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There will be another assignment soon.

# **Equations from Truth Tables**

## Example 1:

Suppose we have 3 inputs (A, B, C) and 1 output F that corresponds to each of the input combinations according to the following truth table:

Α	В	С	F	Minterms
0	0	0	0	m₀
0	0	1	0	m <sub>1</sub>
0	1	0	1	m <sub>2</sub>
0	1	1	1	m <sub>3</sub>
1	0	0	0	$m_4$
1	0	1	0	m <sub>5</sub>
1	1	0	1	m <sub>6</sub>
1	1	1	1	m <sub>7</sub>

Note F is chosen to reflect our requirements—the values for F come from what we're trying to do/solve.

We can get an equation by ORing  $m_2$ ,  $m_3$ ,  $m_6$ , and  $m_7$  together:

$$F = (A'.B.C') + (A'.B.C) + (A.B.C') + (A.B.C)$$

We can optimise this to use fewer gates:

$$F = A'.B.(C' + C) + A.B.(C' + C)$$
  
 $F = A'.B + A.B$   
 $F = B(A' + A)$   
 $F = B$ 

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You can see this is clear if you look at the table. We have now gone from 12 gates to 0 gates required.

### Example 2:

Derive the equation from the following truth table, simplify if possible, and draw the corresponding circuit.

This is how he may write a question in an exam.

Α	В	С	F	Minterms
0	0	0	0	m0
0	0	1	1	m1
0	1	0	0	m2
0	1	1	1	m3
1	0	0	1	m4
1	0	1	0	m5
1	1	0	1	m6
1	1	1	0	m7

$$F = m1 + m3 + m4 + m6$$

$$F = A'.B'.C + A'.B.C + A.B'.C' + A.B.C'$$

$$F = A'.C.(B' + B) + A.C'.(B' + B)$$

$$F = A'.C + A.C' = A \oplus C$$

# **Example 3: Nightclub Bouncer Circuit**

#### Inputs:

- 1.  $\geq$  18? (call this A for age)
- 2. Drunk? (call this D for drunk)
- 3. Male/female? (call this G for gender)
- 4. Well-dressed? (call this C for clothes)

So if A = 1, then you're >= 18.

D = 1, then you're drunk.

G = 1, you're female; G = 0, you're male.

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C = 1, you're well-dressed.

А	D	G	С	Door
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	1
1	0	1	0	1
1	0	1	1	1
1	1	0	0	0
1	1	0	1	1
1	1	1	0	0
1	1	1	1	1

$$\begin{aligned} \text{Door} &= m_3 + m_9 + m_{10} + m_{11} + m_{13} + m_{15} \\ \text{Door} &= \text{A'.D'.G.C} + \text{A.D'.G'.C'} + \text{A.D'.G.C} + \text{A.D.G'.C} + \\ &\quad \text{A.D.G.C} \end{aligned}$$

#### Optimise:

Door = 
$$D'.G.C.(A' + A) + A.D.C.(G + G') + A.D'.(G'.C + G.C')$$

Door = D'.G.C + A.D.C + A.D'.(
$$G \oplus C$$
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