Undirected graphs

(Set of vertices connected pairwise by edges

Path - sequence of vertices connected by edges

Sparse graph: Eval where a issmall.

Cycle - first vertex = last vertex

Discrete Mouth

Connected
Tree: a sychic

Spanning tree

Forest

Spounding forest

Bipartite graph

Graph AP1 (simple)

vertex representation: int O-(V-1), symbol table

Graph (V) / Graph (in)

add Edge (u, w)

Herable adj(v)

int V()

int E()

list of edges: linked list/array adjacency lists; vertex indexed array of lists.

Real world graphs tend to be sparse

Depth-first search

Maze graph: Vertex = intersection, Edge = passage Goal. Explore every intersection

Algorithm

· Unroll a ball of string behind you

· Mark each visited intersection and each visited passage

· Retrace steps when no unvisited options.

To visit a vertex v

• Mark v as visited

* Recursively visit all unmarked vertices adjacent to v. Does not find the shortest path

Désign pattern: Decouple graph borta type from graph processing.

Poths API

Paths (Graph G, int S)

boolean has Path To (v)

Herable path To(v)

(Integer)

DFS marks all vertices connected to 8 in time proportional to the sum of their degrees

After DFS, can check if vertex v is connected to s in constant time and can find v-s path (if exists) in ~

Application: flood-fill

Union find can be faster (because of caching)
Depth-first search: Constant time per query.

Breadth-first search

Put s into a FIFO queue, and mark s as visited Repeat until queue is empty

· Remove vertex v from queue

· Add to queue all unmarked vertices adjacent to u and mark them.

Finds the shortest paths

Connectivity queries

v and w are connected if there is a path between them Connected component: Maximal Set of connected vertices Given connected components, we can answer connectivity queries in constant time.

Find connected component

Initialize all vertices v as unmarked.
For each unmarked vertex v, run DFS to identify all vertices discovered as part of the same component

Challenges!

Is a graph bipartite?
Finding cycles
Euler cycle
Hamilton cycle, NP-complete
Isomorphism, Open problem
Can a graph be represented plane. Difficult