

Multiple Antenna Communications

Lecture 4:

Capacity of Point-to-Point MIMO Channels

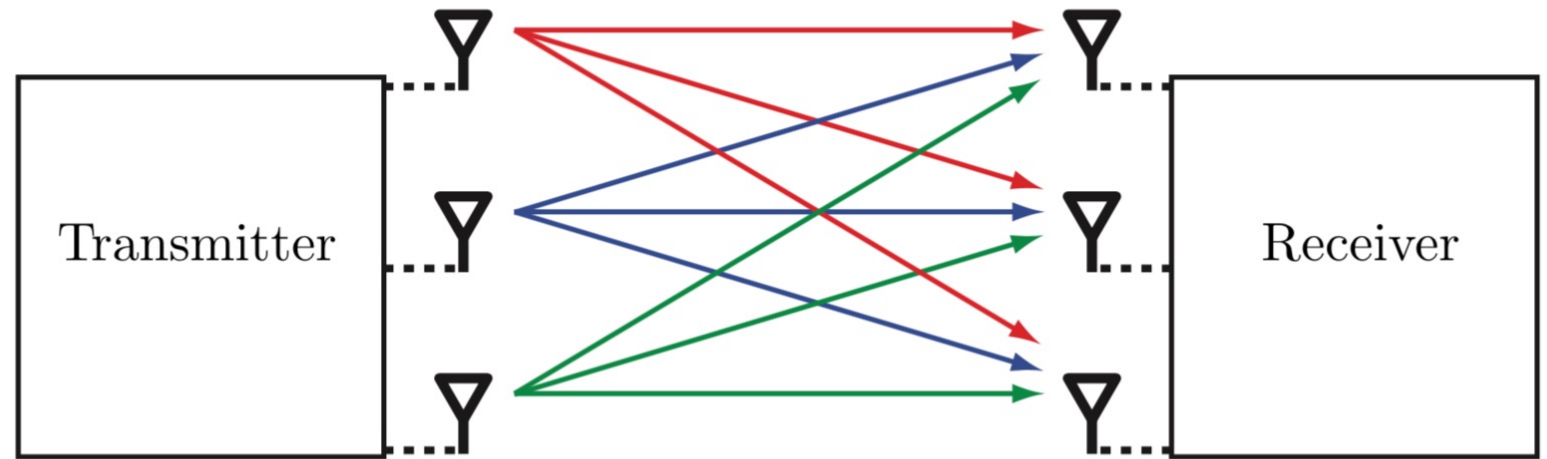
Emil Björnson

Outline

- Point-to-point MIMO channels
 - Basic formulation
 - Some linear algebra results
 - Capacity of MIMO channels
 - Examples of capacity behavior
- Transmitter diversity
 - MISO channels with slow fading

Point-to-point MIMO channel

- K transmit antennas
- M receive antennas
- Deterministic channel



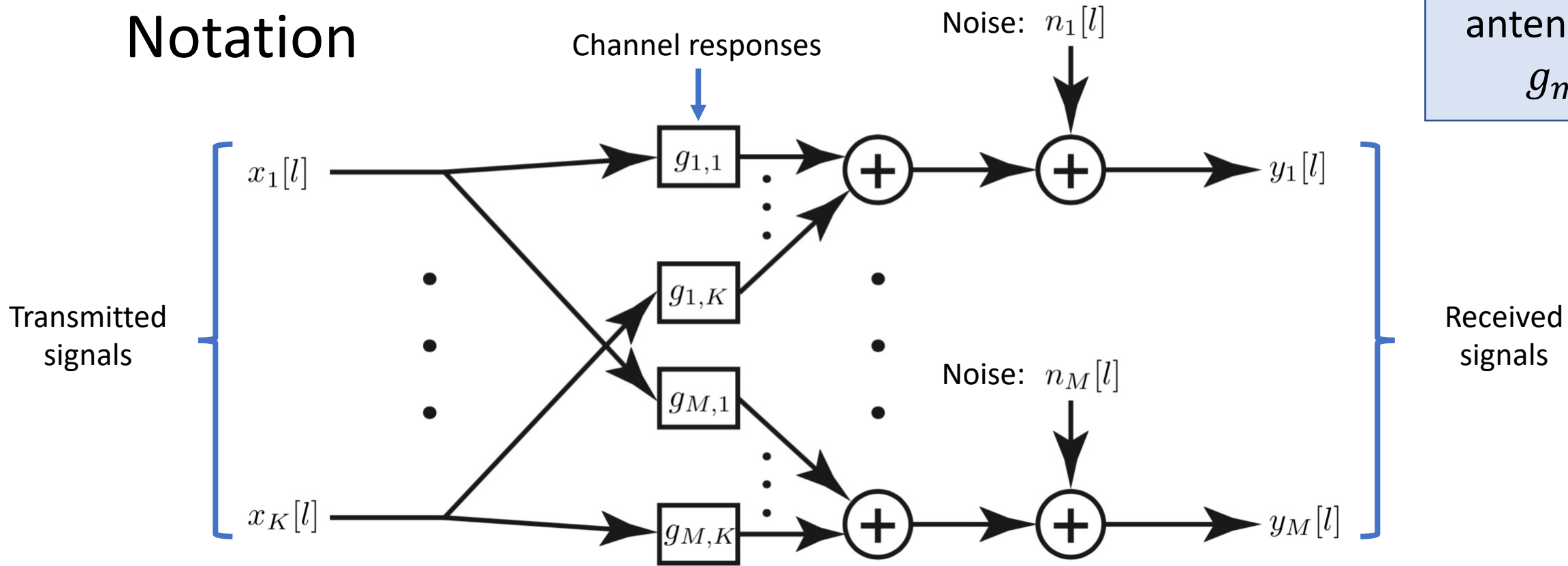
- Generalization of SIMO and MISO channels

What is the capacity of such a channel?

Notation

Channel from
antenna k to
antenna m :

$$g_{m,k}$$



Received signal at antenna m : $y_m[l] = \sum_{k=1}^K g_{m,k} x_k[l] + n_m[l]$

Vector-Matrix Description

- Memoryless channel to antenna m :

$$y_m = \sum_{k=1}^K g_{m,k} x_k + n_m$$

- Matrix form:

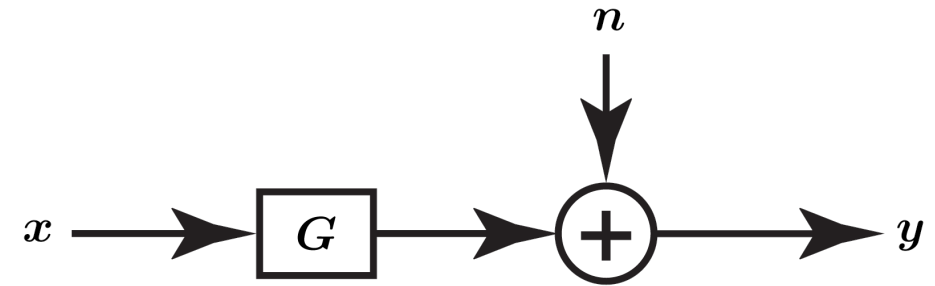
$$\begin{bmatrix} y_1 \\ \vdots \\ y_M \end{bmatrix} = \begin{bmatrix} \sum_{k=1}^K g_{1,k} x_k \\ \vdots \\ \sum_{k=1}^K g_{M,k} x_k \end{bmatrix} + \begin{bmatrix} n_1 \\ \vdots \\ n_M \end{bmatrix} = \underbrace{\begin{bmatrix} g_{1,1} & \dots & g_{1,K} \\ \vdots & \ddots & \vdots \\ g_{M,1} & \dots & g_{M,K} \end{bmatrix}}_{\mathbf{G}} \begin{bmatrix} x_1 \\ \vdots \\ x_K \end{bmatrix} + \begin{bmatrix} n_1 \\ \vdots \\ n_M \end{bmatrix}$$

Short form: $\mathbf{y} = \mathbf{G}\mathbf{x} + \mathbf{n}$

What is the channel capacity?

- Recall: Channel capacity is defined as

$$C = \max_{f(\mathbf{x}): E\{\|\mathbf{x}\|^2\} \leq q} I(\mathbf{x}; \mathbf{y}) \quad \text{bit/s/Hz}$$



- Mutual information: $I(\mathbf{x}; \mathbf{y}) = h(\mathbf{y}) - h(\mathbf{y}|\mathbf{x})$
- Power constraint: $E\{\|\mathbf{x}\|^2\} = \sum_{k=1}^K E\{|x_k|^2\} \leq q$
- Independent $\text{CN}(0, N_0)$ elements in \mathbf{n}

We will not compute $h(\mathbf{y})$ and $h(\mathbf{y}|\mathbf{x})$ directly
but look for a shortcut using linear algebra!

Eigenvalues and eigenvectors

- Consider an $M \times M$ matrix A
 - A non-zero vector \mathbf{u} is an *eigenvector* of A if
$$A\mathbf{u} = \lambda\mathbf{u}$$
where the scalar λ is the *eigenvalue* corresponding to \mathbf{u}

Eigenvalue decomposition

- If \mathbf{A} has M linearly independent eigenvectors, then

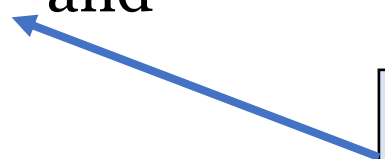
$$\mathbf{A} = \mathbf{U}\mathbf{D}\mathbf{U}^{-1}$$

- \mathbf{U} contains (unit-norm) eigenvectors as columns
- \mathbf{D} is the diagonal matrix with corresponding eigenvalues

- If \mathbf{A} is symmetric ($\mathbf{A} = \mathbf{A}^H$) then $\mathbf{U}^{-1} = \mathbf{U}^H$ and

$$\mathbf{A} = \mathbf{U}\mathbf{D}\mathbf{U}^H$$

- The matrix can be *diagonalized* as
$$\mathbf{U}^H \mathbf{A} \mathbf{U} = \mathbf{D}$$



Unitary matrix:
 $\mathbf{U}\mathbf{U}^H = \mathbf{U}^H\mathbf{U} = \mathbf{I}$

Singular value decomposition

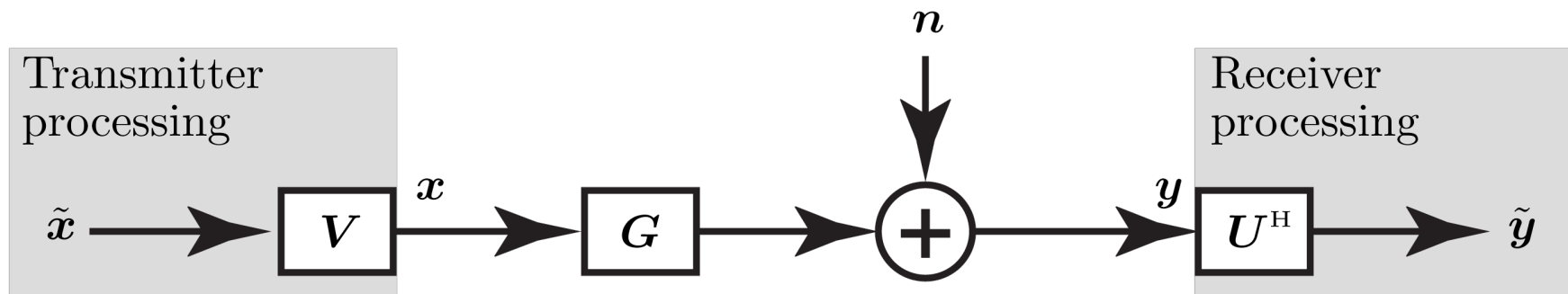
- Every complex $M \times K$ matrix \mathbf{G} can be factorized as

$$\mathbf{G} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^H$$

- Left singular vectors:
 - \mathbf{U} is an $M \times M$ unitary matrix containing the eigenvectors of $\mathbf{G}\mathbf{G}^H$
- Right singular vectors:
 - \mathbf{V} is a $K \times K$ unitary matrix containing the eigenvectors of $\mathbf{G}^H\mathbf{G}$
- Singular values:
 - $\mathbf{\Sigma}$ is an $M \times K$ “diagonal” matrix with $s_1 \geq s_2 \geq \dots \geq s_{\min(M,K)} \geq 0$ on diagonal
 - $s_1^2, s_2^2, \dots, s_{\min(M,K)}^2$ are the non-zero eigenvalues of $\mathbf{G}\mathbf{G}^H$ and $\mathbf{G}^H\mathbf{G}$

Diagonalizing the MIMO channel

- Non-destructive processing:
 - Pre-processing: $\mathbf{x} = \mathbf{V}\tilde{\mathbf{x}}$
 - Post-processing: $\tilde{\mathbf{y}} = \mathbf{U}^H \mathbf{y}$



$$\tilde{\mathbf{y}} = \mathbf{U}^H (\mathbf{G}\mathbf{x} + \mathbf{n}) = \underbrace{\mathbf{U}^H \mathbf{U} \boldsymbol{\Sigma} \mathbf{V}^H \mathbf{V}}_{\text{Singular value decomposition}} \tilde{\mathbf{x}} + \underbrace{\mathbf{U}^H \mathbf{n}}_{\tilde{\mathbf{n}}}$$

Singular value decomposition

$\tilde{\mathbf{n}}$: Rotated noise: i.i.d. $CN(0, N_0)$

Diagonalizing the MIMO channel (2)

- Processed received signal:

$$\tilde{\mathbf{y}} = \mathbf{\Sigma} \tilde{\mathbf{x}} + \tilde{\mathbf{n}}$$

- Let S denote the number of non-zero singular values (rank of \mathbf{G}):

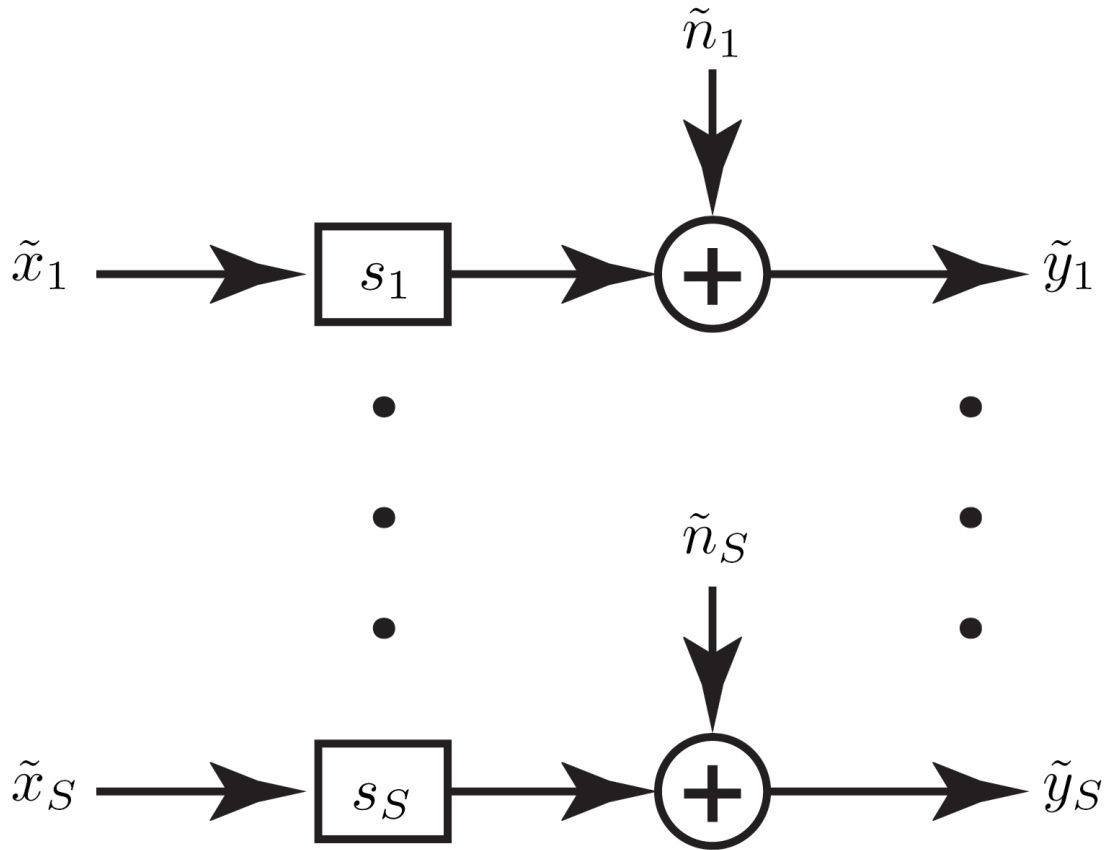
$$\tilde{y}_k = \begin{cases} s_k \tilde{x}_k + \tilde{n}_k, & \text{if } k = 1, \dots, S, \\ \tilde{n}_k, & \text{if } k = S + 1, \dots, M \end{cases}$$

The diagram illustrates the decomposition of the processed received signal \tilde{y}_k into two components based on the index k . The first component, $s_k \tilde{x}_k + \tilde{n}_k$, represents the signal in the S useful subchannels, where s_k is the singular value and \tilde{x}_k is the transmitted signal. The second component, \tilde{n}_k , represents the noise in the $M - S$ useless subchannels. Blue arrows point from the text boxes to the corresponding parts of the piecewise function.

Useful subchannels

Useless subchannels

S parallel channels



Suppose we assign power q_k to subchannel k

- Capacity of subchannel k :

$$R_k = \log_2 \left(1 + s_k^2 \frac{q_k}{N_0} \right)$$

Channel capacity:

$$C = \max_{q_1, \dots, q_S: q_1 + \dots + q_S = q} \sum_{k=1}^S R_k$$

Optimal Power Allocation

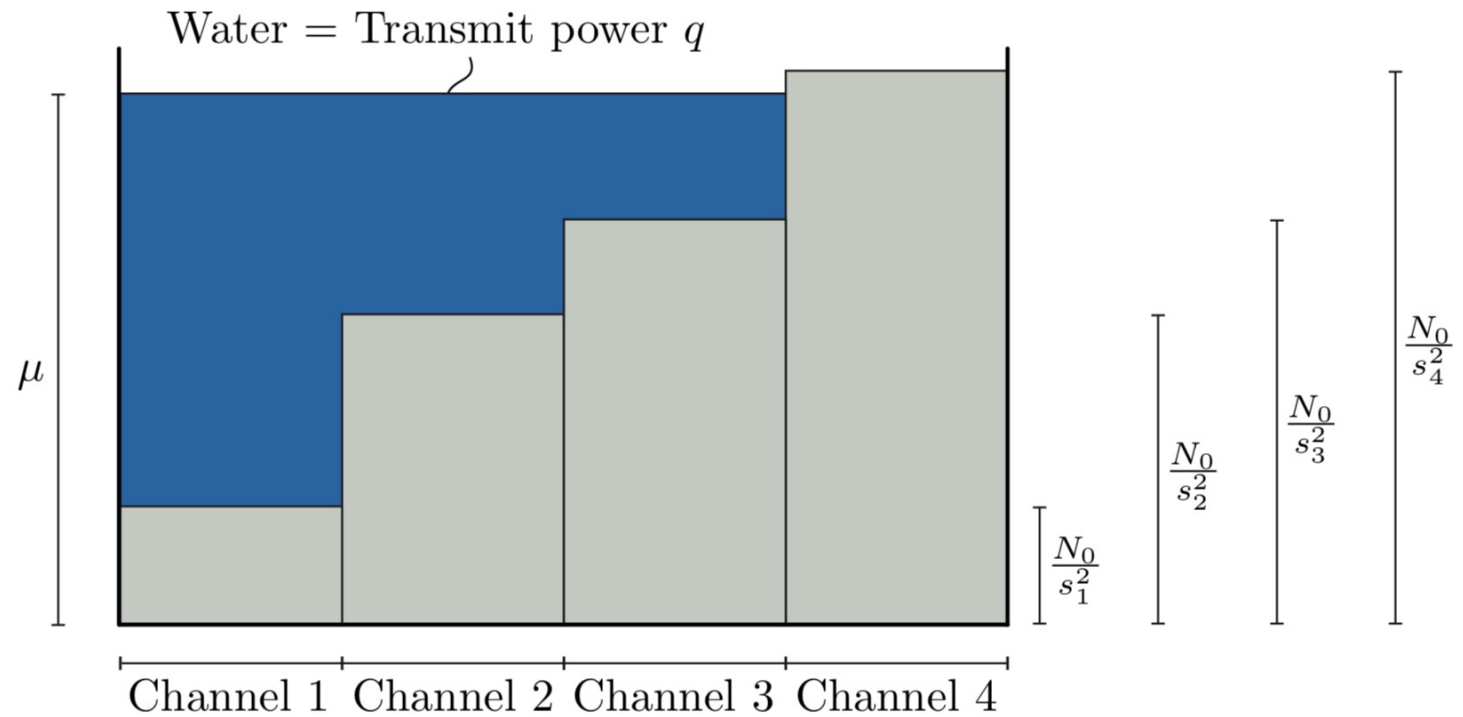
- After some optimization:

$$q_k = \max\left(\mu - \frac{N_0}{s_k^2}, 0\right)$$

where μ is selected such that $q_1 + \dots + q_S = q$

Properties:

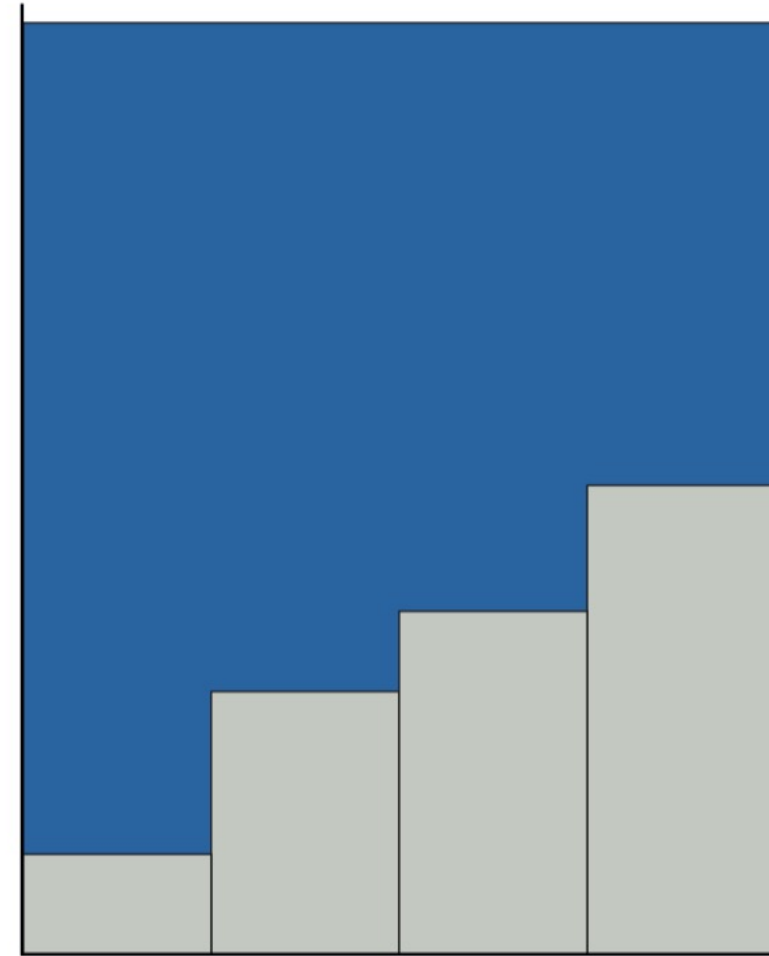
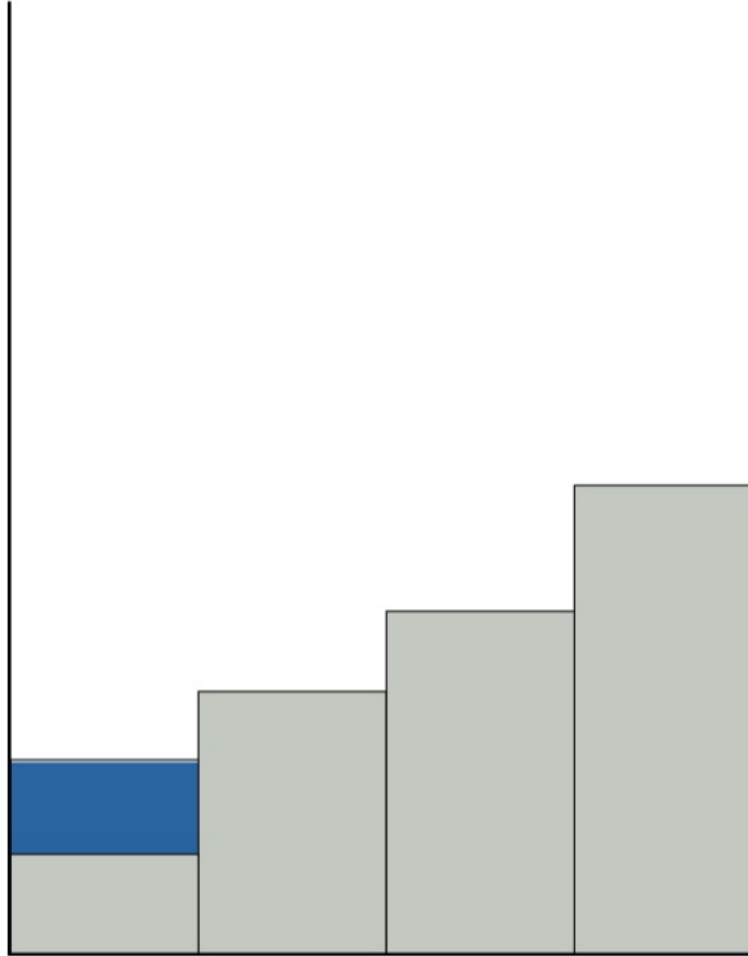
- Larger s_k : More power
- Some subchannels might get zero power



Called waterfilling

Low and high SNR

- Low SNR:
 - Only power to one subchannel
- High SNR:
 - Approximately the same power to all subchannels

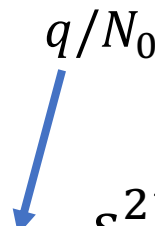


Capacity behavior at high SNR

- Equal power allocation (approximately optimal):

- $q_1 = \dots = q_S = \frac{q}{S}$

- Capacity approximation:

$$C \approx \sum_{k=1}^S \log_2 \left(1 + \text{SNR} \frac{s_k^2}{S} \right) \approx S \log_2(\text{SNR}) + \sum_{k=1}^S \log_2 \left(\frac{s_k^2}{S} \right)$$


Multiplexing gain:

First term proportional to $S = \text{rank}(\mathbf{G}^H \mathbf{G}) \leq \min(M, K)$

Capacity behavior at low SNR

- Assume $s_1 > \dots > s_S$
 - Low SNR: $q_1 = q$, $q_2 = \dots = q_S = 0$
- Capacity approximation:
$$C \approx \log_2(1 + \text{SNR } s_1^2) \approx \log_2(e) \text{SNR } s_1^2$$
- No multiplexing gain

Beamforming gain:
 s_1 is the largest singular value

Capacity comparison with $|g_{m,k}| = 1$

- SISO channel:

$$C = \log_2(1 + \text{SNR})$$

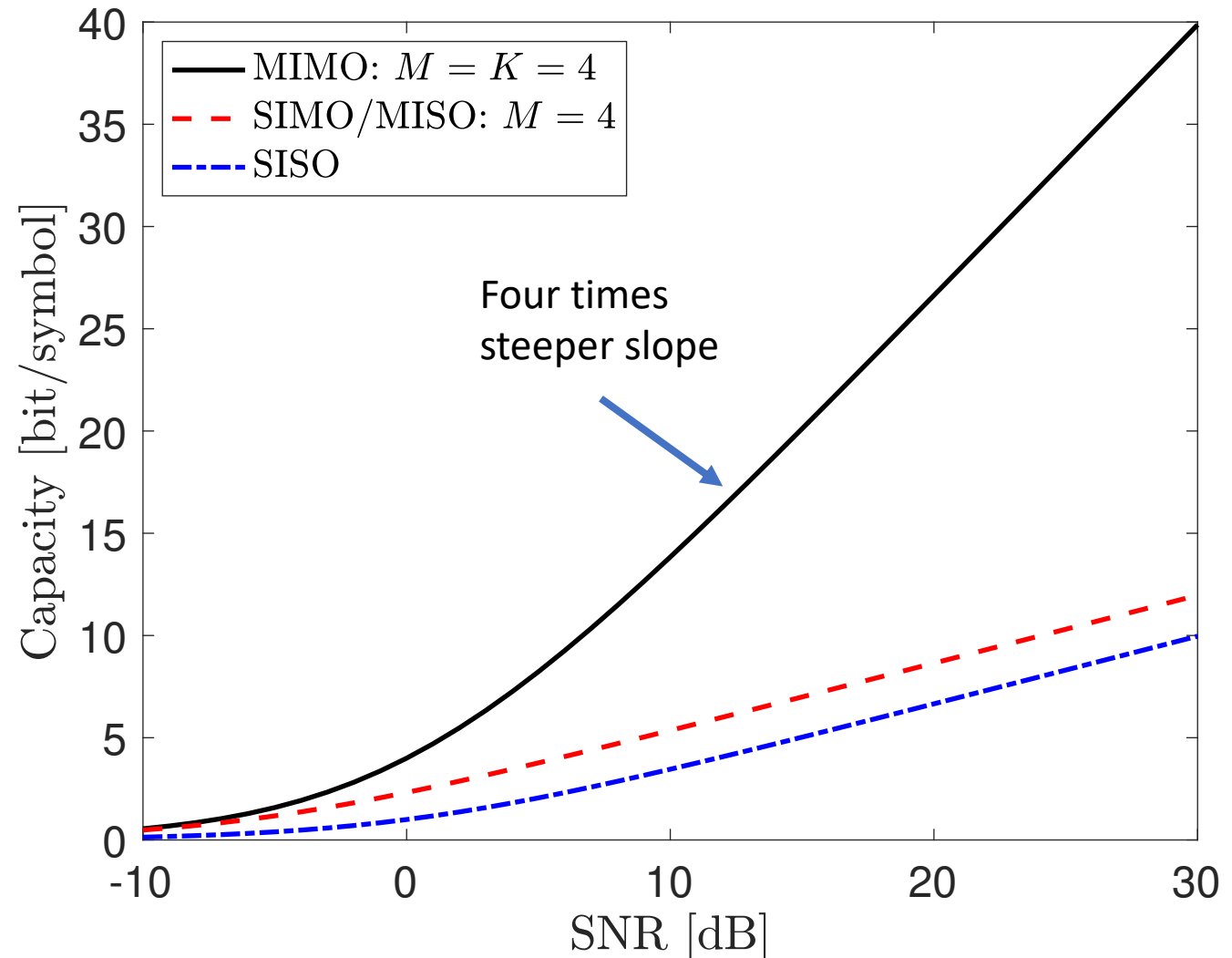
- SIMO/MISO, M antennas:

$$C = \log_2(1 + M \cdot \text{SNR})$$

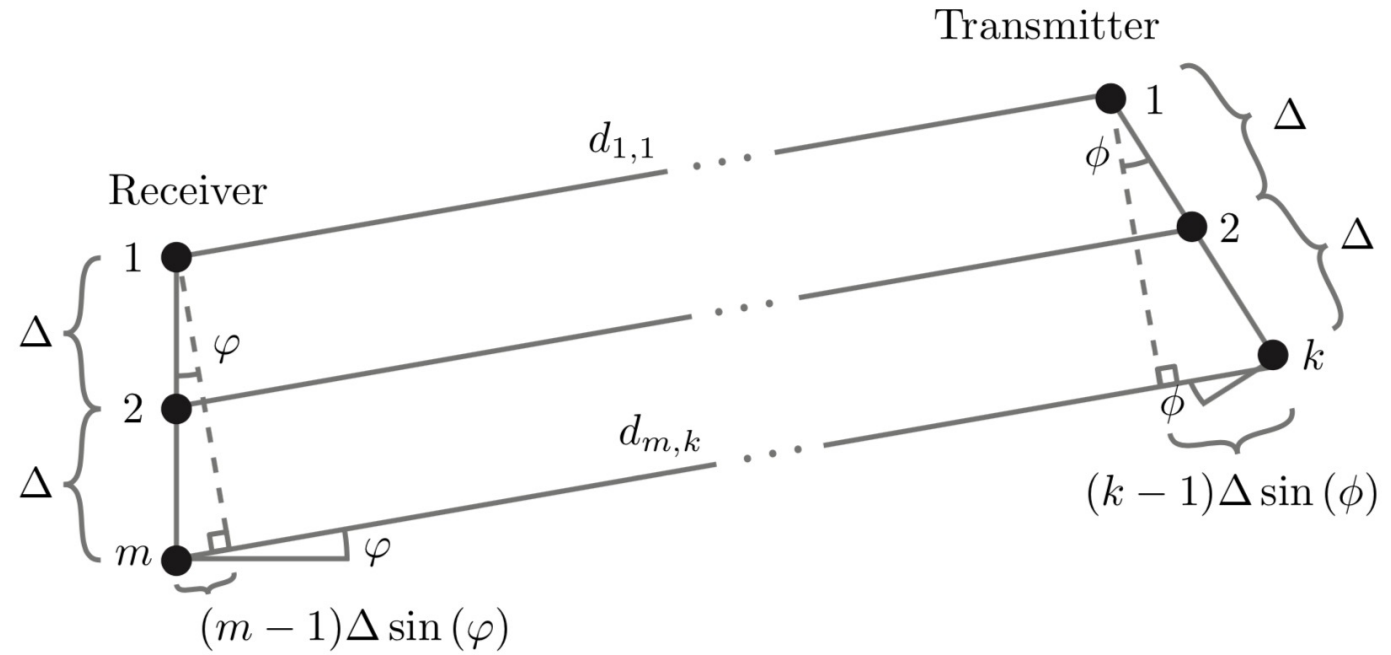
- MIMO, $M = K$ antennas:

- All singular values are \sqrt{M}

$$C = M \cdot \log_2(1 + \text{SNR})$$



Example: Line-of-sight channel



$$\begin{aligned}
 \mathbf{G} &= \begin{bmatrix} g_{1,1} & \dots & g_{1,K} \\ \vdots & \ddots & \vdots \\ g_{M,1} & \dots & g_{M,K} \end{bmatrix} = \sqrt{\beta} \begin{bmatrix} 1 & \dots & e^{-j2\pi \frac{(K-1)\Delta \sin(\phi)}{\lambda}} \\ \vdots & \ddots & \vdots \\ e^{-j2\pi \frac{(M-1)\Delta \sin(\phi)}{\lambda}} & \dots & e^{-j2\pi \frac{(M-1)\Delta \sin(\phi)}{\lambda}} e^{-j2\pi \frac{(K-1)\Delta \sin(\phi)}{\lambda}} \end{bmatrix} \\
 &= \sqrt{\beta} \begin{bmatrix} 1 \\ \vdots \\ e^{-j2\pi \frac{(M-1)\Delta \sin(\phi)}{\lambda}} \end{bmatrix} \begin{bmatrix} 1 & \dots & e^{-j2\pi \frac{(K-1)\Delta \sin(\phi)}{\lambda}} \end{bmatrix}
 \end{aligned}$$

$$\text{Rank 1: } s_1 = \sqrt{\beta MK}$$

Line-of-sight channels: No multiplexing gain

- We have $S = 1$:

- $q_1 = q$

- Capacity:

$$C = \log_2(1 + \beta MK \cdot \text{SNR})$$

- Compare with SIMO and MISO:

$$C = \log_2(1 + \beta M \cdot \text{SNR})$$

Beamforming gain:
Larger in the MIMO case

Slow fading and MISO channels ($M = 2$)

- Received signal

$$y[l] = [g_1 \quad g_2] \begin{bmatrix} x_1[l] \\ x_2[l] \end{bmatrix} + n[l]$$

- Fixed channel g_1, g_2 for the entire transmission
- Assumption: Receiver knows g_1, g_2 , but not the transmitter

- Consider two time slots:
$$\begin{bmatrix} y[1] & y[2] \end{bmatrix} = \underbrace{\begin{bmatrix} g_1 & g_2 \end{bmatrix}}_{=\mathbf{g}^T} \underbrace{\begin{bmatrix} x_1[1] & x_1[2] \\ x_2[1] & x_2[2] \end{bmatrix}}_{=\mathbf{X}} + \begin{bmatrix} n[1] & n[2] \end{bmatrix}$$

Can we select \mathbf{X} in a clever way?

Space-time block coding

- Alamouti code: $\mathbf{X} = \frac{1}{\sqrt{2}} \begin{bmatrix} \tilde{x}[1] & \tilde{x}[2] \\ -\tilde{x}^*[2] & \tilde{x}^*[1] \end{bmatrix}$

Data signals $\tilde{x}[1], \tilde{x}[2]$



- Consider received signal:
$$\begin{aligned} \begin{bmatrix} y[1] \\ y^*[2] \end{bmatrix} &= \frac{1}{\sqrt{2}} \begin{bmatrix} g_1 \tilde{x}[1] - g_2 \tilde{x}^*[2] \\ g_1^* \tilde{x}^*[2] + g_2^* \tilde{x}[1] \end{bmatrix} + \begin{bmatrix} n[1] \\ n^*[2] \end{bmatrix} \\ &= \underbrace{\frac{1}{\sqrt{2}} \begin{bmatrix} g_1 & -g_2 \\ g_2^* & g_1^* \end{bmatrix}}_{=\tilde{\mathbf{G}}} \begin{bmatrix} \tilde{x}[1] \\ \tilde{x}^*[2] \end{bmatrix} + \begin{bmatrix} n[1] \\ n^*[2] \end{bmatrix} \end{aligned}$$

$$\tilde{\mathbf{G}} = \underbrace{\frac{1}{\|\mathbf{g}\|} \begin{bmatrix} g_1 & -g_2 \\ g_2^* & g_1^* \end{bmatrix}}_{=\tilde{\mathbf{U}}} \underbrace{\begin{bmatrix} \frac{\|\mathbf{g}\|}{\sqrt{2}} & 0 \\ 0 & \frac{\|\mathbf{g}\|}{\sqrt{2}} \end{bmatrix}}_{=\tilde{\mathbf{\Sigma}}} \underbrace{\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}}_{=\tilde{\mathbf{V}}^H}$$

This is like a 2×2 MIMO channel!
Transmitter knows $\tilde{\mathbf{V}}$ without
knowing channel

Transmit diversity versus receive diversity

- Ideal capacity with MISO:

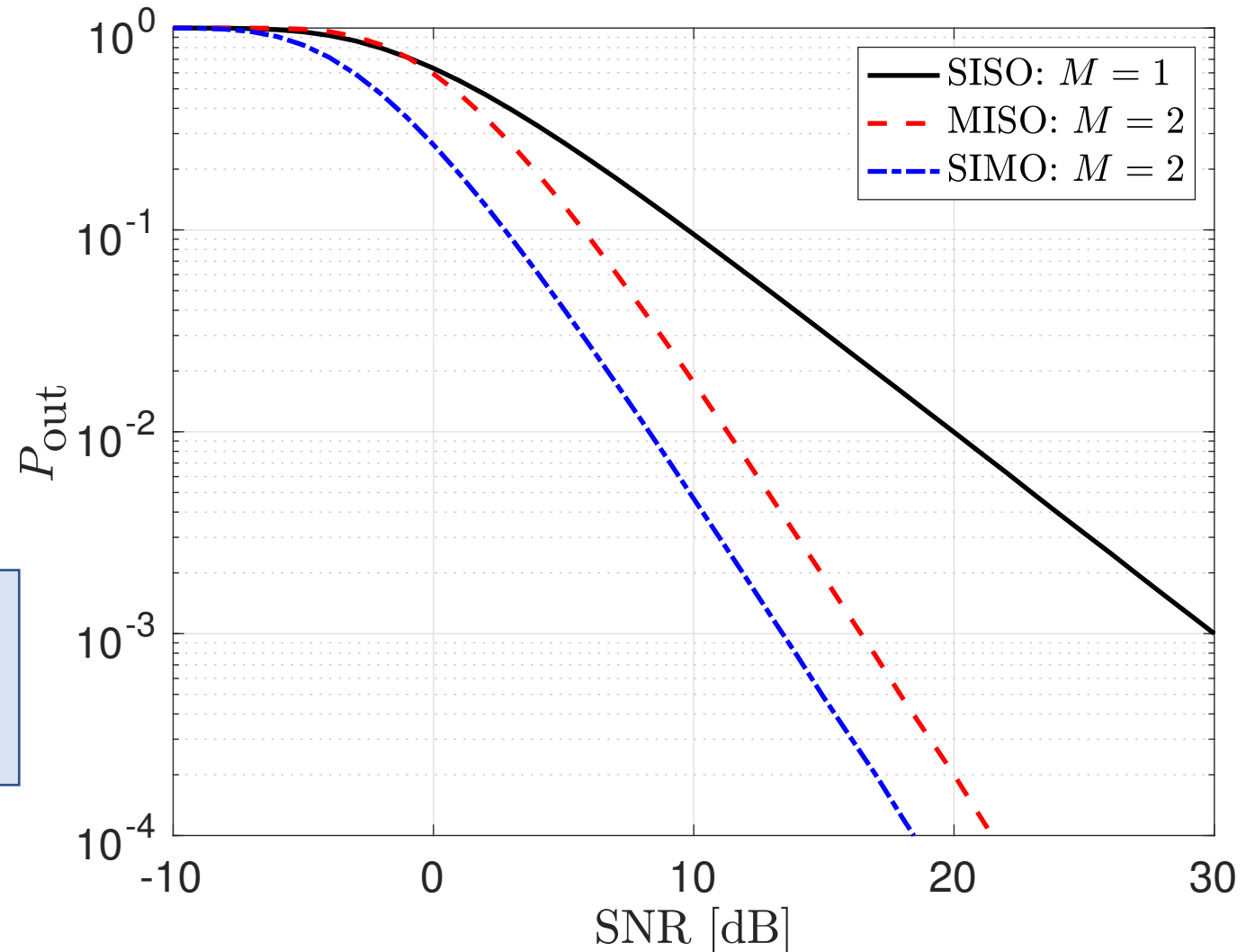
$$C_g = \log_2 \left(1 + \frac{q}{2N_0} \|\mathbf{g}\|^2 \right)$$

- Ideal capacity with SIMO:

$$C_g = \log_2 \left(1 + \frac{q}{N_0} \|\mathbf{g}\|^2 \right)$$

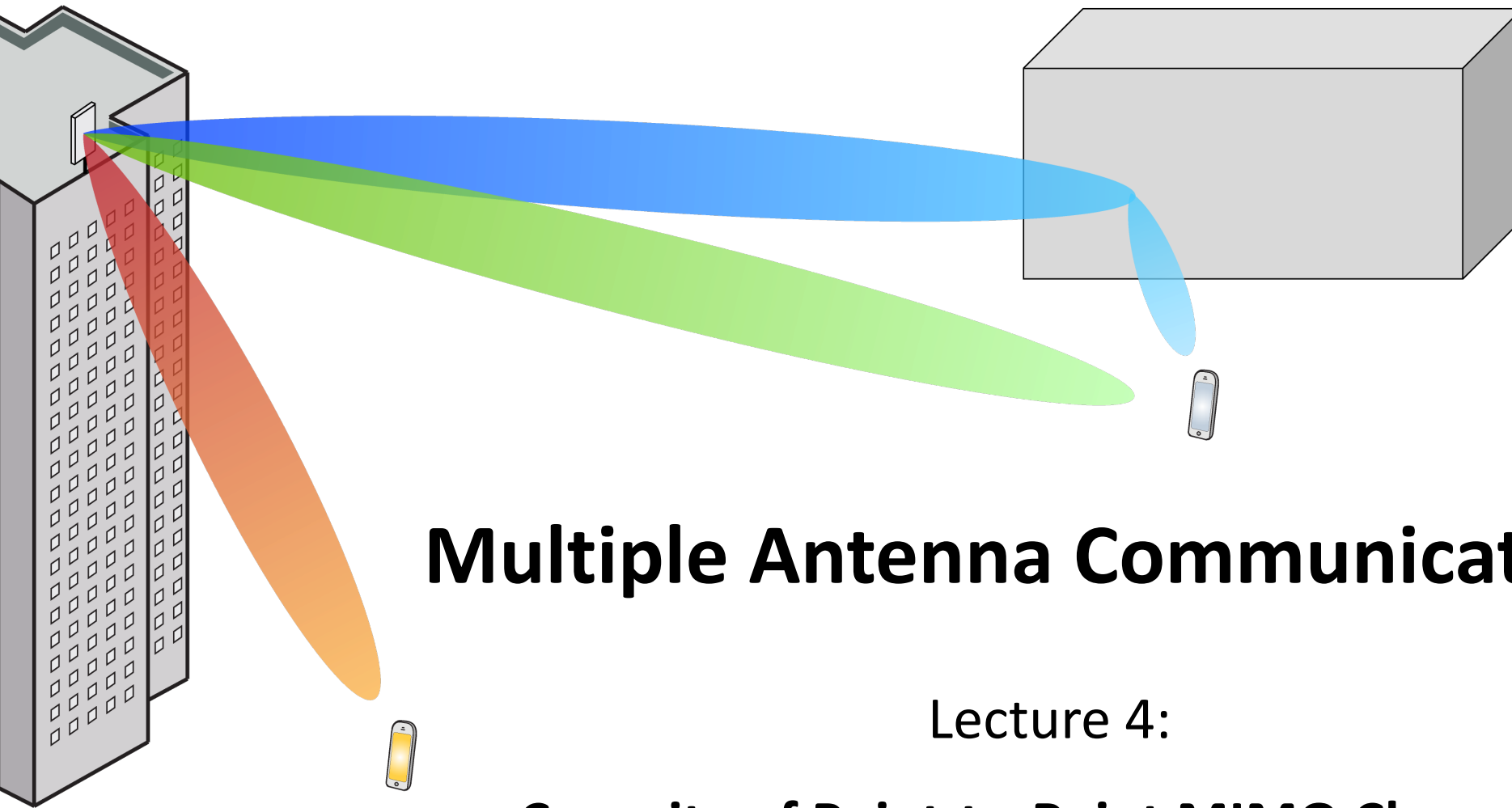
Outage probability:

Same behavior but half the SNR
Diversity gain, but no beamforming gain



Summary

- Capacity of $M \times K$ MIMO channel
 - Send K different messages along eigenvectors of $\mathbf{G}^H \mathbf{G}$
 - This creates K parallel channels
 - Divide power using waterfilling
 - High SNR: Multiplexing gain
 - Low SNR: Beamforming gain
- MISO channels with slow fading
 - No beamforming gain, but diversity gain can be achieved



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Capacity of Point-to-Point MIMO Channels

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