

Cracking the Odd Case of Randomness in Java

*A technique for exploiting insecure randomness in Java
with practical applications*

who



Joseph Surin

[@josep68_](https://twitter.com/josep68_) jsur.in

Security Consultant at elttam

- Organises CTFs (with [DownUnderCTF](#))
- Plays CTFs (with skateboarding dog)
- Loves crypto(graphy)

Randomness is everywhere

Almost any web application uses randomness for something

- Object IDs
- MFA backup codes
- Session tokens
- Access tokens
- Cryptographic keys/nonces
- Email verification links
- Password reset links
- and so on...



Randomness is not done right everywhere

According to [OWASP Top 10 WebApp Security Risks](#)

1. A01:2021-Broken Access Control
2. **A02:2021-Cryptographic Failures**
3. A03:2021-Injection
4. A04:2021-Insecure Design
5. A05:2021-Security Misconfiguration
6. A06:2021-Vulnerable and Outdated Components
7. A07:2021-Identification and Authentication Failures
8. A08:2021-Software and Data Integrity Failures
9. A09:2021-Security Logging and Monitoring Failures
10. A10:2021-Server-Side Request Forgery

Cryptographic failures includes use of insecure randomness!



Randomness is not done right everywhere

```
String password = RandomStringUtils.random(10, true, true);
```

```
private String verificationCode = RandomStringUtils.randomAlphanumeric(16);
```

```
private String accessToken = RandomStringUtils.randomAlphanumeric(ACCESS_TOKEN_LEN);
```

```
const randomPassword = () => {
  function getRandomSpecialChar() {
    const code = Math.round(Math.random() * (38 - 37) + 37);
    return String.fromCharCode(code);
  }
  function getRandomDigit() {
    const code = Math.round(Math.random() * (57 - 48) + 48);
    return String.fromCharCode(code);
  }
  function getRandomLetter() {
    const code = Math.round(Math.random() * (90 - 65) + 65);
    return String.fromCharCode(code);
  }
  let password = '';
  for (let i = 0; i < 6; i += 1) {
    password += getRandomLetter() + getRandomDigit() + getRandomSpecialChar();
  }
  return password;
};
```

```
def genToken(length=64):
    """
    Use this utility function to generate a random string of
    a desired length.
    """
    return ''.join(random.choice(string.letters + string.digits)
                  for x in range(length))
```

Is it actually exploitable?

Our Target: RandomStringUtils

- Part of [Apache Commons Lang 3](#), a popular Java library providing utility classes
- Very easy to use:

```
String password = RandomStringUtils.randomAlphanumeric(32);
```

Insecure by default!

- Java's default (weak) RNG
(`java.util.Random`)



```
36 public class RandomStringUtils {  
37  
38     /**  
39      * <p>Random object used by random method. This has to be not local  
40      * to the random method so as to not return the same value in the  
41      * same millisecond.</p>  
42     */  
43     private static final Random RANDOM = new Random();  
44 }
```

org.apache.commons.lang3

Class RandomStringUtils

java.lang.Object

org.apache.commons.lang3.RandomStringUtils

```
public class RandomStringUtils  
extends Object
```

Generates random Strings.

Caveat: Instances of `Random`, upon which the implementation of this class relies, are not cryptographically secure.

RandomStringUtils under the hood

```
140     public static String randomAlphanumeric(final int count) {  
141         return random(count, true, true);  
142     }  
  
263     public static String random(final int count, final boolean letters, final  
264         boolean numbers) {  
265         return random(count, 0, 0, letters, numbers);  
266     }  
  
283     public static String random(final int count, final int start, final int end,  
284         final boolean letters, final boolean numbers) {  
285         return random(count, start, end, letters, numbers, null, RANDOM);  
286     }
```



RandomStringUtils under the hood

```
348     public static String random(int count, int start, int end, final boolean letters, final boolean numbers,
349                               final char[] chars, final Random random) {
350         // [...]
351
352         if (start == 0 && end == 0) {
353             if (chars != null) {
354                 end = chars.length;
355             } else if (!letters && !numbers) {
356                 end = Character.MAX_CODE_POINT;
357             } else { // chars is null, but letters and numbers are both true, so this branch is taken
358                 end = 'z' + 1;
359                 start = ' ';
360             }
361         } else if (end <= start) {
362             throw new IllegalArgumentException("Parameter end (" + end + ") must be greater than start (" + start + ")");
363         }
364
365         // [...]
366     }
```

RandomStringUtils under the hood

```
367     final StringBuilder builder = new StringBuilder(count);
368     final int gap = end - start; // gap equals 91
369
370     while (count-- != 0) {
371         final int codePoint;
372         if (chars == null) { // chars is null, so this branch is taken
373             // since start is 32 and gap is 91, this set codePoint to a random number between 32 and 122 (inclusive)
374             codePoint = random.nextInt(gap) + start;
375
376             // [...]
377
378         } else {
379             // [...]
380         }
381
382         // [...]
383
384         // codepoints in the range between 32 and 122 contains some non-alphanumeric characters, so this branch is not always taken
385         if (letters && Character.isLetter(codePoint)
386             || numbers && Character.isDigit(codePoint)
387             || !letters && !numbers) {
388             builder.appendCodePoint(codePoint);
389
390             if (numberOfChars == 2) {
391                 count--;
392             }
393
394         } else {
395             count++;
396         }
397     }
398
399     return builder.toString();
400 }
```

**So we just need to crack
random.nextInt(91)...**



how can java random nextInt be cracked if the bound is odd?



It is not possible to crack the randomness of `java.util.Random.nextInt()` if the bound is odd, as the algorithm used to generate the random numbers is considered secure. However, if the bound is even, there could be a potential bias in the generated random numbers, which could be exploited to crack the randomness. Therefore, it is recommended to use an odd bound or use a different random number generator such as `java.security.SecureRandom` for more secure random number generation.

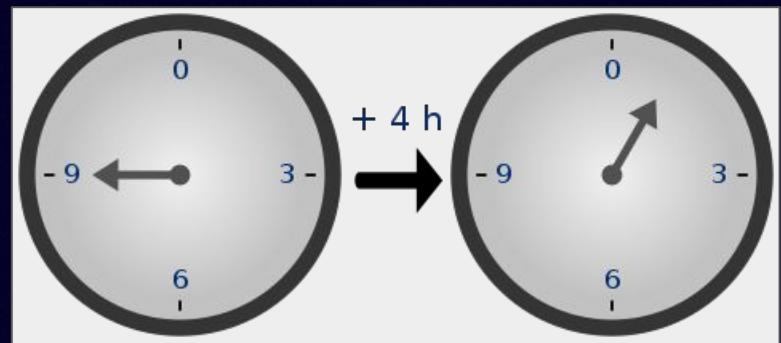
Modular Arithmetic 101

$$a = b \bmod M$$

"a is the remainder after dividing b by M"

$$3 = 11 \bmod 8$$

$$3 = 11 + 8k, \quad k = -1$$



java.util.Random 101

- 48-bit Linear Congruential Generator (LCG)
- x_0 is the seed (initial state)
- $x_{i+1} = (Ax_i + C) \bmod M$
- `nextInt(bound)` returns an integer less than bound
- $y = (x_i \gg 17) \bmod B$



```
protected int next(int bits) {  
    AtomicLong seed = this.seed;  
  
    long oldseed;  
    long nextseed;  
    do {  
        oldseed = seed.get();  
        nextseed = (oldseed * multiplier + addend) & mask;  
        while(!seed.compareAndSet(oldseed, nextseed));  
    }  
    return (int)(nextseed >>> 48 - bits);  
}
```

```
public int nextInt(int bound) {  
    if (bound <= 0) {  
        throw new IllegalArgumentException("bound must be positive");  
    } else {  
        int r = this.next(31);  
        int m = bound - 1;  
        if ((bound & m) == 0) {  
            r = (int)((long)bound * (long)r >> 31);  
        } else {  
            for(int u = r; u - (r = u % bound) + m < 0; ) {  
                System.out.printf("skipped u = %d\n", u);  
                u = this.next(31);  
            }  
        }  
    }  
    return r;  
}
```

Cracking RandomStringUtils - Setup

- Suppose we call `RandomStringUtils.randomAlphanumeric(10)` and get the string "Q6p0ifNqAk"
- Take the ASCII values of each character and subtract 32 to get the outputs of `random.nextInt(91)`
- Call these y_1, y_2, \dots, y_{10} . In this case:

$$y_1 = 49$$

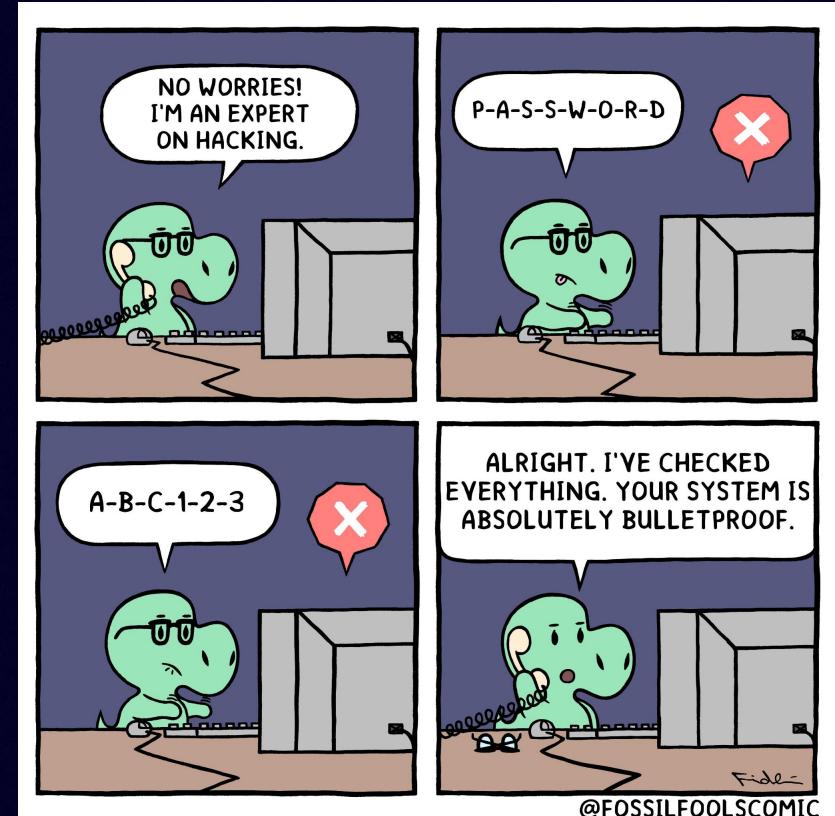
$$y_2 = 22$$

$$y_3 = 80$$

⋮

Cracking RandomStringUtils - Idea #0

- When in doubt, use brute force
- Try all 2^{48} possible Random state values, check which produces a string that agrees with the output
- This will take a while...



Time to stare at equations

Cracking RandomStringUtils - Idea #1

$$y_1 = (x_1 \gg 17) \bmod 91$$

$$\implies y_1 = (x_1 \gg 17) + 91k_1$$

$$\implies y_1 - 91k_1 = (x_1 \gg 17)$$

$$|k_1| < 2^{31}/91$$

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Cracking RandomStringUtils - Idea #1

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Enumerate the $2^{31}/91$ possible values of k_1 to get candidates for $(x_1 \gg 17)$

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Cracking RandomStringUtils - Idea #1

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Enumerate the $2^{31}/91$ possible values of k_1 to get candidates for $(x_1 \gg 17)$

Enumerate the 2^{17} possible values of $x_1 \bmod 2^{17}$ to get candidates for x_1

Complexity: **$2^{41.5}$**

This is the idea used in [alex91ar/randomstringutils](#)

Cracking RandomStringUtils - Idea #2

- Idea #0 uses none of the outputs (to reduce the state search space)
- Idea #1 uses one output
- Idea #2 ...



Time to stare at equations (part 2)

Cracking RandomStringUtils - Idea #2

$$y_2 = (x_2 \gg 17) \bmod 91$$

$$\implies y_2 = (((Ax_1 + C) \bmod M) \gg 17) \bmod 91$$

$$\implies y_2 = (((A(x_{1,U_{31}} + x_{1,L_{17}}) + C) \bmod M) \gg 17) \bmod 91$$

$$x_1 = \boxed{\quad \qquad x_{1,U_{31}} \qquad \quad} \boxed{\quad \qquad x_{1,L_{17}} \qquad \quad}$$

Cracking RandomStringUtils - Idea #2

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Cracking RandomStringUtils - Idea #2

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$$\begin{aligned} \implies y_2 &= (((((Ax_{1,U_{31}} + C) \bmod M) \\ &\quad + (Ax_{1,L_{17}} \bmod M)) \bmod M) \gg 17) \bmod 91 \end{aligned}$$

$$\begin{aligned} \implies y_2 &= (((((Ax_{1,U_{31}} + C) \bmod M) \\ &\quad + (Ax_{1,L_{17}} \bmod M) - b_2M) \gg 17) \bmod 91 \\ &\quad b_2 \in \{0, 1\} \end{aligned}$$

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$$\implies y_2 = (((((Ax_{1,U_{31}} + C) \bmod M) \gg 17) \\ + ((Ax_{1,L_{17}} \bmod M) \gg 17) - b_22^{31} + c_2) \bmod 91 \\ c_2 \in \{0, 1\}$$

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Cracking RandomStringUtils - Idea #2

$$y_2 = (((Ax_{1,U_{31}} + C) \bmod M) \gg 17) \\ + ((Ax_{1,L_{17}} \bmod M) \gg 17) - b_2 2^{31} + c_2 \bmod 91$$

$$\implies (y_2 - (((Ax_{1,U_{31}} + C) \bmod M) \gg 17)) \bmod 91 = \\ (((Ax_{1,L_{17}} \bmod M) \gg 17) - b_2 2^{31} + c_2) \bmod 91$$

From idea #1, we have $2^{31}/91$ candidates for $x_{1,U_{31}} \dots$

Cracking RandomStringUtils - Idea #2

$$\begin{aligned}y_2 = & (((Ax_{1,U_{31}} + C) \bmod M) \gg 17) \\& + ((Ax_{1,L_{17}} \bmod M) \gg 17) - b_2 2^{31} + c_2 \bmod 91\end{aligned}$$

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From idea #1, we have $2^{31}/91$ candidates for $x_{1,U_{31}} \dots$

Cracking RandomStringUtils - Idea #2

Compute a lookup table for left-hand side values

$$(y_2 - (((Ax_{1,U_{31}} + C) \bmod M) \gg 17)) \bmod 91$$

to $x_{1,U_{31}}$ candidates

For each $x_{1,L_{17}}$ candidate, compute the right-hand side value

$$(((Ax_{1,L_{17}} \bmod M) \gg 17) - b_2 2^{31} + c_2) \bmod 91$$

and meet-in-the-middle!



Cracking RandomStringUtils - Idea #2

- Approximately $2^{31}/91^2$ candidates for $x_{1,U_{31}}$ for each left-hand side value

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- Complexity: **2^{37}**

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Cracking RandomStringUtils - Idea #2

- Approximately $2^{31}/91^2$ candidates for $x_{1,U_{31}}$ for each left-hand side value
- 2^{19} right-hand side values to go through
- Complexity: **2^{37}**
- Idea can be generalised, using another output gives a time complexity improvement of 4.5 bits
- Full details: <https://elttam.com/blog/cracking-randomness-in-java/>

That's cool but who cares?



RandomStringUtil language:Java

Filter by

<> Code

45.8k

4



```
nbok.extern.slf4j.Slf4j;  
g.apache.commons.lang3.RandomStringUtils;  
g.springframework.boot.context.event.ApplicationReadyEvent;  
  
// generate password  
password = RandomStringUtils.randomAlphanumeric(16);  
log.info("== Generated random password: {} for super administrator: {} ==",  
import java.util.concurrentresponse,  
import org.apache.commons.lang3.RandomStringUtils;  
import org.apache.commons.lang3.StringUtils;  
  
if (StringUtils.isBlank(cookieSecret)) {  
    cookieSecret = RandomStringUtils.randomAlphanumeric(16);  
}  
  
String rtspConnectionId = kMed...  
+ RandomStringUtils.randomAlphanumeric(16);  
+ uri.getPort() != -1 ? ":" + uri.getPort() : "") + uri.getPath();  
ona > app/models/Email.java  
  
import controllers.routes;  
import org.apache.commons.lang3.RandomStringUtils;  
import play.data.validation.Constraints;  
  
public void sendValidationEmail() {  
    this.token = RandomStringUtils.randomNumeric(50);  
    this.confirmUrl = Url.create(routes.UserApp.confirmEmail(this.id, this.token).url());  
  
566  
567  
if(listOfDisabledFeatures.contains("learningDashboard") == false) {  
    learningDashboardAccessToken = RandomStringUtils.randomAlphanumeric(16);  
}  
import org.apache.commons.lang3.RandomStringUtils;  
commons.lang3.StringUtils;  
} learningDashboardAccessToken = RandomStringUtils.randomAlphanumeric(16);  
-labs/core > ...  
26 import org.apache.commons.lang3.RandomStringUtils;  
...  
56 private static String _getPassword ( String key, int length ) {  
57     return RandomStringUtils.random( length, key );  
58 }  
am;  
193.RandomStringUtils;  
...  
41  
42     private static final String KEY_STORE_PASS = RandomStringUtils.randomAlphanumeric(8);  
43     private static final String KEY_STORE_FILE_PATH = "clientkeystore.jks";  
l-Media-Server > src/main/java/io/antmedia/filter	TokenNameGenerator.java  
import org.apache.commons.lang3.RandomStringUtils;  
import org.apache.commons.lang3.StringUtils;  
if(generatedToken == null) {  
    generatedToken = RandomStringUtils.random(16);  
}
```

Demo

(CVE-2023-24828)

Found with Emily Trau (emily.id.au)

Practical Applications

- Developing exploits quickly
- Detecting insecure randomness in black-box settings

Conclusion

- Insecure randomness ⇒ easily exploitable, high risk issues
- Crypto is fun
- Tool available at <https://github.com/elttam/rsu-cracker>
- Full details + other awesome posts at <https://elttam.com/blog>

Thank You!

Questions?
