Handling of active – inactive CPUs

```
include/linux/cpu.h:
                        * cpu_active mask right after SCHED_ACTIVE. During
include/linux/cpu.h:
                        * This ordering guarantees consistent cpu_active mask and
include/linux/cpumask.h: *
                              cpu active mask - has bit 'cpu' set iff cpu available to migration
include/linux/cpumask.h:extern const struct cpumask *const cpu active mask;
include/linux/cpumask.h:#define num active cpus()
                                                    cpumask weight(cpu active mask)
                                                     cpumask_test_cpu((cpu), cpu_active_mask)
include/linux/cpumask.h:#define cpu_active(cpu)
                                                     ((cpu) == 0)
include/linux/cpumask.h:#define cpu_active(cpu)
include/linux/cpumask.h:void set_cpu_active(unsigned int cpu, bool active);
init/main.c:
                 set_cpu_active(cpu, true);
```

```
kernel/cpu.c:static DECLARE BITMAP(cpu active bits, CONFIG NR CPUS) read mostly;
kernel/cpu.c:const struct cpumask *const cpu active mask = to cpumask(cpu active bits);
kernel/cpu.c:EXPORT SYMBOL(cpu active mask);
kernel/cpu.c:
                       cpumask_set_cpu(cpu, to_cpumask(cpu_active_bits));
kernel/cpu.c:void set cpu active(unsigned int cpu, bool active)
kernel/cpu.c:
                       cpumask set cpu(cpu, to cpumask(cpu active bits));
kernel/cpu.c:
                       cpumask clear cpu(cpu, to cpumask(cpu active bits));
kernel/cpuset.c: if (!cpumask equal(top cpuset.cpus allowed, cpu active mask))
                       if (!cpumask subset(trialcs->cpus allowed, cpu active mask))
kernel/cpuset.c:
kernel/cpuset.c: * synchronized to cpu_active_mask and N_MEMORY, which is necessary in
kernel/cpuset.c: cpumask_copy(&new_cpus, cpu_active mask);
kernel/cpuset.c: /* synchronize cpus allowed to cpu active mask */
kernel/cpuset.c: cpumask copy(top cpuset.cpus allowed, cpu active mask);
kernel/sched/core.c:
                                   if (!cpu active(dest cpu))
kernel/sched/core.c:
                                   if (!cpu_active(dest_cpu))
                       if (likely(cpu_active(dest_cpu))) {
kernel/sched/core.c:
kernel/sched/core.c:
                       dest cpu = cpumask any and(cpu active mask, cpumask);
kernel/sched/core.c:
                       if (!cpumask intersects(new mask, cpu active mask)) {
                       dest_cpu = cpumask_any_and(cpu_active_mask, new_mask);
kernel/sched/core.c:
                       if (unlikely(!cpu active(dest cpu)))
kernel/sched/core.c:
kernel/sched/core.c:static int __cpuinit sched_cpu_active(struct notifier_block *nfb,
kernel/sched/core.c:
                             set_cpu_active((long)hcpu, true);
                             set cpu activecpu active((long)hcpu, false);
kernel/sched/core.c:
                       cpu notifier(sched cpu active, CPU PRI SCHED ACTIVE);
kernel/sched/core.c:
                       if (cpumask_test_cpu(rq->cpu, cpu_active_mask))
kernel/sched/core.c:
kernel/sched/core.c:
                             cpumask_andnot(doms_new[0], cpu_active_mask, cpu_isolated_map);
                       Update cpusets according to cpu active mask. If cpusets are
kernel/sched/core.c: *
kernel/sched/core.c:static int cpuset cpu active(struct notifier block *nfb, unsigned long action,
kernel/sched/core.c:
                        init_sched_domains(cpu_active_mask);
kernel/sched/core.c:
                       hotcpu notifier(cpuset cpu active, CPU PRI CPUSET ACTIVE);
```

```
kernel/sched/fair.c: cpumask_copy(cpus, cpu_active_mask);
```

kernel/sched/fair.c: if (!cpu_active(cpu))

kernel/sched/sched.h: if (!cpu_active(cpu_of(rq)))

kernel/stop_machine.c: BUG_ON(cpu_active(raw_smp_processor_id()));

kernel/stop_machine.c: queue_stop_cpus_work(cpu_active_mask, stop_machine_cpu_stop, &smdata,

```
init/main.c:
     Activate the first processor.
static void __init boot_cpu_init(void)
     int cpu = smp_processor_id();
     /* Mark the boot cpu "present", "online" etc for SMP and UP case */
     set_cpu_online(cpu, true);
     set_cpu_active(cpu, true);
     set_cpu_present(cpu, true);
     set_cpu_possible(cpu, true);
}
kernel/cpu.c:
void set_cpu_active(unsigned int cpu, bool active)
     if (active)
           cpumask_set_cpu(cpu, to_cpumask(cpu_active_bits));
     else
           cpumask_clear_cpu(cpu, to_cpumask(cpu_active_bits));
}
```

```
kernel/cpuset.c:
 * Rebuild scheduler domains.
 * If the flag 'sched load balance' of any cpuset with non-empty
 * 'cpus' changes, or if the 'cpus' allowed changes in any cpuset
 * which has that flag enabled, or if any cpuset with a non-empty
 * 'cpus' is removed, then call this routine to rebuild the
 * scheduler's dynamic sched domains.
 * Call with cpuset mutex held. Takes get online cpus().
static void rebuild sched domains locked(void)
      struct sched_domain_attr *attr;
      cpumask_var_t *doms;
      int ndoms;
      lockdep_assert_held(&cpuset_mutex);
      get_online_cpus();
       * We have raced with CPU hotplug. Don't do anything to avoid
       * passing doms with offlined cpu to partition sched domains().
       * Anyways, hotplug work item will rebuild sched domains.
      if (!cpumask_equal(top_cpuset.cpus_allowed, cpu_active_mask))
            goto out:kernel/cpuset.c:
     /* Generate domain masks and attrs */
      ndoms = generate sched domains(&doms, &attr);
      /* Have scheduler rebuild the domains */
      partition_sched_domains(ndoms, doms, attr);
out:
     put_online_cpus();
}
```

```
/**
 * update cpumask - update the cpus allowed mask of a cpuset and all tasks in it
 * @cs: the cpuset to consider
 * @buf: buffer of cpu numbers written to this cpuset
static int update_cpumask(struct cpuset *cs, struct cpuset *trialcs,
                    const char *buf)
{
      struct ptr heap heap;
      int retval;
      int is load balanced;
      /* top cpuset.cpus allowed tracks cpu online mask: it's read-only */
      if (cs == &top_cpuset)
           return -EACCES;
       * An empty cpus_allowed is ok only if the cpuset has no tasks.
       * Since cpulist_parse() fails on an empty mask, we special case
       * that parsing. The validate_change() call ensures that cpusets
       * with tasks have cpus.
       */
      if (!*buf) {
           cpumask clear(trialcs->cpus allowed);
      } else {
           retval = cpulist_parse(buf, trialcs->cpus_allowed);
           if (retval < 0)
                  return retval:
           if (!cpumask_subset(trialcs->cpus_allowed, cpu_active_mask))
                  return -EINVAL;
      retval = validate_change(cs, trialcs);
      if (retval < 0)
            return retval:
      /* Nothing to do if the cpus didn't change */
      if (cpumask equal(cs->cpus allowed, trialcs->cpus allowed))
           return 0;
      retval = heap_init(&heap, PAGE_SIZE, GFP_KERNEL, NULL);
      if (retval)
```

```
return retval;
is_load_balanced = is_sched_load_balance(trialcs);
mutex_lock(&callback_mutex);
cpumask_copy(cs->cpus_allowed, trialcs->cpus_allowed);
mutex_unlock(&callback_mutex);

/*
    * Scan tasks in the cpuset, and update the cpumasks of any
    * that need an update.
    */
    update_tasks_cpumask(cs, &heap);
heap_free(&heap);
if (is_load_balanced)
        rebuild_sched_domains_locked();
return 0;
}
```

```
kernel/sched/core.c:
static int select fallback rg(int cpu, struct task struct *p)
{
     int nid = cpu to node(cpu);
     const struct cpumask *nodemask = NULL;
     enum { cpuset, possible, fail } state = cpuset;
     int dest_cpu;
      * If the node that the cpu is on has been offlined, cpu to node()
       * will return -1. There is no cpu on the node, and we should
      * select the cpu on the other node.
       */
     if (nid != -1) {
           nodemask = cpumask_of_node(nid);
           /* Look for allowed, online CPU in same node. */
           for_each_cpu(dest_cpu, nodemask) {
                 if (!cpu_online(dest_cpu))
                       continue;
                 if (!cpu_active(dest_cpu))
                       continue;
                 if (cpumask_test_cpu(dest_cpu, tsk_cpus_allowed(p)))
                       return dest cpu;
     }
     for (;;) {
           /* Any allowed, online CPU? */
           for_each_cpu(dest_cpu, tsk_cpus_allowed(p)) {
                 if (!cpu_online(dest_cpu))
                       continue;
                 if (!cpu_active(dest_cpu))
                       continue;
                 goto out;
           }
           switch (state) {
           case cpuset:
                 /* No more Mr. Nice Guy. */
                 cpuset_cpus_allowed_fallback(p);
```

```
state = possible;
                 break;
           case possible:
                 do set cpus allowed(p, cpu possible mask);
                  state = fail:
                 break;
           case fail:
                 printk(KERN ERR "select fallback rg failed\n");
                 printk(KERN_ERR " system_state = %d\n", system_state);
                 printk(KERN_ERR " cpu = %d\n", cpu);
                 printk(KERN_ERR " task = %s\n", p->comm);
                 printk(KERN_ERR " allowed_cpus = %#0101x\n", p->cpus_allowed.bits[0]);
                 printk(KERN_ERR " nr_cpus_allowed = %d\n", p->nr_cpus_allowed);
                 printk(KERN ERR " tsk cpus allowed = %#010lx\n", tsk cpus allowed(p)->bits[0]);
#ifdef CONFIG PREEMPT RT FULL
                 printk(KERN ERR " migrate disable = %x\n" , p->migrate disable);
#endif
#ifdef CONFIG SCHED DEBUG
                 printk(KERN_ERR " migrate_disable_atomic = %x\n" , p->migrate_disable_atomic);
#endif
                  BUG();
                 break;
      }
out:
     if (state != cpuset) {
            * Don't tell them about moving exiting tasks or
             * kernel threads (both mm NULL), since they never
             * leave kernel.
            * /
           if (p->mm && printk_ratelimit()) {
                 printk_deferred("process %d (%s) no longer affine to cpu%d\n",
                             task_pid_nr(p), p->comm, cpu);
           }
      return dest_cpu;
}
```

```
* sched exec - execve() is a valuable balancing opportunity, because at
 * this point the task has the smallest effective memory and cache footprint.
void sched exec(void)
     struct task_struct *p = current;
     unsigned long flags;
     int dest_cpu;
     raw_spin_lock_irqsave(&p->pi_lock, flags);
     dest_cpu = p->sched_class->select_task_rg(p, SD_BALANCE_EXEC, 0);
     if (dest_cpu == smp_processor_id())
           goto unlock;
     if (likely(cpu_active(dest_cpu))) {
           struct migration_arg arg = { p, dest_cpu };
           raw_spin_unlock_irgrestore(&p->pi_lock, flags);
           stop_one_cpu(task_cpu(p), migration_cpu_stop, &arg);
           return;
     }
unlock:
     raw_spin_unlock_irgrestore(&p->pi_lock, flags);
}
```

```
/**
 * migrate me - try to move the current task off this cpu
 * Used by the pin current cpu() code to try to get tasks
 * to move off the current CPU as it is going down.
 * It will only move the task if the task isn't pinned to
 * the CPU (with migrate_disable, affinity or NO_SETAFFINITY)
 * and the task has to be in a RUNNING state. Otherwise the
 * movement of the task will wake it up (change its state
 * to running) when the task did not expect it.
 * Returns 1 if it succeeded in moving the current task
           0 otherwise.
int migrate_me(void)
      struct task struct *p = current;
      struct migration arg arg;
      struct cpumask *cpumask;
      struct cpumask *mask;
      unsigned long flags;
      unsigned int dest_cpu;
      struct rq *rq;
       * We can not migrate tasks bounded to a CPU or tasks not
       * running. The movement of the task will wake it up.
     if (p->flags & PF_NO_SETAFFINITY || p->state)
           return 0;
     mutex lock(&sched down mutex);
      rq = task_rq_lock(p, &flags);
      cpumask = &__get_cpu_var(sched_cpumasks);
     mask = &p->cpus allowed:
     cpumask_andnot(cpumask, mask, &sched_down_cpumask);
      if (!cpumask weight(cpumask)) {
           /* It's only on this CPU? */
           task_rg_unlock(rg, p, &flags);
```

```
mutex_unlock(&sched_down_mutex);
    return 0;
}

dest_cpu = cpumask_any_and(cpu_active_mask, cpumask);

arg.task = p;
    arg.dest_cpu = dest_cpu;

task_rq_unlock(rq, p, &flags);

stop_one_cpu(cpu_of(rq), migration_cpu_stop, &arg);
    tlb_migrate_finish(p->mm);
    mutex_unlock(&sched_down_mutex);

return 1;
}
```

```
* This is how migration works:
 * 1) we invoke migration cpu stop() on the target CPU using
      stop_one_cpu().
 * 2) stopper starts to run (implicitly forcing the migrated thread
      off the CPU)
 * 3) it checks whether the migrated task is still in the wrong rungueue.
 * 4) if it's in the wrong rungueue then the migration thread removes
      it and puts it into the right queue.
 * 5) stopper completes and stop one cpu() returns and the migration
      is done.
 * Change a given task's CPU affinity. Migrate the thread to a
  proper CPU and schedule it away if the CPU it's executing on
 * is removed from the allowed bitmask.
 * NOTE: the caller must have a valid reference to the task, the
 * task must not exit() & deallocate itself prematurely. The
 * call is not atomic; no spinlocks may be held.
int set_cpus_allowed_ptr(struct task_struct *p, const struct cpumask *new_mask)
{
      unsigned long flags;
     struct rq *rq;
     unsigned int dest cpu;
     int ret = 0;
     rq = task_rq_lock(p, &flags);
     if (cpumask_equal(&p->cpus_allowed, new_mask))
            goto out;
     if (!cpumask_intersects(new_mask, cpu_active_mask)) {
           ret = -EINVAL;
           goto out;
      }
      do_set_cpus_allowed(p, new_mask);
```

```
/* Can the task run on the task's current CPU? If so, we're done */
     if (cpumask_test_cpu(task_cpu(p), new_mask) || __migrate_disabled(p))
           goto out;
     dest_cpu = cpumask_any_and(cpu_active_mask, new_mask);
     if (p->on_rq) {
           struct migration_arg arg = { p, dest_cpu };
           /* Need help from migration thread: drop lock and wait. */
           task_rq_unlock(rq, p, &flags);
           stop_one_cpu(cpu_of(rq), migration_cpu_stop, &arq);
           tlb_migrate_finish(p->mm);
           return 0;
     }
out:
     task_rq_unlock(rq, p, &flags);
     return ret;
EXPORT_SYMBOL_GPL(set_cpus_allowed_ptr);
```

```
* Move (not current) task off this cpu, onto dest cpu. We're doing
 * this because either it can't run here any more (set cpus allowed()
 * away from this CPU, or CPU going down), or because we're
 * attempting to rebalance this task on exec (sched exec).
 * So we race with normal scheduler movements, but that's OK, as long
  as the task is no longer on this CPU.
 * Returns non-zero if task was successfully migrated.
static int __migrate_task(struct task_struct *p, int src_cpu, int dest_cpu)
     struct rg *rg_dest, *rg_src;
      int ret = 0:
     if (unlikely(!cpu_active(dest_cpu)))
           return ret;
      rq_src = cpu_rq(src_cpu);
      rq_dest = cpu_rq(dest_cpu);
      raw spin lock(&p->pi lock);
      double_rq_lock(rq_src, rq_dest);
     /* Already moved. */
     if (task_cpu(p) != src_cpu)
           goto done;
      /* Affinity changed (again). */
     if (!cpumask_test_cpu(dest_cpu, tsk_cpus_allowed(p)))
           goto fail;
       * If we're not on a rg, the next wake-up will ensure we're
       * placed properly.
     if (p->on_rq) {
           dequeue_task(rq_src, p, 0);
           set_task_cpu(p, dest_cpu);
           enqueue_task(rq_dest, p, 0);
           check_preempt_curr(rq_dest, p, 0);
done:
```

```
ret = 1;
fail:
    double_rq_unlock(rq_src, rq_dest);
    raw_spin_unlock(&p->pi_lock);
    return ret;
}
```

```
static int cpuinit sched cpu active(struct notifier block *nfb,
                             unsigned long action, void *hcpu)
{
      switch (action & ~CPU TASKS FROZEN) {
     case CPU DOWN FAILED:
           set_cpu_active((long)hcpu, true);
           return NOTIFY_OK;
      default:
           return NOTIFY DONE;
      }
}
static int __cpuinit sched_cpu_inactive(struct notifier_block *nfb,
                             unsigned long action, void *hcpu)
{
      switch (action & ~CPU TASKS FROZEN) {
      case CPU DOWN PREPARE:
           set_cpu_active((long)hcpu, false);
           return NOTIFY_OK;
      default:
           return NOTIFY_DONE;
}
static int init migration init(void)
     void *cpu = (void *)(long)smp_processor_id();
     int err;
     /* Initialize migration for the boot CPU */
      err = migration call(&migration notifier, CPU UP PREPARE, cpu);
      BUG ON(err == NOTIFY BAD);
     migration_call(&migration_notifier, CPU_ONLINE, cpu);
      register_cpu_notifier(&migration_notifier);
     /* Register cpu active notifiers */
      cpu_notifier(sched_cpu_active, CPU_PRI_SCHED_ACTIVE);
     cpu_notifier(sched_cpu_inactive, CPU_PRI_SCHED_INACTIVE);
      return 0;
early_initcall(migration_init);
```

```
static void rq_attach_root(struct rq *rq, struct root_domain *rd)
     struct root_domain *old_rd = NULL;
     unsigned long flags;
     raw_spin_lock_irqsave(&rq->lock, flags);
     if (rq->rd) {
           old rd = rq - > rd;
           if (cpumask_test_cpu(rq->cpu, old_rd->online))
                 set_rq_offline(rq);
           cpumask_clear_cpu(rq->cpu, old_rd->span);
             * If we dont want to free the old rt yet then
             * set old_rd to NULL to skip the freeing later
             * in this function:
             */
           if (!atomic_dec_and_test(&old_rd->refcount))
                 old_rd = NULL;
      }
     atomic inc(&rd->refcount);
      rq->rd = rd;
      cpumask_set_cpu(rq->cpu, rd->span);
     if (cpumask_test_cpu(rq->cpu, cpu_active_mask))
           set_rq_online(rq);
      raw_spin_unlock_irgrestore(&rq->lock, flags);
     if (old_rd)
           call_rcu_sched(&old_rd->rcu, free_rootdomain);
}
```

```
* Partition sched domains as specified by the 'ndoms new'
 * cpumasks in the array doms new[] of cpumasks. This compares
  doms new[] to the current sched domain partitioning, doms cur[].
 * It destroys each deleted domain and builds each new domain.
 * 'doms new' is an array of cpumask var t's of length 'ndoms new'.
 * The masks don't intersect (don't overlap.) We should setup one
 * sched domain for each mask. CPUs not in any of the cpumasks will
 * not be load balanced. If the same cpumask appears both in the
 * current 'doms cur' domains and in the new 'doms new', we can leave
 * it as it is.
 * The passed in 'doms_new' should be allocated using
 * alloc sched domains. This routine takes ownership of it and will
 * free sched domains it when done with it. If the caller failed the
 * alloc call, then it can pass in doms new == NULL && ndoms new == 1,
 * and partition_sched_domains() will fallback to the single partition
 * 'fallback_doms', it also forces the domains to be rebuilt.
 * If doms_new == NULL it will be replaced with cpu_online_mask.
 * ndoms_new == 0 is a special case for destroying existing domains,
 * and it will not create the default domain.
 * Call with hotplug lock held
void partition sched domains(int ndoms new, cpumask var t doms new[],
                       struct sched domain attr *dattr new)
{
     int i, j, n;
      int new topology;
     mutex_lock(&sched_domains_mutex);
      /* always unregister in case we don't destroy any domains */
      unregister sched domain sysctl();
     /* Let architecture update cpu core mappings. */
     new topology = arch update cpu topology();
      n = doms new ? ndoms new : 0;
```

```
/* Destroy deleted domains */
     for (i = 0; i < ndoms cur; i++) {
           for (j = 0; j < n && !new_topology; j++) {
                 if (cpumask equal(doms cur[i], doms new[i])
                      && dattrs equal(dattr cur, i, dattr new, j))
                        goto match1;
           /* no match - a current sched domain not in new doms_new[] */
           detach_destroy_domains(doms_cur[i]);
match1:
      }
     if (doms_new == NULL) {
           ndoms_cur = 0;
           doms new = &fallback doms;
           cpumask andnot(doms new[0], cpu active mask, cpu isolated map);
           WARN ON ONCE(dattr new);
     }
     /* Build new domains */
     for (i = 0; i < ndoms_new; i++) {
           for (j = 0; j < ndoms_cur && !new_topology; j++) {</pre>
                 if (cpumask_equal(doms_new[i], doms_cur[j])
                     && dattrs equal(dattr new, i, dattr cur, j))
                        goto match2;
           /* no match - add a new doms new */
           build_sched_domains(doms_new[i], dattr_new ? dattr_new + i : NULL);
match2:
      }
     /* Remember the new sched domains */
     if (doms_cur != &fallback_doms)
           free_sched_domains(doms_cur, ndoms_cur);
     kfree(dattr_cur); /* kfree(NULL) is safe */
      doms cur = doms new;
      dattr cur = dattr new;
     ndoms cur = ndoms new;
      register_sched_domain_sysctl();
```

```
mutex_unlock(&sched_domains_mutex);
}
```

```
* Update cpusets according to cpu active mask. If cpusets are
 * disabled, cpuset update active cpus() becomes a simple wrapper
 * around partition sched domains().
 * If we come here as part of a suspend/resume, don't touch cpusets because we
 * want to restore it back to its original state upon resume anyway.
static int cpuset cpu active(struct notifier block *nfb, unsigned long action,
                      void *hcpu)
{
      switch (action) {
     case CPU ONLINE FROZEN:
      case CPU_DOWN_FAILED_FROZEN:
             * num cpus frozen tracks how many CPUs are involved in suspend
             * resume sequence. As long as this is not the last online
             * operation in the resume sequence, just build a single sched
             * domain, ignoring cpusets.
             */
           num_cpus_frozen--;
           if (likely(num_cpus_frozen)) {
                 partition sched domains(1, NULL, NULL);
                  break;
            }
             * This is the last CPU online operation. So fall through and
             * restore the original sched domains by considering the
             * cpuset configurations.
      case CPU ONLINE:
      case CPU_DOWN_FAILED:
           cpuset_update_active_cpus(true);
            break;
      default:
           return NOTIFY_DONE;
      return NOTIFY_OK;
}
```

```
void init sched init smp(void)
      cpumask var t non isolated cpus;
      alloc cpumask var(&non isolated cpus, GFP KERNEL);
      alloc_cpumask_var(&fallback_doms, GFP_KERNEL);
      sched init numa();
      get online cpus();
     mutex lock(&sched domains mutex);
      init_sched_domains(cpu_active_mask);
     cpumask_andnot(non_isolated_cpus, cpu_possible_mask, cpu_isolated_map);
      if (cpumask_empty(non_isolated_cpus))
           cpumask_set_cpu(smp_processor_id(), non_isolated_cpus);
     mutex unlock(&sched domains mutex);
     put online cpus();
     hotcpu_notifier(sched_domains_numa_masks_update, CPU_PRI_SCHED_ACTIVE);
     hotcpu_notifier(cpuset_cpu_active, CPU_PRI_CPUSET_ACTIVE);
      hotcpu_notifier(cpuset_cpu_inactive, CPU_PRI_CPUSET_INACTIVE);
     /* RT runtime code needs to handle some hotplug events */
     hotcpu notifier(update runtime, 0);
     init_hrtick();
     /* Move init over to a non-isolated CPU */
     if (set_cpus_allowed_ptr(current, non_isolated_cpus) < 0)</pre>
            BUG();
      sched_init_granularity();
     free cpumask var(non isolated cpus);
      init_sched_rt_class();
#else
void __init sched_init_smp(void)
{
      sched_init_granularity();
}
```

```
kernel/sched/fair.c:
* Check this cpu to ensure it is balanced within domain. Attempt to move
 * tasks if there is an imbalance.
static int load_balance(int this_cpu, struct rg *this_rg,
                  struct sched_domain *sd, enum cpu_idle_type idle,
                 int *balance)
{
     int 1d moved, cur 1d moved, active balance = 0;
     struct sched_group *group;
     struct rg *busiest;
     unsigned long flags;
     struct cpumask *cpus = __get_cpu_var(load_balance_mask);
      struct lb_env env = {
            .sd
                       = sd,
            .dst_cpu
                       = this_cpu,
                             = this_rq,
            .dst_rq
                           = sched_group_cpus(sd->groups),
            .dst_grpmask
            .idle
                       = idle,
            .loop_break = sched_nr_migrate_break,
            .cpus
                       = cpus,
     };
       * For NEWLY IDLE load balancing, we don't need to consider
       * other cpus in our group
      if (idle == CPU_NEWLY_IDLE)
           env.dst_grpmask = NULL;
      cpumask_copy(cpus, cpu_active_mask);
      schedstat_inc(sd, lb_count[idle]);
redo:
     group = find_busiest_group(&env, balance);
     if (*balance == 0)
           goto out_balanced;
```

```
if (!group) {
            schedstat inc(sd, lb nobusyg[idle]);
            goto out balanced;
      }
     busiest = find_busiest_queue(&env, group);
     if (!busiest) {
           schedstat_inc(sd, lb_nobusyg[idle]);
           goto out balanced;
      }
      BUG_ON(busiest == env.dst_rq);
      schedstat_add(sd, lb_imbalance[idle], env.imbalance);
      1d \text{ moved} = 0:
     if (busiest->nr_running > 1) {
             * Attempt to move tasks. If find_busiest_group has found
             * an imbalance but busiest->nr_running <= 1, the group is
             * still unbalanced. ld_moved simply stays zero, so it is
             * correctly treated as an imbalance.
             */
            env.flags |= LBF ALL PINNED;
            env.src cpu = busiest->cpu;
            env.src_rq
                          = busiest:
           env.loop_max = min(sysctl_sched_nr_migrate, busiest->nr_running);
            update_h_load(env.src_cpu);
more_balance:
           local_irq_save(flags);
           double rg lock(env.dst rg, busiest);
             * cur ld moved - load moved in current iteration
                         - cumulative load moved across iterations
             * ld moved
             */
           cur_ld_moved = move_tasks(&env);
            ld moved += cur ld moved;
           double_rq_unlock(env.dst_rq, busiest);
           local_irq_restore(flags);
```

```
* some other cpu did the load balance for us.
if (cur ld moved && env.dst cpu != smp processor id())
      resched cpu(env.dst cpu);
if (env.flags & LBF NEED BREAK) {
      env.flags &= ~LBF NEED BREAK;
      goto more balance;
}
 * Revisit (affine) tasks on src cpu that couldn't be moved to
 * us and move them to an alternate dst cpu in our sched group
 * where they can run. The upper limit on how many times we
 * iterate on same src cpu is dependent on number of cpus in our
 * sched group.
 * This changes load balance semantics a bit on who can move
 * load to a given_cpu. In addition to the given_cpu itself
 * (or a ilb_cpu acting on its behalf where given_cpu is
 * nohz-idle), we now have balance_cpu in a position to move
 * load to given cpu. In rare situations, this may cause
 * conflicts (balance cpu and given cpu/ilb cpu deciding
 * independently and at same time to move some load to
 * given_cpu) causing exceess load to be moved to given_cpu.
 * This however should not happen so much in practice and
 * moreover subsequent load balance cycles should correct the
 * excess load moved.
if ((env.flags & LBF SOME PINNED) && env.imbalance > 0) {
      env.dst_rq = cpu_rq(env.new_dst_cpu);
      env.dst_cpu = env.new_dst_cpu;
      env.flags &= ~LBF_SOME_PINNED;
      env.loop
                  = 0;
                        = sched_nr_migrate_break;
      env.loop_break
      /* Prevent to re-select dst cpu via env's cpus */
      cpumask clear cpu(env.dst cpu, env.cpus);
      /*
```

```
* Go back to "more balance" rather than "redo" since we
             * need to continue with same src cpu.
             */
            goto more balance;
     /* All tasks on this runqueue were pinned by CPU affinity */
     if (unlikely(env.flags & LBF_ALL_PINNED)) {
            cpumask clear cpu(cpu of(busiest), cpus);
            if (!cpumask empty(cpus)) {
                  env.loop = 0;
                  env.loop_break = sched_nr_migrate_break;
                  aoto redo:
            goto out_balanced;
      }
}
if (!ld_moved) {
      schedstat_inc(sd, lb_failed[idle]);
      /*
       * Increment the failure counter only on periodic balance.
       * We do not want newidle balance, which can be very
       * frequent, pollute the failure counter causing
       * excessive cache_hot migrations and active balances.
       */
     if (idle != CPU NEWLY IDLE)
            sd->nr balance failed++:
     if (need_active_balance(&env)) {
            raw_spin_lock_irqsave(&busiest->lock, flags);
            /* don't kick the active_load_balance_cpu_stop,
             * if the curr task on busiest cpu can't be
             * moved to this_cpu
             */
            if (!cpumask_test_cpu(this_cpu,
                       tsk_cpus_allowed(busiest->curr))) {
                  raw_spin_unlock_irgrestore(&busiest->lock,
                                        flags);
                  env.flags |= LBF_ALL_PINNED;
                  goto out_one_pinned;
```

```
}
             * ->active balance synchronizes accesses to
             * ->active balance work. Once set, it's cleared
             * only after active load balance is finished.
            if (!busiest->active_balance) {
                  busiest->active balance = 1;
                  busiest->push cpu = this cpu;
                  active balance = 1;
            raw_spin_unlock_irgrestore(&busiest->lock, flags);
            if (active_balance) {
                  stop_one_cpu_nowait(cpu_of(busiest),
                        active load balance cpu stop, busiest,
                        &busiest->active balance work);
            }
             * We've kicked active balancing, reset the failure
             * counter.
             * /
            sd->nr balance failed = sd->cache nice tries+1;
} else
      sd->nr balance failed = 0:
if (likely(!active_balance)) {
      /* We were unbalanced, so reset the balancing interval */
      sd->balance interval = sd->min interval;
} else {
       * If we've begun active balancing, start to back off. This
       * case may not be covered by the all_pinned logic if there
       * is only 1 task on the busy runqueue (because we don't call
       * move tasks).
      if (sd->balance_interval < sd->max_interval)
            sd->balance_interval *= 2;
}
```

```
kernel/stop_machine.c:
 * stop machine from inactive_cpu - stop_machine() from inactive CPU
 * @fn: the function to run
 * @data: the data ptr for the @fn()
 * @cpus: the cpus to run the @fn() on (NULL = anv online cpu)
 * This is identical to stop machine() but can be called from a CPU which
 * is not active. The local CPU is in the process of hotplug (so no other
 * CPU hotplug can start) and not marked active and doesn't have enough
 * context to sleep.
 * This function provides stop_machine() functionality for such state by
 * using busy-wait for synchronization and executing @fn directly for local
 * CPU.
 * CONTEXT:
 * Local CPU is inactive. Temporarily stops all active CPUs.
 * RETURNS:
 * 0 if all executions of @fn returned 0, any non zero return value if any
 * returned non zero.
 */
int stop machine from inactive cpu(int (*fn)(void *), void *data,
                         const struct cpumask *cpus)
{
      struct stop machine data smdata = { .fn = fn, .data = data,
                                  .active cpus = cpus }:
     struct cpu_stop_done done;
      int ret;
      /* Local CPU must be inactive and CPU hotplug in progress. */
      BUG_ON(cpu_active(raw_smp_processor_id()));
      smdata.num_threads = num_active_cpus() + 1;
                                                     /* +1 for local */
     /* No proper task established and can't sleep - busy wait for lock. */
     while (!mutex_trylock(&stop_cpus_mutex))
           cpu relax();
      /* Schedule work on other CPUs and execute directly for local CPU */
      set_state(&smdata, STOPMACHINE_PREPARE);
      cpu_stop_init_done(&done, num_active_cpus());
```