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Riding the Binlog: an in Deep Dissection of the Replication Stream

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- Based in Amsterdam since 1996
- Online Hotel and Accommodation Agent:
 - 170 offices worldwide
 - +795.000 properties in 221 countries
 - 42 languages (website and customer service)
- Part of the Priceline Group
- And we use MySQL:
 - Thousands (1000s) of servers, ~85% replicating
 - >110 masters: ~25 >50 slaves & ~8 >100 slaves

Riding the Binlog: Replication

- One master / one or more slaves
- The master records all writes in a journal: the binary logs
- Each slave:
 - Downloads the journal and saves it locally (IO thread): relay logs
 - Executes the relay logs on the local database (SQL thread)
 - Could produce binary logs to be itself a master (log-slave-updates)
- Understood replication features:
 - Asynchronous
 - Single threaded

Riding the Binlog: Replication'

Riding the Binlog: Replication"

- More understood binlog/replication features:
 - Master side binary log filtering
 - Binlog_Do_DB / Binlog_Ignore_DB (avoid those in favor of slave side filtering)
 - sql_log_bin
 - Slave side filtering:
 - Replicate_Do_DB / Replicate_Ignore_DB
 - Replicate_Do_Table / Replicate_Ignore_Table
 - Replicate_Wild_Do_Table / Replicate_Wild_Ignore_Table
 - Circular replication
- New innovations in replication land:
 - Global Transaction Identifiers (GTIDs)
 - Multi-Source and Multi-Threaded Replication (many types)
 - Binlog Servers

Riding: the Binary Logs (on master)

```
> ls -1 # Dots (.) in the file sizes are added for readability.
[...]
-rw-rw---- 1 mysql mysql 537.362.886 Sep 13 10:19 binlog.000878
-rw-rw---- 1 mysql mysql 537.341.120 Sep 13 10:22 binlog.000879
-rw-rw---- 1 mysql mysql 6431.760.382 Sep 13 11:13 binlog.000880
-rw-rw---- 1 mysql mysql
                         537.394.773 Sep 13 11:26 binlog.000881
                         537.288.989 Sep 13 11:29 binlog.000882
-rw-rw---- 1 mysql mysql
-rw-rw---- 1 mysql mysql
                         111.419.262 Sep 13 11:30 binlog.000883
-rw-rw---- 1 mysql mysql
                         10.411.193 Sep 13 11:30 binlog.000884
-rw-rw---- 1 mysql mysql
                          537.271.616 Sep 13 11:34 binlog.000885
-rw-rw---- 1 mysql mysql
                          537.290.232 Sep 13 11:38 binlog.000886
[...]
```

• Let's use mysqlbinlog to see what is in those files.

```
mysql plams <<< "CREATE TABLE test1 (\
 my pk BIGINT AUTO INCREMENT, \
my bigint BIGINT DEFAULT NULL, \
 my datetime DATETIME DEFAULT NULL, \
 PRIMARY KEY(my pk));"
# at 217
#150919 12:04:27 server id 1 end log pos 428 CRC32 0xebef2e40\
Query thread id=3100 exec time=0 error code=0
use `plams`
SET TIMESTAMP=1442657067
CREATE TABLE test1 ( my pk BIGINT AUTO INCREMENT, my bigint
BIGINT DEFAULT NULL, my datetime DATETIME DEFAULT NULL,
PRIMARY KEY (my pk))
```

```
mysql plams <<< "INSERT INTO test1(my pk) VALUES (1);"
# (Omitting "SET TIMESTAMP=...".)
# at 428
#150919 12:04:30 server id 1 end log pos 509 ... Query ...
BEGIN
# at 509
#150919 12:04:30 server id 1 end log_pos 620 ... Query ...
INSERT INTO test1 (my pk) VALUES (1)
# at 620
#150919 12:04:30 server id 1 end log_pos 651 ... Xid = ...
COMMIT
```

```
# PK (auto increment) not specified.
mysql plams <<< "INSERT INTO test1 (my_bigint, my_datetime) \</pre>
VALUES (floor(1024*1024*rand()), now());"
# (Omitting BEGIN/COMMIT.)
# at 740 / # at 772 / # at 811
#150919 12:04:38 server id 1 end log pos 772 ... Intvar
SET INSERT ID=2
#150919 12:04:38 server id 1 end log pos 811 ... Rand
SET @@RAND SEED1=80977572, @@RAND SEED2=154388004
#150919 12:04:38 server id 1 end log pos 976 ... Query ...
SET TIMESTAMP=1442657078
INSERT INTO test1 (my bigint, my datetime) VALUES
(floor(1024*1024*rand()), now())
                                                Booking.com
```

```
mysql plams <<< "
BEGIN;
INSERT INTO test1 (my datetime) VALUES (now());
SELECT sleep(5);
INSERT INTO test1 (my datetime) VALUES (now());
COMMIT; " &
mysql plams <<< "
SELECT sleep(1);
INSERT INTO test1(my datetime) VALUES (now());"
```

```
# (Omitting BEGIN/COMMIT/"SET TIMESTAMP=...".)
# at 1096 / # at 1128 / SET INSERT_ID=4
#150919 12:05:00 server id 1 end log pos 1257 ... Query ...
INSERT INTO test1(my datetime) VALUES (now())
# at 1377 / # at 1409 / SET INSERT ID=3
#150919 12:04:58 server id 1 end log pos 1538 ... Query ...
INSERT INTO test1(my datetime) VALUES (now())
# at 1538 / # at 1570 / SET INSERT ID=5
#150919 12:05:03 server id 1 end log pos 1699 ... Query ...
INSERT INTO test1(my_datetime) VALUES (now())
```

```
mysql plams <<< "SET binlog_format = ROW;
INSERT INTO test1(my_datetime) VALUES (now());"

# (Omitting BEGIN/COMMIT.)

# at 1811
#... Table_map: `plams`.`test1` mapped to number 143
# at 1863
#... Write_rows: table id 143 flags: STMT_END_F</pre>
```

BINLOG '

```
# mysqlbinlog -v
# (Omitting BEGIN/COMMIT.)
# at 1811
#... Table_map: `plams`.`test1` mapped to number 143
# at. 1863
#... Write rows: table id 143 flags: STMT END F
BINLOG '
VTP9VRMBAAAANAAAEcHAAAAAI8AAAAAAEABXBsYW1zAAV0ZXN0MQADCAqSAQAGVr3KSw==
### INSERT INTO `plams`.`test1`
### SET
### @1=6
### @2=NULL
### @3='2015-09-19 12:05:09'
```

```
# mysqlbinlog -vv
# (Omitting BEGIN/COMMIT.)
# at 1811
#... Table map: `plams`.`test1` mapped to number 143
# at 1863
#... Write rows: table id 143 flags: STMT END F
BINLOG '
VTP9VRMBAAAANAAAEcHAAAAAI8AAAAAAAAABABXBsYW1zAAV0ZXN0MQADCAqSAQAGVr3KSw==
### INSERT INTO `plams`.`test1`
### SET
### @1=6 /* LONGINT meta=0 nullable=0 is null=0 */
### @2=NULL /* LONGINT meta=0 nullable=1 is null=1 */
### @3='2015-09-19 12:05:09' /* DATETIME(0) meta=0 nullable=1 is null=0 */
```

- Binary Logs: Journal for MySQL Replication
 - Files with a basename and 6 digit index
 - Stream of transactions in commit order
 - A transaction is never split in 2 binlog files
 - But a transaction is split in many events
- Want to know more:
 - Read the documentation
 - Or source dive: sql/log_event.{h,cc}

Riding: some Event Types

```
sql/log event.h
enum Log event type {
  UNKNOWN EVENT= 0,
 QUERY EVENT= 2,
  STOP EVENT= 3,
 ROTATE EVENT= 4,
  INTVAR EVENT= 5,
 LOAD EVENT= 6,
  CREATE FILE EVENT= 8,
 APPEND BLOCK_EVENT= 9,
 EXEC LOAD EVENT= 10,
  DELETE FILE EVENT= 11,
 NEW LOAD EVENT= 12,
 RAND EVENT= 13,
 USER VAR EVENT= 14,
```

```
FORMAT DESCRIPTION EVENT= 15,
XID EVENT= 16,
TABLE MAP EVENT = 19,
INCIDENT EVENT= 26,
HEARTBEAT LOG EVENT= 27,
IGNORABLE LOG EVENT= 28,
ROWS QUERY LOG EVENT= 29,
WRITE ROWS EVENT = 30,
UPDATE ROWS EVENT = 31,
DELETE ROWS EVENT = 32,
GTID LOG EVENT= 33,
ANONYMOUS GTID LOG EVENT= 34,
PREVIOUS GTIDS LOG EVENT= 35,
```

Riding: the Binlogs on Slaves

- Depends on log-slave-updates (lsu):
 - If Isu disabled: no trx from the master in the binlogs of the slave
 - else: SQL Thread logs its transactions
- binlogs on a slave != binlogs on the master:
 - FLUSH BINARY LOGS
 - Restart of mysqld on the master or slave
 - RESET MASTER on slave
 - Different binlog filenames on master and slave
 - ...

Riding: the Binlogs on Slaves'

Binlogs on the master:

```
-rw-rw---- 1 mysql mysql 40720631 Sep 19 13:16 binlog.000001
-rw-rw---- 1 mysql mysql 16777476 Sep 19 13:17 binlog.000002
-rw-rw---- 1 mysql mysql 16777476 Sep 19 13:18 binlog.000003
-rw-rw---- 1 mysql mysql 10124412 Sep 19 13:18 binlog.000004
-rw-rw---- 1 mysql mysql 8388820 Sep 19 13:18 binlog.000005
-rw-rw---- 1 mysql mysql 8388820 Sep 19 13:19 binlog.000006
-rw-rw---- 1 mysql mysql 6964108 Sep 19 13:19 binlog.000007
```

And on a slave:

```
-rw-rw---- 1 mysql mysql 67109007 Sep 19 13:21 binlog.000001
-rw-rw---- 1 mysql mysql 67108948 Sep 19 13:22 binlog.000002
-rw-rw---- 1 mysql mysql 52518911 Sep 19 13:22 binlog.000003
-rw-rw---- 1 mysql mysql 45279097 Sep 19 13:30 binlog.000004
```

The content of the binary logs on slaves is (hopefully) the same as on the master, but the transport is different.

Riding: the Binlogs on Slaves"

- Depends on slave-local transactions (and Isu):
 - If Isu disabled: the slave behaves like a master
 - The slave has its own transaction/binlog stream
 - This stream is independent from the master stream (Isu disabled)
 - else: local trx UNION master trx (multiplexing)
 - The slave transaction stream merges with the master stream
 - Those might be independent, or might depend on each other

Riding: the Binlogs on Slaves"

- If Isu enabled, also depends on replication filters:
 - If no filter: binlogs on slave contain local and master trx
 - else: depends of MySQL 5.6 GTIDs:
 - Without 5.6 GTIDs, filtered transaction are skipped (exclusion)
 - With 5.6 GTIDs: an empty transaction is logged (mutation)

(Transactions can be filtered fully or partially : also mutation)

Riding: the Binlogs on Slaves" '

- Circular-replication (needs log-slave-updates):
 - Combination of slave local transactions and filters
 - local trx UNION (master trx SUB "same server_id" trx) (multiplexing with exclusion)

Riding: the Binlogs on Slaves" "

- Depends on Multi-Source Replication:
 - If Isu disabled: only slave local transactions
 - else: local trx UNION M1 trx UNION M2 trx ... (multiplexing)
 (Identically configured slave might multiplex in different ways)
 (Filters can be added to multi-source → mutation)

Riding: the Binlogs on Slaves" "

- Depends on Multi-Threaded Replication (MySQL 5.6 schema based):
 - If Isu disabled: only slave local transactions
 - else: reordering of transactions (morphing)
- On the master, some transactions in 3 schemas (A, B, C):
 - Order in the Binary Logs: A1, A2, B1, B2, C1, A3, C2, B3, C3
- On the slave, transactions in different schema run in //:
 "A1, A2, A3" in // with "B1, B2, B3" in // with "C1, C2, C3"
 - One possible order: A1, B1, C1, A2, B2, C2, A3, B3, C3
 - Another: A1, C1, A2, C2, A3, C3, B1, B2, B3 (if B1 is big)
 - Many others...

Riding: the Binlogs on Slaves" "

- Depends on MySQL 5.7 Logical Clock Parallel Replication:
 - If Isu disabled: only slave local transactions
 - else if slave_preserve_commit_order = OFF (default): reordering of transactions (morphing)
 - else (slave_preserve_commit_order = ON)
 modification of parallelism information (still morphing)

Riding: the Binlogs on Slaves" ""

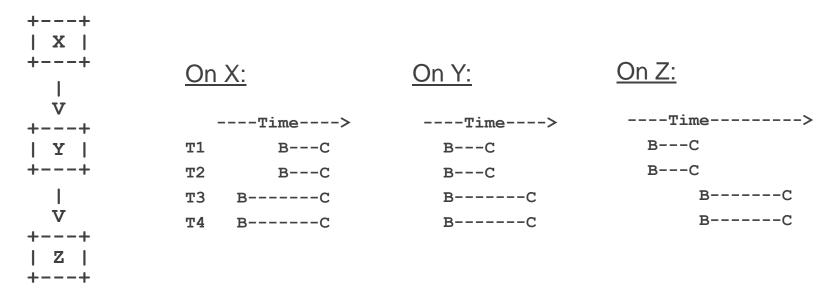
Four transactions on X and Y:

```
On Y
| X |
                                         (with slave_preserve_commit_order):
          On X:
                          On Y:
            ----Time--->
                         ----Time--->
                                        ----Time--->
                                         В----С
          T1 B----C
                          В----С
                                         B----. C
                           B----C
          T2 B----C
| Y |
                         B---C
                                        B---. . C
          T3 B---C
                                         B---- . C
          T4 B----C
                           В----С
+---+
```

- Transaction order in the binary logs of Y: (without slave_preserve_commit_order)
 - T3, T4, T2, T1

Riding: the Binlogs on Slaves" "" "

Four other transactions on X and Y:



- IM (Y) might stall the parallel replication pipeline
- Replicating with IM will slow down parallel replication

Riding: the Binlogs on Slaves" "" "

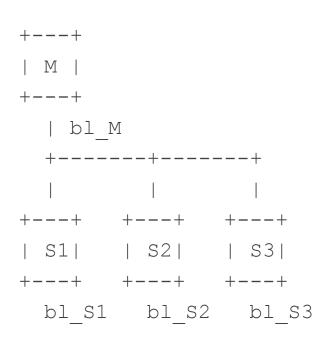
- Let's try to summarise:
 - Multiplexing:
 - Local transactions + log-slave-updates
 - Circular Replication (needs log-slave-updates)
 - Multi-Source + log-slave-updates
 - Morphing:
 - All types of Multi-Threaded replication + log-slave-updates
 - Mutation:
 - Filters + MySQL 5.6 GTIDs + log-slave-updates
- See the pattern...
 - log-slave-updates

Riding the Binlog: why Isu?

- Lsu is needed for:
 - master promotion with GTIDs
 - fan-out
 - circular replication
 - external trigger
 - converting SBR to RBR
 - new version/feature testing
 - multiplexing with a side effect

Riding the Binlog: Isu with GTIDs

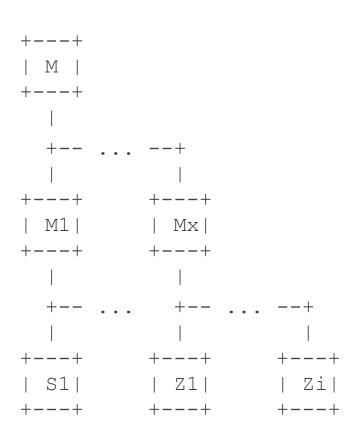
 After the failure of M, new master promotion needs levelling the remaining slaves



- If S2 is the most up to date slave:
 - bl_S2 will be used to level S1 and S3
- Ok if bl_S2 is a different transport of bl_M
- Ok-ish if bl_S2 is a simple morphing of bl_M
- NOT-Ok if bl_S2 is a mutation of bl_M
- If you must replicate using Isu put Intermediate Master between the master and mutation (filtering) slave (or you might loose those slave during promotion)
- Or use Binlog Servers

Riding the Binlog: Isu for Fan-Out

Remote site or too many slaves: needs a fan-out solution



- Intermediate Masters slow down replication (including parallel replication)
- Their failure needs to be managed
- Rogue transactions need to be managed
- And we still have the risks/problems associated with morphing and mutations
- IMs come with many drawbacks
- Avoiding them might be better
- Binlog Servers allow that

Riding the Binlog: Circular Replication

Circular replication is a form of multiplexing

- But Multi-Source replication is also multiplexing
- Replace Circular replication by Multi-Source
 - High Availability is tricky
 - But it was not simple either with Circular Replication
 - And it is very simple with <u>Binlog Servers</u>

Riding the Binlog: Circular Replication'

Replace the following:

```
|S11| <---+ +---+ +---> |S21|
+---+ +- | M1 | ----> | M2 | -+ +---+
 +---+ | +-+-+ | +---+
 +---+ | | +---+ | | +---+
+---+ | +-- | M3 | <-+ | +---+
|S13| <--+ +-+-+ +--> |S23|
+---+
         /----\
       +---+ V +---+
       |S31| +---+ |S33|
       +---+ IS32| +---+
```

Riding the Binlog: Circular Replication"

By this:

```
+---+ +---+
     Δ~~ Ι~~ Ι~~
+----|+|----+
   |+----|+----|+
| | +---- | | +----+-----| | +
VVV VVV
                VVV
                   VVV
+---+
| S1| | S2|
               |...| | Sn|
```

Riding the Binlog: Binlog Server

- Binlog Server (BLS): is a daemon that:
 - Downloads the binary logs from the master
 - Saves them identically as on the master
 - Serves them to slaves

- A or X are the same for B and C:
 - By design, the binary logs served by A and X are the same

Riding the Binlog: BLS – Fan-out

Replace Intermediate Master by Binlog Servers:

• If BLS fails, repoint slaves to other BLSs (easy by design)

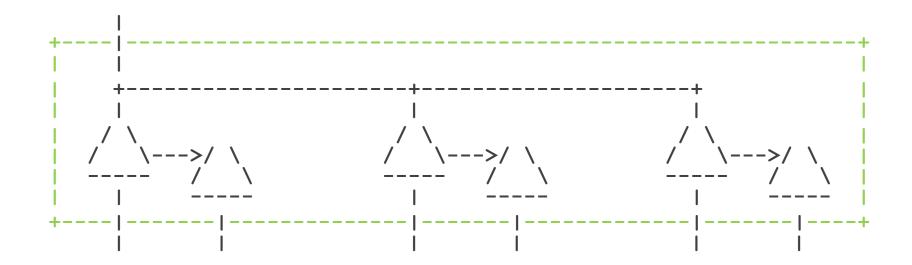
Riding the Binlog: BLS – HA

Distributed Binlog Serving Service (DBSS):

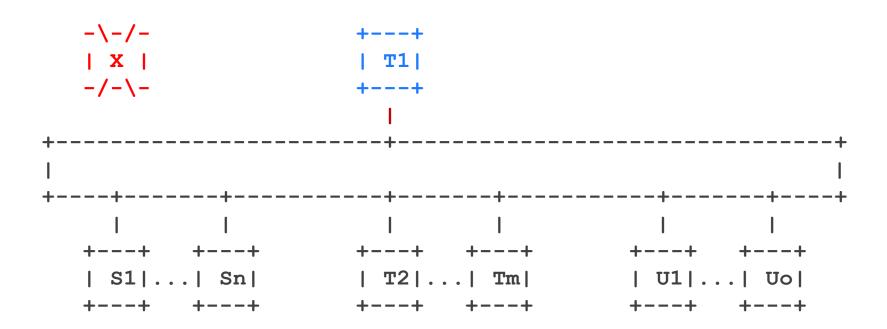
```
+---+
| M |
+---+
|
|
+---+
|
|
+---+
|
|
+---+
|
| | | | | |
|
+---+ +---+ +---+ +---+ +---+
| S1|...| Sn| | T1|...| Tm| | U1|...| Uo|
+---+ +---+
```

Riding the Binlog: BLS – HA'

Zoom in DBSS:



Riding the Binlog: BLS – HA"

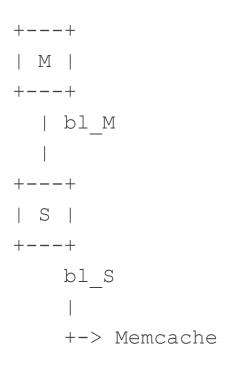


Riding the Binlog: why Isu?'

- Lsu is needed for:
 - master promotion with GTIDs: use with care or Binlog Servers
 - fan-out: Binlog Servers
 - circular replication: replaced by Multi-Source + Binlog Servers
 - external trigger
 - converting SBR to RBR
 - new version/feature testing
 - multiplexing with a side effect

Riding the Binlog: external trigger

 If we need to trigger an event after commit on a slave, log-slave-updates is a good solution



- An example is cache invalidation
- But morphing and mutation is still there
- Remember above when tempted to use the slave binlogs for other use-cases (HA)
- Cache invalidation does not need full binlogs
- A custom binary log format dedicated to that might be useful

Riding the Binlog: SBR to RBR

- When a master is SBR and you need Row-Based Events
 - log-slave-updates is currently your only solution
- But consuming Row-Based Events on a slave introduces morphing, mutation and a single point of failure
- Also, migrating from SBR to RBR is complex
- Wouldn't it be great if:
 - the master logs in both SBR and RBR format,
 - And each slave chooses to follow the SBR or RBR stream.
- I suggest a new binary log format: multiplexed
 - (with the corresponding configuration on slaves)

Riding the Binlog: why Isu?"

- Lsu is needed for:
 - master promotion with GTIDs: use with care or Binlog Servers
 - fan-out: Binlog Server
 - circular replication: replaced by Multi-Source + Binlog Servers
 - external trigger (cache invalidation): OK
 - converting SBR to RBR: new binary log format
 - new version/feature testing
 - multiplexing with a side effect

Riding the Binlog: why Isu?"

- Lsu is needed for:
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 - new version/feature testing: ok, not really production
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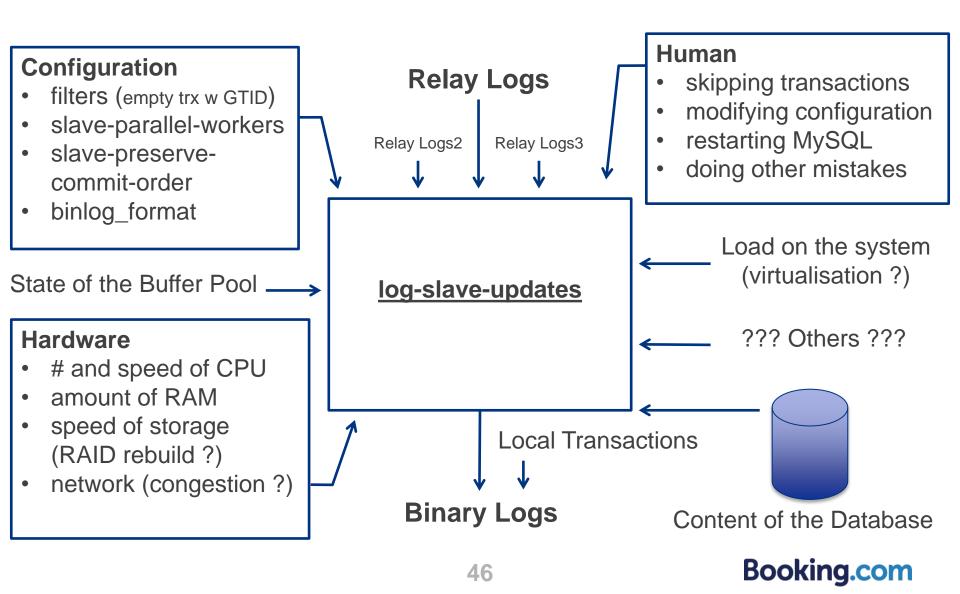
Riding the Binlog: why Isu?"

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 - fan-out: Binlog Server
 - circular replication: replaced by Multi-Source + Binlog Servers
 - external trigger (cache invalidation): OK
 - converting SBR to RBR: new binary log format
 - new version/feature testing: ok, not really production
 - multiplexing with a side effect:
 very complex use-case, not covered in the talk

Riding the Binlog: Conclusion

- The Binary logs are a stream of transactions
- Enabling log-slave-updates makes this stream evolves
- Those evolutions are complex
- Tips to manage those evolutions in combination with HA:
 - A filtering slave should replicate through an intermediate master
 - A slave with local transaction should also replicate through an IM
 - Avoid using log-slave-updates for replication (use Binlog Servers)
- Learn more tomorrow: Binlog Server at Booking.com
 - Wed. 23 September, 2:10PM 3:00PM @ Matterhorn 1
 - https://www.percona.com/live/europe-amsterdam-2015/sessions/binlog-servers-bookingcom

Riding: the Binlog Function on Slaves



Riding: the Binlog Function on Slaves

Without log-slave-updates

Local Transactions

Binary Logs

Riding the Binlog: Links

Binlog Servers:

- http://blog.booking.com/mysql_slave_scaling_and_more.html
- http://blog.booking.com/abstracting_binlog_servers_and_mysql_maste_r_promotion_wo_reconfiguring_slaves.html

IM Morphing Loosing Parallelism Information:

- http://blog.booking.com/better_parallel_replication_for_mysql.html
- (http://blog.booking.com/evaluating_mysql_parallel_replication_2-slave_group_commit.html)

Others

 https://www.percona.com/blog/2009/05/14/why-mysqls-binlog-do-dboption-is-dangerous/

Thanks

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