

Directed graphs (Digraphs)

- The edges has direction: outdegree and indegree
- Can be seen as intersections directed path and one way streets for example. directed cycle
- Has many applications

Problems:

- $s \rightarrow t$ path, shortest $s \rightarrow t$ path
- directed cycle, topological sort
- strong connectivity
- transitive closure
- PageRank

API:

```

Digraph(int V)
Digraph(In in)
addEdge(int v, int w)
Iterable<Integer> adj(int v)
int V()
int E()
Digraph reverse()
String toString()
    
```

Adjacency lists

- vertex indexed array of lists
- Real world problems tend to be sparse

Reachability problem: which vertices can be reached from v ?

DFS (Directed)

To visit a vertex v :

- Mark v as visited
- Recursively visit all unmarked vertices pointing from v .

Every program is a digraph

- Dead-code elimination
- Infinite loop detection

Mark-sweep garbage collection

BFS is same as for undirected graph

- Put s (source) onto FIFO queue, and mark s as visited.
- Repeat until the queue is empty:
 - remove the least recently added vertex v
 - for each unmarked vertex pointing from v :
 - add to queue and mark as visited.

Finds the shortest directed path.

Multiple-source shortest path

- Enqueue \leftarrow

Topological sort (ex. Precedence scheduling)

DAG - Directed acyclic graph

If cycles exist, topological sort is not possible (\Leftrightarrow)

Run depth-first search

- Return vertices in reverse postorder

Non-connected vertices may be placed in any order

- First vertex in postorder has outdegree 0.
- Second-to-last vertex in postorder can only point to last vertex.
- ...

Proof in presentation

DFS visits each vertex exactly once. The order can be important.

- Preorder
- Postorder
- Reverse postorder

Strongly connected components.

- Two vertices are strongly connected if a directed path exists in both way. This is an equivalence relation.
- A strong component is a maximal subset of strongly-connected vertices.

Kosaraju-Sharir algorithm

Strong components in G is same as in G^R

Kernel DAG: Contract each strong component into a single vertex.

- Compute topological order (reverse postorder) in kernel DAG
- Run DFS, considering vertices in reverse topological order

Kosaraju-Sharir algorithm computes the strong components of a digraph in time proportional to $E+V$.