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## Contractibility as Uniqueness

Hardy Lecture Tour, Manchester

# An analogy



contractibility :: uniqueness

1. Contractibility as Uniqueness
2. Categorifying Uniqueness
3.  $\infty$ -Categorifying Uniqueness



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## Contractibility as Uniqueness

# The algebra of paths

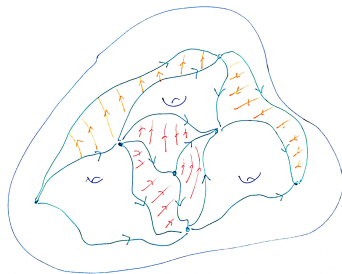
The standard technique used to distinguish your favorite space  $A$  from other spaces is to compute an algebraic invariant of the space.

The “algebra of paths” of a space is described in increasing precision by:

- the fundamental group  $\pi_1(A, x)$  of loops in  $A$  based at  $x$  up to homotopy
- the fundamental groupoid  $\pi_1 A$  of paths in  $A$  up to homotopy
- the fundamental  $\infty$ -groupoid  $\pi_\infty A$  of paths in  $A$

$\pi_\infty A$  has:

- points of  $A$  as objects
- paths of  $A$  as 1-arrows
- paths between paths in  $A$  as 2-arrows
- paths between paths between paths in  $A$  as 3-arrows, and so on ...

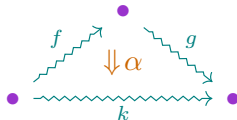


# Witnesses to composition

Q: How do we define the **composite** of two paths?

A: We don't!

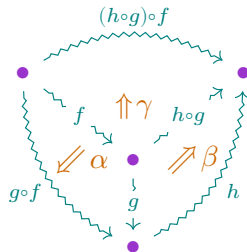
Instead of a composition **operation**, composites of paths are **witnessed** by higher paths.



Q: How unique is path composition?

Partial A: Unique enough for **associativity**.

Given composable paths  $f, g, h$  and specified higher paths  $\alpha, \beta, \gamma$  witnessing composition relations, these higher paths compose. More precisely, a 3-arrow expresses a coherence between compositions witnessed by 2-arrows.



# Homotopical uniqueness of path composition



**Theorem.** The space of composites of two paths  $f$  and  $g$  in  $A$  is contractible.

**Proof:** The space of composites of paths  $f$  and  $g$  in  $A$  is defined by the pullback:

$$\begin{array}{ccc}
 S^{n-1} & \longrightarrow & \text{Comp}(f, g) \hookrightarrow A^\Delta \\
 \downarrow & \nearrow & \downarrow \text{restrict} \\
 D^n & \xrightarrow{\quad} & * \xrightarrow{f \wedge g} A^\Lambda
 \end{array}
 \quad \Leftrightarrow \quad
 \begin{array}{ccc}
 S^{n-1} \times \Delta \cup_{S^{n-1} \times \Lambda} D^n \times \Lambda & \longrightarrow & A \\
 \downarrow & \nearrow & \\
 D^n \times \Delta & &
 \end{array}$$

A space is **contractible** just when any sphere  $S^{n-1}$  can be filled to a disk  $D^n$  for  $n \geq 0$ . This filling problem transposes to an extension problem, and the extension exists since the inclusion admits a continuous deformation retract.  $\square$

# Summary



- In a **group(oid)**, composable arrows have a **unique** composite.
- In a  **$\infty$ -group(oid)**, composable arrows have a **contractible space** of composites.

The analogy

set-based mathematics	::	“higher” mathematics
uniqueness	::	contractibility

can be made even tighter.

**Aim:** Express the classical notion of uniqueness more categorically.



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## Categorifying Uniqueness



# Uniqueness



To say  $C$  has a **unique** element means

$$\exists x \in C, \forall y \in C, x = y$$

Here “ $x = y$ ” is a **predicate** — a mathematical statement that is either true or false, depending on two free variables  $x, y \in C$ .

In **proof-relevant mathematics**, we interpret “ $x = y$ ” as the set of all **proofs** that  $x$  equals  $y$  (which is empty if  $x$  and  $y$  are not equal).

Then we can form the set  $\sum_{x \in C} \prod_{y \in C} x = y$

inspired by a notational analogy with the sentence  $\exists x \in C, \forall y \in C, x = y$ .

The set  $\sum_{x \in C} \prod_{y \in C} x = y$  is also a set of proofs

— but proofs of what?

# Proofs of uniqueness



The set  $\sum_{x \in C} \prod_{y \in C} x = y$  is also a set of proofs

— but proofs of what?

An element of  $\sum_{x \in C} \prod_{y \in C} x = y$  is

- the choice of some element  $c \in C$
- together with a proof, for all  $z \in C$ , that  $c$  equals  $z$ .

Thus  $\sum_{x \in C} \prod_{y \in C} x = y$  is the set of proofs of the sentence  $\exists x \in C, \forall y \in C, x = y$

asserting that  $C$  has a **unique** element.

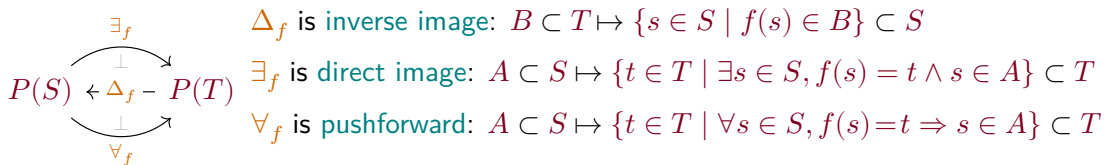
It remains to explain the analogy:

logic	$\exists$	$\forall$
sets	$\sum$	$\prod$

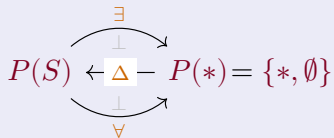
## Digression: quantifiers as adjoints



A set function  $f : S \rightarrow T$  induces order-preserving functions between their powersets:



For the unique function  $! : S \rightarrow *$  these reduce to



The set  $P(S) = \{A \subset S\}$  can be identified with the set of **predicates**  $p(s)$  with one free variable  $s \in S$  — the corresponding subset is  $\{s \in S \mid p(s) \text{ is true}\}$ . If we interpret the two elements of  $P(*)$  by  $* =: \text{“true”}$  and  $\emptyset =: \text{“false”}$  then

- $\exists$  is the function that sends the predicate  $p(s)$  to the sentence  $\exists s \in S, p(s)$
- $\forall$  is the function that sends the predicate  $p(s)$  to the sentence  $\forall s \in S, p(s)$

## Digression: locally cartesian closed categories



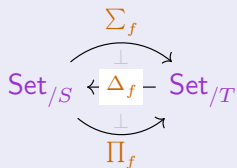
For any function  $f: S \rightarrow T$  there are functors:

$$\begin{array}{c} \begin{array}{ccc} & \exists_f & \\ \curvearrowright & \perp & \curvearrowleft \\ P(S) & \xleftarrow{\Delta_f} & P(T) \\ \curvearrowleft & \perp & \curvearrowright \\ & \forall_f & \end{array} \end{array} \quad \begin{array}{l} \Delta_f \text{ is inverse image: } B \subset T \mapsto \{s \in S \mid f(s) \in B\} \subset S \\ \exists_f \text{ is direct image: } A \subset S \mapsto \{t \in T \mid \exists s \in S, f(s) = t \wedge s \in A\} \subset T \\ \forall_f \text{ is pushforward: } A \subset S \mapsto \{t \in T \mid \forall s \in S, f(s) = t \Rightarrow s \in A\} \subset T \end{array}$$

In proof-relevant mathematics, it is natural to replace the poset  $P(S)$  by the category  $\mathbf{Set}_S$  of  $S$ -indexed sets. An object  $\{P(s)\}_{s \in S}$  is a family of sets where  $P(s)$  can be thought of as the set of proofs of some predicate  $p(s)$  on  $s \in S$ .

$$\begin{array}{c} \begin{array}{ccc} & \Sigma_f & \\ \curvearrowright & \perp & \curvearrowleft \\ \mathbf{Set}_S & \xleftarrow{\Delta_f} & \mathbf{Set}_T \\ \curvearrowleft & \perp & \curvearrowright \\ & \Pi_f & \end{array} \end{array} \quad \begin{array}{l} \Delta_f \text{ is substitution: } \{Q(t)\}_{t \in T} \mapsto \{Q(f(s))\}_{s \in S} \\ \Sigma_f \text{ is sum: } \{P(s)\}_{s \in S} \mapsto \{\sum_{s \in f^{-1}(t)} P(s)\}_{t \in T} \\ \Pi_f \text{ is product: } \{P(s)\}_{s \in S} \mapsto \{\prod_{s \in f^{-1}(t)} P(s)\}_{t \in T} \end{array}$$

# Summary



$\Delta_f$  is **substitution**:  $\{Q(t)\}_{t \in T} \mapsto \{Q(f(s))\}_{s \in S}$

$\Sigma_f$  is **sum**:  $\{P(s)\}_{s \in S} \mapsto \{\sum_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

$\Pi_f$  is **product**:  $\{P(s)\}_{s \in S} \mapsto \{\prod_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

The triple of adjoint functors gives a more formal way to understand the set

$$\sum_{x \in C} \prod_{y \in C} x = y$$

- ~ The set of proofs “ $x = y$ ” defines an indexed set  $\{x = y\}_{x, y \in C} \in \mathbf{Set}_{/C \times C}$
- ~ Product along the projection  $\pi_1 : C \times C \rightarrow C$  gives  $\{\prod_{y \in C} x = y\}_{x \in C} \in \mathbf{Set}_{/C}$
- ~ Sum along  $! : C \rightarrow *$  gives the set  $\sum_{x \in C} \prod_{y \in C} x = y \in \mathbf{Set}_{/*} \simeq \mathbf{Set}$



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$\infty$ -Categorifying Uniqueness

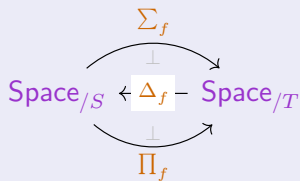
# A convenient category of spaces



Now replace **Set** by a convenient category **Space** of spaces and continuous maps.

A family of spaces  $\{E_b\}_{b \in B} \in \mathbf{Space}/_B$  is a continuous map  $\pi: E \rightarrow B$ , where the space  $E_b$  is the fiber over a point  $b \in B$  in the base space while the total space  $E \simeq \sum_{b \in B} E_b$ .

Any continuous  $f: S \rightarrow T$  gives rise to an adjoint triple:



$\Delta_f$  is pullback

$\Sigma_f$  is composition

$\Pi_f$  is pushforward

# Identifications as paths



Q: For a space  $C$ , how to interpret the family of spaces  $\{x = y\}_{x,y \in C} \in \mathbf{Space}_{/C \times C}$ ?

First guess:

$$\begin{array}{c} C \\ \Delta \downarrow \\ C \times C \end{array} \in \mathbf{Space}_{/C \times C} \text{ — but a better choice is the path space } \begin{array}{c} C^I \\ \Downarrow \\ C \times C \end{array} \in \mathbf{Space}_{/C \times C}.$$

New idea:

A point  $p \in x = y$  is a **path** from  $x$  to  $y$  in  $C$ , providing a **proof** that  $x$  equals  $y$ .



# A space of proofs



What is a point in the space  $\sum_{x \in C} \prod_{y \in C} x = y$ ?

The functor  $\sum_B : \mathbf{Space}/B \rightarrow \mathbf{Space}$  takes  $\{E_b\}_{b \in B}$  to the total space  $\sum_{b \in B} E_b$ .

$\leadsto$  a point in  $\sum_{b \in B} E_b$  is a pair  $(a, e_a)$  of a point  $a \in B$  and a point  $e_a \in E_a$

The functor  $\prod_B : \mathbf{Space}/B \rightarrow \mathbf{Space}$  takes  $\{E_b\}_{b \in B}$  to the space of sections  $\prod_{b \in B} E_b$ .

$\leadsto$  a point in  $\prod_{b \in B} E_b$  is section  $s : B \rightarrow \sum_{b \in B} E_b$  of the projection to  $B$

- So a point in  $\sum_{x \in C} \prod_{y \in C} x = y$  is a pair  $(c, h)$  where  $c \in C$  and  $h \in \prod_{y \in C} c = y$ .
- The point  $h \in \prod_{y \in C} c = y$  is a section  $h : C \rightarrow \sum_{y \in C} c = y$  to the projection.

Together  $(c, h) \in \sum_{x \in C} \prod_{y \in C} x = y$  defines:

- a center of contraction  $c$  and
- a contracting homotopy  $h$ ,

proving that the space  $C$  is contractible!

# Contractibility as uniqueness



In summary, a point in the set

$$\sum_{x \in C} \prod_{y \in C} x = y$$

is a proof that  $C$  is **unique**, while a point in the space

$$\sum_{x \in C} \prod_{y \in C} x = y$$

is a proof that  $C$  is **contractible**.

**The point:** this gives a glimpse of how uniqueness in **homotopy type theory** in fact expresses a contractibility condition — thus, uniqueness is “homotopical uniqueness.”

Thank you!