

Chebyshev polynomial of the first kind

$T_n(x)$ (degree n)

1) definition $T_n(x)$ is solution of the following diff. eq.

$$(1-x^2) y'' - x y' + n^2 y = 0$$

Solution $T_n(x) = \cos(n \arccos x) \quad \therefore -1 \leq T_n(x) \leq 1$

Set $x = \cos \theta$

then $T_n(\cos \theta) = \cos n\theta$

e.g.

$$n=0 \quad T_0 = 1$$

$$\left. \begin{array}{l} n=1 \end{array} \right\} T_1 = \cos \theta = x$$

$$\left. \begin{array}{l} n=2 \end{array} \right\} T_2 = \cos 2\theta = 2\cos^2 \theta - 1 = 2x^2 - 1$$

$$T_3 = 4\cos^3 \theta - 3\cos \theta \Rightarrow 4x^3 - 3x$$

using $\cos(n+1)\theta + \cos(n-1)\theta = 2\cos \theta \cos n\theta, \quad n \geq 1$

We have $T_{n+1}(x) + T_{n-1}(x) = 2x T_n(x)$

recursion relation.

(2) orthogonal relation (continuous)

$$A = \int_{-1}^1 \frac{T_n(x) T_m(x)}{\sqrt{1-x^2}} dx = ?$$

$$= \int_{-1}^1 \frac{\cos n\theta \cos m\theta d(\cos\theta)}{\sin\theta}$$

$$= - \int_{\pi}^0 \cos n\theta \cos m\theta d\theta$$

$$a) \quad m=n=0, \quad A = \int_0^{\pi} d\theta = \pi$$

$$b) \quad m=n \neq 0, \quad A = \int_0^{\pi} \cos^2 n\theta d\theta = \frac{\pi}{2}$$

$$c) \quad m \neq n, \quad A = \int_0^{\pi} \frac{1}{2} (\cos(m+n)\theta + \cos(n-m)\theta) d\theta =$$

$$\therefore A = \begin{cases} 0 & m \neq n \\ \frac{\pi}{2} & m = n \neq 0 \\ \pi & m = n = 0 \end{cases}$$

* $T_n(x)$ is a polynomial of order n ,

it has n zeros in $[-1, 1]$ at

$$x = \cos \frac{\pi(k-\frac{1}{2})}{n}, \quad k=1, 2, \dots, n$$

Proof

$$A = \sum_{k=1}^m \cos i \theta_k \cos j \theta_k = \frac{1}{2} \sum_k [\cos(i+j)\theta_k + \cos(i-j)\theta_k]$$

a). if $i=j=0$, $A = m$.

b). if $i=j \neq 0$, $A = \frac{1}{2} \sum_k (\cos 2i \theta_k + 1)$

We will show

$$B = \sum_{k=1}^m \cos 2n \theta_k = 0, \quad \theta_k = \frac{\pi}{m} (k - \frac{1}{2})$$

$$B = \operatorname{Re} \left[\sum_k e^{i 2n \theta_k} \right]$$

$$= \operatorname{Re} \left[e^{-\frac{2\pi n}{2m} i} \sum_k e^{\frac{2\pi n}{m} i k} \right]$$

$$= \operatorname{Re} \left[e^{-\frac{\pi n}{m} i} e^{\frac{2\pi n}{m} i} \frac{1 - e^{\frac{2\pi n i}{m} m}}{1 - e^{\frac{2\pi i n}{m}}} \right]$$

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0

$$\therefore A = \frac{1}{2} \sum_k 1 = \frac{m}{2}$$

c). $i \neq j$

We will show for $m \neq n$

$$A = \frac{1}{2} \sum_k [\cos(m+n)\theta_k + \cos(m-n)\theta_k] = 0, \quad \theta_k = \frac{\pi}{N} (k - \frac{1}{2})$$

1st term

$$\theta_k = \frac{\pi}{N} (k - \frac{1}{2})$$

$$\operatorname{Re} \left[\sum_k e^{i(m+n)\theta_k} \right] = \operatorname{Re} \left[e^{-\frac{\pi}{2N}(m+n)i} \sum_k e^{\frac{\pi}{N}(m+n)ik} \right]$$

$$= \operatorname{Re} \left[e^{-\frac{\pi}{2N}(m+n)i} e^{\frac{\pi}{N}(m+n)i} \frac{1 - e^{\frac{\pi}{N}(m+n)i \times N}}{1 - e^{\frac{\pi}{N}(m+n)i}} \right]$$

$$= \operatorname{Re} \left[\frac{e^{\frac{\pi}{2N}(m+n)i}}{1 - e^{\frac{\pi}{N}(m+n)i}} (1 - (-1)^{m+n}) \right]$$

\downarrow
real

$$= \operatorname{Re} \left[\frac{(1 - e^{-\frac{\pi}{N}(m+n)i}) e^{\frac{\pi}{2N}(m+n)i}}{(1 - e^{\frac{\pi}{N}(m+n)i})(1 - e^{-\frac{\pi}{N}(m+n)i})} \right] (1 - (-1)^{m+n})$$

denominator is real

$$= \operatorname{Re} \left[e^{\frac{\pi}{2N}(m+n)i} - e^{-\frac{\pi}{2N}(m+n)i} \right] \frac{1 - (-1)^{m+n}}{|1 - e^{\frac{\pi}{N}(m+n)i}|^2}$$

\downarrow

$2i \sin \frac{\pi}{2N}(m+n)$ no real part

$$= 0$$

Similarly 2nd term = 0.

Proof $T_n(\cos \theta) = \cos n\theta = 0$

$$\therefore n\theta_k = k\pi - \frac{1}{2}\pi, \quad k=1, 2, \dots, n$$

$$\therefore \theta_k = \frac{\pi(k - \frac{1}{2})}{n}, \quad x_k = \cos \theta_k.$$

* $T_n(x)$ has $n+1$ extrema (maxima, minima)

$$\text{at } x_k = \cos \frac{\pi k}{n}, \quad k=0, 1, \dots, n$$

$$\therefore \cos n\theta = \pm 1 \Rightarrow n\theta_k = k\pi$$

$$\therefore \theta_k = \frac{k\pi}{n}$$

* This is a very useful property of T_n .

(3) discrete orthogonality relation.

for $i, j < m$

$$\sum_{k=1}^m T_i(x_k) T_j(x_k) = \begin{cases} 0 & i \neq j \\ \frac{m}{2} & i=j \neq 0 \\ m & i=j=0 \end{cases}$$

where x_k ($k=1, 2, \dots, m$) are zeros of $T_m(x)$.

Chebyshev approximation.

$$f(x) \approx \sum_{k=1}^N C_k T_{k-1}(x) - \frac{1}{2} C_1 \quad \text{in } [-1, 1] \quad \textcircled{A}$$

* better approximation for large N .

* to find coefficient C_k , we use zeros of T_N

assume $T_N(x_k) = 0$, $k=1, 2, \dots, N$

$$\text{then } f(x_k) = \sum_{j=1}^N C_j T_{j-1}(x_k) - \frac{1}{2} C_1$$

$$B_i = \sum_{k=1}^N T_i(x_k) f(x_k) = \sum_k \sum_j C_j T_i(x_k) T_{j-1}(x_k) - \frac{1}{2} C_1 \sum_k T_i(x_k) T_0(x_k)$$

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$$a). i=0, B_0 = N C_1 - \frac{1}{2} N C_1 = \frac{N}{2} C_1$$

$$b). i \neq 0, B_i = \frac{N}{2} C_{i+1}$$

$$\therefore B_i = \frac{N}{2} C_{i+1} \quad \text{for general } i$$

$$\therefore C_j = \frac{2}{N} B_{j-1} = \frac{2}{N} \sum_{k=1}^N T_{j-1}(x_k) f(x_k)$$

Note that eq. \textcircled{A} is exact when $x = x_k$.

Why Chebyshev approximation?

* minimax polynomial p_L .

Among all polynomials g_L of the same degree, p_L has the smallest maximum deviation from the true function $f(x)$.

$$\begin{aligned} f(x) &= \sum_L g_L a_L && a_L \text{ coefficient.} \\ &= \sum_L p_L a_L && \leftarrow \text{minimax.} \end{aligned}$$

* although minimax polynomial is very difficult to find, Chebyshev approx. polynomial is almost minimax.

* generating function of $T_n(x)$

$$\sum_{n=0}^{\infty} T_n(x) t^n = \frac{1 - tx}{1 - 2tx + t^2}$$