W (a) B) Cascadeless (a) Cascadeless (a) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless (c) Cascadeless	ssified as (x), W3(y S C) R1(C) ssified as (, R2(y), F	Strict (), R2(y), W2(y) Strict Strict R3(y), W1(y)	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1 Not Recoverable z), W3(y), W2(y), C3 Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable	B) Cascadeless (z), R3(x), R3(y), W (z), R3(x), R3(y), W (z), R3(x), R3(y), W (z) (z), R3(x), R3(y), W (z) (z), R3(x), R3(y), W (z) (z), R3(x), R3(z) (z), R3(x), R1(z) (z) (z), R3(x), R1(z)	ssified as 1(x), W3(y C) R1(c) R1(c) R2(y), R	Strict (), R2(y), W2(x), R2(x), W2(x) Strict Strict Strict P1(x)	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1 Not Recoverable z), W3(y), W2(y), C3
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F	B) Cascadeless (z), R3(x), R3(y), W (z), R3(x), R3(y), W (z), R3(x), R3(y), W (z) (z), R3(x), R3(y), W (z) (z), R3(x), R3(y), W (z) (z), R3(x), R3(z) (z), R3(x), R1(z) (z) (z), R3(x), R1(z)	ssified as (x), W3(y) (x), W3(y) (x) (x) (x) (x) (x) (x) (x) (x) (x) (x	Strict (), R2(y), W2(x), R2(x), W2(x) Strict Strict Strict P1(x)	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1 Not Recoverable z), W3(y), W2(y), C3 Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class	ssified as (x), W3(y (x),	Strict (), R2(y), W2(x), R2(x), W2(x) Strict R3(y), W1(x) Strict R1(x)	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1 Not Recoverable z), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class	ssified as (x), W3(y (x), W3(y (x), W3(y (x), C) (x) (x) (x) (x) (x) (x) (x) (x) (x) (x	Strict (x), R2(y), W2(x) Strict (x), R2(x), W2(x) Strict R3(y), W1(x) Strict R1(x)	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(D) y, W1(x), R2	Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable (x), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c R1(X), F A) Recoverable	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class can be classified as Cascadeless can be classified as	ssified as (x), W3(y (x), W3(y (x), W3(y (x), C) (x) (x) (x) (x) (x) (x) (x) (Strict Strict Strict Strict Strict R1(x Strict R1(x Strict C1, W2(X Strict	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(D) y, W1(x), R2	Not Recoverable C1, C2, C3 Not Recoverable R1(y), W1(y), C1 Not Recoverable z), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c R1(X), F A) Recoverable 38- The following s	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class can be classified as Cascadeless chedule can be class can be classified as Cascadeless can be classified as	ssified as (x), W3(y (x), W3(y (x), W3(y (x), C) (x) (x) (x) (x) (x) (x) (x) (Strict (x), R2(y), W2(x) Strict (x), R2(x), W2(x) Strict R1(x) Strict (x), C1, W2(X) Strict	D) 2(z), W2(y), x), W1(x), C2, (D) (x), C1, W2(D) (x), C2 (D)	Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable (x), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c R1(X), F A) Recoverable 38- The following s R1(x), W1(x),	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class can be classified as (Z(X), W1(X), R1(Y) B) Cascadeless (X2(X), W1(X), R1(Y) Cascadeless	ssified as (x), W3(y (x),	Strict (x), R2(y), W2(x) Strict R3(y), W1(x) Strict R1(x) Strict (x), C1, W2(X) Strict (x), W1(y), A	D) 2(z), W2(y), x), W1(x), C2, I D) x), C1, W2(D) y), W1(x), R2 D) 1, C2 D)	Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable (x), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A Not Recoverable Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c R1(X), F A) Recoverable 38- The following s	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class can be classified as Cascadeless chedule can be class can be classified as Cascadeless can be classified as	ssified as (x), W3(y (x),	Strict Strict Strict Strict Strict R1(x Strict N, C1, W2(X Strict N, W1(y), A Strict	D) (z(z), W2(y), (x), W1(x), C2, (D) (x), C1, W2((D)) (x), C2 (D) (x), C2 (D) (x), C2 (D)	Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable (x), W1(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A Not Recoverable Not Recoverable
A) Recoverable 33- The following s R1(x), R2(x), R1 A) Recoverable 34- This schedule c A) Recoverable 35- The following s R1(x), F A) Recoverable 36- The following s A) Recoverable 37- This schedule c R1(X), F A) Recoverable 38- The following s R1(x), W1(x), A) Recoverable	B) Cascadeless chedule can be class (z), R3(x), R3(y), W B) Cascadeless can be classified as Cascadeless chedule can be class can be classified as (Z(X), W1(X), R1(Y) B) Cascadeless (X2(X), W1(X), R1(Y) Cascadeless	ssified as (x), W3(y (x), W3(y (x), W3(y (x), W3(y (x), C) (x) (x) (x) (x) (x) (x) (x) (x) (x) (x	Strict Strict Strict Strict R1(x Strict R1(x Strict N, C1, W2(X Strict N, W1(y), A Strict	D) (z(z), W2(y), (x), W1(x), C2, (D) (x), C1, W2((D)) (x), C2 (D) (x), C2 (D) (x), C2 (D)	Not Recoverable Not Recoverable Not Recoverable Not Recoverable Not Recoverable (x), W3(y), W2(y), C3 Not Recoverable (x), R1(y), W2(x), C2, A Not Recoverable Not Recoverable

Examination Committee

Dr.Hosam Refaat