

$$\begin{array}{ccc}
 \text{Hom}(A, A) & \xrightarrow{\text{Hom}(A, f)} & \text{Hom}(A, X) \\
 \downarrow \Phi_A & & \downarrow \Phi_X \\
 F(A) & \xrightarrow{Ff} & F(X)
 \end{array}$$

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What you needa know about Yoneda

Emma Bach (she/her)

Seminar on Functional Programming and Logic, Summer Semester 2025

Motivation

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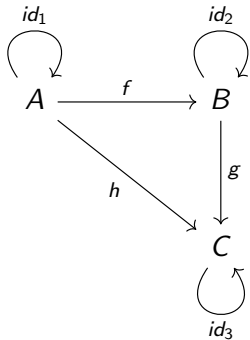
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- ▶ “*Tell me your company, and I will tell you what you are.*”¹
- ▶ The Yoneda lemma is the result of applying this way of thinking to mathematical objects in category theory.
- ▶ As a result, a category \mathbb{C} is often best understood by instead studying functors from that category into \mathbf{Set} .

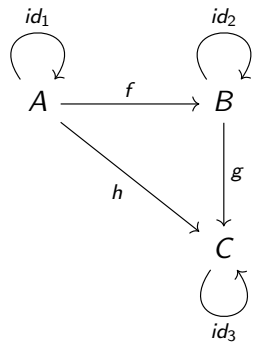
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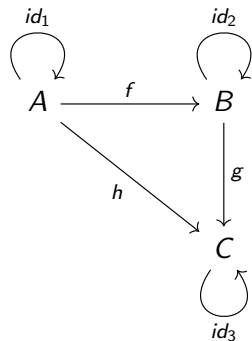


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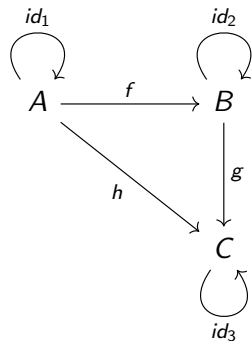
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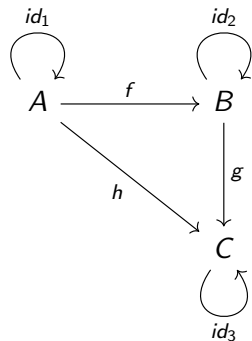
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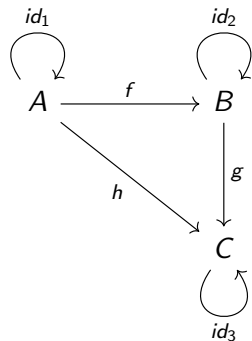
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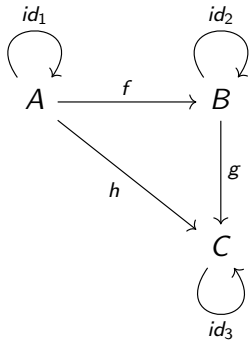
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- ▶ For every category \mathbb{C} , the *opposite category* \mathbb{C}^{op} is a category.
- ▶ For every pair of categories \mathbb{C}, \mathbb{D} , the *product category* $\mathbb{C} \times \mathbb{D}$ is a category.

Functors

A *functor* $F : \mathbb{C} \rightarrow \mathbb{D}$ is a structure-preserving map between two categories:

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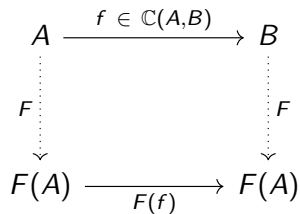
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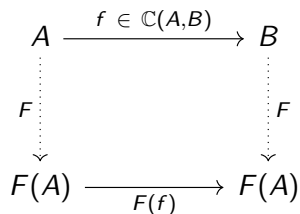
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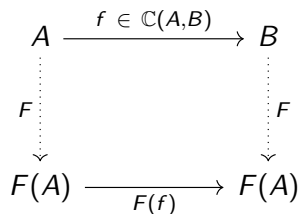
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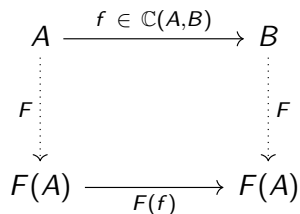
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Functors from a category into itself are known as *endofunctors*.

Natural Transformations

$$A \xrightarrow{f} B$$

► Structure-preserving maps between functors.

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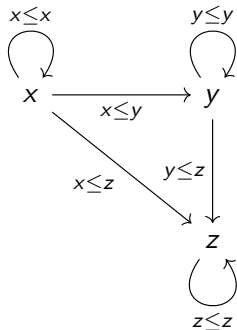
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 - ▶ These morphisms satisfy the *naturality condition*:

$$\forall f \in \mathbb{C}(A, B) : \phi_B \circ F(f) = G(f) \circ \phi_A$$

Exercise 1 : Order Categories

- a) Let \leq be a reflexive, transitive order (a *preorder*) on a set M . Show that if we define objects by $|\mathbb{P}re(M, \leq)| = M$ and morphisms by $\exists! f_{x \leq y} \in \mathbb{P}re(x, y) \Leftrightarrow x \leq y$, then $\mathbb{P}re(M, \leq)$ forms a category.
- b) Let $F : \mathbb{M} \rightarrow \mathbb{M}$ be an endofunctor on \mathbb{M} . Show that F defines a monotonic function $M \rightarrow M$, i.e. $\forall x, y : x \leq y \implies F(x) \leq F(y)$.
- c) Let $F, G : M \rightarrow M$ be monotonic functions. Let ϕ be a natural transformation $F \rightarrow G$. Show that $\forall x \in M : F(x) \leq G(x)$.

Exercise 1 : Order Categories, Solution a)



- ▶ \leq is reflexive, so we have
 $\forall x : x \leq x \implies \exists id_{x \leq x} \in \mathbb{P}re(x, x)$.
- ▶ Because of transitivity, for every pair of morphisms $f_{x \leq y}$ and $g_{y \leq z}$, we have a composed morphism $(g \circ f)_{x \leq z}$.
- ▶ Since our morphisms are just witnesses of an ordering, they don't care about the order of function application, so composition is associative.

Exercise 1 : Order Categories, Solution b)

$$\begin{array}{ccc} x & \xrightarrow{x \leq y} & y \\ \downarrow F & & \downarrow F \\ F(x) & \xrightarrow{F(x) \leq F(y)} & F(y) \end{array}$$

- By the definition of functors, F must take each morphism $f_{x \leq y} \in \mathbb{P}re(x, y)$ to a morphism $F(f)_{F(x) \leq F(y)} \in \mathbb{P}re(F(x), F(y))$.

Exercise 1 : Order Categories, Solution c)

$$\begin{array}{c} F(x) \\ \downarrow \phi_x : F(x) \leq G(x) \\ G(x) \end{array}$$

- By the definition of natural transformations, for every object x , ϕ_x is a morphism $F(x) \rightarrow G(x)$. If such a morphism exists, we have $F(x) \leq G(x)$.

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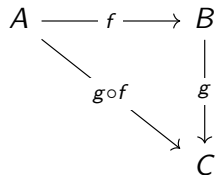
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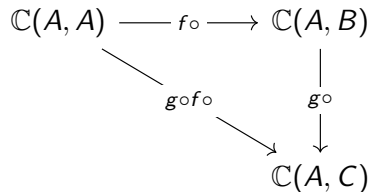
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- ▶ So our free theorem is a proof that any parametrically polymorphic function r is a natural transformation!
- ▶ It turns out that parametrically polymorphic functions correspond exactly to natural transformations between endofunctors $\mathbb{H}ask \rightarrow \mathbb{H}ask$.

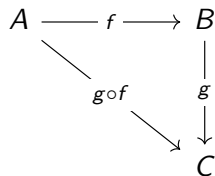
Homfunctors



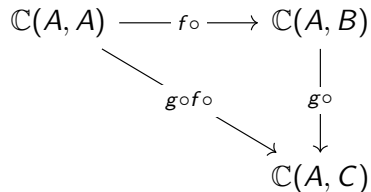
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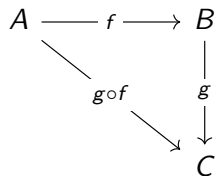
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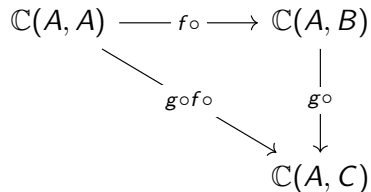
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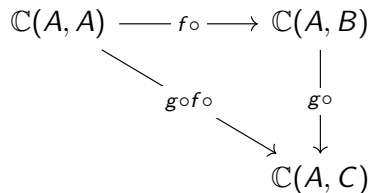
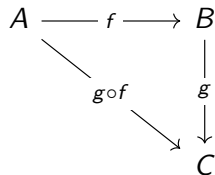
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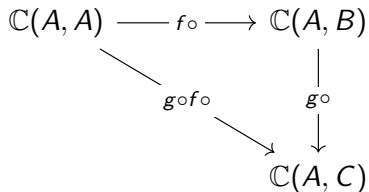
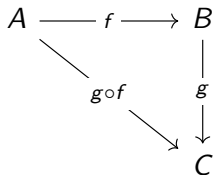


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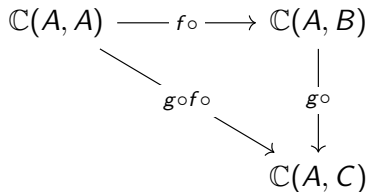
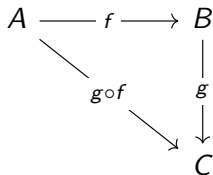
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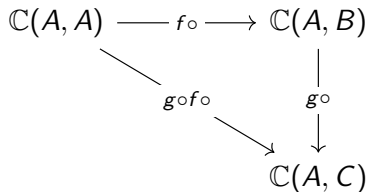
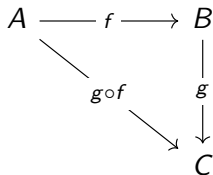
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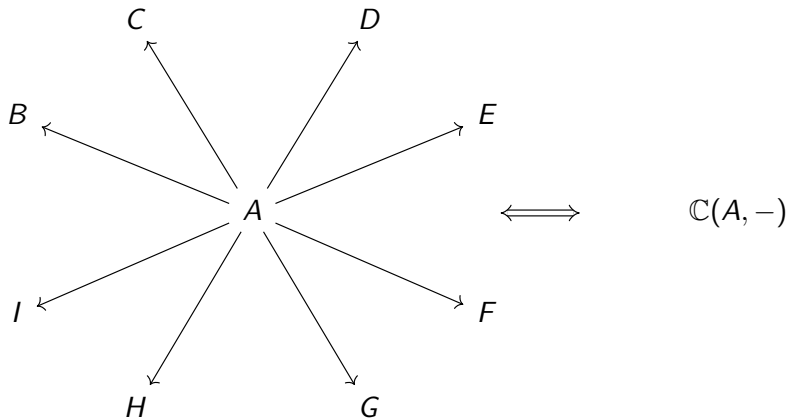
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- ▶ A morphism $\phi \in \mathbb{D}^{\mathbb{C}}(F, G)$ is a natural transformation $F \rightarrow G$.

The Yoneda Embedding



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- Formally, we want a bijective functor

$$\mathcal{Y} : \mathbb{C} \rightarrow \mathbf{Set}^{\mathbb{C}}$$

$$A \mapsto \mathbb{C}(A, -)$$

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- Is it actually possible to construct all of the necessary natural transformations?

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- ▶ Furthermore, this isomorphism is a natural transformation.
- ▶ So we can construct the Yoneda embedding \mathcal{Y} from the set $F(A)$.
- ▶ Vice versa, if we know all natural transformations $\text{Nat}(\mathbb{C}(A, -), F)$, we can construct the set $F(A)$.

Constructing the bijection

$$A \xrightarrow{f} B$$

► Let $\phi \in \text{Nat}(\mathbb{C}(A, -), F)$. Since ϕ is natural transformation, we have

$$\begin{array}{ccc} \mathbb{C}(A, A) & \xrightarrow{\mathbb{C}(A, f) = f \circ} & \mathbb{C}(A, B) \\ \downarrow \phi_A & & \downarrow \phi_B \\ F(A) & \xrightarrow{F(f)} & F(B) \end{array}$$

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- ▶ Remember that these functors are $\mathbb{C} \rightarrow \text{Set}$.
- ▶ This means our morphisms are just regular set functions.

Constructing the bijection

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ id_A & \xrightarrow{\mathbb{C}(A,f)=f \circ} & f \\ \downarrow \phi_A & & \downarrow \phi_B \\ u \in F(A) & \xrightarrow{F(f)} & \phi_B(f) = F(f)(u) \end{array}$$

- If we apply these functions to the identity morphism id_A , we get:

$$\phi_B(f \circ id_A) = F(f)(\phi_A(id_A))$$

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- ▶ The Yoneda lemma is often viewed as a generalization of Cayley's theorem.

Exercise 2 - Cayley's Theorem for Monoids

Use the Yoneda embedding to show that every monoid M is isomorphic to a monoid of functions $M \rightarrow M$.

Hint 1: The Yoneda embedding gives an isomorphism between objects and their homfunctors.

Hint 2: Two weeks ago we saw that every monoid M defines a category \mathbb{M} with a single object $*$ and a morphism m for each element $m \in M$, where we define morphism composition to be the monoid operation.

Exercise 2 - Cayley's Theorem for Monoids

- The Yoneda embedding is an isomorphism mapping each object to its homfunctor.

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- ▶ Thus, the Yoneda embedding on M is an isomorphism between monoid objects and a set of functions $M \rightarrow M$. These functions form a monoid under composition.



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- ▶ *Profunctor optics* are neat and flexible representations of optics as individual polymorphic function.
- ▶ In particular, profunctor optics make composition of optics trivial.
- ▶ Equivalence between optics and their profunctor representations comes down to the Yoneda lemma.

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data Adapter a b s t = Adapter { from :: s -> a, to :: b -> t }
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- ▶ This lets us view the category $\mathbb{H}ask^{op} \times \mathbb{H}ask$ as the category $\mathbb{A}da$ of adapters.

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- ▶ We define the category \mathbb{Prof} of Profunctors on Haskell types to be the functor category $\mathbb{Set}^{(\mathbb{Hask}^{op} \times \mathbb{Hask})} = \mathbb{Set}^{Ada}$

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$$\begin{array}{ccc} F & \xrightarrow{\eta} & G \\ \text{\tiny $-(A)$} \downarrow \text{\tiny \vee} & & \downarrow \text{\tiny $-(A)$} \text{\tiny \vee} \\ F(A) & \xrightarrow{\eta_A} & G(A) \end{array}$$

- Given a category \mathbb{C} , the operation of applying a functor $F : \mathbb{C} \rightarrow \mathbf{Set}$ to an object $A \in |\mathbb{C}|$ is itself a functor from $\mathbf{Set}^{\mathbb{C}}$ to \mathbb{C} .

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- ▶ We write $-(A)$ for this functor.

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- The profunctor representation of an adapter is given by:

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type AdapterP a b s t =  
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- ▶ Specifically, the homsets in $\mathbb{A}daP$ are:

$$\mathbb{A}daP((A, B), (S, T)) = \mathbb{S}et^{\mathbb{P}rof}(-(A, B), -(S, T))$$

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- ▶ Equivalence of adapters and profunctor adapters can be shown by applying the Yoneda embedding twice.
- ▶ Similar techniques can be used to show the equivalence of any optic and its profunctor representation.