

1D1020 2018-09-17 #7

## Quicksort (Tony Hoare)

- Shuffle
- Partition so that for some value  $j$ 
  - The element  $a[j]$  is correctly placed
  - No element greater than  $a[j]$  is to the left of  $j$
  - No element less than  $a[j]$  is to the right of  $j$
- Sort all sub-arrays recursively

Hoare invented Quicksort to translate from Russian to English.

"There are two ways of constructing a software design: One way is to make it so simple that there are obviously no deficiencies, and the other way is to make it so complicated that there are no obvious deficiencies. The first method is far more difficult."

Quicksort Partitioning { Repeat until the indices  $i$  and  $j$  cross

- Scan by  $i$  from left to right while  $(a[i] < a[j])$
- Scan by  $j$  from right to left while  $(a[j] > a[i])$
- Swap  $a[i]$  and  $a[j]$

Not stable. Can be made stable with extra array, but then merge sort is probably better anyway.

Can be 30% faster than merge sort. Worst case:  $O(N^2)$   
Harder to implement. Extremely improbable  
Average case:  $O(N \log N)$

Randomized divide-and-conquer algorithm.  
In-place logarithmic stack space.

Improvements:

- Insertion sort on  $\sim 10$  elements.
- Partitioning element: Median-of-three

Problem: Duplicate keys

Stop scanning when equal keys are found

3-way partitioning (Dijkstra)

Dijkstra's method not good with few equal keys

Sedgewick-Bentley

3-way partitioning can make sorting possible in almost linear time.

## When to choose different sorting algorithms?

Things to take into account:

- Stable?
- Parallel?
- In-place?
- Deterministic?
- Duplicate keys?
- Different types of keys?
- Linked list or array?
- Large or small elements? (copy/swap)
- Order of the input?
- Performance guarantees?

"System sort" is often fast enough for common cases

## Priority Queues

Often: low value  $\rightarrow$  high priority

Enqueue: Insert element

Dequeue: Remove and return element with highest priority

Stability: Preserve FIFO-order among elements with equal priority

Lazy: Priority queues need only be partially sorted

Linked list implementation

- Keep elements ordered in descending priority order
- Enqueue: Sort the new element into correct position
- Dequeue: Remove the first (highest priority) element.

Binary heap - A balanced binary tree with heap property

- All nodes have higher or equal priority than any of its two children
- The highest priority element is found in the root.

Array heap representation 

1	2	3	4
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 ...

Bottom-up reheapify (swim)

Usage: Enqueue a new element to the first empty place in the lowest level, restore heap order by bottom-up reheapify.

Swim: If the child has higher priority than its parent - swap them. Continue until no swap or root is reached.

Top-down reheapify (sink)

Usage: Dequeue - remove the top element and replace it with the last (rightmost) element in the bottom level - do a top-down reheapify (sink)

Sink: If the node has lower priority than any child - swap it with highest priority child. Continue until no swap or no children.

Heap: Simple to implement, Theoretical  $O(\log N)$  performance for enqueue & dequeue

Heap sort

Bad cache performance - store elements in different order to solve

## Special purpose priority queues

Process/thread scheduling in OS

- Fixed number of priorities
- Array of linked lists - one per priority.
- Enqueue and Dequeue is  $O(1)$ .

Calendar queue

- Array with  $\sim 365$  elements - one for each day
- One linked list per day. Modulo operations for different years

#days  $\frac{N}{2}$ , daylength  $\sim 2 \Rightarrow O(1)$

Fast general purpose priority queue

- Splay tree, Tarjan & Sleator

Priority queues are also used for many other algorithms.