Question 1 [30%]

- (i) Give both a regular expression and a deterministic finite automaton that captures Give both a regular expression and a telegraph of letters and digits) that contain the string all alphanumeric strings (i.e. consisting of letters and digits) that contain the string all alphanumeric strings (You may make use of the following shorthands) all alphanumeric strings (i.e. consisting of the following shorthand: (dot) "august" as a substring. (You may make use of the following shorthand: (dot)for "any symbol" and $[\wedge x]$ for "any symbol other than x".) (ii) Draw a deterministic finite automaton that captures the strings over the alphabet
- Draw a deterministic finite attended and the each consecutive pair of as there is at least $\{a,b,c\}$ in the following set: Between each consecutive pair of as there is at least one b and and least one c.
- (iii) Give a regular expression that captures the set of binary strings such that every consecutive block of five symbols contains at most three zeros. (10%)

Question 2 [25%]

(i) Consider the following simple grammar for the programming language PLO.

```
(prog) \rightarrow (list)
\langle list \rangle \rightarrow \langle assignment \rangle; \langle list \rangle \mid \epsilon
(assignment) \rightarrow id := (expr)
(expr) \rightarrow (expr) + (term) \mid (expr) - (term) \mid (term)
(term) → (term) * (factor) | (term) div (factor) | (term) mod
(factor) | (factor)
(factor) \rightarrow ((expr)) \mid id \mid num
```

Give a complete parse tree for the following program.

```
x := (17);
y := x + 1;
z := y;
```

(10%)(ii) Modify the grammar so that semicolons act as statement separators rather than statement terminators (i.e. so that a three-statement program would appear as

```
x := 17;
y := x + 1;
z := y
```

i.e. without a semicolon after the last statement.)

(iii) Modify the grammar of Part (i) to incorporate a simple while-loop construct with

```
while <expr> do
  begin
  t>
  end
```

(10%)

Question 3 [30 %]

(i) The following grammar is intended to capture the syntax of simple context-free grammars. The symbols \Rightarrow . • , N , \parallel and T are all terminals. State which nonterminals are nullable.

$$\begin{array}{c} \langle G \rangle \rightarrow \langle P \rangle \langle G' \rangle \\ \langle G' \rangle \rightarrow \langle P \rangle \langle G' \rangle \mid \epsilon \\ \langle P \rangle \rightarrow \langle L \rangle \Rightarrow \langle R \rangle \bullet \\ \langle L \rangle \rightarrow N \\ \langle R \rangle \rightarrow \langle A \rangle \langle A' \rangle \\ \langle A \rangle \rightarrow \langle S \rangle \langle S' \rangle \\ \langle A' \rangle \rightarrow \parallel \langle A \rangle \langle A' \rangle \mid \epsilon \\ \langle S \rangle \rightarrow N \mid T \\ \langle S' \rangle \rightarrow \langle S \rangle \langle S' \rangle \mid \epsilon \end{array}$$

(5%)

- (ii) Analyze the grammar to calculate the FIRST sets for each of the nonterminals in the grammar. (10%)
- (iii) Give a succinct description of a recursive descent parser this grammar. Give pseudocode for the parsing methods for nonterminals (G), (G') and (P). Give a crisp summary (in words) of the behaviour of the remaining parsing methods. Summarise carefully any assumptions regarding the behaviour of the scanner or any ancillary methods you rely upon.

Question 4 [15 %]

(i) Consider the code fragment based on the PL0 grammar introduced in Question 2 as augmented with the while loop of Part (iii) of that question.

```
sum := 0;
k := 1;
while k <= 100 do
   begin
   sum := sum + k;
end;</pre>
```

Give a translation of the while loop shown into three-address code (TAC), showin clearly the correspondence between the PL0 code and the TAC. (10%

(ii) Write a brief note on the relationship between the concepts of regular expression and finite automata, specifically on the expressive power of both to capture pattern in strings. (5%)