







LEARNING to CONTEST ARGUMENTATIVE CLAIMS

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- Assumption-based Argumentation ABA frameworks
- Learning ABA frameworks
- Contesting argumentative claims
- Redressing as a way of learning



- Assumption-based Argumentation
 - **ABA frameworks**
- Learning ABA frameworks
- Contesting argumentative claims
- Redressing as a way of learning



Rule-based systems

for non-monotonic reasoning formalisms, can be used for **explainable AI** by providing **arguments** for **claims**

Arguments are **structured: derivations** built from rules supported by **assumptions**

Rules are **defeasible** by deriving arguments for **contraries** of assumptions

 Assumption-based Argumentation ABA frameworks

- Learning ABA frameworks
- Contesting argumentative claims
- Redressing as a way of learning

Automated logic-based learning
of ABA frameworks from
background knowledge
+
positive & negative examples

Algorithm based on **transformation rules**, using **Answer Set Programming**

- Assumption-based Argumentation ABA frameworks
- Learning ABA frameworks
- Contesting argumentative claims
- Redressing as a way of learning

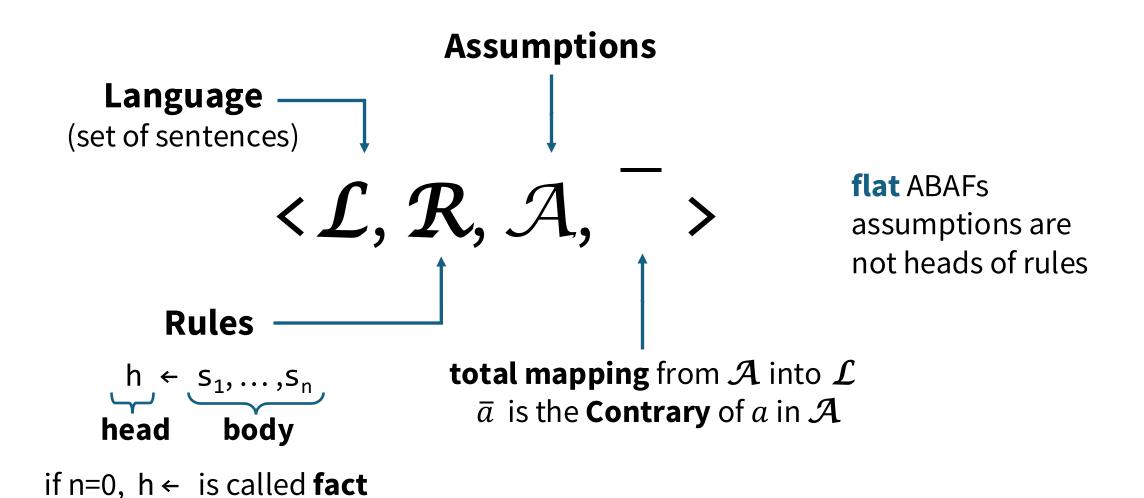
highly desirable property for **human-centric Al**

claims are subject to **contestation**:

- rejected claims may be desirable
- accepted claims may be undesirable

modify rules incrementally to **reconcile** contestations

ABA FRAMEWORKS



ABA FRAMEWORKS

an example ...

```
employed(jo) ← onleave(jo) ← maternity(jo) ←
employed(bob) ← onleave(bob) ← maternity(diana) ←
employed(claudia) ← facts

loan(X) ← employed(X), nobreaks(X)

breaks(X) ← onleave(X)

rules

contrary of
nobreaks
```

nobreaks(X) renders the rule defeasible:
it can be applied only if breaks(X) cannot be derived

ABA FRAMEWORKS - SEMANTICS

"acceptable" extensions: sets of arguments able to "defend" themselves from "attacks" (as determined by the chosen semantics)

```
employed(jo) ← onleave(jo) ← maternity(jo) ←
employed(bob) ← onleave(bob) ← maternity(diana) ←
employed(claudia) ←
loan(X) ← employed(X), nobreaks(X)
breaks(X) ← onleave(X)
```

- Arguments are deductions of claims using rules and supported by assumptions
- Attacks are directed at the assumptions in the support of arguments

```
{
    arg1: { nobreaks(jo) } ⊢ loan(jo)
    arg2: { nobreaks(bob) } ⊢ loan(bob)
    arg3: { nobreaks(claudia) } ⊢ loan(claudia)
    arg4: { } ⊢ breaks(jo)
    arg5: { } ⊢ breaks(bob)
}
```

We focus on stable extensions

any set of arguments S that

- 1. do not attack each other (conflict-free)
- 2. S attacks all arguments it does not contain

Accepted claims: loan(claudia) Rejected claims: loan(jo) loan(bob)

Given

- 1. ABA framework $\mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle$ (background knowledge) with at least one stable extension
- 2. **Ep** = { **positive examples** }
- 3. En = { negative examples }
- 4. T = { learnable predicates }

find $\mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'}, \mathcal{A'}, - \rangle$ with a stable extension S such that

- i. **F** ⊆ **F**'
- ii. **positive** are **covered**: every positive has an argument in S
- iii. negative are not covered: no negative has an argument in S

Given

- 1. ABA framework $\mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle$ (background knowledge) with at least one stable extension
- 2. **Ep** = { **positive examples** }
- 3. En = { negative examples }
- 4. T = { lea loan es }

find $\mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'}, \mathcal{A'}, - \rangle$ with a stable extension S such that

- i. **F** ⊆ **F**'
- ii. **positive** are **covered**: every positive has an argument in S
- iii. negative are not covered: no negative has an argument in S

```
Given
```

```
1. ABA framework \mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle (background knowledge)
      with at least one stable extension
  2. Ep = { positive examples }
  3. En = { negative examples }
  4. T = \{ lea \}
                                    les }
                       loan
find \mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'} |
                    employed(jo) ←
                                                 onleave(jo) ←
                                                                        maternity(jo) ←
      F ⊆ F' employed(bob) ←
                                                 onleave(bob) ←
                                                                        maternity(diana) ←
       positive ar employed(claudia) ←
  ii.
```

iii. negative are not covered: no negative has an argument in S

```
Given
```

```
1. ABA framework \mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle (background knowledge)
      with at least one stable extension
  2. Ep = { | loan(claudia) | s }
  3. En = { negative examples }
  4. T = \{ lea \}
                                    les }
                       loan
find \mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'} |
                     employed(jo) ←
                                                  onleave(jo) ←
                                                                         maternity(jo) ←
 i. F \subseteq F' employed(bob) \leftarrow
                                                  onleave(bob) ←
                                                                         maternity(diana) ←
      positive ar employed(claudia) ←
  iii.
       negative are not covered: no negative has an argument in S
```

```
Given
```

```
1. ABA framework \mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle (background knowledge)
      with at least one stable extension
  2. Ep = { loan(claudia) s }
  3. En = \{ nell loan(bob) es \}
  4. T = \{ led \}
                                    les }
                       loan
find \mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'} |
                    employed(jo) ←
                                                 onleave(jo) ←
                                                                        maternity(jo) ←
                   employed(bob) ←
                                                 onleave(bob) ←
                                                                        maternity(diana) ←
      F⊆F'
      positive ar employed(claudia) ←
  ii.
```

iii. **negative** are **not covered**: no negative has an argument in S

```
Given
```

```
1. ABA framework \mathbf{F} = \langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \rangle (background knowledge)
         with at least one stable extension
   2. Ep = { | loan(claudia) | s }
   3. En = \{ ne \mid loan(bob) \mid es \}
   4. T = \{ led \}
                                                les }
                               loan
find \mathbf{F'} = \langle \mathcal{L'}, \mathcal{R'} | \text{employed(jo)} \leftarrow \text{onleave(jo)} \leftarrow \text{maternity(jo)} \leftarrow
i. \mathbf{F} \subseteq \mathbf{F'} employed(bob) \leftarrow onleave(bob) \leftarrow maternity(diana
                                                                                                maternity(diana) ←
        postive a employed(claudia) ←
  iii. negative a loan(X) ← employed(X), nobreaks(X)
breaks(X) ← onleave(X)
F' is a solution to the ABA tearning problem
```

ABA LEARNING VIA TRANSFORMATION RULES

Learning ABA frameworks relies upon a set of transformation rules

$$<\mathcal{L}_1,\mathcal{R}_1,\mathcal{A}_1,\overset{-\imath}{>}\longrightarrow <\mathcal{L}_2,\mathcal{R}_2,\mathcal{A}_2,\overset{-\imath}{>}\longrightarrow \dots\longrightarrow <\mathcal{L}_n,\mathcal{R}_n,\mathcal{A}_n,\overset{-\imath}{>}$$
 background knowledge intensional solution

→ ∈{ Rote Learning, Folding, Assumption Introduction, Subsumption }

learnt rules **do not**make explicit
reference to specific
values in the universe

A **strategy** controls the order of application of the transformation rules

TRANSFORMATION RULES at work ROTE LEARNING

Add facts

- from **positive** examples
- for contraries of assumptions

to get a (non-intensional) solution

It's enough to learn

TRANSFORMATION RULES at work FOLDING

Towards an **intensional** solution ...

Generalise

loan(X) ← X=claudia

to

 $loan(X) \leftarrow employed(X)$

by using

 $employed(X) \leftarrow X=claudia$

WARNING

It also constructs an argument for a negative example: loan(bob)

ABA LEARNING is **parametric** w.r.t. the folding strategy
It includes a portfolio of strategies, such as "nondeterministic" and "greedy" folding

TRANSFORMATION RULES at work ASSUMPTION INTRODUCTION

Repairing the ABA framework to get a solution ...

Add an **assumption** to avoid

- rejecting a positive example
- accepting a negative example

```
loan(X) ← employed(X), nobreaks(X)
```

with contrary breaks(X)

AND REPEAT!

Rote Learning breaks(X) ← X=bob

Folding breaks(X) ← onleave(X)

No more rules to learn: LEARNING COMPLETED!

```
employed(jo) ← onleave(jo) ← maternity(jo) ←
employed(bob) ← onleave(bob) ← maternity(diana) ←
employed(claudia) ←

loan(X) ← employed(X), nobreaks(X)
breaks(X) ← onleave(X)
```

CONTESTATION

Given

1. an ABA framework F' want of **solution** of the ABA learning problem (**F**, <**Ep**, **En**>, **T**) not c 2. a new claim (example) $c \notin Ep \cup En$ want of c **c** ? (predicted) claims F' Ep, En, F_{BK}

CONTESTATION

Given

- an ABA framework F' solution of the ABA learning problem (F, <Ep, En>, T)
- 2. a new claim (example) *c* ∉ Ep ∪ En

F' is **contested** by

want of <i>c</i> want of <i>not c</i>

iff ∄ stable extension S of **F**'s.t.

Ep U { c } are covered in S Ep are covered in S & En are not covered in S En U { c } are not covered in S

REDRESS

When a solution **F**' is **contested** by either **want of** c or **want of not** c, **redressing** consists in deriving a **solution F**" of the ABA learning problem

```
• (want of c) (F, ⟨Ep∪{c}, En⟩, T)
```

• (want of *not c*) (F, $\langle Ep, En \cup \{c\} \rangle, T$)

How to redress

from scratch

trivial form

- 1. forgetting F'
- 2. solving the <u>original</u> problem w/c
 either (F, ⟨Ep ∪ {c}, En ⟩, T)
 or (F, ⟨Ep, En ∪ {c}⟩, T)

Incremental

modifies F' as little as possible

- selecting some of the **learnt** rules & making them **defeasible** by assumption introduction, to get the <u>new</u> problem (F'_{ai}, 〈 Ep, En 〉, T'_{ai})
- solving the new problem w/c
 either (F'_{ai}, ⟨Ep ∪ {c}, En ⟩, T'_{ai})
 or (F'_{ai}, ⟨Ep, En ∪ {c}⟩, T'_{ai})

INCREMENTAL REDRESS w/ABA Learning

Suppose (the current solution) F'

```
employed(jo) ← onleave(jo) ← maternity(jo) ←
employed(bob) ← onleave(bob) ← maternity(diana) ←
employed(claudia) ←
loan(X) ← employed(X), nobreaks(X)
- breaks(X) ← onleave(X)
```

is contested by the **want of** loan(jo), which is **not covered** by any stable extension We can **incrementally modify** F" by applying the transformation rules

```
→ breaks(X) ← onleave(X), alpha(X) (1) by assumption introduction rule
c_alpha(X) ← X=jo (2) by rote learning rule
c_alpha(X) ← maternity(X) (3) by folding rule
```

to get a new solution **F**" with at least one stable extension covering loan(jo)

Algorithm 1: RASP-ABAlearn

```
Input: (\langle \mathcal{R}_0, \mathcal{A}_0, \overline{\phantom{a}}^0 \rangle, \langle \mathcal{E}^+, \mathcal{E}^- \rangle, \langle \mathcal{E}_C^+, \mathcal{E}_C^- \rangle, \mathcal{T}): redress problem
      Output: \langle \mathcal{R}, \mathcal{A}, \overline{\phantom{a}} \rangle: incremental redress relative to \langle \mathcal{E}_{\mathcal{C}}^+, \mathcal{E}_{\mathcal{C}}^- \rangle
 1 \mathcal{R} := \mathcal{R}_0; \mathcal{A} := \mathcal{A}_0; \overline{\phantom{a}} := \overline{\phantom{a}}^0; \mathcal{R}_l := \emptyset;
 2 RoLe(); Gen(); return \langle \mathcal{R}, \mathcal{A}, \overline{\phantom{A}} \rangle;
  3 Procedure RoLe()
              P := ASP(\langle \mathcal{R}, \mathcal{A}, \overline{\phantom{A}} \rangle, \langle \mathcal{E}^+, \mathcal{E}^- \rangle, \langle \mathcal{E}_C^+, \mathcal{E}_C^- \rangle, \mathcal{T});
              if \neg sat(P) then
                      fail:
              _{
m else}
                      S := as(P);
                      // R1. Rote Learning
                      for each newp(t) \in S do
                             \mathcal{R}_l := \mathcal{R}_l \cup \{p(X) \leftarrow X = t\};
11
                      \mathbf{end}
12
              \mathbf{end}
13 Procedure Gen()
              foreach \rho: (p(X) \leftarrow X = t) \in \mathcal{R}_l do
                      \mathcal{R}_l := \mathcal{R}_l \setminus \{\rho\};
                      // R4. Fact Subsumption
                      if \neg sat(ASP(\langle \mathcal{R} \cup \mathcal{R}_l, \mathcal{A}, \overline{\ }), \langle \mathcal{E}^+, \mathcal{E}^- \rangle, \langle \mathcal{E}_G^+, \mathcal{E}_G^- \rangle, \emptyset)) then
16
                              // R2 w/ R3. Folding with Assumption Introduction
                             \langle \rho_{\alpha}, \alpha(X), C_{\alpha} \rangle := Folding WAsmIntro(\rho);
17
                             \mathcal{R}\!:=\!\mathcal{R}\!\cup\!\{
ho_g\};
18
                             \mathcal{A} := \mathcal{A} \cup \{\alpha(X)\};
                             \overline{\alpha(X)} := c \quad \alpha(X);
20
                              // R1. Rote Learning
                             for each c \alpha(t) \in C_{\alpha} do
21
                                    \mathcal{R}_l := \mathcal{R}_l \cup \{c \mid \alpha(X) \leftarrow X = t\};
23
                             end
                      \mathbf{end}
25
              \mathbf{end}
26 Function Folding WAsmIntro(ρ)
               // R2. Folding
              while foldable(\rho, \mathcal{R}) do
27

\rho := fold(\rho, \mathcal{R});

28
              // R3. Assumption Introduction
              Let \rho be H \leftarrow B; X := vars(B);
31
              if there exists \alpha(X) \in \mathcal{A} relative to B then
                      \rho_q := H \leftarrow B, \alpha(X); \quad C_\alpha := \emptyset;
                     if \neg sat(ASP(\langle \mathcal{R} \cup \{\rho\}, \mathcal{A}, \overline{\ \ \ }), \langle \mathcal{E}^+, \mathcal{E}^- \rangle, \langle \mathcal{E}_C^+, \mathcal{E}_C^- \rangle, \emptyset)) then
34
                             fail;
                      \mathbf{end}
              else // introduce an assumption \alpha(X), with a new predicate \alpha
                      \rho_g := H \leftarrow B, \alpha(X);
                      F := \langle \mathcal{R} \cup \{\rho\}, \mathcal{A} \cup \{\alpha(X)\}, \overline{\phantom{A}} \cup \{\alpha(X) \mapsto c \ \alpha(X)\} \rangle;
                     C_{\alpha} := \{c_{\alpha}(X) \mid c_{\alpha}(X) \in as(ASP(F, \langle \mathcal{E}^+, \mathcal{E}^- \rangle, \langle \mathcal{E}^+_C, \mathcal{E}^-_C \rangle, \{c_{\alpha}\}))\};
39
              return \langle \rho_g, \alpha(X), C_\alpha \rangle;
```

A GLIMPSE OF IMPLEMENTATION via ASP

ASP encoding

```
loan(X) :- employed(X), nobreaks(X). ...
nobreaks(X) :- employed(X), not breaks(X).
breaks(X) :- onleave(X).
```

```
{ breaksP(X) } :- onleave(X).
breaks(X) :- breaksP(X).
#minimize{1,X: breaksP(X)}.
:- not loan(claudia).
:- loan(bob).
```

learning facts for contraries

Answer sets

(1-to-1 correspondence with **Stable extensions**)

```
{ breaks(bob), ...}, ...
```

Rote learning



https://github.com/ABALearn/aba_asp

EXPERIMENTS

RASP-ABAlearn

https://github.com/ABALearn/aba_asp

$|T_S| 13524$ $|S_S| 6953$ S_R 6953 $|T_R|$ 8482 $|S_S|$ 6519 42250 42173 42357 42673 | 4232 $\langle 1503, 1374 \rangle |S_S| 33409$ 33409 33410 33410 33410 autism33409 33410 33410 33410 33410 **≈** 6568 $|T_R|471191|$ 11680 | 11409 | 279 $\langle 171, 464 \rangle$ $\langle 214, 1587 \rangle \mid S_S \mid 34762 \mid 34762 \mid 34763 \mid 34763 \mid 34764 \mid 34763 \mid 34763 \mid 34763 \mid$

Learning problems

six standard datasets of the UCI ML Repo as ABA learning problems

Experimental processes

Redress from scratch (S) vs. Incremental redress (R)

11 runs of RASP-ABALearn:

: run (S) and (R) using 90% of positive and negative examples

1-10: run (S) and (R) each using a randomly selected **new example** either **positive** or **negative**

												The second secon
		+	+	+	+	_	+	_	_	_	_	+
	$ T_S $	13524	14427	14552	15004	14985	15252	15048	15114	15246	15097	15329
	$ T_R $	12741	970	905	999	44	1126	47	46	44	44	1144
>	$ S_S $	6953	6954	6955	6956	6956	6958	6958	6958	6958	6958	6961
	$ S_R $	6953	6954	6955	6956	6956	6957	6957	6957	6957	6957	6958

size of the background knowledge (# rules)

 $S_R \mid 34762 \mid 34762 \mid 34763 \mid 34763 \mid 34764 \mid 34765 \mid 34765$

 $\langle |Ep|, |En| \rangle$

 $\frac{T_S}{T_R}$ time in ms

 S_S

of learn rules

CONCLUSIONS

- Contestability for ABA frameworks learnt from background knowledge + positive and negative examples
- Method for incremental redress
- Redressing as a way of learning from additional positive or negative examples
- experiments show that incremental redress is more efficient than re-learning from scratch

