

$$R = 20 \text{ M}$$

$$R = 100$$

$$B = 40\%$$

unpolar

$$V = 14$$

$$P = 6$$

$$P_R = 7$$

$$a) \text{bit}_R = \left\lfloor \frac{B}{R} \right\rfloor = 4$$

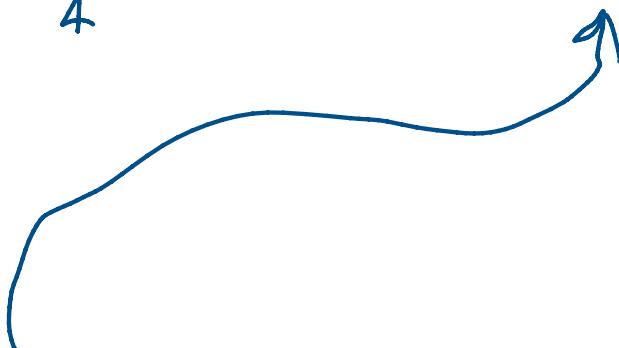
$$\text{bit}_i = \left\lfloor \frac{B}{V+P} \right\rfloor = 204$$

40%  
14 + 6 > 20

$$m_0 = \left\lceil \frac{R}{\text{bit}_R} \right\rceil = 5 \text{ M}$$

$\frac{20 \text{ M}}{\text{bit}_R}$   
 $\frac{1}{4}$

$$m_{0,i} = \left\lceil \frac{m_0}{\text{bit}_i} \right\rceil = 2451 \Omega$$



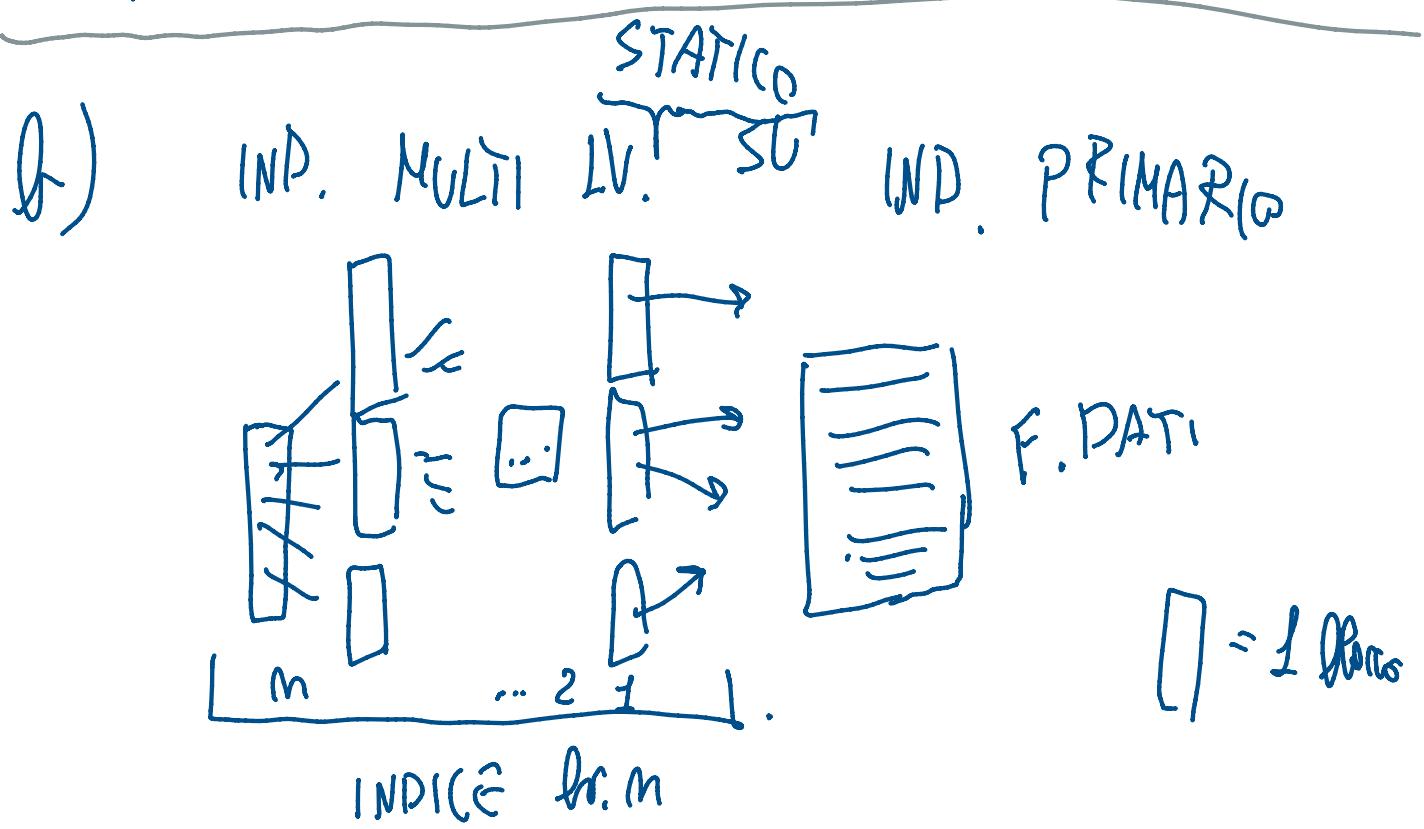
$$m_{acc \text{ IP.}} = m_{acc \text{ indice}} + 1 =$$

$$= \log_2(m_{0,i}) + 1 =$$

=  $D_{nn} (1111 \dots)$

$$\begin{aligned} &= \log_2(24510) + 1 = \\ &= 15 + 1 = 16 \end{aligned}$$

(neste indice no f. multidimensional)

$$\log_2(m_b) = 23$$


$$lfp_2 = (\text{some prime}) = 4$$

$$lfp_2 = (\text{some prime}) = 204 [\text{n bloco}]$$

$m_{\text{bloco}} = 1 \quad \dots \quad 5M \quad \dots \quad n \quad \dots$

$$M_{lb} = (\dots) = 5M \quad M_{lb_i} = (\dots) = 24510$$

$f_0 = lb_i \pi_i \leftarrow$  FAN OUT / F. DI RIDUZIONE

- prende un indice
- si conta IP finché non ottiene 2 blocchi

$$M_{lb_1} = M_{lb_i} = 24510$$

$$M_{lb_2} = \left\lceil \frac{M_{lb_1}}{f_0} \right\rceil = 121$$

204

$$M_{lb_3} = \left\lceil \frac{M_{lb_2}}{f_0} \right\rceil = 1$$

$$\begin{aligned} \text{dim Totali ind MLV} &= M_{lb_1} + M_{lb_2} + M_{lb_3} = \\ &= 24510 + 121 + 1 = 24632 \text{ blocchi} \end{aligned}$$

$$\text{Maccioni } \text{MLV} \leq M_{lb_1} + 1 = 3 + 1 = 4$$

c)  $\beta$ , oltre ...  $\rightarrow$  piena al 70%.

c) D'avorio

$$P \leq \left( \text{Cap. max di un nodo} \right)$$

$$P \cdot P + (P-1)(V + Pr) \leq B$$

(num. a modi figli)

come (dove, P. additi)

ORDINE DEL' ALBERO

a reward

$$P \cdot 6 + (P-1)(14 + 7) \leq 4096$$

$$6P + 21P - 21 \leq 4096$$

$$27P \leq 4117 \quad P \leq \frac{4117}{27} \approx 152,48$$

$$P = 152$$

$$\Rightarrow m_a \text{ media } P \cdot \text{a figli} = 106$$

$$m_a \text{ media } P \cdot \text{a altri} = 105$$

OBIETTIVO < ouvre  $\geq 20M$  entry

Inv.	m modi	m entry	m. Punt. figli
I	1	105	106

$$\text{II} \quad 106 \quad 105 \cdot 106 \approx \quad 106 \cdot 106 = \\ 11130 \quad 11236$$

$$\text{III} \quad 11236 \quad 105 \cdot 11236 = \quad 106 \cdot 11236 = \\ 1179780 \quad = 1191016$$

$$\text{IV} \quad 1191016 \quad 105 \cdot \\ 1191016 = \\ = 125056680$$

$$n_{\text{way}} = 105 + 11130 + 1179780 + 125056680 = \\ = 126247195 \quad (> 20M)$$

$$n_{\text{blochi}} = 1 + 106 + 11236 + 1191016 = \\ = 1202359 \text{ blochi}$$

$$n_{\text{occhi}} = \text{occhi indici} + 1 = \\ = n_{\text{hr}} + 1 = \\ = 4 + 1 = 5$$

s'è fatto  $B^+$

2 sq.

modi non:  $P \cdot P + (P-1) \cdot V \leq B$

" foggia:  $P_{bef} (V + P_k) + P \leq B$

