

# A simple example

A simple design for DB for managing a company's orders for spare parts:

- Keeps track of suppliers we use and items we require
- Also track details of outstanding orders

Three tables:

- suppliers (snum, sname, status, city)
- parts (pnum, pname, colour, weight)
- ordered\_from (snum, pnum, quantity)

We assume status is city-dependent, e.g. a delivery charge.

## A worse design

A worse design (because of more redundancy), would be to move (status) column into the ordered\_from table.

Now the same city-dependent numbers are being repeated many times.

# Design principle: avoid redundancy

We want to avoid redundancy wherever possible.

We try to represent one fact in one place, rather than in multiple places.

## Issues arising from redundancy

- Duplicating info in several places wastes space.
- If we change one copy of a piece of info, we need to make sure we change them all, which is inconvenient.

# Normalisation

## Normal forms

Normal forms capture desirable traits that reduce risk of certain DB problems.

They are ordered by number, with higher numbers denoting stricter conditions.

## First normal form

R is in First Normal Form (1NF) if and only if it contains atomic values only (no multi-valued attributes). We took this as given when we started looking at DBs.

## A poor design

first (snum, status, city, pnum, quantity) and parts (pnum, pname, colour, weight)

Here we've combined the info on suppliers and orders together.

### Problems:

- The city value for suppliers is replicated many times.
- Can't record suppliers details until that supplier supplies at least one part.
- If we delete the last tuple for a supplier, we lose all the information about that supplier.

# Functional Dependencies

Dependencies capture how different attributes in table relate to one another.

Definition: Given relation R, attribute Y is functionally dependent on attribute X if and only if, whenever two tuples of R agree on their X-value, they are also guaranteed to agree on their Y-value. Notated ( $X \rightarrow Y$ ).

By analysing these dependencies with normalisation, we can reduce redundancy problems.

## Functional dependencies in our poor design

parts:

- $pnum \rightarrow pname$
- $pnum \rightarrow weight$
- $pnum \rightarrow colour$

first:

- $pnum, snum \rightarrow status, city, quantity$

- $\text{snun} \rightarrow \text{status, city}$
- $\text{city} \rightarrow \text{status}$

In this example, pnum and snun together form the key for the 'first' table, but some attributes are determined only by part of this key (snun). This is called a non-full dependency.

## Second normal form

A DB without any non-full dependencies is in Second Normal Form (2NF).

### A second (now better) design

If we split first(...) into second(snun, status, city) and ordered\_from(snun, pnum, quantity), then we have a DB in 2NF.

There are still some problems:

- The status value for each city is replicated many times.
- Can't enter status values for a city until we have a supplier located in that city.
- If we delete the only tuple for a particular city, we lose the status info for that city.

### Dependencies in the second design

- $\text{snun} \rightarrow \text{city}$
- $\text{city} \rightarrow \text{status}$
- $\text{snun} \rightarrow \text{status}$

Status has a transitive dependence on snun via city.

We can fix this by splitting the table into suppliers(snun, city) and cities(snun, status).

## Third normal form

DBs without any transitive dependencies are in Third Normal Form.