

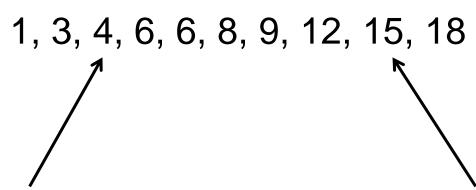
Selecting the kth ordered item from an unordered list

Now the journalist who was successfully treated for breast cancer is fronting a campaign to promote awareness of clinical trials in Ireland for the treatment of cancer.

Ireland has a clinical trial participation rate of 6pc, which is high internationally and puts the country in the top quartile globally, but ICORG is seeking to increase this number.

A common task in data analytics, scientific computing, social sciences, statistics, economics, etc is identifying *percentiles* in a collection of items.

Given a set of items, the kth percentile is the value below which k% of the items are ranked.



70th percentile, because 70% of values come after this point,

10th percentile, because 10% of values come after this point, The *median* is a measure of the average of a collection.

less influenced by extreme values than the mean

The median is the 50th percentile.

The upper *quartile* is the subset of values that are 'better' than the 75th percentile.

The lower *quartile* is the subset of values that are 'poorer' than the 25th percentile.

Given an unsorted list of items, how should we compute the kth percentile? how should we compute the median?

6 4 12 9 6 15 3 18 1 8

Method 1:

repeat k times:
find the biggest item in the list
replace that item with None
report the most recent item found

Analysis:

Each iteration is O(n), so O(k*n)

For small $k \ll \log n$, this gives k*O(n) = O(n)

For $k \sim = \log n$, this is $O(n \log n)$

For large $k \gg \log n$, this is $O(n^2)$

Method 2:

for each element in list:

add element to a sorted linked list
step through linked list to position k

Analysis:

Worst case is for a sorted list, additions require $O(1 + 2 + 3 + ... + n) = O(n^2)$ followed by k steps, so $O(n^2)$

Similar analysis for adding to a sorted array-based list, but instead of $O(n^2)$ comparisons, need $O(n^2)$ assignments for a reverse sorted list to shift up each item

Method 3:

sort the list by preferred sorting algorithm read item in position k

Analysis:

For merge or heap sort, O(n log n)
For quicksort, O(n²), but O(n log n) on average

Method 4:

apply quicksort idea to identify element in position k

```
pick a pivot
sweep through list to gather all items less than, equal to or greater than pivot
if more items 'less than pivot' than k
recurse on all items less than pivot
else if more items 'less than or equal to pivot' than k
return pivot
else
recurse on all items greater than pivot
```

'decrease and conquer'

Note: don't sort every sublist – all we are doing is picking a pivot from some part of the list, dividing elements from that part, choosing *one* of the divisions, and recursing.

Analysis of method 4:

- worst case always choose a pivot that decreases the list by just 1 item, so n + (n-1) + (n-2) + ... + (n-(n-(k+1))), which is $O(n^2)$ for small k
 - average case

By a similar argument to the one for quicksort, it is possible to show that the average time for a list of length n is O(n).

Practical considerations:

We can implement quickselect in-place, in the same style as in-place quicksort.

Efficient implementation in python is not so obvious ...

Exercise:

- develop versions of quickselect which:
- (i) creates lists less, equal and more using python lists and append
- (ii) creates lists less, equal and more using 3 list comprehensions
- (iii) implements in place
- and algorithms which
- (iv) pre-sort with mergesort, heapsort or quicksort, then select
- (v) pre-sort with python list sort, then select

Which is fastest on average?

Next lecture

Graphs and Graph Algorithms