CS179F: Projects in Operating System

Memory mapped files

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Virtual Address Space

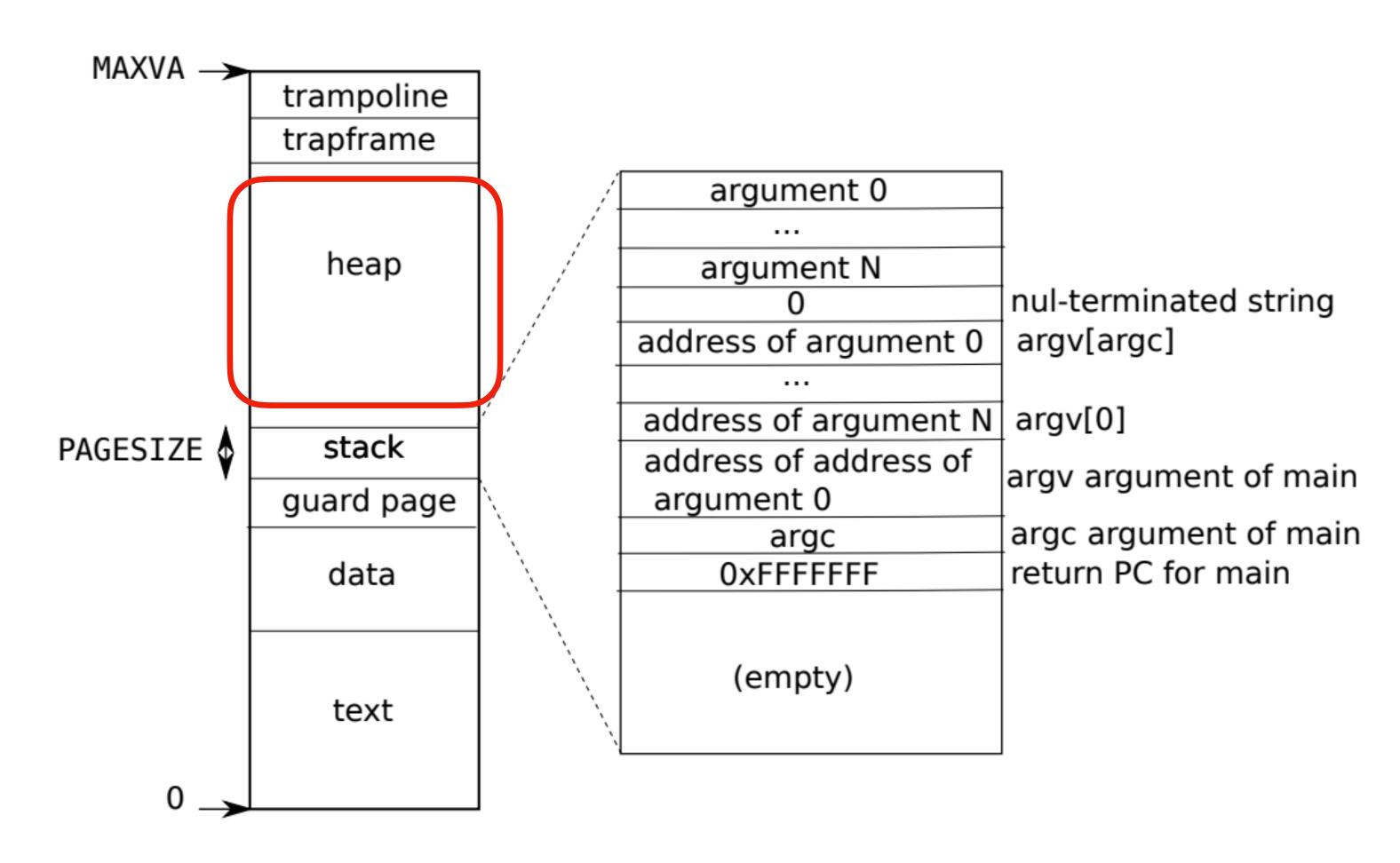


Figure 3.4: A process's user address space, with its initial stack.

Virtual Memory Allocation

- sbrk() a single memory region
- mmap(void *addr, size_t length, int prot, int flags, int fd, off t offset)
 - Multiple (segments) of memory regions
- munmap(void *addr, size t length)
 - Can unmap partially of a region

- Question: how to keep track of memory regions?
- VMA: virtual memory area
 - What metadata is necessary?
 - mmap(void *addr, size_t length, int prot, int flags, int fd, off t offset)

- Question: how to keep track of memory regions?
- VMA: virtual memory area
 - Q1: what metadata is necessary?
 - [start, end) [addr, addr + length)
 - permissions prot
 - behaviors flags
 - file object, offset fd, offset

- Q2: how to manage vmas (e.g., how to find a vma)?
 - Array (fixed size)
 - Linked list

- Q3: how to handle lazy allocation (when a page fault happens, how do we know it's due to lazy allocation)?
 - How to improve the performance?
 - Heap (sorted tree) red-black tree

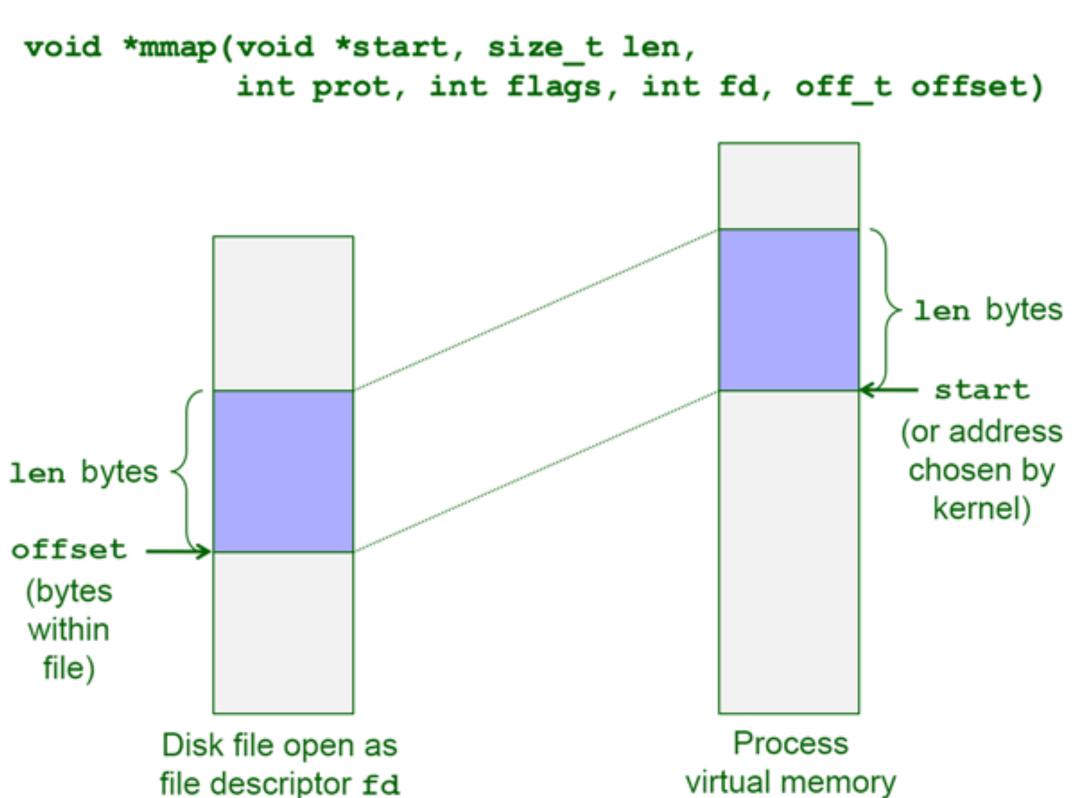
- Q4: how to allocate vma?
 - Find an unallocated virtual space region that is large enough
 - Recall heap allocation
- Q5: how to handle munmap?
 - Shrink, split, or free

- Q6: how to handle a memory mapped file?
 - SHARED vs PRIVATE
 - offset?
 - lazy allocation
 - buffer?

Step 0: Pull the code for Lab 5

- In your xv6-riscv code base, run
 - \$ git fetch
 - \$ git checkout mmap
- Make sure this checkout to this branch new branch, and then
 - \$ make clean
 - \$ make qemu
- Read lab instructions carefully
 - https://github.com/emidec/cs179f-fall23/blob/xv6-riscv-fall23/doc/lab5.md
 - You will use the knowledge/skills you learned in Lab2 and Lab4
- Important dates
 - 12/15/2023 1:59PM (NO EXTENSIONS)

mmap() and munmap() system call



Reference: https://www.clear.rice.edu/comp321/html/laboratories/lab10/

- mmap() maps files (or devices) into virtual memory of the calling process where start specify the starting address of the mapping and len specifies the length.
- In this lab, we assume that:
 - start will always be zero, meaning that the kernel will decide the address at which to map the file
 - prot indicates if the memory mapped should be readable and(or) writable, it can be PROT_WRITE, PROT_READ or both.
 - o flags can be
 - MAP SHARED: modifications should be written back
 - MAP PRIVATE: modifications should not be written back
 - offset always be 0. The mapping always starts from the beginning of the file.
- munmap(void* start, size_t len)
 - Remove the mapping. If **MAP_SHARED**, write back the modifications
 - You can assume that it will either unmap at the start, or at the end, or the whole region (but not punch a hole in the middle of a region).
- Return value:
 - On success, mmap returns the pointer to the mapped area, on error, it returns -1
 - On success, munmap returns 0, on error, returns -1

Your Goal

You should implement enough mmap and munmap functionality to make the mmaptest test program work.

If mmaptest doesn't use a mmap feature, you don't need to implement that feature.

Test File

```
AAAA...AAA AA...A 000...0

1 Page 1/2 Page 1/2 Page

99 makefile(f);
100 if ((fd = open(f, 0_RDONLY)) == -1)
101 err("open");
```

Test 1: Basic checks

```
char *p = mmap(0, PGSIZE*2, PROT_READ, MAP_PRIVATE, fd, 0);
if (p == MAP_FAILED)
   err("mmap (1)");
_v1(p);
if (munmap(p, PGSIZE*2) == -1)
   err("munmap (1)");
```

Test 2: Private writable mapping for a file opened read-only is allowed

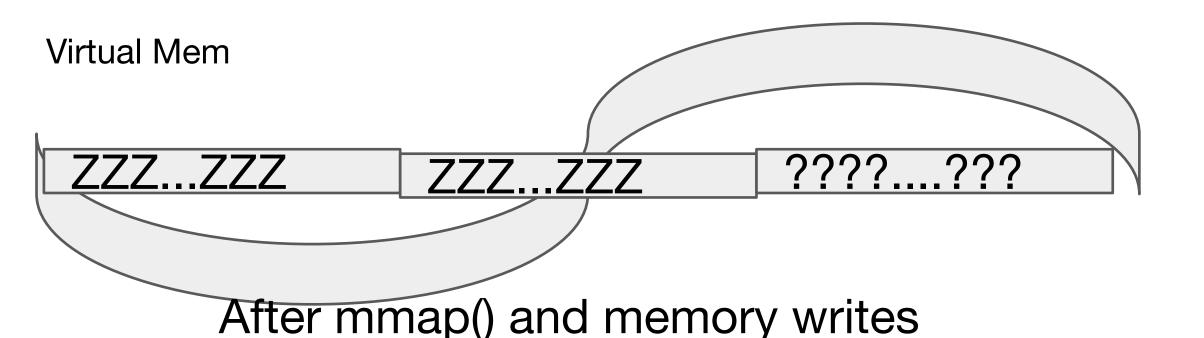
```
// should be able to map file opened read-only with private writable
// mapping
p = mmap(0, PGSIZE*2, PROT_READ | PROT_WRITE, MAP_PRIVATE, fd, 0);
if (p == MAP_FAILED)
    err("mmap (2)");
if (close(fd) == -1)
    err("close");
_v1(p);
for (i = 0; i < PGSIZE*2; i++)
    p[i] = 'Z';
if (munmap(p, PGSIZE*2) == -1)
    err("munmap (2)");</pre>
```

Test 3: Reject a shared readable/writable mapping of a file opened read-only

```
// check that mmap doesn't allow read/write mapping of a
// file opened read-only.
if ((fd = open(f, 0_RDONLY)) == -1)
   err("open");
p = mmap(0, PGSIZE*3, PROT_READ | PROT_WRITE, MAP_SHARED, fd, 0);
if (p != MAP_FAILED)
   err("mmap call should have failed");
if (close(fd) == -1)
   err("close");
```

Test File





Test File

ZZZZZZ	ZZZZZZ
1 Page	1 Page

After munmap()

Test 4: munmap

```
// check that mmap does allow read/write mapping of a
     // file opened read/write.
    if ((fd = open(f, 0_RDWR)) == -1)
       err("open");
     p = mmap(0, PGSIZE*3, PROT_READ | PROT_WRITE, MAP_SHARED, fd, 0);
     if (p == MAP_FAILED)
       err("mmap (3)");
     if (close(fd) == -1)
       err("close");
     // check that the mapping still works after close(fd).
158
     _v1(p);
     // write the mapped memory.
     for (i = 0; i < PGSIZE*2; i++)
       p[i] = 'Z';
     // unmap just the first two of three pages of mapped memory.
     if (munmap(p, PGSIZE*2) == -1)
       err("munmap (3)");
166
     // check that the writes to the mapped memory were
     // written to the file.
     if ((fd = open(f, 0_RDWR)) == -1)
       err("open");
     for (i = 0; i < PGSIZE + (PGSIZE/2); i++){
173
       char b;
       if (read(fd, &b, 1) != 1)
         err("read (1)");
       if (b != 'Z')
         err("file does not contain modifications");
177
     if (close(fd) == -1)
       err("close");
181
     // unmap the rest of the mapped memory.
     if (munmap(p+PGSIZE*2, PGSIZE) == -1)
       err("munmap (4)");
```

Test File 1

12345

Test File 2

67890

Virtual Mem

12345

67890

Test 5: map two files

```
// mmap two files at the same time.
     int fd1;
     if((fd1 = open("mmap1", O_RDWR|O_CREATE)) < 0)
       err("open mmap1");
     if(write(fd1, "12345", 5) != 5)
       err("write mmap1");
     char *p1 = mmap(0, PGSIZE, PROT_READ, MAP_PRIVATE, fd1, 0);
    if(p1 == MAP_FAILED)
       err("mmap mmap1");
     close(fd1);
     unlink("mmap1");
199
     int fd2;
     if((fd2 = open("mmap2", O_RDWR|O_CREATE)) < 0)
       err("open mmap2");
     if(write(fd2, "67890", 5) != 5)
       err("write mmap2");
     char *p2 = mmap(0, PGSIZE, PROT_READ, MAP_PRIVATE, fd2, 0);
    if(p2 == MAP_FAILED)
       err("mmap mmap2");
     close(fd2);
     unlink("mmap2");
210
     if(memcmp(p1, "12345", 5) != 0)
212
       err("mmap1 mismatch");
     if(memcmp(p2, "67890", 5) != 0)
214
       err("mmap2 mismatch");
     munmap(p1, PGSIZE);
     if(memcmp(p2, "67890", 5) != 0)
       err("mmap2 mismatch (2)");
     munmap(p2, PGSIZE);
220
     printf("mmap_test OK\n");
```

Parent's process Virtual Mem ????????????? ???????????? After mmap() Duplicate Child's process Virtual Mem **p1** ??????????? p2

After fork()

Test 6: forktest

```
225 // mmap a file, then fork.
226 // check that the child sees the mapped file.
227 //
228 void
229 fork_test(void)
     int fd;
     int pid;
     const char * const f = "mmap.dur";
     printf("fork_test starting\n");
     testname = "fork_test";
     // mmap the file twice.
     makefile(f);
     if ((fd = open(f, O_RDONLY)) == -1)
       err("open");
     unlink(f);
     char *p1 = mmap(0, PGSIZE*2, PROT_READ, MAP_SHARED, fd, 0);
     if (p1 == MAP_FAILED)
       err("mmap (4)");
     char *p2 = mmap(0, PGSIZE*2, PROT_READ, MAP_SHARED, fd, 0);
     if (p2 == MAP_FAILED)
       err("mmap (5)");
      // read just 2nd page.
     if(*(p1+PGSIZE) != 'A')
       err("fork mismatch (1)");
     if((pid = fork()) < 0)
       err("fork");
     if (pid == 0) {
       _v1(p1);
       munmap(p1, PGSIZE); // just the first page
       exit(0); // tell the parent that the mapping looks OK.
      int status = -1;
     wait(&status);
     if(status != 0){
       printf("fork_test failed\n");
       exit(1);
     // check that the parent's mappings are still there.
      _v1(p2);
     printf("fork_test OK\n");
```

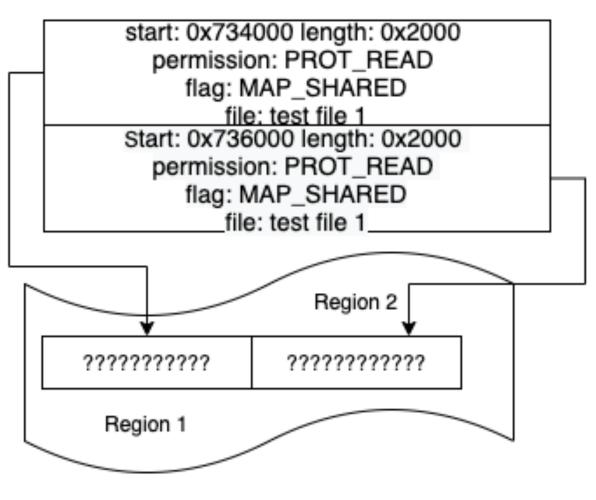
Design (1/2)

- You must use lazy allocation (lab 2). That is, mmap should not allocate physical memory or read the file. Instead, do that in page fault handling code in (or called by) usertrap. Reasons:
 1) performance 2) mapping large file (larger than physical memory)
- Keep track of what mmap has mapped for each process. Define a structure (referred as VMA in the following discussions) recording address, length, permission, file, etc. for a virtual memory region created
- In mmap(), find an unused region in the process's address space in which to map the file, and add a VMA to the process's table of mapped regions. The VMA should contain a pointer to a struct file for the file being mapped; mmap should increase the file's reference count so that the structure doesn't disappear when the file is closed

Design (2/2)

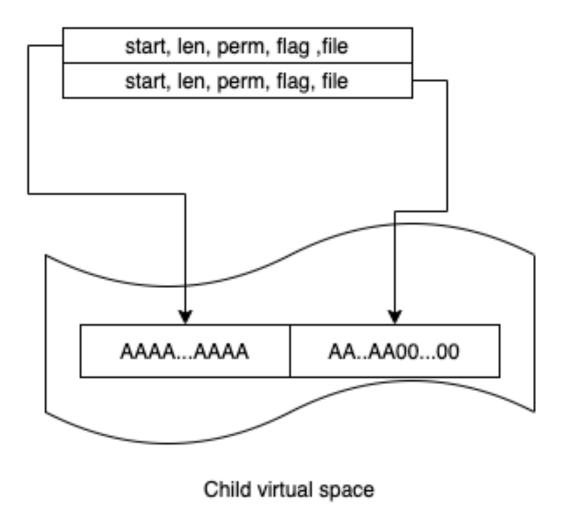
- In usertrap(), read 1 page of the relevant file into that page, and map it into the user address space
- Implement munmap (): find the VMA for the address range and unmap the specified pages. If munmap removes all pages of a previous mmap, it should decrement the reference count of the corresponding struct file. If an unmapped page has been modified and the file is mapped MAP_SHARED, write the page back to the file.
- Modify exit() to unmap the process's mapped regions as if munmap had been called
- Modify fork() to ensure that the child has the same mapped regions as the parent. Don't
 forget to increment the reference count for a VMA's struct file

Running example based on fork test()

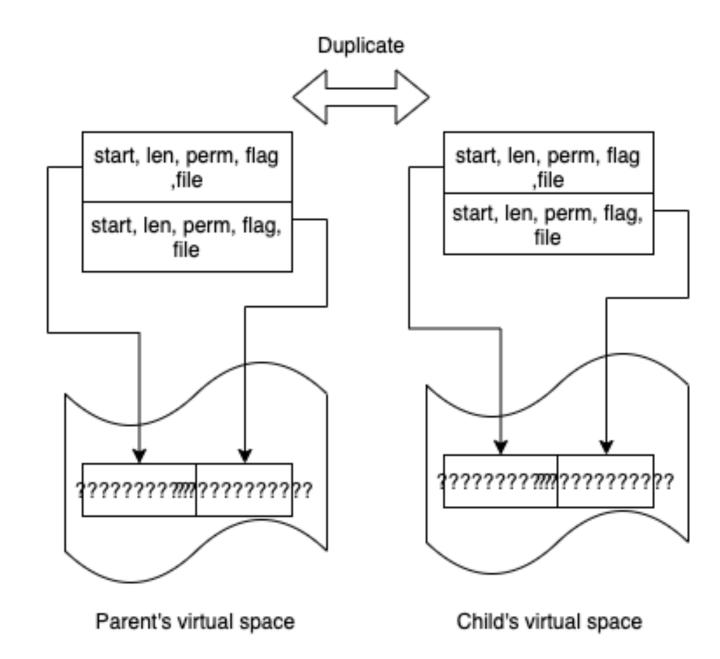


Parent's virtual space

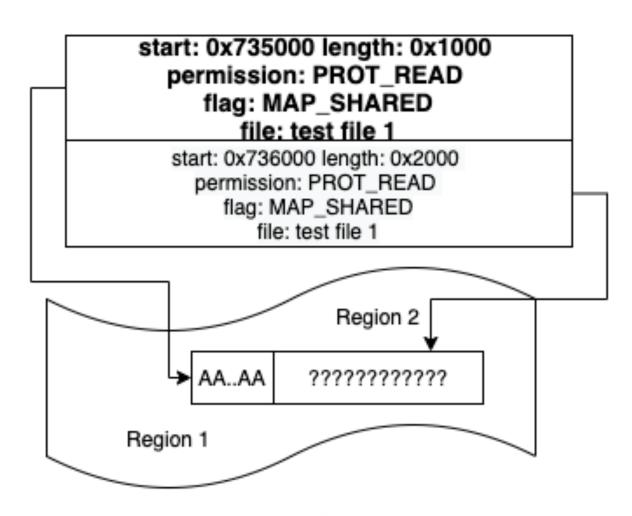
After call mmap()



After usertrap()



After call fork()



Child's virtual space

After munmap() for the first page

Implementation step 0: Add flags/definitions for the syscalls

• Add \$U mmaptest in Makefile to add mmaptest test. Add entry ("mmap"); entry ("munmap"); in user/usys.pl Add the following function definitions in user/user.h void* mmap(void*, unsigned int, int, int, int, unsigned int); int munmap(void*, unsigned int); • Add the following definitions in kernel/syscall.h #define SYS mmap #define SYS munmap 24 • Add the following mapping in kernel/syscall.c [SYS mmap] sys_mmap,

[SYS munmap] sys_munmap,

Implementation step 0: Add flags/definitions for the syscalls

Add the following definitions in kernel/syscall.c

```
o extern uint64 sys_mmap(void);
o extern uint64 sys_munmap(void);

• Add the following definitions in kernel/fcntl.h
o #define PROT_READ 0x000
o #define PROT_WRITE 0x001
o #define MAP_PRIVATE 0x000
o #define MAP_SHARED 0x001

• Your implementations can go to kernel/sysfile.c
o uint64_t mmap() { // your implementation }
```

Uint64 t munmap() { // your implementation }

Implementation: defining and allocating structure vma

- Define your **VMA** structure (recording start, length, permission, flags and file for each mapped memory range). You can define it in kernel/proc.h
- Add a table in struct proc of all the VMAs for a process, for example

```
o struct vma* vma_table[4];  //vma regions
```

- Since the xv6 kernel doesn't have a memory allocator in the kernel, it's OK to declare a fixed-size array of VMAs and allocate from that array as needed. A size of 16 should be sufficient
 - Hint: look at how file structure is allocated in kernel/file.h

Implementation: mmap()

- Check and respond to different permissions, flags.
- Allocate a vma
- Update the vma with addr, length, file, permissions, etc.
- Add the vma to the process's vma table
- Call filedup to increase the reference count of the file

Implementation: usertrap()

- Given a faulty address, check if it belongs to a VMA region, and find the VMA region if it does
- Deciding the offset in the file at which the data copy should start by looking at the faulty address and the VMA's start address
- Allocate a physical page
- Call readi () to transfer the data from file to the physical page
- Map the physical map to user's virtual memory space by calling mappagers ()

Implementation: munmap() exit() and fork()

- Find the VMA for the address range and unmap the specified pages (hint: use uvmunmap()).
- If all pages are removed, decrease the reference count of the corresponding struct file (hint: look at fileclose)
- If an unmapped page has been modified and the file is mapped MAP_SHARED, write the
 page back to the file (hint: look at filewrite())
- Modify exit(), which removes all mapped memory regions for a process
- Modify fork() to ensure that the child has the same mapped regions as the parent. Don't forget to increment the reference count for a VMA's struct file
- You will probably need to modify uvmcopy() and uvmunmap() in a way similar to what you've done in Lab2.

Self-Assessment

```
$ make qemu-gdb
(3.6s)
  mmaptest: mmap_test: OK
  mmaptest: fork_test: OK
usertests:
$ make qemu-gdb
OK (31.7s)
Score: 100/100
```

- You can run each test individually
 - \$ make qemu
 - \$ mmaptest
 - \$ usertests
- Or use
 - \$ make grade