## Algolab Graph and BGL Introduction

Martin Raszyk<sup>1</sup>

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<sup>&</sup>lt;sup>1</sup>using material from Andreas Bärtschi, Petar Ivanov, Chih-Hung Liu, and Daniel Wolleb

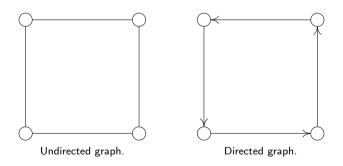
### Roadmap

- definition of a graph and its representations
- declaring and initializing a graph in BGL
- examples of standard graph algorithms in BGL
- ▶ tutorial problem: from a problem statement to a full solution

## Graph: Definition

#### Definition

A graph G = (V, E) consists of a set of *n* vertices (nodes) V and m edges E.



### Directed Graph: Representation

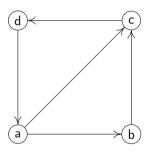
Adjacency matrix:

Space complexity:  $\Theta(n^2)$ .

Adjacency list: USE THIS REPRESENTATION!

Vertex	List of neighbors
a	[b, c]
b	[c]
С	[d]
d	[a]

Space complexity:  $\Theta(n+m)$ .

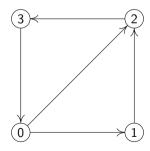


# Adjacency List in C++ Standard Library

```
#include <vector>

typedef std::vector<int> neighbor_list;
typedef std::vector<neighbor_list> cpp_graph;
```

vertex	neighbor_list
0	[1, 2]
1	[2]
2	[3]
3	[0]



## Adjacency List in BGL

► C++ Standard Library

```
#include <vector>

typedef std::vector<int> neighbor_list;
typedef std::vector<neighbor_list> cpp_graph;
```

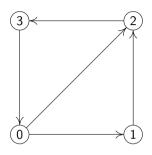
► BGL

```
#include <boost/graph/adjacency_list.hpp>
typedef boost::adjacency_list<boost::vecS, boost::directedS> graph;
```

## Warm-up: Initialize a Graph...

```
void init_graph()
{
  graph G(4);

  boost::add_edge(0, 1, G);
  boost::add_edge(1, 2, G);
  boost::add_edge(2, 3, G);
  boost::add_edge(3, 0, G);
  boost::add_edge(0, 2, G);
```



#### WARNING

boost::add\_edge(0, 7, G); would extend the vertex set of G to eight vertices!

#### Warm-up: ... and Iterate over its Edges

► all edges:

```
typedef boost::graph_traits<graph>::edge_iterator edge_it;
edge_it e_beg, e_end;
for (boost::tie(e_beg, e_end) = boost::edges(G); e_beg != e_end; ++e_beg) {
   std::cout << boost::source(*e_beg, G) << " " << boost::target(*e_beg, G) << "\n";
}</pre>
```

neighbors of a vertex:

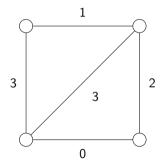
```
typedef boost::graph_traits<graph>::out_edge_iterator out_edge_it;

out_edge_it oe_beg, oe_end;
for (boost::tie(oe_beg, oe_end) = boost::out_edges(0, G); oe_beg != oe_end; ++oe_beg) {
   assert(boost::source(*oe_beg, G) == 0);
   std::cout << boost::target(*oe_beg, G) << "\n";
}
} /* end of function init_graph */</pre>
```

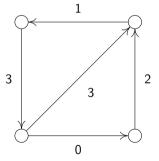
# Weighted Graph: Definition

#### **Definition**

A weighted graph G = (V, E, w) consists of a set of vertices V, edges E, and a weight function  $w : E \to \mathbb{Z}$ .



Undirected weighted graph.



Directed weighted graph.

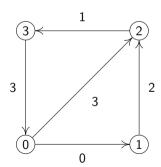
#### Weighted Graph in BGL

#### Definition

A weighted graph G = (V, E, w) consists of a set of vertices V, edges E, and a weight function  $w : E \to \mathbb{Z}$ .

### Initialize a Weighted Graph

```
typedef boost::property_map<weighted_graph, boost::edge_weight_t>::type weight_map;
typedef boost::graph_traits<weighted_graph>::edge_descriptor edge_desc;
void init_weighted_graph()
  weighted_graph G(4);
  weight_map weights = boost::get(boost::edge_weight, G);
  edge_desc e:
     boost::add_edge(0, 1, G).first; weights[e]=0;
     boost::add_edge(1, 2, G).first; weights[e]=2;
     boost::add_edge(2, 3, G).first; weights[e]=1;
     boost::add_edge(3, 0, G).first; weights[e]=3;
  e = boost::add_edge(0, 2, G).first; weights[e]=3;
 edge_weight_t: a C++ type
 edge_weight: a C++ value
```



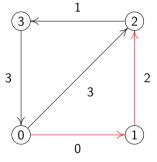
### Predefined Vertex and Edge Properties

Some *predefined* vertex and edge properties:

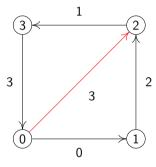
- vertex\_degree\_t
- vertex\_distance\_t
- edge\_weight\_t
- edge\_capacity\_t
- edge\_residual\_capacity\_t
- edge\_reverse\_t

All property maps must be initialized and maintained manually!

#### Shortest Path between Two Vertices



Shortest path from 0 to 2.



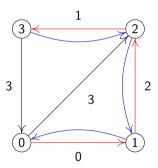
Longer path from 0 to 2.

#### Distance between Two Vertices: Dijkstra's Algorithm

Time complexity of boost::dijkstra\_shortest\_paths is  $O(n \log n + m)$ .

# Shortest Path between Two Vertices: Dijkstra's Algorithm

```
#include <boost/graph/dijkstra_shortest_paths.hpp>
typedef boost::graph_traits<weighted_graph>::vertex_descriptor vertex_desc;
int dijkstra_path(const weighted_graph &G, int s, int t, std::vector<vertex_desc> &path) {
 int n = boost::num_vertices(G);
 std::vector<int>
                           dist_map(n);
 std::vector<vertex_desc> pred_map(n):
 boost::dijkstra_shortest_paths(G, s,
   boost::distance_map(boost::make_iterator_property_map(
     dist_map.begin(), boost::get(boost::vertex_index, G)))
/* dot! */ .predecessor_map(boost::make_iterator_property_map(
     pred_map.begin(), boost::get(boost::vertex_index, G))));
 int cur = t:
 path.clear(); path.push_back(cur);
 while (s != cur) {
   cur = pred_map[cur];
   path.push_back(cur);
 std::reverse(path.begin(), path.end());
 return dist_map[t];
```

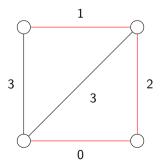


Shortest path edges are red. Predecessor edges are blue.

# Minimum Spanning Tree

#### **Definition**

A minimum spanning tree of a connected undirected weighted graph G = (V, E) is an acyclic subgraph of G connecting all vertices in V and having the minimum sum of edge weights.



### Minimum Spanning Tree: Kruskal's Algorithm

```
#include <boost/graph/kruskal_min_spanning_tree.hpp>
typedef boost::adjacency_list<br/>boost::vecS, boost::vecS, boost::undirectedS,
 boost::no_property, boost::property<boost::edge_weight_t, int> > weighted_graph;
typedef boost::graph_traits<weighted_graph>::edge_descriptor
                                                                   edge desc:
void kruskal(const weighted_graph &G) {
 std::vector<edge desc> mst: // vector to store MST edges (not a property map!)
 boost::kruskal_minimum_spanning_tree(G, std::back_inserter(mst));
 for (std::vector<edge_desc>::iterator it = mst.begin(); it != mst.end(); ++it) {
    std::cout << boost::source(*it, G) << " " << boost::target(*it, G) << "\n":
```

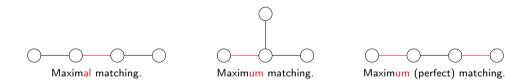
Time complexity of boost:: $kruskal_minimum_spanning_tree$  is  $O(m \log m)$ .

## Maximum Matching

#### Definition

A maximum matching in an undirected graph G = (V, E) is a subset of its edges with maximum cardinality such that no two edges in the matching share any endpoint.

Remark. A maximal matching (i.e., one which cannot be further extended) can be obtained by a simple greedy algorithm. It is not necessarily a maximum matching!



# Maximum Matching: Edmond's Algorithm

```
#include <boost/graph/max_cardinality_matching.hpp>
typedef boost::adjacency_list<boost::vecS, boost::vecS, boost::undirectedS> graph;
typedef boost::graph_traits<graph>::vertex_descriptor
                                                                            vertex desc:
void maximum_matching(const graph &G) {
 int n = boost::num vertices(G):
 std::vector<vertex_desc> mate_map(n); // exterior property map
 const vertex_desc NULL_VERTEX = boost::graph_traits<graph>::null_vertex();
 boost::edmonds_maximum_cardinality_matching(G.
    boost::make_iterator_property_map(mate_map.begin(), boost::get(boost::vertex_index, G)));
  int matching_size = boost::matching_size(G,
    boost::make_iterator_property_map(mate_map.begin(), boost::get(boost::vertex_index, G)));
 for (int i = 0; i < n; ++i) {
   // mate_map[i] != NULL_VERTEX: the vertex is matched
   // i < mate_map[i]: visit each edge in the matching only once</pre>
    if (mate_map[i] != NULL_VERTEX && i < mate_map[i]) std::cout << i << " " << mate_map[i] << "\n";
Time complexity of boost::edmonds_maximum_cardinality_matching is O(mn \cdot \alpha(m, n)).
```

### BGL Outlook - More Algorithms Available!

To be covered in the remainder of this tutorial:

- breadth-first search (breadth\_first\_search)
- strongly connected components in directed graphs (strong\_components)

To be covered in upcoming tutorials:

- maximum-flow algorithms
- minimum-cost maximum-flow algorithms

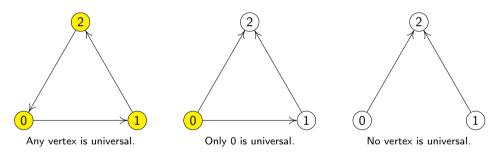
#### Exercise:

- connected components in undirected graphs (connected\_components)
- check if a graph is bipartite (is\_bipartite)
- minimum spanning tree (kruskal\_minimum\_spanning\_tree, prim\_minimum\_spanning\_tree)

#### Tutorial Problem: Universal Vertex

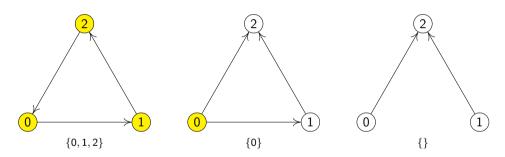
#### Definition

Let G = (V, E) be a directed graph. We call a vertex  $u \in V$  universal if every other vertex  $v \in V$  can be reached from u via a directed path.



#### Tutorial Problem: The Problem Statement

**Input** A directed unweighted graph G = (V, E). **Output** All universal vertices in G.



## Tutorial Problem: Checking if a Vertex is Universal

```
#include <boost/graph/breadth_first_search.hpp>
#include <boost/graph/properties.hpp>
typedef boost::adjacency_list<boost::vecS, boost::vecS, boost::directedS> graph;
typedef boost::default_color_type
                                                                          color:
const color black = boost::color_traits<color>::black(); // visited by BFS
const color white = boost::color_traits<color>::white(); // not visited by BFS
bool is_universal(const graph &G, int u) { // Is u universal in G?
 int n = boost::num vertices(G):
 std::vector<color> vertex_color(n); // exterior property map
 boost::breadth_first_search(G, u,
    boost::color_map(boost::make_iterator_property_map(
      vertex_color.begin(), boost::get(boost::vertex_index, G))));
 // u is universal iff no vertex is white (i.e., all vertices are reachable from u)
 return (std::find(vertex_color.begin(), vertex_color.end(), white) == vertex_color.end());
```

Time complexity of boost::breadth\_first\_search is O(n+m).

#### Tutorial Problem: A Slow Solution

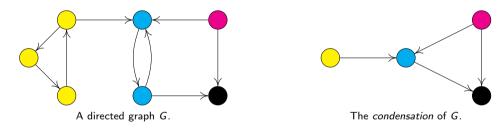
```
void testcase() {
  int n. m:
  std::cin >> n >> m;
  graph G(n);
  for (int i = 0; i < m; ++i) {</pre>
    int u, v;
    std::cin >> u >> v;
    boost::add_edge(u, v, G);
  for (int i = 0: i < n: ++i) {
    if (is_universal(G, i)) std::cout << i << " ";</pre>
  std::cout << "\n";
```

Time complexity of testcase is  $O(n^2 + nm)$ .

### Tutorial Problem: Strongly Connected Components

#### **Definition**

A strongly connected component (SCC) of a directed graph G = (V, E) is any maximal subset of vertices  $C \subseteq V$  such that all vertices in C are pairwise reachable (via directed paths).



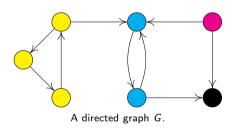
#### Lemma

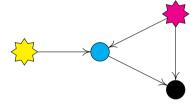
The condensation of a directed graph is acyclic.

#### Tutorial Problem: Source SCC

#### **Definition**

A strongly connected component is called a *source* if it has no incoming edges from other SCCs.





The condensation of G.

#### Lemma

Any non-empty directed graph has a source SCC.

#### Lemma

If a directed graph has a single source SCC, then it is a universal vertex in its condensation.

# Tutorial Problem: Full Solution (I)

```
#include <boost/graph/adjacency_list.hpp>
#include <boost/graph/strong_components.hpp>
typedef boost::adjacency_list<boost::vecS, boost::vecS, boost::directedS> graph;
typedef boost::graph_traits<graph>::edge_iterator
                                                                           edge_it;
void testcase() {
  int n. m:
  std::cin >> n >> m;
  graph G(n);
  for (int i = 0: i < m: ++i) {
    int u. v:
    std::cin >> u >> v:
    boost::add_edge(u, v, G);
```

# Tutorial Problem: Full Solution (II) — Strongly Connected Components

```
// scc_map[i]: index of SCC containing i-th vertex
std::vector<int> scc_map(n); // exterior property map
// nscc: total number of SCCs
int nscc = boost::strong_components(G,
   boost::make_iterator_property_map(scc_map.begin(), boost::get(boost::vertex_index, G)));
```

Time complexity of boost::strong\_components is O(n+m).

# Tutorial Problem: Full Solution (III) – Source SCCs

```
// is_src[i]: is i-th SCC a source?
std::vector<bool> is_src(nscc, true);
edge_it ebeg, eend;

for (boost::tie(ebeg, eend) = boost::edges(G); ebeg != eend; ++ebeg) {
   int u = boost::source(*ebeg, G), v = boost::target(*ebeg, G);
   // edge (u, v) in G implies that component scc_map[v] is not a source
   if (scc_map[u] != scc_map[v]) is_src[scc_map[v]] = false;
}
```

## Tutorial Problem: Full Solution (IV) – Finding All Universal Vertices

```
int src_count = std::count(is_src.begin(), is_src.end(), true);
if (src_count > 1) { // no universal vertex among multiple SCCs
    std::cout << "\n";
    return;
}
assert(src_count == 1); // recall property of the condensation DAG (directed acyclic graph)
// all vertices in the single source SCC are universal
for (int v = 0; v < n; ++v) {
    if (is_src[scc_map[v]]) std::cout << v << " ";
}
std::cout << "\n";
} /* end of function testcase */</pre>
```

Time complexity of testcase is O(n+m).

# Tutorial Problem: Full Solution (V) – Main Function

```
int main()
{
    int T;
    std::cin >> T;
    while(T--) testcase();
    return 0;
}
```