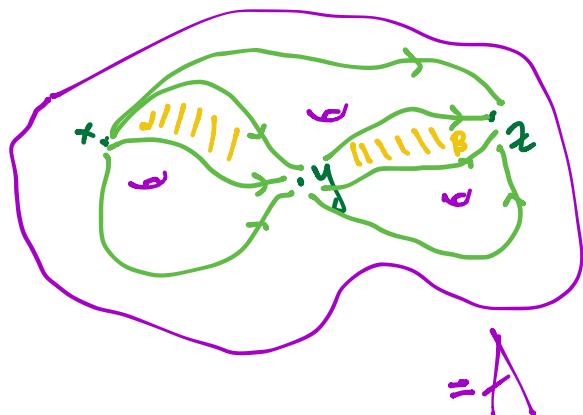


CONTRACTIBILITY as UNIQUENESS

The standard technique used to distinguish your favorite space A from other spaces is to compute an algebraic invariant of the space.



The "algebra of paths" in a space is described in increasing precision by

- the fundamental group $\pi_1(A, x)$ of loops in A based at x up to homotopy
- the fundamental groupoid $\Pi_1 A$ of paths in A up to homotopy
- the fundamental ∞ -groupoid $\mathrm{Tot} \mathcal{A}$ of paths in A

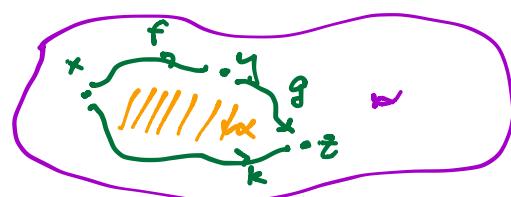
The A has

- points of A as objects
- paths in A as 1-dimensional arrows
- homotopies between paths in A as 2-dimensional arrows
- homotopies between homotopies between paths in A as 3-dimensional arrows, and so on...

Q: How do we define the **composite** of two paths? **A:** We don't!

Instead of a composition operation, composites of paths are **witnessed** by homotopies:

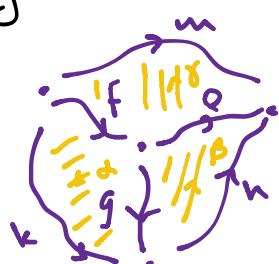
α is a witness that k is a composite of



Q: How unique is this? **Partial A:** Unique enough for **associativity**:

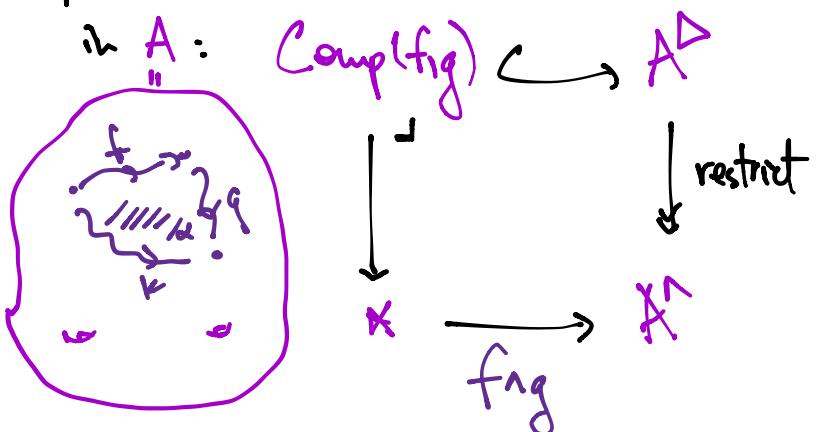
Given composable paths f, g, h and specified homotopies witnessing composition relations, these homotopies **compose**.

More precisely, a 3-arrow expresses a coherence between compositions witnessed by 2-arrows.

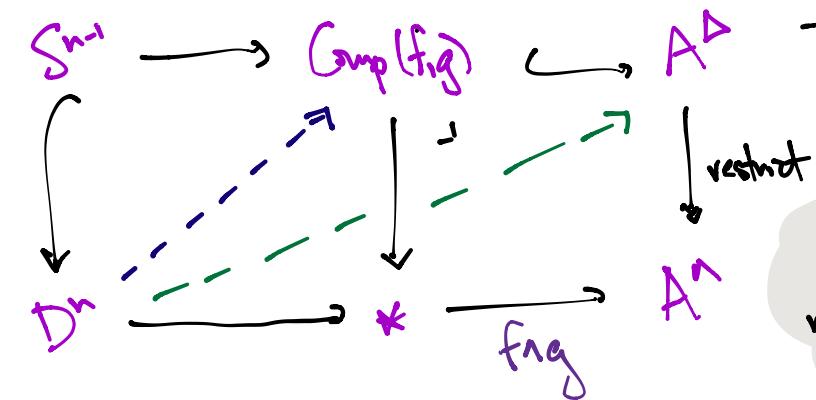


THEOREM: The space of composites of two paths f and g in A is contractible.

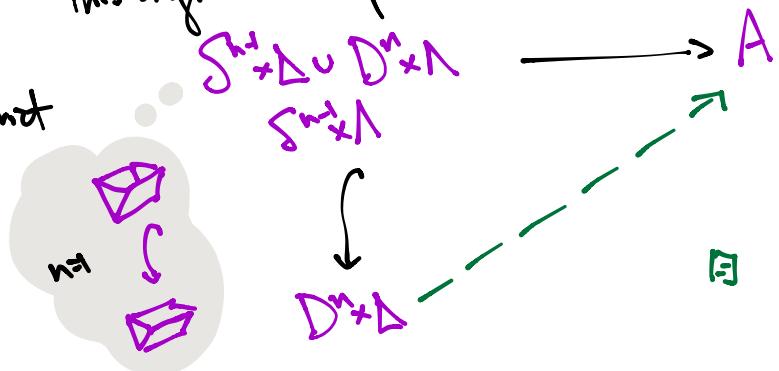
proof: The first step is to define the space of composites of paths f and g



A space is **Contractible** just when any sphere S^{n-1} can be filled in to a disk $D^n \quad \forall n \geq 1$, so we need to define the blue map, for which it suffices to define the green map.



This diagram transposes to:



The extension exists since the inclusion admits a continuous (deformation) retraction. \square

SUMMARY

In a group(oid) any composable pair of arrows has a **unique** composite.

In an ∞ -group(oid) any composable pair of arrows has a **contractible space** of composites.

The ANALOGY ordinary mathematics :: higher mathematics Can be made even tighter.
 Uniqueness :: Contractibility

To say a set C has a unique element means $\exists x \in C, \forall y \in C, x = y$

Here " $x = y$ " is a predicate — a mathematical statement that is either true or false that depends on two free variables $x, y \in C$.

In proof relevant mathematics, we instead interpret " $x = y$ " as the set of all proofs that x equals y (which is empty if x and y are not equal).

Then we can form the set $\sum_{x \in C} \prod_{y \in C} x = y$ inspired by a notational analogy

with the sentence $\exists x \in C, \forall y \in C, x = y$.

In proof relevant mathematics, $\sum_{x \in C} \prod_{y \in C} x = y$ is also a set of proofs, but proofs of what?

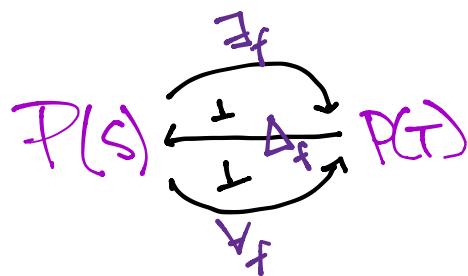
An element of this set is a choice of some element $c \in C$ together with a proof for all $z \in C$, that c equals z . In other words $\sum_{x \in C} \prod_{y \in C} x = y$ is the set of proofs that C contains a unique element (which is empty if this isn't true).

Similarly, we will see that if C is a space, then $\sum_{x \in C} \prod_{y \in C} x = y$ can be interpreted as a code that defines another space. Once again we can interpret this as a space of proofs.... but proofs of what?

To explain this requires a digression to explain the analogy

$$\begin{array}{lll} \text{logic} & \exists & \vee \\ \text{Set theory} & \sum & \prod \end{array}$$

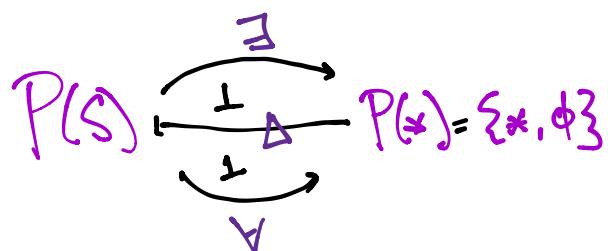
DIGRESSION A set function $S \xrightarrow{f} T$ induces order-preserving functions between their powersets



Δ_f is the inverse image: $B \subseteq T \mapsto \{s \in S \mid f(s) \in B\} \subseteq S$
 \exists_f is the direct image: $A \subseteq S \mapsto \{t \in T \mid \exists s \in S, f(s) = t \wedge s \in A\} \subseteq T$
 \forall_f is the preimage: $A \subseteq S \mapsto \{t \in T \mid \forall s \in S, f(s) = t \Rightarrow s \in A\} \subseteq T$

For the unique function $S \xrightarrow{\star}$ these reduce to

"Quantifiers as adjoints"



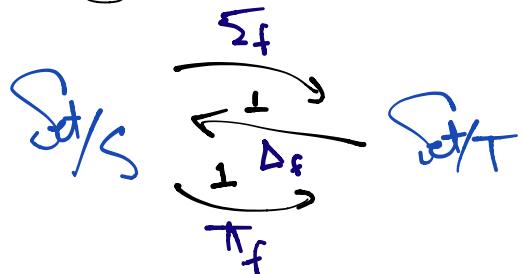
The powerset $P(S)$ can be identified with the set of predicates $p(s)$ with one free variable $s \in S$ — the corresponding sheet is $\{s \in S \mid p(s) \text{ is true}\}$.

If we interpret the two elements of $P(\star)$ by declaring \star means "true" and ϕ means "false"

then \exists is the functor that sends a predicate $p(s)$ to the sentence $\exists s \in S, p(s)$
while \forall is the functor that sends a predicate $p(s)$ to the sentence $\forall s \in S, p(s)$.

In proof relevant mathematics, it's better to replace the poset $P(S)$ by the category Set/S of S -indexed sets. An object in Set/S is a family of sets $\{P(s)\}_{s \in S}$ where $P(s)$ can be thought of as the set of proofs of some predicate $p(s)$ on $s \in S$.

For any function $S \xrightarrow{f} T$ there are functors



Δ_f is substitution: $\{\mathcal{Q}(t)\}_{t \in T} \mapsto \{\mathcal{Q}(f(s))\}_{s \in S}$

\sum_f is sum: $\{P(s)\}_{s \in S} \mapsto \{\sum_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

\prod_f is product: $\{P(s)\}_{s \in S} \mapsto \{\prod_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

This gives a more formal way to understand the set

$$\sum_{x \in C} \prod_{y \in C} x=y$$

Recall " $x=y$ " is the set of proofs that x equals y where $x, y \in C$.

This should be thought of as an indexed set $\{x=y\}_{x,y \in C} \in \text{Set}/C \times C$.

The product along the projection functor $C \times C \xrightarrow{\Pi} C$ gives $\{\prod_{y \in C} x=y\}_{x \in C} \in \text{Set}/C$.

Then the sum along $C \xrightarrow{!} *$ gives $\sum_{x \in C} \prod_{y \in C} x=y \in \text{Set}/*$ = Set.

EXTENDED

ANALOGY

In logic, $\exists + \forall$ are constructions on predicates

In Set theory, $\sum - \prod$ are constructions on indexed sets of proofs

$$\begin{array}{ccc} \exists_f & & \text{P}(S) \xleftarrow{\perp} \Delta_f \xrightarrow{\perp} \text{P}(T) \\ & \swarrow \perp \quad \downarrow \perp \quad \uparrow \perp \quad \searrow \perp & \\ \text{Set}/S & & \text{Set}/T \end{array}$$

Now replace Set by a suitably nice category of topological spaces and continuous functions.

$$\begin{array}{ccc} \Sigma_f & & \text{Space}/S \xleftarrow{\perp} \Delta_f \xrightarrow{\perp} \text{Space}/T \\ & \swarrow \perp \quad \downarrow \perp \quad \uparrow \perp \quad \searrow \perp & \\ & & \prod_f \end{array}$$

For a space C this constructs a new space $\sum_{x \in C} \prod_{y \in C} x=y$,

where an important new idea tells us to interpret " $x=y$ "

as the space of paths in C from the point x to the point y .

Once more, $\sum_{x \in C} \prod_{y \in C} x=y$ is a space of proofs.... but proofs of what?

Q: What is a point in the space $\sum_{x \in C} \prod_{y \in C} x = y$?

A point in $\sum_{x \in C} \prod_{y \in C} x = y$ is given by the choice of a basepoint $c \in C$ together with a point in the space $\prod_{y \in C} c = y$.

This latter point in $\prod_{y \in C} c = y$ is given by a continuous function γ from \mathbb{Z}^C to the space of paths in C from c to y .

Together this data defines a basepoint c and a contracting homotopy γ .

SUMMARY $\sum_{x \in C} \prod_{y \in C} x = y$ is the space of proofs that C is contractible!

This gives a glimpse of the meaning of uniqueness in a new proposed foundation system for mathematics called HOMOTOPY TYPE THEORY.