



# Computational Framework for Optimal Robust Beamforming in Coordinated Multicell Systems

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### Introduction

- Downlink Coordinated Beamforming
  - N cells with  $N_t$ -antenna base stations
  - Each serves *K* single-antenna users
  - Common narrowband frequency resource
  - Limited by co-user interference

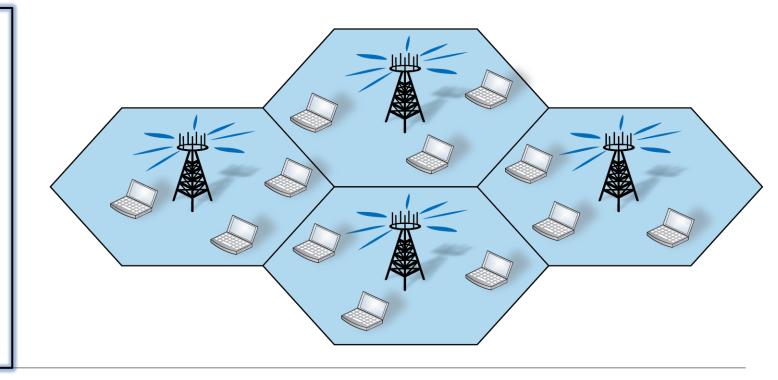
### **Problem**

Compute Optimal Linear Beamforming

**General Conditions** 

Robustness to Channel Uncertainty

Generally NP-hard: Systematic Algorithm





### System Model

- Parameters for User j in Cell i

  - Linear beamforming vector:  $\mathbf{w}_{i,j} \in \mathbb{C}^{N_t \times 1}$  Channel from cell m:  $\mathbf{h}_{m,i,j} \in \mathbb{C}^{N_t \times 1}$
  - Signal-to-interference-and-noise ratio (SINR):

$$\text{SINR}_{i,j} = \underbrace{\frac{\left|\mathbf{h}_{i,i,j}^H \mathbf{w}_{i,j}\right|^2}{\left|\mathbf{h}_{i,i,j}^H \mathbf{w}_{i,l}\right|^2 + \sum\limits_{m \neq i} \lVert \mathbf{h}_{m,i,j}^H \mathbf{W}_m \rVert_2^2 + \sigma^2}_{\text{Intra-cell interference}}$$

- Notation:  $\mathbf{W}_i = [\mathbf{w}_{i,1} \ ... \ \mathbf{w}_{i,K}]$ 



# System Model (2)

- Arbitrary Power Constraints in Cell i
  - Constraints:

$$\mathcal{W}_i = \left\{ \mathbf{W}_i \colon \operatorname{tr} \left\{ \mathbf{W}_i^H \mathbf{Q}_{i,k} \mathbf{W}_i \right\} \le q_{i,k} \ \forall k \right\}$$

Positive - Notation:  $\mathbf{W}_i = [\mathbf{w}_{i,1} \ \ldots \ \mathbf{w}_{i,K}]$  semi-definite

- Ex: Per-antenna and per-cell

- Channel State Information (CSI)
  - Perfect CSI within each cell
  - Uncertain inter-cell CSI: Ellipsoidal uncertainty set

$$\mathbf{h}_{m,i,j} \in \mathcal{U}_{m,i,j} = \left\{ \widehat{\mathbf{h}}_{m,i,j} + \mathbf{B}_{m,i,j} \boldsymbol{\epsilon}_{m,i,j} : \ \|\boldsymbol{\epsilon}_{m,i,j}\|_2 \leq 1 \right\}$$
 Uncertainty Known Shape of Unknown set estimate ellipsoid error vector

### **Error Sources**

Estimation Feedback Delays

More severe between cells

Positive



### Measures of User Performance

- General Measure of User Performance
  - Arbitrary strictly increasing function:

$$g_{i,j}(SINR_{i,j})$$
 with  $g_{i,j}(0) = 0$ 

- Ex: Mutual information, Bit error rate, Mean squared error

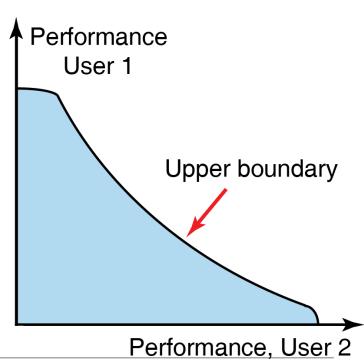
- Worst-Case Robust User Performance
  - We will try to maximize:

$$g_{i,j}(\widetilde{\text{SINR}}_{i,j})$$
 where  $\widetilde{\text{SINR}}_{i,j} = \min_{\substack{\mathbf{h}_{m,i,j} \in \mathcal{U}_{m,i,j} \\ \forall m \neq i}} \text{SINR}_{i,j}$ 



# Measures of User Performance (2)

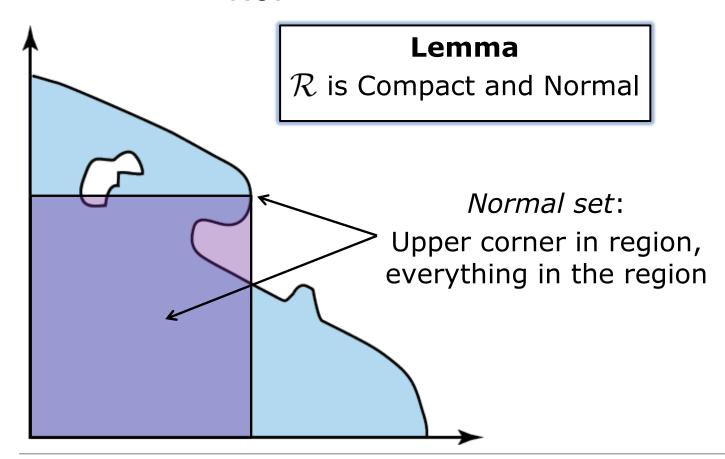
- Many Users
  - One Performance Measure  $g_{i,j}(\mathrm{SINR}_{i,j})$  per User
- Fairness Dimension
  - Divide power and control co-user interference
- Robust Performance Region  $\mathcal R$ 
  - NK users  $\leftrightarrow NK$  dimensions
  - All possible combinations
  - Good points: On upper boundary
  - Unknown shape





# Measures of User Performance (3)

- Can it have any shape?
  - Can be non-convex!
- No!





# System Performance

- Which Point in R to Select?
- ullet System Performance Function  $f:\mathcal{R} o\mathbb{R}$ 
  - Strictly increasing and Lipschitz continuous

### Examples

- Sum performance:  $f(\mathbf{g}) = \sum_{i,j} g_{i,j}$ 

- Proportional fairness:  $f(\mathbf{g}) = \prod_{i,j} g_{i,j}$ 

- Harmonic mean:  $f(\mathbf{g}) = NK(\sum_{i,j} g_{i,j}^{-1})^{-1}$ 

- Max-min fairness:  $f(\mathbf{g}) = \min_{i,j} g_{i,j}$ 

- Can be modified with weights

# Accumulated performance





### **Problem Formulation**

Optimize System Performance

$$\underset{\mathbf{W}_{i} \in \mathcal{W}_{i} \ \forall i}{\operatorname{maximize}} \ f\left(g_{1,1}(\widetilde{\mathbf{SINR}}_{1,1}), \dots, g_{N,K}(\widetilde{\mathbf{SINR}}_{N,K})\right)$$
 (1)

- Lemma: Optimum on upper boundary of  ${\cal R}$
- Generally NP-hard: Exponential complexity
- Only suboptimal strategies in practice
- Goal: Computational Framework for Solving (1)
  - Enable benchmarking and study properties
- Approach
  - Solve a special case of f()
  - Exploit it to solve (1) for general f()



- Maximize Performance with Fairness Constraints
  - Generalization of classic max-min fairness:

$$\max_{\mathbf{W}_{i} \in \mathcal{W}_{i} \ \forall i} \min_{i,j} \frac{g_{i,j}(\widetilde{\mathbf{SINR}}_{i,j}) - a_{i,j}}{\alpha_{i,j}} \\
\text{s.t.} \quad g_{i,j}(\widetilde{\mathbf{SINR}}_{i,j}) \ge a_{i,j} \quad \forall i, j$$

- Lowest acceptable performance level:  $g_{i,j} \geq a_{i,j} \geq 0$
- Users get a portion of exceeding resources:  $\alpha_{i,j} \geq 0$

#### Lemma

- Solved by line-search in  ${\mathcal R}$  (bisection)
- Exploiting that  ${\mathcal R}$  is normal and compact



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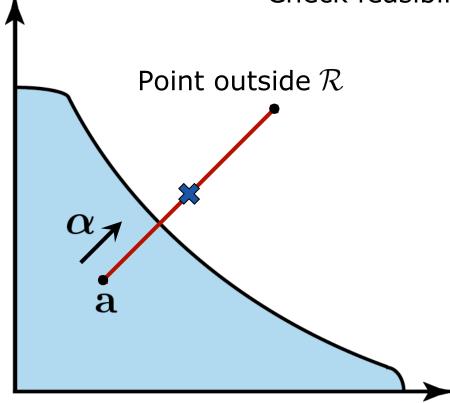
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- Geometrical Interpretation
  - Bisection: Fast convergence
  - Check feasibility at midpoint **c**:



#### **Theorem**

- Feasibility checked as convex problem
- CSI uncertainty handled using S-lemma

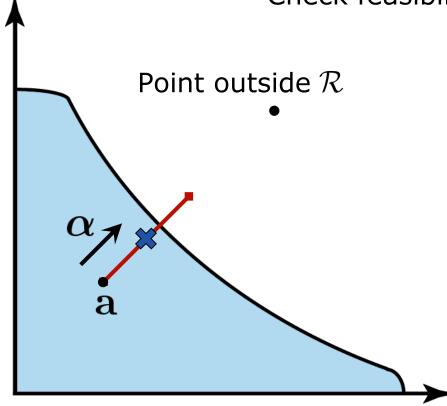
find 
$$\mathbf{W}_{i} \in \mathcal{W}_{i}, \lambda_{m,i,j} \geq 0, b_{m,i,j} \geq 0 \ \forall i, j, m \neq i$$
  
s.t. 
$$\begin{bmatrix} b_{m,i,j} - \lambda_{m,i,j} \ \mathbf{\hat{h}}_{m,i,j}^{H} \mathbf{W}_{m} & \mathbf{0} \\ \mathbf{W}_{m}^{H} \mathbf{\hat{h}}_{m,i,j} & b_{m,i,j} \mathbf{I}_{K} \ \mathbf{W}_{m}^{H} \mathbf{B}_{m,i,j} \end{bmatrix} \succeq \mathbf{0} \ \forall i, j, m \neq i$$

$$\mathbf{0} \quad \mathbf{B}_{m,i,j}^{H} \mathbf{W}_{m} \ \lambda_{m,i,j} \mathbf{I}_{N_{t}} \end{bmatrix}$$

$$\sqrt{1 + \frac{1}{C_{i,j}} \mathbf{h}_{i,i,j}^{H} \mathbf{w}_{i,j}} \ge \sqrt{\|\mathbf{h}_{i,i,j}^{H} \mathbf{W}_{i}\|_{2}^{2} + \sum_{m \neq i} b_{m,i,j}^{2} + \sigma^{2}} \quad \forall i, j$$



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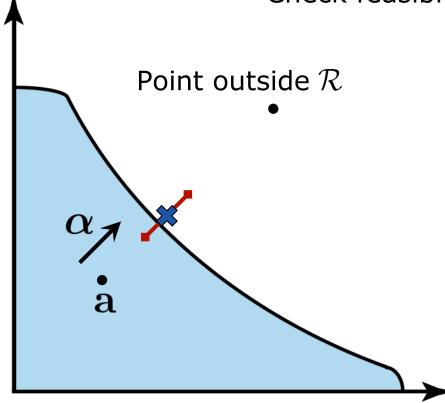
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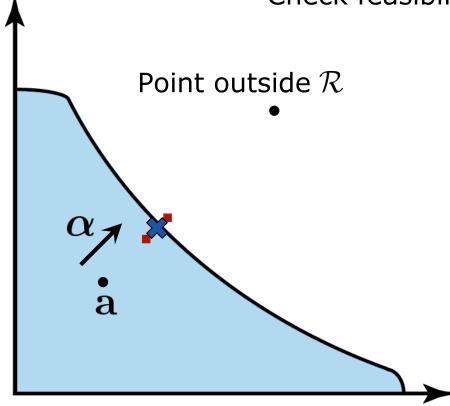
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### Framework for General Case

- Systematic Algorithm with Minimal Search Space
  - Search in  ${\mathcal R}$  and concentrate on important parts
  - Improve lower/upper bounds on optimum:

$$f_{\min} \le f_{\mathrm{opt}} \le f_{\max}$$

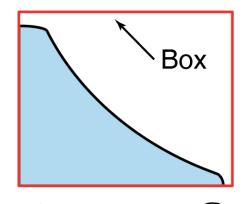
- Continue until  $f_{\mathrm{max}} - f_{\mathrm{min}} < arepsilon$ 

- Iterations in Polynomial Time
  - Fairness-profile optimization

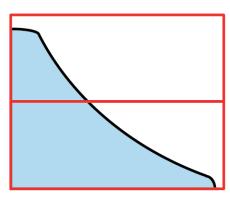


## Framework for General Case (2)

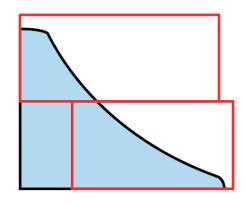
- Branch-Reduce-Bound (BRB) Algorithm
  - 1. Cover  $\mathcal{R}$  with a box
  - 2. Divide the box into two sub-boxes
  - 3. Remove parts with no solutions in  $[f_{\min}, f_{\max}]$
  - 4. Search for solutions to improve bounds (Fairness-profile optimization)
  - 5. Continue with sub-box with largest value



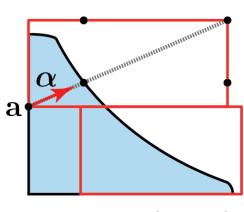
1) Cover region  ${\cal R}$ 



2) Branch (Divide)



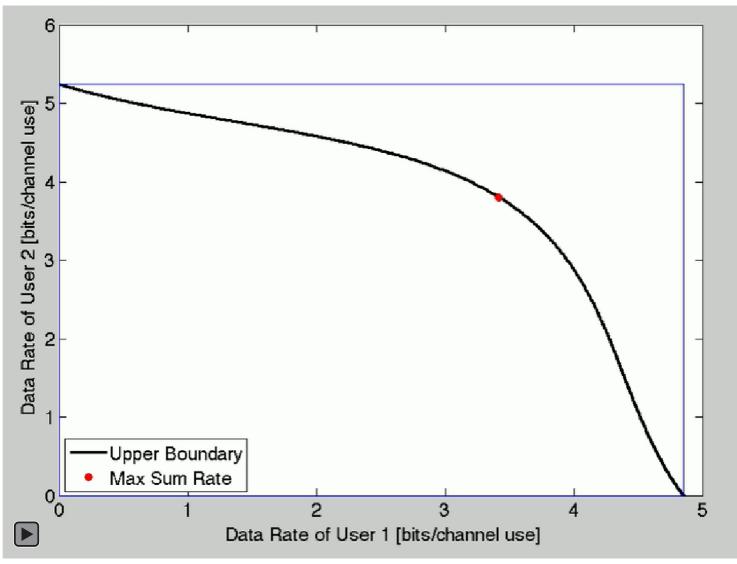
3) Reduce using bounds



4) Improve bounds



# Framework for General Case (3)



### **Theorem**

- Global Convergence
- Accuracy ε>0 in finitely many iterations
- Exponential complexity only in NK
- Polynomial complexity in  $N_t$  and #constraints
- Any accuracy of fairness-profile opt



### Numerical Illustrations

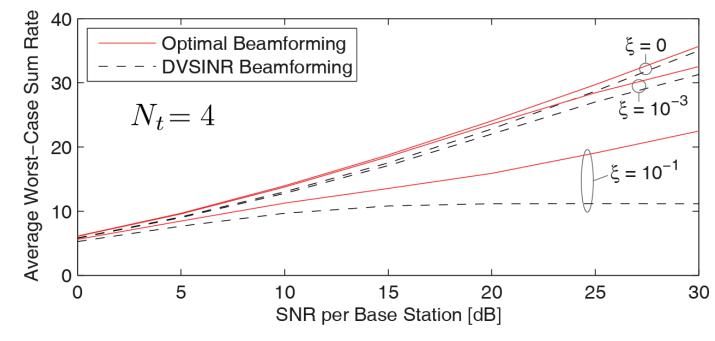
### 2 Cells and 2 User/Cell

Per-base station constraints

Uncorrelated Rayleigh fading:  $\mathbb{E}\{\|\mathbf{h}_{i,i,j}\|^2\} = 2\mathbb{E}\{\|\widehat{\mathbf{h}}_{m,i,j}\|^2\}$ 

Spherical uncertainty sets:  $\mathbf{B}_{m,i,j} = \sqrt{\xi} \mathbf{I}_{N_t}$ 

- Robustness of Heuristic Beamforming
  - DVSINR beamforming [Björnson et al., 2010]
  - Robust to small intercell uncertainties
  - Highly suboptimal at higher uncertainties





# Numerical Illustrations (2)

### 2 Cells and 2 User/Cell

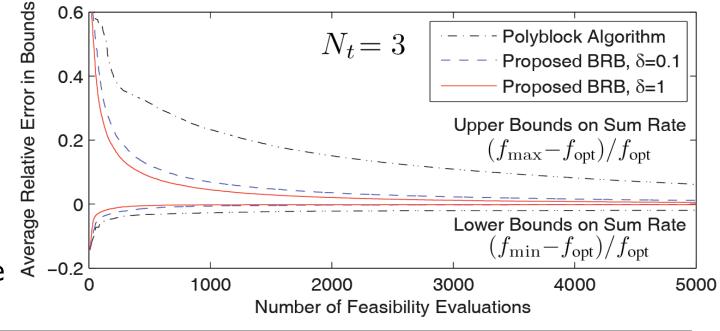
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### Convergence of Lower/Upper Bounds

- Compared with Polyblock algorithm
- Plot relative error of lower/upper bounds
- BRB algorithm has faster convergence
- Accurate fairness-profile not necessary





### Summary

- Robust Coordinated Beamforming
  - Generally NP-hard ↔ Suboptimal strategies in practice
- Contribution: Computational Framework
  - Enables benchmarking and analysis
  - Robustness and general performance measures/constraints
- Fairness-Profile Optimization
  - Special case solved in polynomial time: Even with robustness
  - Subproblem of general algorithm
- Branch-Reduce-and-Bound Algorithm
  - Systematic algorithm for the general problem
  - Guaranteed to find global solution
  - More general and better convergence than previous work



### Extensions

### Journal Article

- E. Björnson, G. Zheng, M. Bengtsson, B. Ottersten, "Robust Monotonic Optimization Framework for Multicell MISO Systems," IEEE Transactions on Signal Processing, Under Minor Revision, arXiv:1104.5240v2.
- Contains all mathematical details
- Extension 1: All channels can be uncertain
- Extension 2: Applicable whenever subproblem can be solved efficiently



### Thank You for Listening!

### **Questions?**

Papers and Presentations Available:

http://www.ee.kth.se/~emilbjo