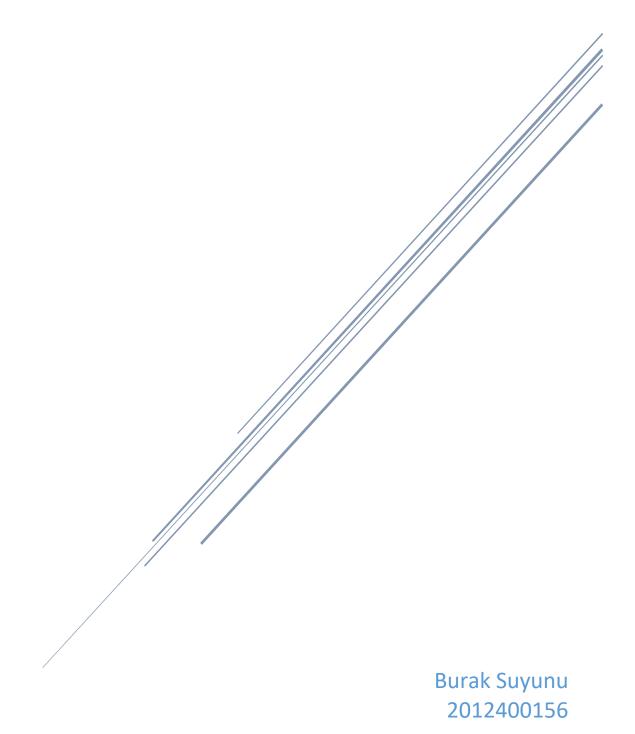
FINAL PROJECT – TABU SEARCH ON TRAVELLING SALESPERSON PROBLEMS WITH PROFITS

IE 517 - Spring17



Important Note: This report is a summary of my detailed work. So, please check the Jupyter Notebook that I have created in https://github.com/suyunu/TSPs-with-Profit repository. You can directly view the code from here: https://github.com/suyunu/TSPs-with-Profit/blob/master/ts-tspp.ipynb . You can find the detailed process of the code with heavy commenting.

Introduction

In this project, we tried to solve Travelling Salesperson Problems with Profits (TSPs with profits) with Tabu Search (TS). Before I start doing anything on the problem, I made a literature survey. There are lots of papers in the literature about TSPs with profits but those papers are generally tries to solve it with some constraints. So actually, I couldn't find a good paper to pointing out our problem which has no constraint. But the following paper has some good ideas about the general structure of the problem even it has a constraint on the tour length:

 Gendreau, Michel, Gilbert Laporte, and Frédéric Semet. "A tabu search heuristic for the undirected selective travelling salesman problem." European Journal of Operational Research 106.2-3 (1998): 539-545.

Travelling Salesperson Problems with Profits

Traveling Salesperson problems with profits (TSPs with profits) are a generalization of the traveling salesman problem (TSP), where it is not necessary to visit all vertices. A profit is associated with each vertex. The overall goal is the simultaneous optimization of the collected profit and the travel costs.

(http://pubsonline.informs.org/doi/abs/10.1287/trsc.1030.0079?journalCode=trsc)

Solution Representation

I used a simple permutation representation. The list [1, 2, 3, 4, 5, 1] represents the route of the salesperson. All the routes should start with "1" and end with "1" which is the depot.

Tabu Search

In this part I will explain the steps of tabu search for the travelling salesperson problems with profits.

Pseudocode

- 1. (Initialization) Construct an initial tour by means of a construction heuristic.
- 2. (Insertion Partitions) Determine all insertion partitions according to proximity measure and retain 10 of them.
- 3. Repeat Step **4-10** for *10.000* iterations:
- 4. (Insertion Candidate) Randomly choose one insertion partition and determine the best insertion candidate from this partition
- 5. (**Deletion Chains**) Determine the deletion chains.
- 6. (Deletion Candidate) Determine the best deletion candidate from deletion chains
- 7. (Insertion or Deletion) Compare the results of the insertion and deletion then apply the best one. If the best move is deletion, then declare all vertices of deletion tabu for θ iteration
- 8. **(Tour Improvement)** If the iteration count is multiple of 5, apply 2-opt
- 9. **(Best Solution Update)** If newly generated solution has a better objective than the incumbent solution then apply *3-opt* to the newly generated solution to improve the tour quality and make it the incumbent solution.
- 10. **(Shuffle to Reset)** If there hasn't been an improvement in γ iteration, then assign incumbent solution to the current solution and shuffle the route. Also, resets the tabu list.

Initialization

- 1. Determine a tour length V and start building a tour
- 2. Until length of route T reaches V, repeat:
- 3. Randomly determine a spot j in the tour T and add the city $v_j \notin T$ having the minimal ratio $(d_{ij} + d_{jk} d_{ik})/p_j$
- 4. Apply 2-opt to the generated list

After lots of test initial tour length V=N/2 gave the best results. So function runs 5 time and constructs 5 route with length N/2, then chooses the route with best objective.

Insertion Partitions

Dispersion Index

Dispersion index of a non-empty list R,

$$\Gamma(R) = \begin{cases} \frac{1}{|R|(|R|-1)} \sum_{v_i, v_j \in R} d_{-}ij \ , & if |R| > 1 \\ 0 & , & if |R| = 1 \end{cases}$$

Proximity Measure

Proximity measure between two non-empty lists R and S,

$$\Delta(R,S) = \frac{2}{|R||S|} \left(\sum_{v_i \in R, v_j \in S} d_i j \right) - \Gamma(R) - \Gamma(S)$$

Note that, if $R = \{v_i\}$ and $S = \{v_j\}$, then $\Delta(R, S) = d_{ij}$

Procedure

Using this proximity measure, we define several partitions of $V \setminus \{v_0\}$ within a preprocessing step of the algorithm. Each of these partitions contains clusters of vertices

- 1. (First Partition) Set r := 1 and $P_r := \{\{v_1\}, ..., \{v_n\}\}$
- 2. **(Next Partitions)** If r=N stop. Otherwise, define P_{r+1} form P_r by merging the two clusters C_{ri^*} and C_{rk^*} of P_r yielding $\min_{i\neq j}\{\Delta(C_{ri},C_{rk})\}$ set r:=r+1 and repeat this step.

We calculate all possible N-1 partitions but we only retain partitions P_r corresponding to $r=\{1,int(N/2),int(2N/3),int(3N/4),int(4N/5),int(5N/6),int(6N/7),int(7N/8),int(8N/9),int(9N/10)\}$ at the begining. One reason for keeping at most 10 partitions is to save memory. Moreover, removing partitions that are very similar to one another will create a diversification effect in the search process. This will become clearer later.

Insertion Candidates

The value of insertion of a cluster C'_{rk} from the partition P_r is measured by the ratio of added profit over added distance.

The gravity centre $\overline{v_k}$ of C'_{rk} is first computed for all clusters of P_r , and a preliminary move evaluation is made according to the formula

$$\bar{g}(C'_{sk}) = \frac{\sum_{v_h \in C'_{sk}} p_h}{l(T \cup \{\overline{v_k}\}) - l(T)}$$

The cluster C'_{sk^*} corresponding to $max_k\{\bar{g}(C'_{sk})\}$ is then selected. The exact move evaluation associated with C'_{sk^*} is

$$\bar{g}(C'_{sk^*}) = \frac{\sum_{v_h \in C'_{sk^*}} p_h}{l(T \cup C'_{sk^*}) - l(T)}$$

Deletion Chains

The sets of vertices H_{ij} candidate for removal are defined as follows. Consider a solution $T=\{v_0,\ldots,v_{j_0},v_{i_1},\ldots,v_{j_1},v_{i_2},\ldots,v_{j_{\lambda-1}},v_{i_0},\ldots,v_0\}$ are the λ longest edges of the tour and λ is an input parameter randomly selected in the interval $[2,\delta/2]$, and δ is the maximum between 4 and the number of vertices appearing on the initial tour. Then the sets H_{ij} are simply $H_{i_1j_1},\ldots$, $H_{i_{\lambda-1}j_{\lambda-1}}$

Deletion Candidates

The value of a move associated with the removal of a chain H_{ij} is measured by the ratio of saved distance over lost profit, and is computed as

$$\bar{g}(H_{ij}) = \frac{l(T) - l(T \setminus H_{ij})}{\sum_{v_k \in H_{ij}} p_k}$$

Insertion or Deletion

Compare the results of the insertion and deletion then apply the best one. If the best move is deletion, then declare all vertices of deletion tabu for θ iteration where θ is a random number between (5,25)

Tour Improvement

If the iteration count is the multiple of 5, apply 2-opt.

2-opt

The main idea behind it is to take a route that crosses over itself and reorder it so that it does not. To speed up the 2-opt while comparing the new route to the old one, we just compare the added and removed edges length. Also, at each iteration of 2-opt we are searching for the best update move and apply it.

Best Solution Update

If newly generated solution has a better objective than the incumbent solution then apply 3-opt to the newly generated solution to improve the tour quality and make it the incumbent solution.

3-opt

3-opt analysis involves deleting 3 edges in a tour, reconnecting the tour in all other possible ways, and then evaluating each reconnection method to find the optimum one. This process is then repeated for all different set of 3 connections. To speed up 3-opt process, unlike our 2-opt implementation, we don't search for the best move in all the edge pairs, but we take the first move that results in a better tour. To make things stochastic, we select edges randomly.

Shuffle to Reset

If there hasn't been an improvement in γ iteration, then assign incumbent solution to the current solution and shuffle the route. I chose γ as 10000. Also, it resets the tabu list. This is a magic reset step which enables different solutions by shuffling the route. Shuffling is important because, tabu search itself and especially 2-opt and 3-opt methods are not guaranteed to find optimal paths. For example, the order of given nodes of a route may change the final 2-opt route result. By shuffling, we are increasing chances to find different routes which may have better objective values.

Failed Extensions

We have also tried some other extensions to the tabu search to improve the results. However, some of them just failed.

Intermediate-Term Memory

Intensification rules intended to bias the search towards promising areas of the search space. If a node is always showing up in the solution then we made this node forbid to be in tabu list.

Long-Term Memory

Diversification rules that drive the search into new regions. If a node is always showing up in the solution, then we made this node forbid to enter the solutions for a longer time than a normal tabu.

Shuffling

After completion of tabu search heuristic, to try to make an improvement in the route length, we shuffle and optimize the route several times.

Extra - Visualization of 2-Opt

I have added an extra property to my 2-opt function. You can track the edge changes in the 2-opt algorithm visually via sending the function some parameters.

Sample Output

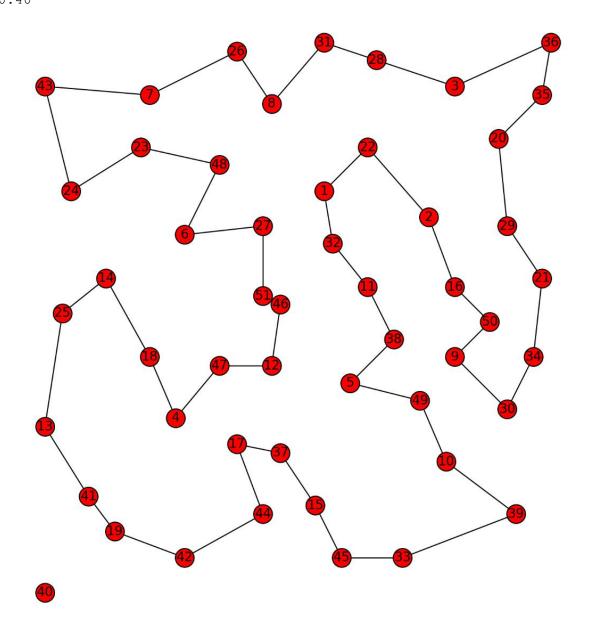
```
Instance:
ei151-HP

Best Objective Value:
704.73

Number of Customers Visited (Depot Excluded):
49

Sequence of Customers Visited:
[1, 32, 11, 38, 5, 49, 10, 39, 33, 45, 15, 37, 17, 44, 42, 19, 41, 13, 25, 14, 18, 4, 47, 12, 46, 51, 27, 6, 48, 23, 24, 43, 7, 26, 8, 31, 28, 3, 36, 35, 20, 29, 21, 34, 30, 9, 50, 16, 2, 22, 1]

CPU Time (s):
40.40
```



Results

Instance	Best	No. of	Sequence of customers visited	CPU Time
	Objective Value	customers visited		(s)
eil51-LP	49.98	19	[1, 32, 11, 38, 9, 16, 50, 34, 30, 10, 33, 45, 15, 44, 17, 4, 18, 14, 6, 48, 1]	27.34
eil51-HP	704.73	49	[1, 32, 11, 38, 5, 49, 10, 39, 33, 45, 15, 37, 17, 44, 42, 19, 41, 13, 25, 14, 18, 4, 47, 12, 46, 51, 27, 6, 48, 23, 24, 43, 7, 26, 8, 31, 28, 3, 36, 35, 20, 29, 21, 34, 30, 9, 50, 16, 2, 22, 1]	39.85
eil76-LP	160.05	53	[1, 43, 42, 41, 56, 23, 49, 16, 3, 44, 32, 9, 39, 72, 58, 10, 38, 11, 53, 14, 19, 35, 7, 8, 46, 34, 52, 27, 13, 57, 15, 37, 20, 70, 60, 71, 69, 36, 47, 21, 74, 30, 68, 75, 76, 67, 26, 12, 40, 17, 51, 6, 33, 73, 1]	98.92
eil76-HP	1241.40	74	[1, 33, 63, 16, 3, 44, 32, 40, 12, 17, 51, 6, 68, 4, 75, 76, 26, 67, 34, 46, 52, 27, 45, 29, 48, 30, 2, 74, 28, 61, 21, 47, 36, 69, 71, 60, 70, 20, 37, 5, 15, 57, 13, 54, 19, 8, 35, 7, 53, 14, 59, 11, 66, 65, 38, 10, 58, 72, 39, 9, 25, 55, 50, 18, 24, 49, 23, 56, 41, 43, 42, 64, 22, 62, 73, 1]	101.60
eil101-LP	262.15	75	[1, 69, 31, 88, 62, 10, 32, 90, 63, 11, 19, 47, 48, 82, 7, 18, 83, 60, 5, 84, 61, 16, 86, 44, 91, 100, 85, 93, 98, 37, 92, 59, 99, 96, 6, 94, 13, 95, 97, 87, 42, 43, 15, 57, 41, 22, 74, 75, 56, 23, 39, 4, 72, 73, 21, 40, 58, 53, 101, 28, 12, 80, 68, 24, 29, 78, 34, 9, 51, 81, 33, 79, 3, 77, 76, 50, 1]	196.62

eil101-HP	1639.82	98	[1, 69, 27, 101, 53, 28, 26, 12, 80, 68,	227.62
			29, 24, 55, 25, 4, 39, 67, 23, 56, 75,	
			41, 22, 74, 72, 73, 21, 40, 58, 13, 87,	
			57, 15, 43, 42, 14, 44, 38, 86, 16, 61,	
			85, 91, 100, 37, 98, 93, 99, 96, 59, 92,	
			97, 95, 94, 6, 89, 52, 18, 83, 60, 5, 84,	
			17, 45, 46, 8, 82, 7, 48, 19, 47, 36, 49,	
			64, 11, 63, 90, 32, 10, 62, 88, 31, 70,	
			30, 20, 66, 65, 71, 35, 34, 78, 81, 9,	
			51, 33, 79, 3, 77, 76, 50, 1]	

Evaluation of Results and Observations

In this part I just want to make a few comments about results and the general structure of the algorithm.

First, to achieve this version of the code, I have run hundreds of different combinations of parameters and methodologies. I've tried different term memories, shuffling, different tour improvement approaches. But in the end, this layout gave the best results in terms of both objective value and time.

Second, I also run this version of the code several times to test different characteristics of the algorithm, especially I've tried to find the optimal iteration count. Results showed that, with the reset shuffling idea, a better solution can be found at any stage of the iteration. In my test, once it achieved the highest objective value in 100 iterations, in another test, it found a better route at the 9500 iteration. However, tests showed that it is unlikely to find a better solution after 10000 iterations.

My point is that, one can also use 1000 iterations rather than 10000 iterations. (In the original paper that I mentioned at the beginning 1000 iterations was used.) With 1000 one can still achieve very good results but with less confidence. However, time spent will be way more less than 10000. So obviously, like all the heuristics, and other methods here we have a trade off again. I chose 10000 iterations, because after some optimization in the code, it gave results in a tolerable time with a better confidence interval.

Code

Importing required libraries

```
iimport numpy as np
import math
import time
import random
import itertools
import queue
import pandas as pd
from IPython.display import display, Markdown
import networkx as nx
import matplotlib.pyplot as plt
```

Reading data

```
# HP or LP
dataset = "HP"
#51,76 or 101
N = 51
filename = "dataset-" + dataset + ".xls"
df = pd.read_excel(filename, sheetname = "eil"+str(N), header = None, index_col = 0)
df.columns = ['x', 'y', 'prof']
display(df[0:10])
distances = [-1]
prof = [-1]
for lab, row in df.iterrows():
  tempDist = [-1]
  prof.append(row['prof'])
  for lab2, row2 in df.iterrows():
    dist = math.sqrt( math.pow(row['x']-row2['x'], 2) + math.pow(row['y']-row2['y'], 2) )
    tempDist.append(dist)
  distances.append(tempDist)
# dff holds the main data as given from the xls
# Started the indices from 1
dff = [[0,0,0]]
for lab, row in df.iterrows():
  dff.append([row['x'],row['y'],row['prof']])
```

Tabu Search Functions

```
def calculateObj(route):
```

```
if len(route) == 0:
    return -99999999
  objVal = 0
  for i in range(1,len(route)):
    objVal = objVal + dff[route[i]][2] - distances[route[i-1]][route[i]]
  return objVal
def calculateTour(route):
  objVal = 0
  for i in range(1,len(route)):
    objVal = objVal + distances[route[i-1]][route[i]]
  return objVal
def updateGraph(G, old_route, route, se, visualize):
  G.remove edge(old route[se[0]-1], old route[se[0]])
  G.remove_edge(old_route[se[1]], old_route[se[1]+1])
  G.add_edge(route[se[0]-1], route[se[0]])
  G.add_edge(route[se[1]], route[se[1]+1])
  if visualize:
    nx.draw(G,pos,with_labels = True)
    plt.show()
    print(str(old_route[se[0]-1]) + ',' + str(old_route[se[0]]) + ' - ' + str(old_route[se[1]]) + ',' +
str(old_route[se[1]+1]))
  return G
def twoOpt(route, G=None, visualize = False):
  if G!= None:
    pos=nx.get node attributes(G,'pos')
  if visualize and G!= None:
    nx.draw(G,pos,with_labels = True)
    plt.show()
  se = (0,0)
  xx = 0
  while(True):
    xx = xx + 1
    temp_route = list(route)
    old route = list(route)
    route distance = -999999999
    for i in range(1, len(route)-2):
       for j in range(i+1, len(route)-1):
         new route = route[:i] + list(reversed(route[i:j+1])) + route[j+1:]
         diff distance = distances[route[i-1]][route[i]] + distances[route[j]][route[j+1]]
         diff_distance
                                 diff_distance
                                                        distances[new_route[i-1]][new_route[i]]
distances[new_route[j]][new_route[j+1]]
         if diff_distance > route_distance:
```

```
temp_route = list(new_route)
           route_distance = diff_distance
           se = (i,j)
    if route_distance > 0.01:
      route = list(temp_route)
      if G != None:
        G = updateGraph(G, old_route, route, se, visualize)
    else:
      break
  return route, G
def threeOptSwap(route, i, j, k):
  bestRoute = list(route)
  best diff = 0
  a = i
  b = j+1
  c = k+2
  nRoute = route[:a] + list(reversed(route[a:b])) + list(reversed(route[b:c])) + route[c:]
  diff = distances[route[a-1]][route[a]] + distances[route[b-1]][route[b]] + distances[route[c-1]]
1]][route[c]]
  diff
         = diff
                         distances[route[a-1]][route[b-1]] - distances[route[a]][route[c-1]]
distances[route[b]][route[c]]
  if diff > best diff:
    best diff = diff
    bestRoute = list(nRoute)
  nRoute = route[:a] + route[b:c] + route[a:b] + route[c:]
  diff = distances[route[a-1]][route[a]] + distances[route[b-1]][route[b]] + distances[route[c-
1]][route[c]]
  diff = diff - distances[route[a-1]][route[b]] - distances[route[c-1]][route[a]] - distances[route[b-
1]][route[c]]
  if diff > best diff:
    best diff = diff
    bestRoute = list(nRoute)
  nRoute = route[:a] + route[b:c] + list(reversed(route[a:b])) + route[c:]
  diff = distances[route[a-1]][route[a]] + distances[route[b-1]][route[b]] + distances[route[c-
1]][route[c]]
  diff
                         distances[route[a-1]][route[b]] -
                                                                distances[route[c-1]][route[b-1]]
              diff -
distances[route[a]][route[c]]
  if diff > best_diff:
    best diff = diff
    bestRoute = list(nRoute)
  nRoute = route[:a] + list(reversed(route[b:c])) + route[a:b] + route[c:]
  diff = distances[route[a-1]][route[a]] + distances[route[b-1]][route[b]] + distances[route[c-1]]
1]][route[c]]
  diff = diff - distances[route[a-1]][route[c-1]] - distances[route[b]][route[a]] - distances[route[b-
1]][route[c]]
```

```
if diff > best_diff:
    best_diff = diff
    bestRoute = list(nRoute)
  return bestRoute, best_diff
def threeOpt(route):
  xx = 0
  while(True):
    xx += 1
    temp route = list(route)
    old_route = list(route)
    best_diff = 0.01
    brk = False
    li = list(range(1, len(route)-2))
    random.shuffle(li)
    for i in li:
       lj = list(range(i, len(route)-2))
      random.shuffle(lj)
      for j in lj:
         lk = list(range(j, len(route)-2))
         random.shuffle(lk)
         for k in lk:
           new_route, new_diff = threeOptSwap(route, i, j, k)
           if new_diff > best_diff:
             temp_route = list(new_route)
             best_diff = new_diff
             brk = True
             break
         if brk:
           break
      if brk:
         break
    if not brk:
       break
    if best_diff > 0.01:
       route = list(temp_route)
    else:
       break
  return route
def initialization():
  "Construction Heuristic"
  best_objs = []
  best_routes = []
  for i in [int(N/2)]:
    local_route = []
    for t in range(5):
       route = [1,1]
```

```
for j in range(i):
         min_obj = 99999999
         k = random.randint(0, len(route)-2)
         temp route = list(route)
         for lab in range(1,N+1):
           if lab not in route:
             new_route = route[:k+1] + [lab] + route[k+1:]
             diff obj
                                 (distances[route[k]][lab]
                                                                      distances[lab][route[k+1]]
                          =
distances[route[k]][route[k+1]]) / prof[lab]
             if diff obj < min obj:
                temp_route = list(new_route)
                min obj = diff obj
         route = list(temp_route)
      temp_route = twoOpt(route)[0]
      temp_obj = calculateObj(temp_route)
      if temp_obj > local_obj:
         local_obj = temp_obj
         local_route = list(temp_route)
    best_routes.append(local_route)
    best objs.append(local obj)
  route = list(best_routes[0])
  rat = 0
  for i in range(len(best_routes)):
    if best_objs[i]/len(best_routes[i]) > rat:
       rat = best_objs[i]/len(best_routes[i])
       route = list(best_routes[i])
  return route
def dispersionIndex(cluster):
  if len(cluster) == 1:
    return 0
  else:
    sm = 0
    for c1 in cluster:
      for c2 in cluster:
         sm = sm + distances[c1][c2]
    return sm / (len(cluster)*(len(cluster)-1))
def proximityMeasure(cluster1, cluster2):
  sm = 0
  for c1 in cluster1:
    for c2 in cluster2:
       sm = sm + distances[c1][c2]
  return (2/(len(cluster1)*len(cluster2)))*sm - dispersionIndex(cluster1) - dispersionIndex(cluster2)
def insertionCandidates():
```

```
candidates = []
      rList = [1, int(N/2), int(2*N/3), int(3*N/4), int(4*N/5), int(5*N/6), int(6*N/7), int(7*N/8), int(8*N/9), int(8*
int(9*N/10)]
       Pr = []
       Pr = [[x] \text{ for } x \text{ in range}(2,N+1)]
      candidates.append(list(Pr))
      for r in range(2,N):
             minProx = 999999999
             minProxInd = []
             for i in range(len(Pr)):
                   for j in range(i+1, len(Pr)):
                          pM = proximityMeasure(Pr[i], Pr[j])
                          if pM < minProx:
                                minProx = pM
                                minProxInd = [i, j]
             Pr.append(Pr[minProxInd[0]]+Pr[minProxInd[1]])
              del(Pr[minProxInd[1]])
             del(Pr[minProxInd[0]])
             if r in rList:
                   candidates.append(list(Pr))
       return candidates
def deletionCandidates(route):
      candidates = []
      edges = []
       K = random.randint(2,int(max(4,len(route))/2))
      for i in range(len(route)-1):
             edges.append([distances[route[i]][route[i+1]], i, i+1])
      edges = list(reversed(sorted(edges)))[:K]
       edges.sort(key=lambda x: x[1])
      for i in range(K-1):
             tempList = []
             for j in range(edges[i][2], edges[i+1][1]+1):
                   tempList.append(route[j])
             candidates.append(tempList)
       return candidates
def findBestInsertionCandidate(route, tabuList, insCandidates):
       bestInsCandidate = []
       bestInsObj = -99999999
```

```
for iC in insCandidates:
    profitSum = 0
    gCenter = [0,0]
    for c in iC:
      if c not in route and c not in tabuList:
         gCenter[0] = gCenter[0] + dff[c][0]/len(iC)
         gCenter[1] = gCenter[1] + dff[c][1]/len(iC)
         profitSum = profitSum + dff[c][2]
    minDist = 99999999
    for j in range(len(route)-1):
       distAdd1 = calculateDist(dff[route[j]][0],dff[route[j]][1],gCenter[0],gCenter[1])
       distAdd2 = calculateDist(gCenter[0],gCenter[1],dff[route[j+1]][0],dff[route[j+1]][1])
       distRem = calculateDist(dff[route[j]][0],dff[route[j]][1],dff[route[j+1]][0],dff[route[j+1]][1])
       dist = distAdd1 + distAdd2 - distRem
       if dist < minDist:
         minDist = dist
    if profitSum/minDist > bestInsObj:
       bestInsObj = profitSum/minDist
       bestInsCandidate = list(iC)
  return bestInsCandidate
def calculateDist(x1,y1,x2,y2):
  return math.sqrt( math.pow(x1-x2, 2) + math.pow(y1-y2, 2) )
Main Solver
# Iteration Count
ITER = 10000
# Start the timer
t1 = time.clock()
# Create the initial route
route = initialization()
# Determine all possible insertion partitions
insCandidatesAll = insertionCandidates()
tabuList = {}
solutionIndex = [0]
bestRoute = list(route)
bestObj = calculateObj(bestRoute)
# Start tabu search
for i in range(ITER):
```

```
# Choose one insertion partition ramdompy
  insCandidates = list(insCandidatesAll[random.randint(0,len(insCandidatesAll)-1)])
  # Determine deletion candidates
  if len(route) < 3:
    delCandidates = []
  else:
    delCandidates = deletionCandidates(route)
  candidateRoute = []
  tabuAddition = []
  # Find best insertion candidate from the selected partition
  bestInsCandidate = findBestInsertionCandidate(route, tabuList, insCandidates)
  # Calculate the gain of inserting the insertion candidate to the route
  insertedRoute = list(route)
  profitSum = 0
  distSum = 0
  random.shuffle(bestInsCandidate)
  for c in bestInsCandidate:
    if c not in insertedRoute and c not in tabuList:
      profitSum = profitSum + dff[c][2]
      minDist = 99999999
      temp_route = list(insertedRoute)
      for j in range(len(insertedRoute)-1):
        new_route = insertedRoute[:j+1] + [c] + insertedRoute[j+1:]
                         distances[insertedRoute[j]][c] +
                                                               distances[c][insertedRoute[j+1]]
distances[insertedRoute[j]][insertedRoute[j+1]]
        if diffDist < minDist:
          temp_route = list(new_route)
           minDist = diffDist
      insertedRoute = list(temp route)
      distSum = distSum + minDist
  if distSum == 0:
    distSum = 99999999
  insertedObj = profitSum / distSum
  # Choose the best deletion candidate from the selected ones, then calculate its gain
  deletedRoute = list(route)
  maxDeletedObj = -99999999
  for dC in delCandidates:
    tempRoute = list(route)
    profitSum = 0
    distSum = 0
    for c in dC:
      if c in tempRoute:
        cPrev = tempRoute[tempRoute.index(c)-1]
        cNext = tempRoute[tempRoute.index(c)+1]
```

```
profitSum = profitSum + dff[c][2]
      distSum = distances[cPrev][c] + distances[c][cNext] - distances[cPrev][cNext]
      tempRoute.remove(c)
  if profitSum != 0 and distSum/profitSum > maxDeletedObj:
    maxDeletedObj = distSum/profitSum
    deletedRoute = list(tempRoute)
    tabuAddition = list(dC)
deletedObj = maxDeletedObj
# Compare the insertion and deletion gains, and apply the better one
if insertedObj > deletedObj:
  candidateRoute = list(insertedRoute)
  chosen = ['I', len(insertedRoute)-len(route)]
else:
  candidateRoute = list(deletedRoute)
  chosen = ['D', len(route)-len(deletedRoute)]
# Update the tabu list
for key, value in list(tabuList.items()):
  tabuList[key] = tabuList[key] - 1
  if tabuList[key] == 0:
    del(tabuList[key])
# If deletion action is performed then add the chosen deletion candidates to the tabu list.
if chosen[0] == 'D':
  for tA in tabuAddition:
    if tA in route:
      tabuList[tA] = random.randint(5,25)
route = list(candidateRoute)
# Improve the route
if i % 5 == 0:
  route = twoOpt(route)[0]
# Best solution update
if calculateObj(route) > bestObj:
  solutionIndex.append(i)
  route = threeOpt(route)
  bestRoute = list(route)
  bestObj = calculateObj(route)
# Shuffle to Reset
if i - solutionIndex[-1] >= 1000:
  tabuList.clear()
  tempRoute = bestRoute[1:-1]
  random.shuffle(tempRoute)
  tempRoute = [1] + tempRoute + [1]
  route = list(tempRoute)
  solutionIndex.append(i)
```

```
# Stop the timer
t2 = time.clock()
```

Results

```
print("Instance: ")
print("eil" + str(N) + "-" + str(dataset))
print()
print("Best Objective Value:")
print("%.2f" %calculateObj(bestRoute))
print()
print("Number of Customers Visited (Depot Excluded):")
print(len(bestRoute)-2)
print()
print("Sequence of Customers Visited:")
print(bestRoute)
print()
print("CPU Time (s):")
timePassed = (t2-t1)
print("%.2f" %timePassed)
%config InlineBackend.figure_format = 'retina'
plt.figure(figsize=(9,9))
G=nx.Graph()
for lab, row in df.iterrows():
  G.add_node(lab, pos = (row['x'], row['y']))
for i in range(1,len(bestRoute)):
  G.add_edge(bestRoute[i-1], bestRoute[i])
pos=nx.get_node_attributes(G,'pos')
nx.draw(G,pos,with_labels = True)
plt.show()
```