# **Wasserstein-Distance**

This is an brief introduction to Wasserstein-Distance, including its formulation, computation and application.

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### **Tutorials**

- 1. Optimal Transport for Applied Mathematicians
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## Introduction

We will start from some some intuitive examples.

### **Existing metrics**

- 1. metrics

  - 1. KL divergence:  $D_{\mathrm{KL}}(P\|Q) = -\sum_i P(i) \ln \frac{Q(i)}{P(i)}$ 2. JS divergence:  $D_{\mathrm{JS}}(P,Q) = \frac{1}{2} \left( D_{\mathrm{KL}} \left( P \| \frac{P+Q}{2} \right) + D_{\mathrm{KL}} \left( Q \| \frac{P+Q}{2} \right) \right)$

3. F divergence:  $D_f(p\|q) = \int q(x) f\left(rac{p(x)}{q(x)}
ight) dx$ , where f is a convex function.

#### 2. Issues

- 1. can not evaluate 2 distributions with different support set.
  - 1. example:  $\{p(x)|x\in[0,1]\}$  and  $\{q(y)|y\in[2,3]\}$
- 2. use KL/JS divergence as loss function -> gradient vanishing!
- 3. need well-defined distance metric

### **Transportation problem**

	$D_1$	$D_2$	• • •	$D_n$	Supply
$O_1$	$c_{11}$	$c_{12}$		$c_{1n}$	$a_1$
$O_2$	$c_{21}$	$c_{22}$		$c_{2n}$	$a_2$
:	•	•	٠	:	•
$O_m$	$c_{m1}$	$c_{m2}$	• • •	$c_{mn}$	$a_m$
Demand	$b_1$	$b_2$		$b_n$	

#### 1. Problem:

- 1. origin:  $\{O_1, ..., O_m\}$ , destination:  $\{D_1, ..., D_n\}$
- 2. supply:  $\{a_1, ..., a_m\}$ , demand:  $\{b_1, ..., b_n\}$
- 3. transport supply goods in origin locations to satisfy demand in destinations
- 4. transport plan/matrix:  $x_{ij}$
- 5. transport cost:  $c_{ij}$

#### 2. Formulation

1. Primal formation

Minimize 
$$\sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij}$$
  
Subject to:

$$egin{array}{l} \sum_{j=1}^n x_{ij} = a_i & ext{for } i = 1, 2, \ldots, m \ \sum_{i=1}^m x_{ij} = b_j & ext{for } j = 1, 2, \ldots, n \ x_{ij} \geq 0 & ext{for } i = 1, 2, \ldots, m ext{ and } j = 1, 2, \ldots, n \end{array}$$

2. Dual formulation

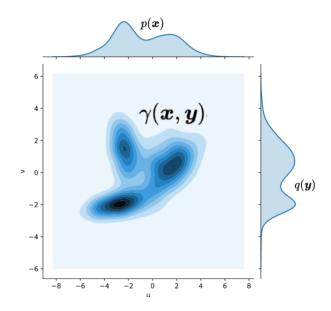
Maximize 
$$\sum_{i=1}^m a_i f_i + \sum_{i=1}^n b_i g_i$$
 Subject to:

$$f_i + g_j \le c_{ij}$$
 for  $i = 1, 2, ..., m$ , for  $j = 1, 2, ..., n$ 

### **Formulation**

### Wasserstein distance (Kantorovich formulation)

1. view it as a continuous version of transportation problem



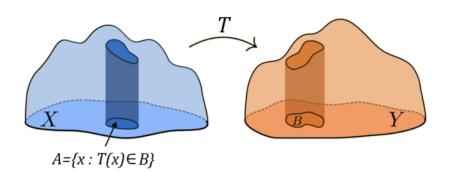
Minimize  $\mathcal{W}[p,q]=\inf_{\gamma\in\Pi[p,q]}\iint\gamma(x,y)c(x,y)dxdy$  Subject to:

$$\int \gamma(x,y) dy = p(x)$$
$$\int \gamma(x,y) dx = q(y)$$

- 2. cost function c(x, y):
  - 1. any norm,  $\|x-y\|_1, \ \|x-y\|_2, \ \|x-y\|_2^2$
- 3. joint distribution  $\gamma(x,y)$ :
  - 1. with marginal distribution  $\gamma(x)=p(x), \gamma(y)=q(y)$

## **Optimal transport and Wasserstein distance**

1. Optimal transport (Monge formulation)



$$C_M(T) = \int_{\Omega} c(x, T(x)) p(x) d$$
  
 $q = T(p)$ 

- 1. transport map: q(y) = T(p(x))
  - 1. non-linear constraint

### **Several dual formulations**

#### Kantorovich-Rubinstein Duality

$$\mathcal{W}[p,q] = \max_{f,g} \left\{ \int [p(x)f(x) + q(x)g(x)] dx \mid f(x) + g(y) \leq c(x,y) 
ight\}$$

- 1. primal-dual optimality condition
- 1. f(x) + g(y) = c(x, y)
- 2. Proof:
- 1. forward: if primal and dual reach optimality, then

$$\underbrace{\iint \gamma(x,y)c(x,y)dxdy}_{\text{primal formulation}} \tag{1}$$

$$= \underbrace{\int [p(x)f(x) + q(x)g(x)]dx} \tag{2}$$

primal=dual when reaching optimality

$$= \underbrace{\iint [f(x) + g(y)]\gamma(x, y)dxdy}_{\text{marginal distribution}}$$
(3)

$$\to f(x) + g(y) = c(x, y) \tag{4}$$

2. backward: if f(x) + g(y) = c(x, y) holds, then:

$$\underbrace{\int [p(x)f(x) + q(x)g(x)]dx} \tag{5}$$

$$\underbrace{\int [p(x)f(x) + q(x)g(x)]dx}_{\text{dual formulation}}$$

$$= \underbrace{\int \int [f(x) + g(y)]\gamma(x,y)dxdy}_{\text{marginal distribution}}$$
(5)

$$= \underbrace{\iint \gamma(x,y)c(x,y)dxdy}_{f(x)+g(y)=c(x,y)} \tag{7}$$

$$\rightarrow$$
 primal = dual when reaching optimality (8)

#### Lipschitz constrained formulation

$$\mathcal{W}[p,q] = \max_f \left\{ \int [p(x)f(oldsymbol{x}) - q(oldsymbol{x})f(oldsymbol{x})]dx \mid \|f\|_L \leq 1 
ight\}$$

1. consider the optimality condition when x=y

1. 
$$f(y) + g(y) = c(y, y) = 0 \rightarrow g(y) = -f(y)$$

- 2. take g(y) = -f(y) into the  $\mathcal{W}[p,q]$
- 1. objective function:  $\max_f \left\{ \int [p(x)f(m{x}) q(m{x})f(m{x})] dx \right\}$
- 2. constraints:  $||f||_L \leq 1$
- 1.  $f(x)-f(y)\leq c(x,y)$  and  $f(y)-f(x)\leq c(y,x)$  2.  $\|f\|_L=rac{|f(x)-f(y)|}{c(x,y)}\leq 1$

#### **Unconstrained formulation**

$$\mathcal{W}[p,q] = \max_f \int f(x) dp(x) + \int \min_x [c(x,y) - f(x)] dq(y)$$

1. C-transform:

For  $f \in C(\Omega)$  define its c-transform  $f^c \in C(\Omega)$  by

$$f^c(y) = \inf\{c(x,y) - f(x) \mid x \in \Omega\}$$

and its  $ar{c}$ -transform  $g^{ar{c}} \in C(\Omega)$  by

$$g^{\bar{c}}(x) = \inf\{c(x,y) - g(y) \mid y \in \Omega\}$$

1.  $f^{c\hat{c}}(x) > f(x)$ , "=" holds when f is concave

- 2. Consider g(y) is the C-transform of f(x)
- 1.  $f^c(y) = \inf_x \{c(x,y) f(x)\}$
- 2. Proof of such a transform will not affect the optimality
- 1. prove f(x) and  $f^c(y)$  satisfy the constraint

$$f(x) + \inf\{c(x,y) - f(x)\}\tag{9}$$

$$\leq f(x) + c(x,y) - f(x) \tag{10}$$

$$=c(x,y) \tag{11}$$

The constraint is always be satisfied under C-transform

2. prove f(x) and  $f^{c}(y)$  reach optimality condition

$$f(x) = g^c(x) \tag{12}$$

$$\to f^c(y) = g^{c\hat{c}}(y) \ge g(y) \tag{13}$$

$$\to f(x) + f^c(y) \ge f(x) + g(y) \tag{14}$$

when f(x)+g(y)=c(x,y),  $c(x,y)\leq f(x)+f^{c}(y) \leq c(x,y)$ Therefore  $f(x)+f^{c}(y)=c(x,y)$  and reaches optimality.

#### Quadratic cost function

- 1. quadratic cost function:  $c(x,y)=\frac{1}{2}\sqrt{x-y}^2$
- 2. The C-transform can be simplified as:

$$f(x) = \inf_{y} \left\{ \frac{1}{2} \|x - y\|^2 - g(y) \right\}$$

$$= \frac{1}{2} \|x\|^2 + \inf_{y} \left\{ -\langle x, y \rangle + \frac{1}{2} \|y\|^2 - g(y) \right\}$$

$$= \frac{1}{2} \|x\|^2 - \sup_{y} \left\{ \langle x, y \rangle - \left[ \frac{1}{2} \|y\|^2 - g(y) \right] \right\}$$

$$:= \phi(x) : \text{convex}$$

- 1.  $\phi(x)$  is the convex conjugate of  $rac{1}{2}\|y\|^2-g(y)$
- 1. Brenier theorem:
- 1. Under quadratic case, optimal transport map T(x) is equivalent with transport plan  $\gamma(x,y)$

$$T(x) = x - \nabla f(x) = x - (x - \nabla \phi(x)) = \nabla \phi(x)$$

#### **Convex formulation**

under quadratic case:

$$egin{aligned} f(x)+g(y) &\leq rac{1}{2}\|x-y\|_2^2 &\Longleftrightarrow \ \left[rac{1}{2}\|x\|_2^2-f(x)
ight]+\left[rac{1}{2}\|y\|_2^2-g(y)
ight] \geq \langle x,y
angle \end{aligned}$$

2. define:

1. 
$$f'(x) = \frac{1}{2} ||x||_2^2 - f(x)$$

2. 
$$g'(y) = rac{1}{2} \|y\|_2^2 - g(y)$$

3. The objective function becomes:

$$\mathcal{W}[p,q] = C_{p,q} - \min_{f',g'} \left\{ \mathbb{E}_p[f'(x)] + \mathbb{E}_q[g'(y)] \mid f'(x) + g'(y) \geq \langle x,y 
angle 
ight\} \ C_{p,q} = rac{1}{2} \mathbb{E}_p[\|X\|_2^2] + \mathbb{E}_q[\|Y\|_2^2]$$

4. apply the conjugate transformation

1. 
$$g'(y)=f^{'*}(y)=\sup_{x}\left\{\langle x,y
angle-\underbrace{\left[rac{1}{2}\|x\|^2-f(x)
ight]}_{f'(x)}
ight\}$$

5. unconstrained optimization

$$\mathcal{W}[p,q] = C_{p,q} - \min_{f',g'} \left\{ \mathbb{E}_p[f'(x)] + \mathbb{E}_q[f^{'*}(y)] 
ight\}$$

- 1. similar proof as C-transform
- 1. constraint:
- 1.  $f'(x) + f^{'*}(y) \geq \langle x,y 
  angle$
- 2. optimality
- 1.  $f^{**} \leq f$ , "=" holds when f is convex
- 2. f' and g' are convex
- 1. According to the Brenier theorem, when reach optimality
- 1. abla g'(y) = T(y) is the optimal transport map

2. 
$$f^{'*}(y) = \sup_x \left\{ \langle x,y \rangle - \left\lceil \frac{1}{2} \|x\|^2 - f(x) \right\rceil \right\} = \langle T(y),y \rangle - \left\lceil \frac{1}{2} \|T(y)\|^2 - f(T(y)) \right\rceil$$

2. convex formulation:

$$\mathcal{W}[p,q] = C_{p,q} - \min_{f' \in ext{cvx}} \max_{g' \in ext{cvx}} \left\{ \mathbb{E}_p[f'(x)] + \mathbb{E}_q[f^{'*}(y)] 
ight\}$$

# Computation

# **Application**