



NATIONAL TEXTILE
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DEPARTMENT OF COMPUTER SCIENCE

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SECTION SE: 5th (A)

Operating System-Hometask lab

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Operating Systems – COC 3071

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After-mid Homework -1

Part 1: Semaphore theory

1. A counting semaphore is initialized to 7. If 10 wait() and 4 signal() operations are performed, find the final value of the semaphore.

Initial value of semaphore=7

Then 10 times wait, for 7 times, the value of semaphore is decremented by 1

So semaphore =0, blocked=3

Then 4 signal():

1:wake a process

2:wake a process

3:wake a process

4:No process to wake so increment by 1.

Semaphore value=1

2. A semaphore starts with value 3. If 5 wait() and 6 signal() operations occur, calculate the resulting semaphore value.

When 5 wait() called, semaphore has value 3. So, $5-3=2$, 2 are **BLOCKED**.

Value =0.

When 6 times **Signal()** occur:

1st signal: 1st block wake.

2nd signal: 2nd block wake

$6-2=4$. No process to wake anymore so increment by 4.

Value=0+4

3. A semaphore is initialized to 0. If 8 signal() followed by 3 wait() operations are executed, find the final value.

When semaphore starts with value 0. When 8 **signal()** called, semaphore increases by 8.

Value = 8. Then 3 **wait()** are called:

$8 - 3 = 5$, no process is blocked. **Value = 5.**

4. A semaphore is initialized to 2. If 5 wait() operations are executed:

a) **How many processes enter the critical section?**

2 processes will enter the critical section

b) **How many processes are blocked?**

3 processes are blocked

5. A semaphore starts at 1. If 3 wait() and 1 signal() operations are performed:

a) **How many processes remain blocked?**

1 process remains blocked.

b) **What is the final semaphore value?**

When semaphore starts with value 1. When 3 **wait()** called, semaphore has value 1.

So, $3 - 1 = 2$, 2 are **BLOCKED**.

Value = -2. When 1 **signal()** occurs:

1st signal: 1 **blocked process wakes up**. Value = -1

No increment happens. Value = -1.

6.

semaphore S = 3;

wait(S); wait(S);

signal(S);

wait(S); wait(S);

a) **How many processes enter the critical section?**

4 processes enter the wait section

b) **What is the final value of S?**

The final value is: $3 - 2 + 1 - 2 = 0$

7.

semaphore S = 1;

wait(S); wait(S);

signal(S);

signal(S);

a) **How many processes are blocked?**

0 processes blocked

b) **What is the final value of S?**

Semaphore S starts with value 1.

wait(S) called first time:

$S = 1 - 1 = 0$, no process blocked.

wait(S) called second time:

$S = 0$, 1 process is **BLOCKED**.

Value = -1.

signal(S) called first time:

1 **blocked process wakes up**.

Value becomes 0.

signal(S) called second time:

No blocked process to wake, so semaphore increments.

Value = 1.

8. A binary semaphore is initialized to 1. Five **wait()** operations are executed without any **signal()**. How many processes enter the critical section and how many are blocked?

1 process enters a critical section. The other 4 are blocked.

9. A counting semaphore is initialized to 4. If 6 processes execute wait() simultaneously, how many proceed and how many are blocked?

6-4=2, 2 processes are blocked and 4 processes proceed.

10. A semaphore S is initialized to 2. wait(S); wait(S); wait(S); signal(S); signal(S); wait(S);

a) Track the semaphore value after each operation.

For first wait(S): value=2-1=1

For second wait(S): value=1-1=0

For third wait(S): value=0-1=-1

For First Signal: value=-1+1=0

For Second signal: 0+1=1

For 4th wait: 1-1=0

b) How many processes were blocked at any time?

1 process

11. A semaphore is initialized to 0. Three processes execute wait() before any signal(). Later, 5 signal() operations are executed.

a) How many processes wake up?

All processes wake up

b) What is the final semaphore value?

Final semaphore value: 0-3+5=2

Part 2: Semaphore Coding

Consider the Producer–Consumer problem using semaphores as implemented in Lab-10 (Lab-plan attached). Rewrite the program in your own coding style, compile and execute it successfully, and explain the working of the code in your own words.

Submission Requirements:

- Your rewritten source code
- A brief description of how the code works
- Screenshots of the program output showing successful execution

```
#include <stdio.h>
#include <pthread.h>
#include <semaphore.h>
```

```

#include <unistd.h>

sem_t slots;

void* vehicle(void* arg) {
    int num = *(int*)arg;

    printf("Vehicle %d wants to enter.\n", num);

    sem_wait(&slots);

    printf("Vehicle %d entered successfully.\n", num);

    sleep(2);

    printf("Vehicle %d exited.\n", num);

    sem_post(&slots);

    return NULL;
}

int main() {
    pthread_t threads[10];
    int numbers[10];

    sem_init(&slots, 0, 3);

    for (int i = 0; i < 10; i++) {
        numbers[i] = i + 1;
        pthread_create(&threads[i], NULL, vehicle, &numbers[i]);
    }

    for (int i = 0; i < 10; i++) {
        pthread_join(threads[i], NULL);
    }
    sem_destroy(&slots);
    return 0;
}

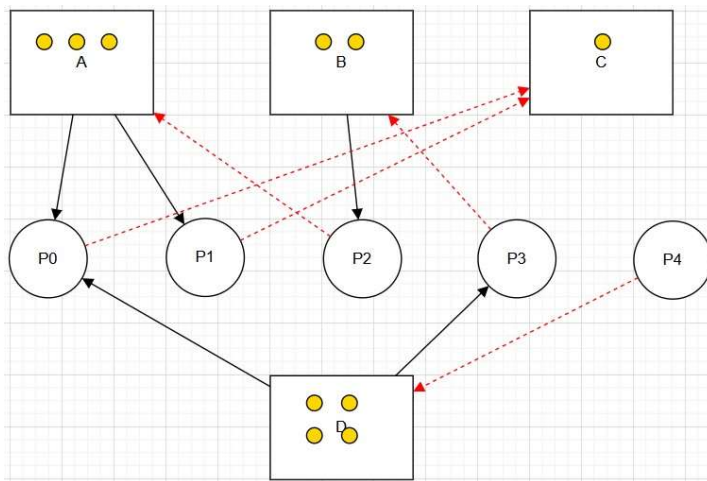
```

EXPALNATION:

- A semaphore 'slots' represents the available parking spaces, its value is 3.
- Multiple threads are created, where each thread represents a vehicle.
- Each vehicle tries to enter by calling `sem_wait()`.
- If a slot is available, the vehicle enters; otherwise, it waits (blocks) until a slot is free.
- The slot is free, when semaphore has value greater than 0.
- Once inside, the vehicle stays for some time using `sleep()`.
- After leaving, the vehicle calls `sem_post()` to free the slot.
- This ensures that only 3 vehicles can enter at a time, even though 10 threads are running.

Part 3: RAG (Recurse Allocation Graph)

- Convert the following graph into matrix table ,



Allocation matrix:

Process	A	B	C	D
P0	1	0	0	1
P1	1	0	0	0
P2	0	1	0	0
P3	0	0	0	1
P4	0	0	0	0

Resource Matrix:

Process	A	B	C	D
P0	0	0	1	0

P1	0	0	1	0
P2	1	0	0	0
P3	0	1	0	0
P4	0	0	0	1

Part 4: Banker's Algorithm

System Description:

- The system comprises five processes (P0–P3) and four resources (A,B,C,D).
- Total Existing Resources:

Total			
A	B	C	D
6	4	4	2

- Snapshot at the initial time stage:

	Allocation				Max				Need			
	A	B	C	D	A	B	C	D	A	B	C	D
P0	2	0	1	1	3	2	1	1				
P1	1	1	0	0	1	2	0	2				
P2	1	0	1	0	3	2	1	0				
P3	0	1	0	1	2	1	0	1				

Questions:

1. Compute the Available Vector:

- Calculate the available resources for each type of resource.

2. Compute the Need Matrix:

- Determine the need matrix by subtracting the allocation matrix from the maximum matrix.

3. Safety Check:

- Determine if the current allocation state is safe. If so, provide a safe sequence of the processes.

- Show how the Available (working array) changes as each process terminates.

Allocation	Max	Need	Available:-
A B C D	A B C D	A B C D	A B C D
P ₀ 2 6 1 1	3 2 1 1	1 2 0 0	6 4 4 2
P ₁ 1 1 0 0	1 2 0 2	0 1 0 2	
P ₂ 1 0 1 0	3 2 1 0	2 2 0 0	
P ₃ 0 1 0 1	2 1 0 1	2 0 0 0	

⇒ P₃ executes.

Allocation	Max	Need	Available:-
A B C D	A B C D	A B C D	A B C D
P ₀ 2 0 1 1	3 2 1 1	1 2 0 0	4 5 4 3
P ₁ 1 1 0 0	1 2 0 2	0 1 0 2	
P ₂ 1 0 1 0	3 2 1 0	2 2 0 0	
P ₃ 0 0 0 0	0 0 0 0	0 0 0 0	

⇒ P₂ executes :-

Allocation					Max					Need					<u>Available</u>
	A	B	C	D		A	B	C	D		A	B	C	D	
P ₀	2	0	1	1		3	2	1	1		1	2	0	0	<u>5 5 5 2</u>
P ₁	1	1	0	0		1	2	0	2		0	1	0	2	
P ₂	0	0	0	0		0	0	0	0		0	0	0	0	
P ₃	0	0	0	0		0	0	0	0		0	0	0	0	

⇒ P₁ executes :-

Allocation					Max					Need					<u>Available</u>
	A	B	C	D		A	B	C	D		A	B	C	D	
P ₀	2	0	1	1		3	2	1	1		1	2	0	0	<u>5 5 5 2</u> 6 6 5 3
P ₁	0	0	0	0		0	0	0	0		0	0	0	0	
P ₂	0	0	0	0		0	0	0	0		0	0	0	0	
P ₃	0	0	0	0		0	0	0	0		0	0	0	0	

⇒ P₀ Executes :-

Allocation					Max					Need					<u>Available</u>
	A	B	C	D		A	B	C	D		A	B	C	D	
P ₀	0	0	0	0		0	0	0	0		0	0	0	0	<u>7 4 6 4</u>
P ₁	0	0	0	0		0	0	0	0		0	0	0	0	
P ₂	0	0	0	0		0	0	0	0		0	0	0	0	
P ₃	0	0	0	0		0	0	0	0		0	0	0	0	

P₃ → P₂ → P₁ → P₀ Safe ✓

Submission Guidelines:

- Ensure all answers are well-explained and calculations are shown step-by-step.
- Submit your assignment on MS Team and GitHub in a PDF format.
- VIVA based Evaluation so Develop your own solution after getting help.