

## Problem Set 7, Part I

### Problem 1: Working with stacks and queues

1-1)

```
public static void doubleAllStack(Stack<Object> stack, Object item)
{
    if(item==null){
        throw new NullPointerException();
    }
    Stack<Object> stack1 = new Stack<Object>();
    while(!stack.isEmpty()){
        Object top=stack.pop();
        if(item.equals(top)){
            stack1.push(top);
        }
        stack1.push(top);
    }
    while(!stack1.isEmpty()){
        stack.push(stack1.pop());
    }
}
```

1-2)

```
public static void doubleAllQueue(Queue<Object> stack, Object item)
{
    if(item==null){
        throw new NullPointerException();
    }
    Queue<Object> q= new Queue<Object>();
    while(!stack.isEmpty()){
        Object top=stack.remove();
        if(item.equals(top)){
            q.insert(top);
        }
        q.insert(top);
    }
    while(!q.isEmpty()){
        stack.insert(q.remove());
    }
}
```

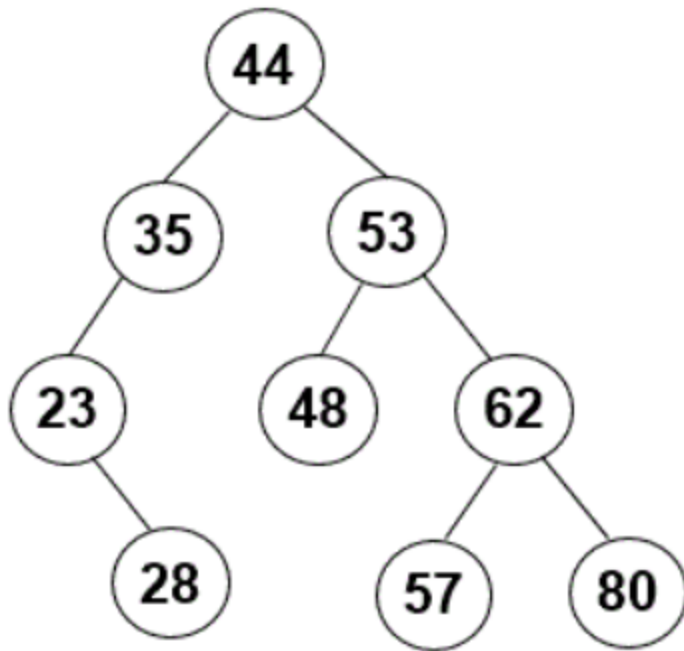
## Problem 2: Using a queue to search a stack

```
boolean found=false;
//checks if item I is in the stack by removing all of the items in the
//stack and placing them in the queue Q
while(!S.isEmpty()){
    Object item = S.pop();
    //if the item is found, the boolean variable found is set to true
    if(I.equals(item)){
        found=true;
    }
    Q.insert(item);
}
//adds the items from the queue back into the stack
while(!Q.isEmpty()){
    S.push(Q.remove());
}

//the items are in backwards order so we add the items from the stack
//back into the queue
while(!S.isEmpty()){
    Q.insert(S.pop());
}

//finally, we add the items from the queue back into the stack and
//they are in the correct order
while(!Q.isEmpty()){
    S.push(Q.remove());
}

//returns the boolean variable found, indicating that item I was found
return found;
```



**Problem 3: Binary tree basics**

3-1) Height: 3

3-2) Leaf Nodes: 4

Interior Nodes: 5

3-3) 44 35 23 28 53 48 62 57 80

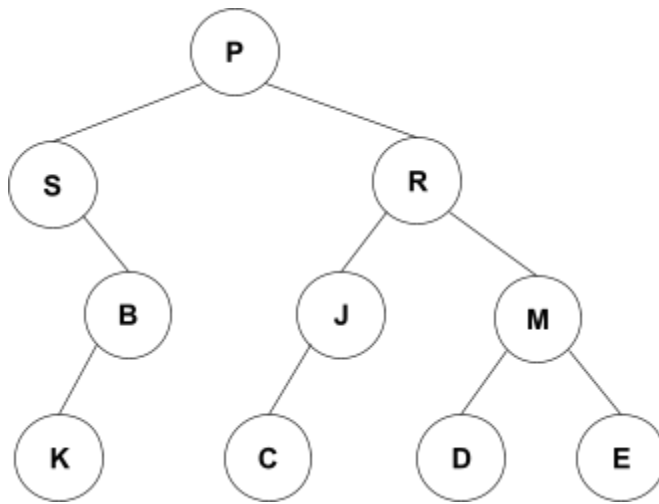
3-4) 28 23 35 48 57 80 62 53 44

3-5) 44 35 53 23 48 62 28 57 80

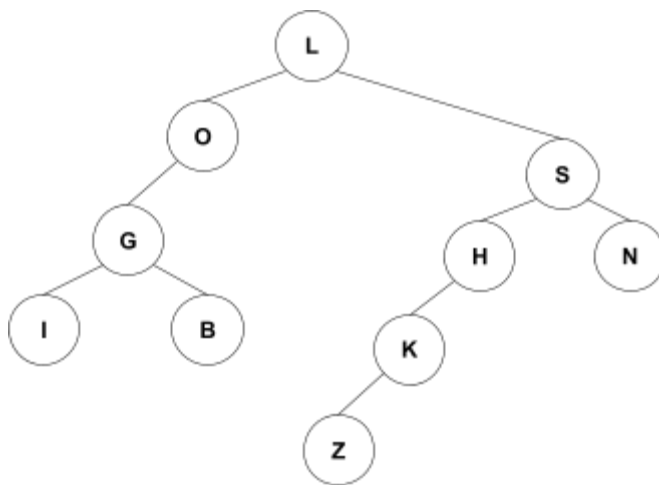
3-6) Yes, this is a search tree as for each node the left subtree is less than the node and the right subtree is greater than the node

3-7) Yes, this is a balanced search tree as for each node its subtree has either the same height or its height - 1 .

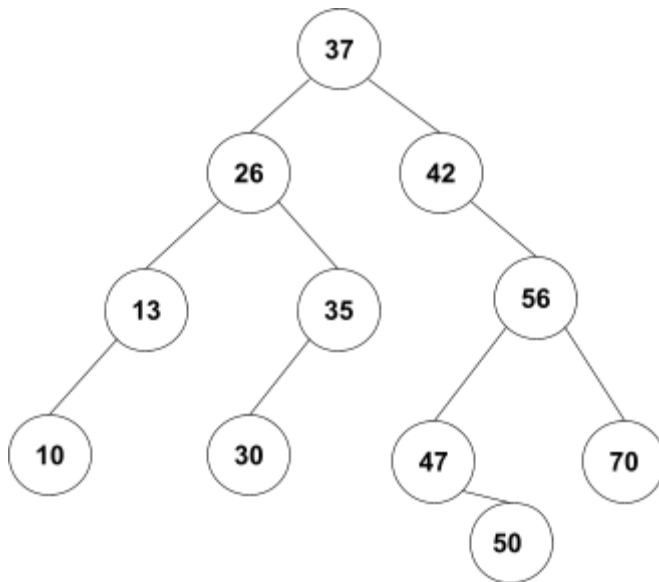
**Problem 4: Tree traversal puzzles**  
**4-1)**



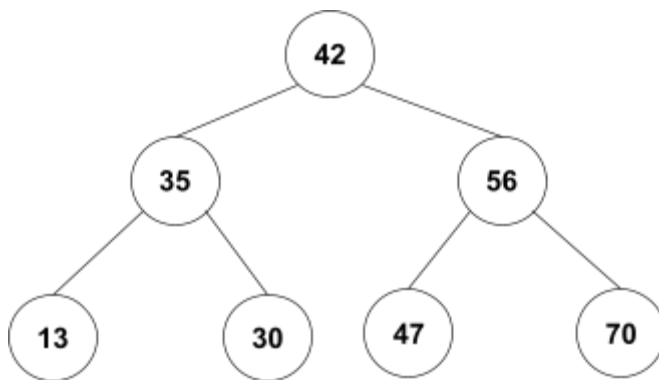
**4-2)**



**Problem 5: Binary search trees**  
**5-1)**



**5-2)**



## Problem 6: Counting keys below a threshold

### 6-1)

Best case:  $O(1)$  The best case occurs when the tree only consists of a root node, in this case the root key will only be compared to  $t$  and therefore no recursive calls will be made. There would only be a call to the `numSmallerInTree` method that would check if the root key is greater than or less than  $t$ . This does not traverse the tree therefore in the best case it would be  $O(1)$ .

Worst case:  $O(n)$  In the worst case the entire tree has to be searched and it is unbalanced. If it is unbalanced this gives a worst case time complexity of  $O(n)$  to search it, as the entire tree has to be searched. The Tree contains  $n$  items and you must process all  $n$  of the nodes, performing  $O(1)$  operations on each, which gives  $O(n)$  time complexity.

### 6-2)

```
private static int numSmallerInTree(Node root, int t) {
    int count=0;
    if ( root.left != null && root.key < t) {
        count += numSmallerInTree(root.left, t);
    }
    if (root.right!= null && root.key< t) {
        count += numSmallerInTree(root.right, t);
    }
    if(root.left!=null && root.key>t){
        count += numSmallerInTree(root.left, t);
    }

    if (root.key < t) {
        return 1 + count;
    } else {
        return count;
    }
}
```

**6-3)** Best Case:  $O(1)$ , if the tree only consists of a root node, it does not have to traverse the tree therefore it has a time efficiency of  $O(1)$

Worst Case:  $O(n)$ , This method has an  $O(n)$  time efficiency as if every node was less than  $t$  it would search the whole tree giving an  $O(n)$  efficiency.

**Problem 7: Balanced search trees**

