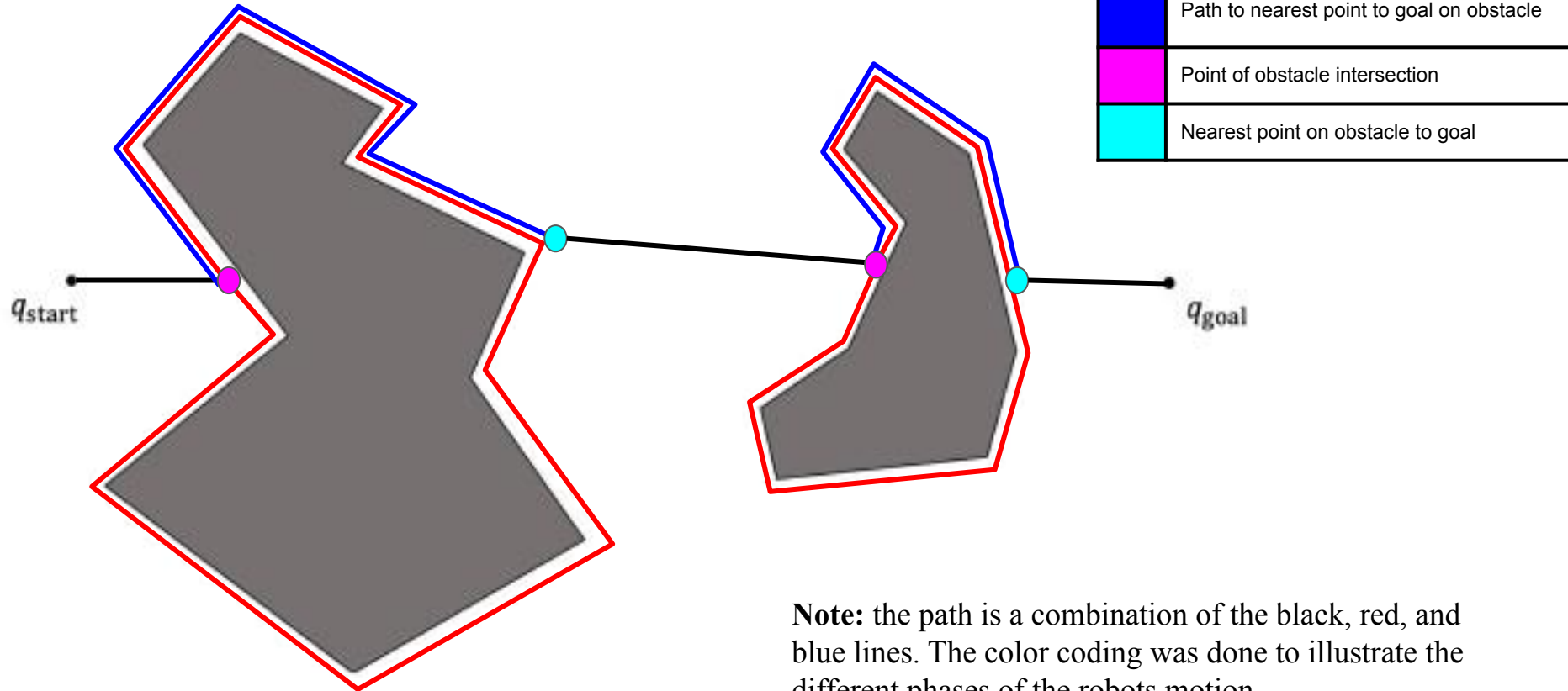
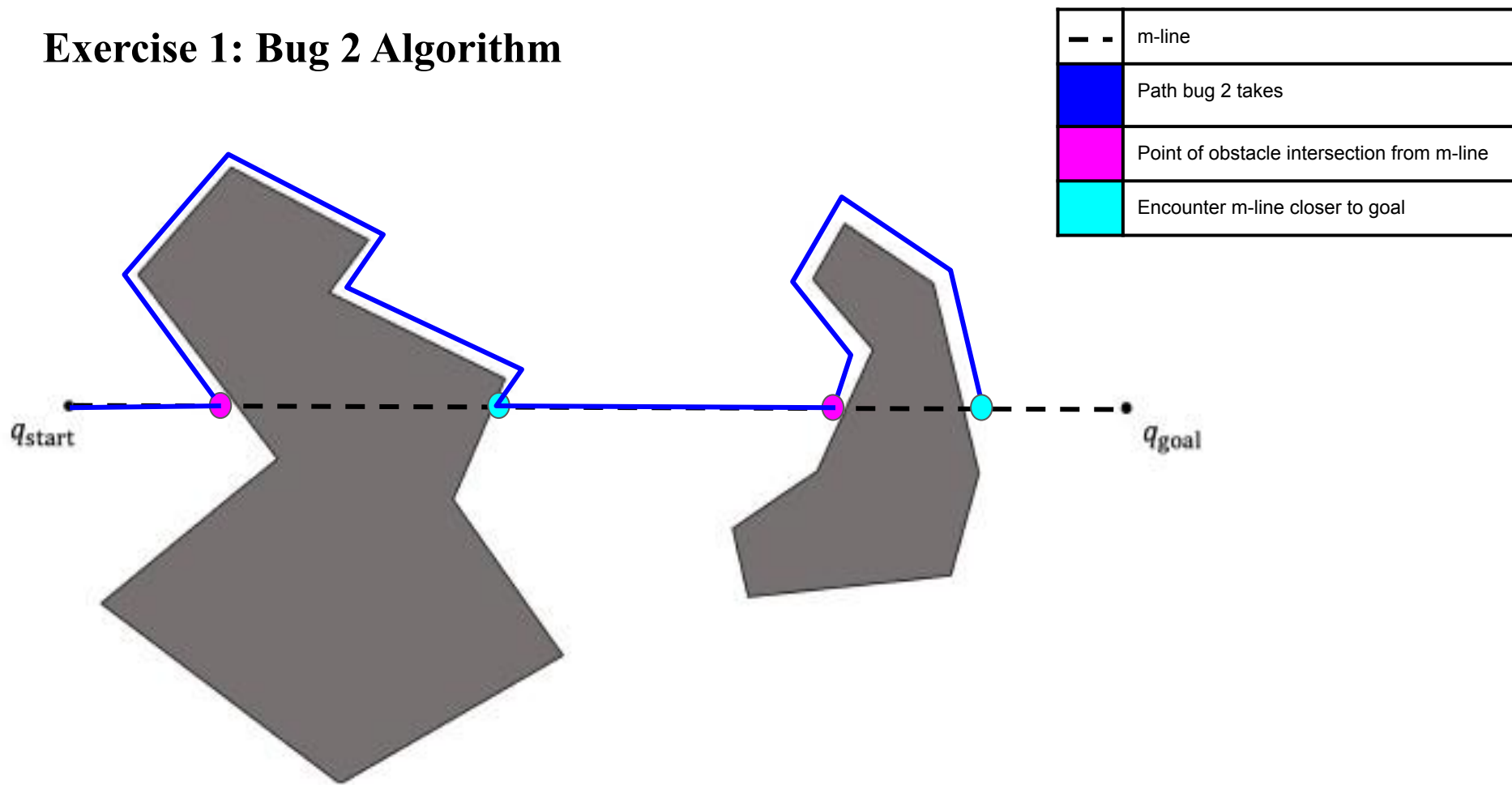


Exercise 1: Bug 1 Algorithm

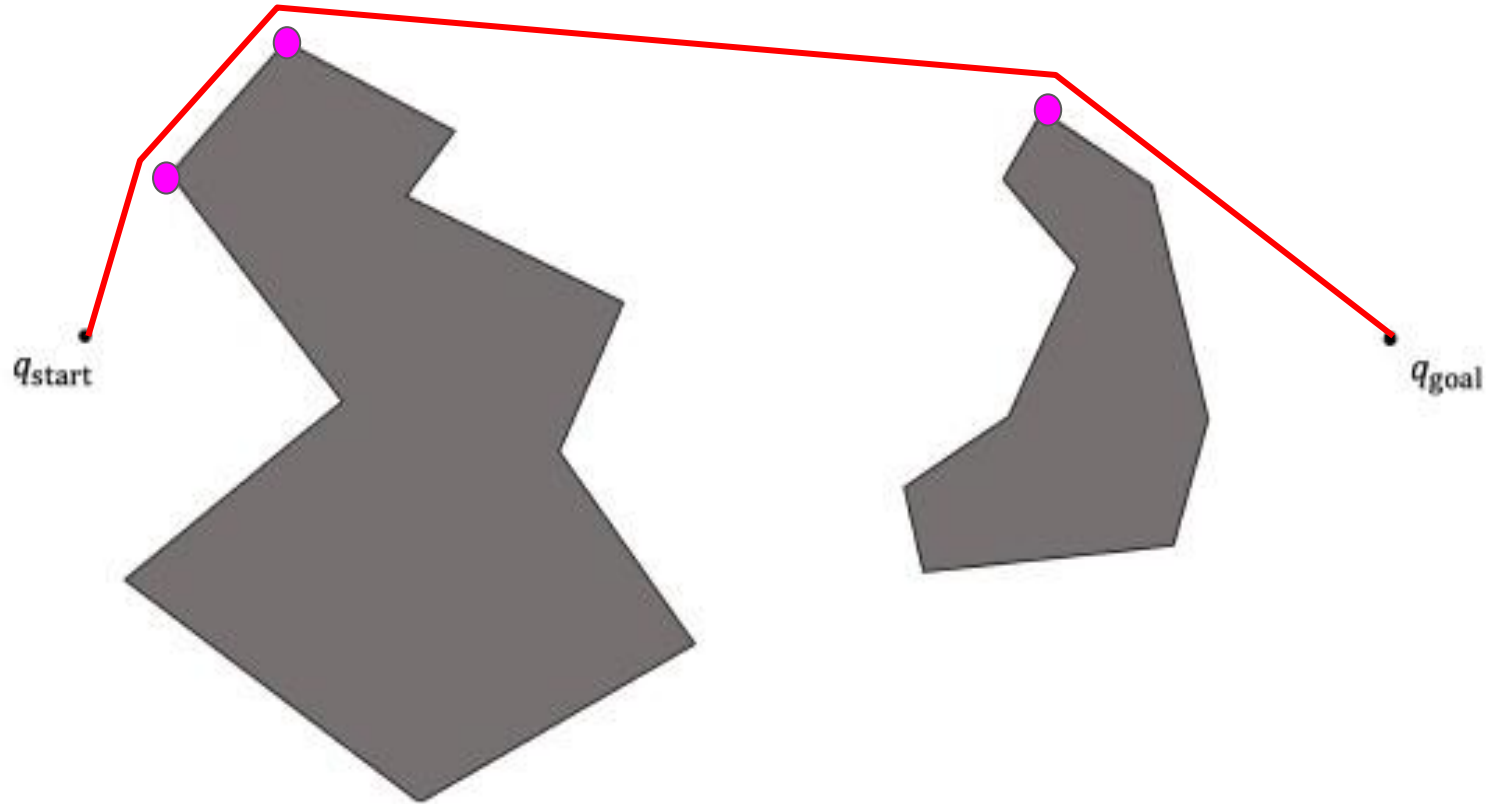


Exercise 1: Bug 2 Algorithm



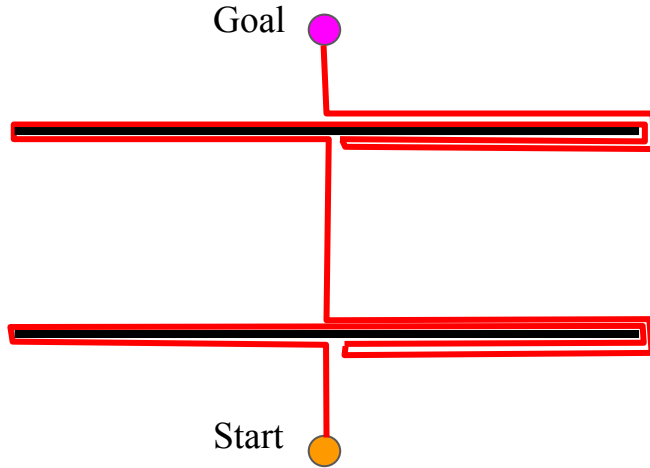
Exercise 1: Tangent Bug Algorithm

	Tangent bug path
	Discontinuities that minimize heuristic

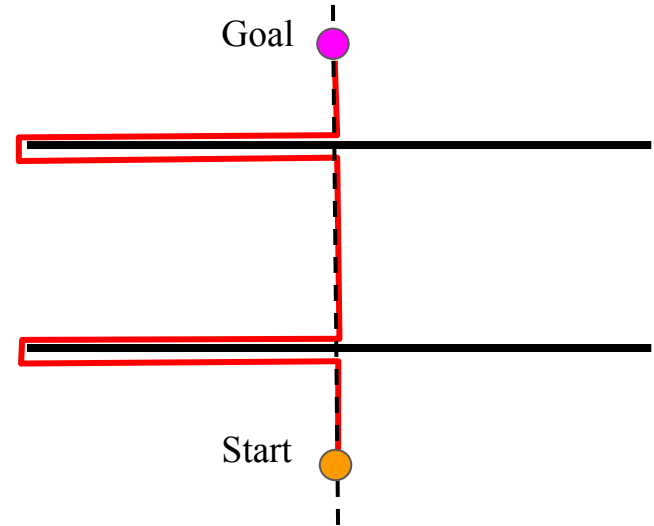


Exercise 2

Bug 1



Bug 2

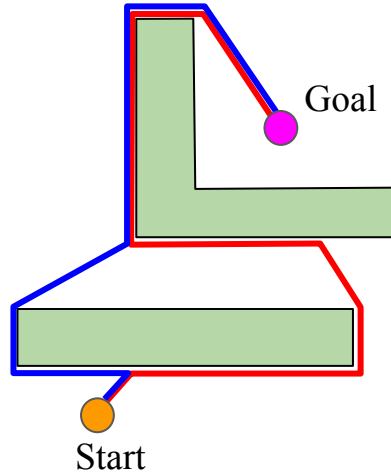


We can image obstacles of 1 dimension (a line of distance w) with perimeters equal to $2w(*)$. Here the Bug 1 algorithm takes a path of $D + 1.5\sum P_i$ - easily verified by inspection. Bug 2 outperforms Bug 1 by twice the addition of the obstacles' perimeters ($2\sum P_i$).

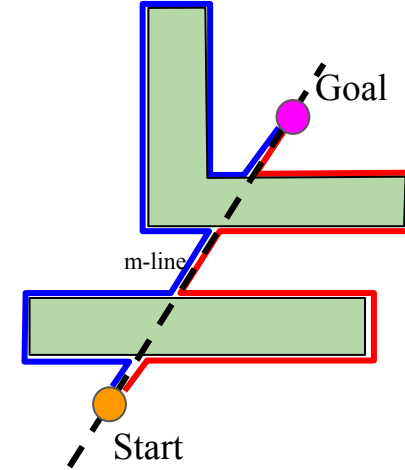
(*) Perimeter = $2w + 2h$, and we take the limit as $h \rightarrow 0$, we get that the perimeter of a line is $2w$

Exercise 3

Tangent Bug



Bug 2

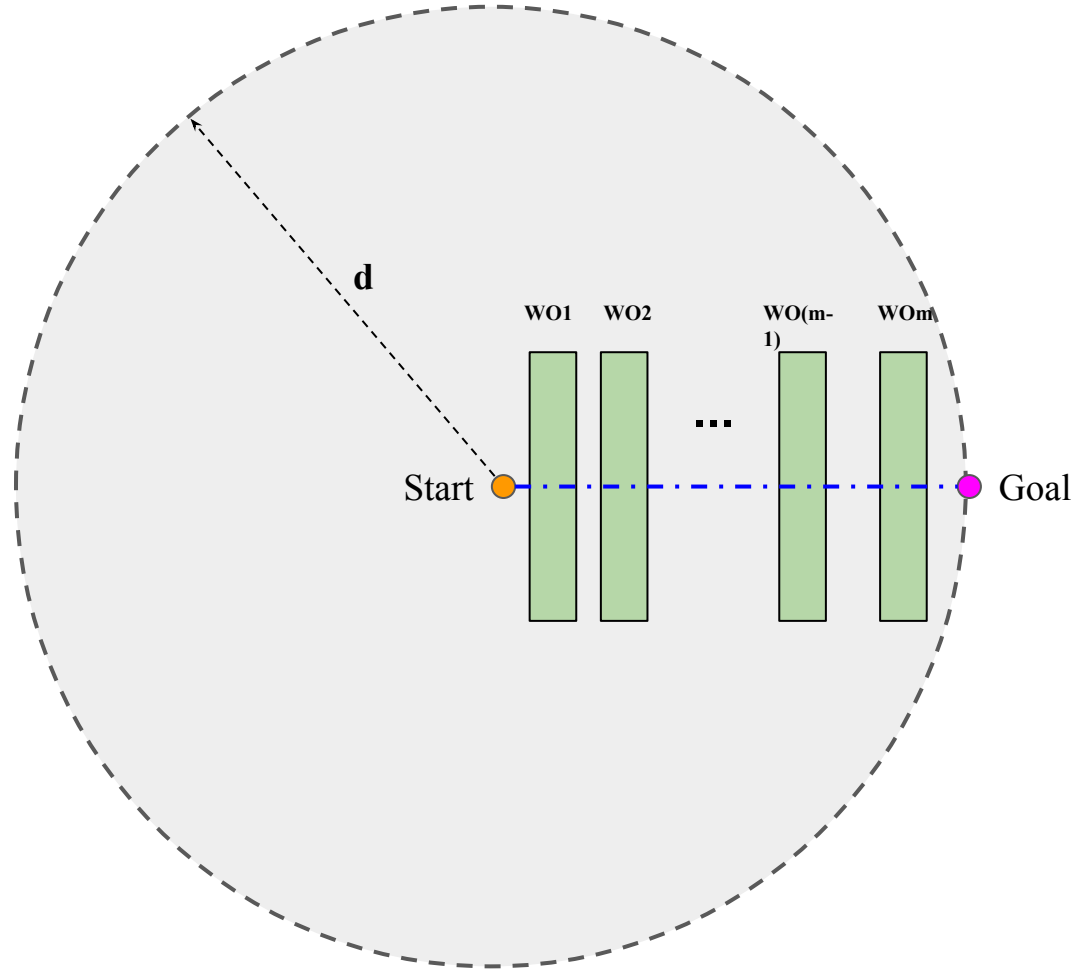


The path the Bug 2 algorithm takes is largely dependent on whether it goes **CW** or **CCW** when it encounters an obstacle, this similar for the Tangent Bug. We can see that because the Tangent Bug is allowed to move toward the goal once $d_{reach} < d_{follow}$, it goes back into motion-to-goal once the bug has found a way around the obstacle. Bug 2 always travels on the m-line (if it finds a point on the m-line closer to the goal).

Exercise 4

Suppose all of the obstacles are rectangles and the rectangles are all the same size and are parallel to each other. Arrange all of the interior obstacles WO_i , $i \in [1, m]$, between start and goal such that the line between start and goal intersects each object twice - seen on the right. This means that every hit point and leave point q_{hi} and q_{li} , respectively, will lie on the m-line*. This means that q_{li} will necessarily lead to $q_{h(i+1)}$, until we reach q_{lm} , and then there should be a clear path to goal. For this reason, the maximum number of obstacles Bug 1 will encounter is m obstacles.

*** I understand that this definition of the m-line is used for Bug 2, but is a very convenient definition for this example**



Exercise 5

Is the Tangent Bug algorithm complete?

Proof:

- 1) Suppose that the Tangent Bug Algorithm (TBA) is NOT complete (proof by **contradiction**)
- 2) Therefore, there exists a path from start to goal
 - a) By assumption it is finite length, and encounters obstacles a finite number of times
- 3) Let's say that TBA does not find a path to goal
 - a) Either it **never terminates** or it **terminates**
 - b) Suppose it **never terminates**
 - i) If it does not terminate, then the robot keeps finding leave(L)/depart(D) points that are not the corresponding hit points(H)
 - ii) Each L/D is closer to goal than the corresponding H. This means we are approaching the goal steadily but surely with each step
 - iii) Because each step takes us closer, there are a finite number of H/D/L pairings we will exhaust before reaching goal
 - iv) A goal will be found (**contradiction**)
 - c) Suppose it **terminates** (incorrectly)
 - i) The closest point after a H point must be a L point where it moves into the corresponding obstacle
 - (1) Between this L point and the goal is another point on the obstacle close to the goal (Jordan Curve theorem)
 - (2) Because we assumed that a path exists, we should have encountered this point, **contradicting** the definition of a leave point