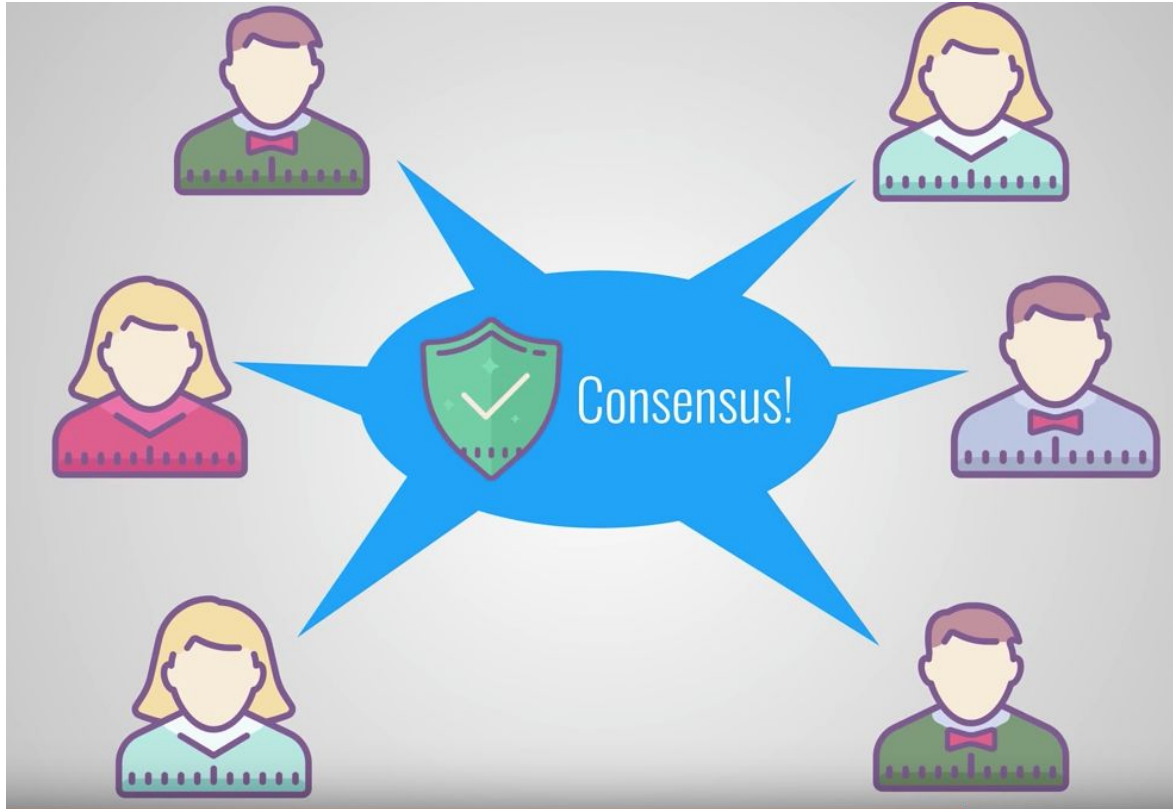
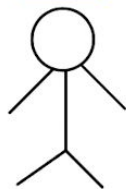

Transport Paxos To Cloud — Environment —

Paxos Algorithm

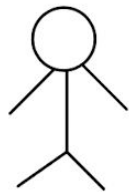
What is Paxos?



Roles in Paxos

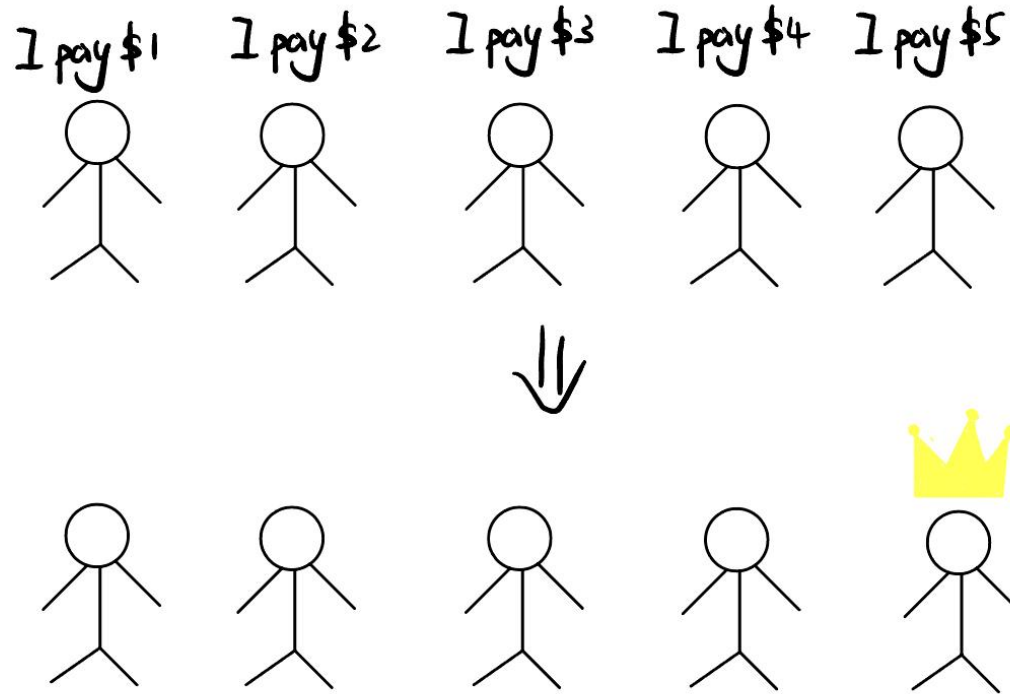


Leader: collects information from learners and makes decision on which value

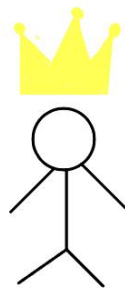
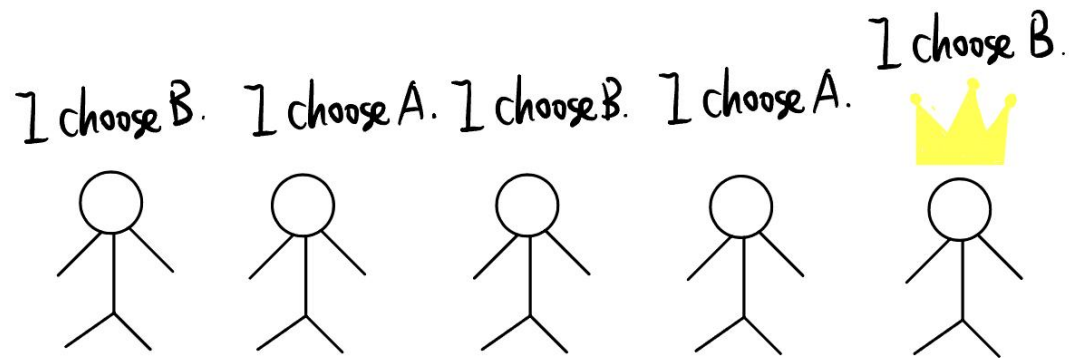


Learner: vote for values and learn the decision from the leader

Leader Election(first consensus)



Leader Makes Decision(second consensus)

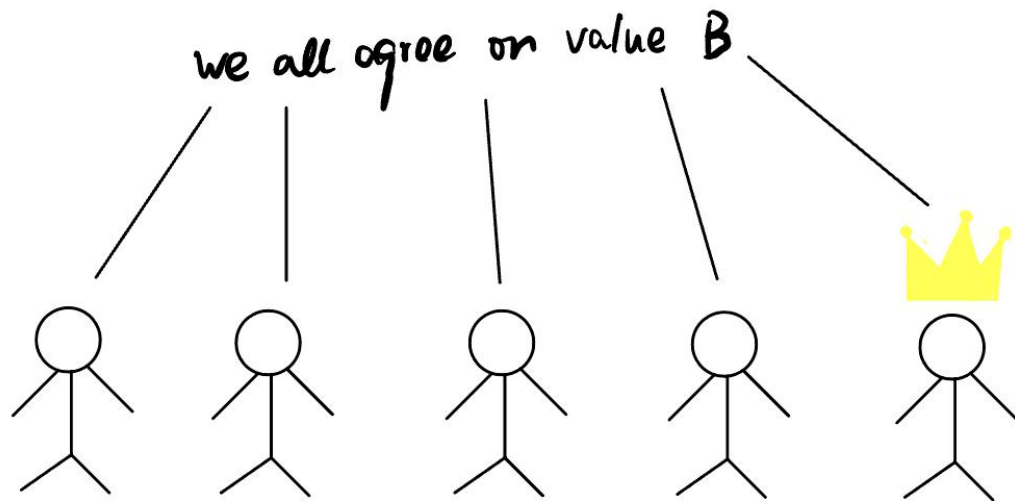


value B = 3

value A = 2

\Rightarrow value B is chosen

Learner learn the decision



My Implementation

Application: key/value store between client and servers

Client: put(key, value) and get(key)

Server: reply putOK or value of the key

Learner: vote for client request

Leader: decide which request execute first and store changes to database

Why servers vote for different client request?

The answer is unreliable network.

Example: delayed message

Round 1: every server votes for request 1, however, one vote message is delayed.

Round 2: other servers vote for request 2, however, the delayed message is arrived.

Leader will consider delayed message as a vote in round 2

Paxos Client

1. **broadcast** command to all server
2. only **accept** the first correct result from one server, and ignore repeated result

Paxos Server

Phase 1: leader election

1. every server start with a ballot number
2. every server **broadcasts** their own ballot numbers to all servers
3. server keep track of received ballot number
4. server only **vote** for the one with ballot number that is higher than its current recorded ballot number
5. server with >50% votes become the leader(majority calculation)

Paxos Server

Phase 2: stable leader and decision notification

1. leader **broadcast** command to all servers
2. leader and other servers execute the command from the client
3. leader **send** the result back to the client

Unreliable Network

How paxos system handle delayed or lost message?

Answer: resend mechanism

1. For almost every type of message, there is a timer.
2. When time up, the message will be **resended** until it get desired reply

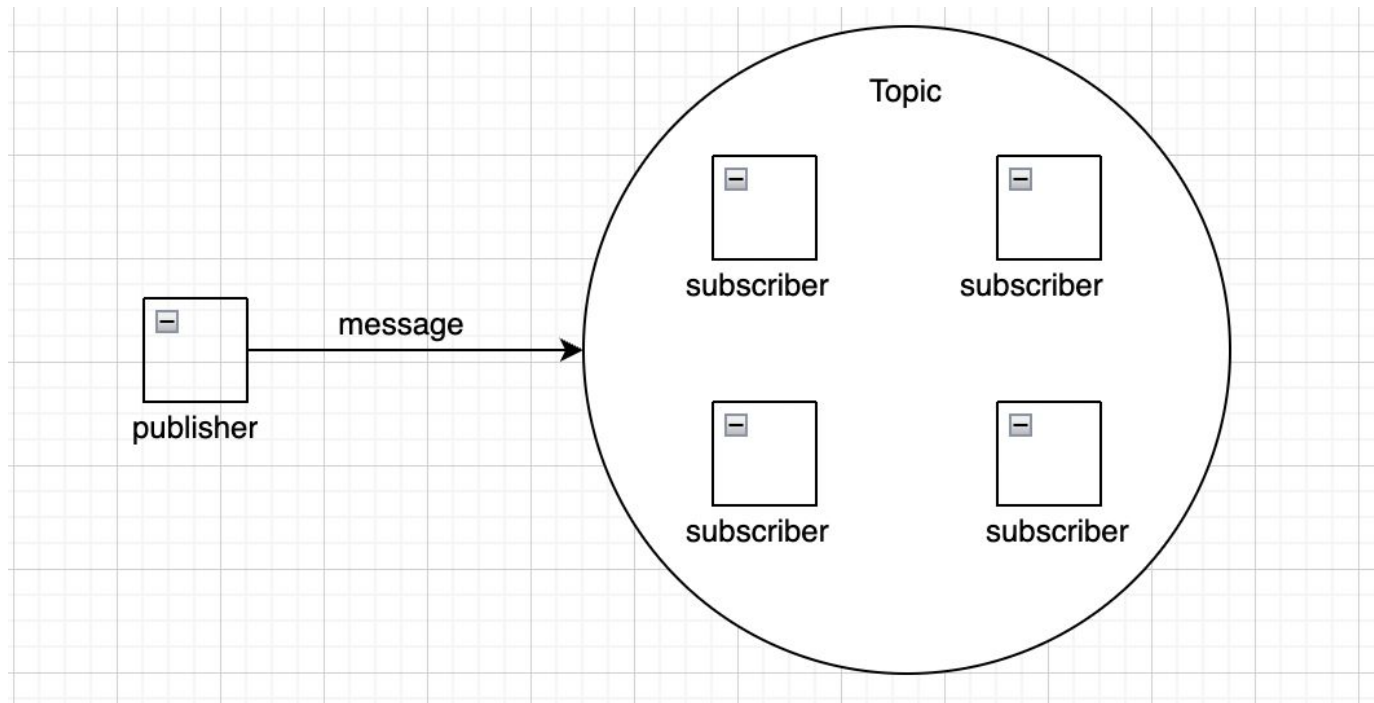
EMOX

— Paxos Communication —

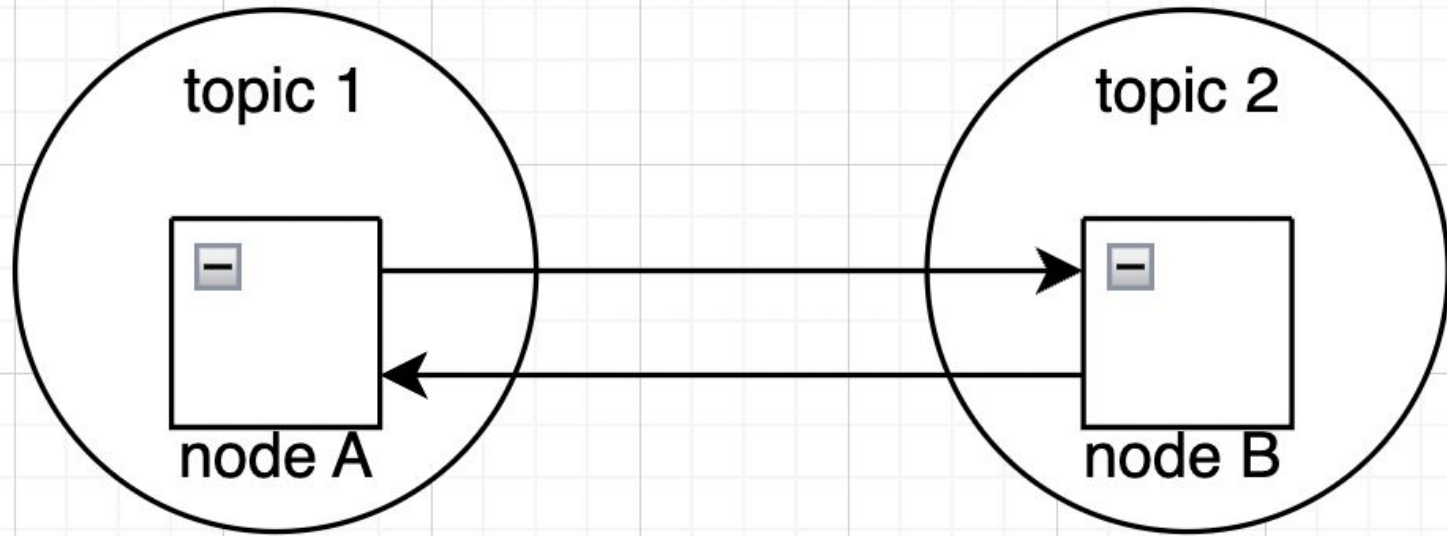
EMQX (MQTT-Message Queuing Telemetry Transport)



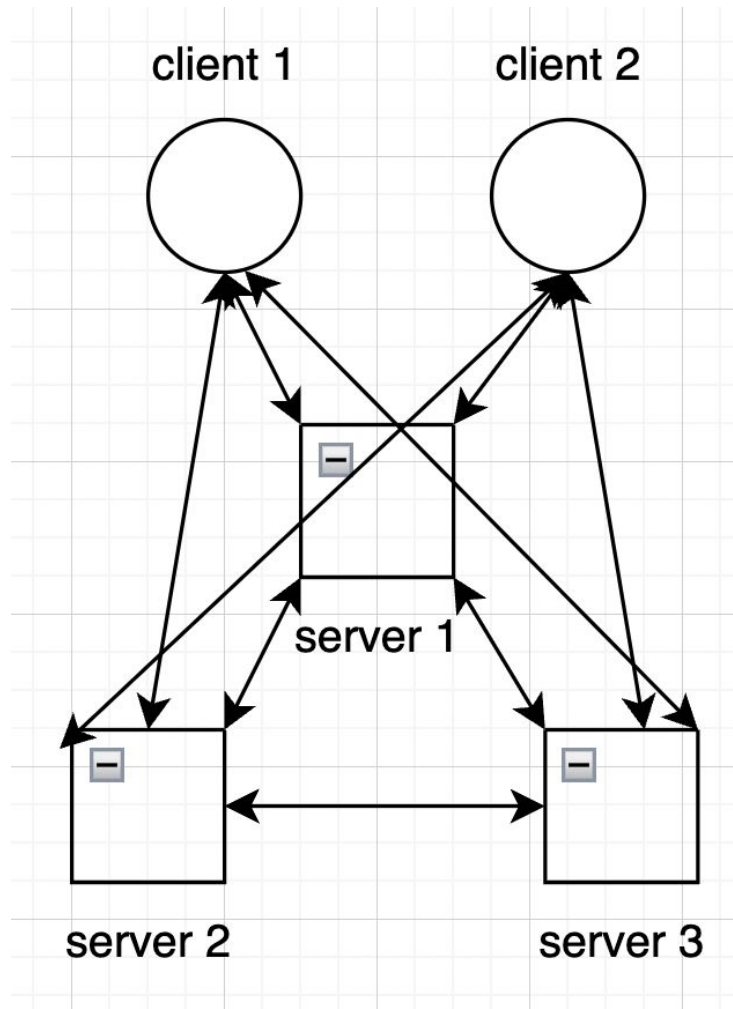
EMQX topic



My Design of Node Communication



Example of a simple paxos system





Kubernetes



EMQX in K8s

```
zhenhuansu@Zhenhuans-MacBook-Pro cs551_project_demo % kubectl get pods
```

NAME	READY	STATUS	RESTARTS	AGE
balanced-54c864d9c-zfrv6	1/1	Running	1 (30d ago)	30d
client-deployment-7cbcc54c86-84jcd	1/1	Running	0	9m27s
client-deployment-7cbcc54c86-j5mwj	1/1	Running	0	9m27s
client-deployment-7cbcc54c86-x7npv	1/1	Running	0	9m27s
emqx-core-69bd466b69-0	1/1	Running	0	16h
emqx-core-69bd466b69-1	1/1	Running	0	16h

Deployment of client and server

1. every client or server node is inside of one pod
2. All of pods communicate based on multiple topics through EMQX cluster

EMQX dashboard

EMQX

Upgrade →

admin

Monitoring

Access Control

Authentication

Authorization

Banned Clients

Flapping Detect

Integration

Management

Diagnose

System

Authentication

< Back | Built-in Database

Connected password_based

Overview Settings **User Management**

Username Is superuser

Username	Is superuser	Actions
client1	No	Edit Delete
client2	No	Edit Delete
server1	No	Edit Delete
server2	No	Edit Delete
server3	No	Edit Delete

v5.3.2

Shared Files among Pods

1. I have multiple testcases with different configuration of node size
2. every server needs to know the the size of server to calculate majority number
3. every server also need to know if other servers connect to the EMQX cluster so it can start to run Paxos algorithm

Persistent Volume

Persistent Volume

```
apiVersion: v1
kind: PersistentVolume
metadata:
  name: example-pv
spec:
  capacity:
    storage: 5Gi
  accessModes:
    - ReadWriteMany
  hostPath:
    path: /Users/zhenhuansu/D
```

Persistent Volume Claim

```
apiVersion: v1
kind: PersistentVolumeClaim
metadata:
  name: shared-volume-claim
spec:
  accessModes:
    - ReadWriteMany
  storageClassName: ""
  resources:
    requests:
      storage: 1Gi
```


Pod Volume

```
spec:
  containers:
    - name: client-container
      image: emma188/client:latest
      ports:
        - containerPort: 8080
      volumeMounts:
        - name: config-volume
          mountPath: /app/files
  volumes:
    - name: shared-storage
      persistentVolumeClaim:
        claimName: shared-volume-claim
```

Chaos Mesh

unreliable network

Chaos Mesh experiment



Schedules



Experiments



Events



Archives



Settings



IO Injection

HTTP

HTTP Fault



Kernel Fault



Network Attack



Pod Fault



Stress Test



Clock Skew



JVM Fault



Partition



Loss



Delay

What is the principle of network failure in chaos mesh?

Chaos Controller: Chaos Mesh defines a custom resource called NetworkChaos that allows users to specify chaos experiments related to network behavior.

Chaos Daemon: Chaos Daemon running on each targeted node receives the instructions from the Chaos Controller. It uses iptables rules to manipulate network traffic.

Complexity of Paxos in terms of number of messages

N is the number of servers

Phase 1: $O(N \times N)$ every node send its ballot number to each other for leader election

Phase 2: $O(N)$ leader send command to other servers.

This is only under the reliable network condition

My Implementation

I have 6 types of message and 5 timers

Example:

HeartBeatMessage: the message that leader broadcast to other servers to let them know that leader is still working.

HeartBeatTimer: leader broadcast HeartBeatMessage to all the servers every 100 millisecond.

My test case

Test goal: test paxos system under bad network condition

Configuration: 7 servers and 5 client

Network condition: Message loss: 20%, Message delay: 10 ms, Multiple network partitions

Details:

client: constantly send their command

server: randomly split servers into 2 partition groups in size of 4 and 3. Do the repartition 5 times.

Result

Incoming Messages



Outgoing Messages



Question?
