## Oblivious Fransfer

- Sender has m., ..., mn & H.
- · Received has i a It. n ]
- · Goal. (1) Receiver learns mi, and no other mj.

## Or from El Bamal

· Group & ; < G> = q ; hash fune 1. 62 > Mxk

6

6

· CPA secure (Es, Ds) channel.

Receiver (Bob)	R	Sender (News)	
Q - Z - Z	L V + 9"	BZ Zg	
	- uegave >		El Barrel ph
Take c.		forjin 1. m.	4; = 8 2 v 2 2
· we vo		uje uv	7
ka- M(U, W)		with the	
m = Ds (k, ci)	(G,,:,, Cn)	kj = le(v, wi)	El famales
		C; L & (tj. m;)	J (4, 6) = (8 , 9)
	-		

- . The article is which Bob want to read in energyted with Bob's public key.
- · Other articles everypted with some other public lays (waknown).
- · Bob can decrypt and relact ui, the article

Lealwel Digital Signatures Clients (pk) Vendor (8K) Software Client Software update Vendor secure boy sh) - 12 Signing algorithm -> Client . Whiley s - Huslid, Security Game. contall upa Adversary A Challenger m2, 11 mg m, el -> (m, o) (pk, ex) - Gen 0,4 8(st, m.) 0, ,, 0g Adu wins of V(ph, m, s) = 'secept' and to e & m, ..., mg 5 Def. 88 = (Gen, 8, 4) is secure if for all eff. A: dignations Chame Adv [A, SS] = Pr[Awins] < negl. 85 = ( Sen, S. 4). Attaches can find mo 7 m, 3. %. V(pt, mo, o) = V(pt, m., o) 40, (pt st) + Ben D. Can this 85 be seeme? Lo No, signatures con be forged: (1) Ast to sign ma Bruss 5. (a) Forge (m, to).

Hilloy

Euler Theorem.

· (Zp) is called a gydic group, that is

I g = (Zp) \* such that 2', g, g, g, g, g, g, g' f = (Zp) \*

g is called a generator of (Zp)+.

Exemple. p=4

· Not every dement in a generator:

f, d, d, 2, 2, 255 = f1, d, 45.

Solving Quadratic Equations (mod p)

· Solve: and + Box + C in Zp

· Solution: 2= (-B + 582-4ac) /2a in Zp

1) Find (2a) in Zp using Enclid

a) Find square root of 62-4ac in Zp (if exists) wing a square root algorithm.