

Natural Language Processing - Problem Set #4

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Problem 1

a

Since $f_1(x, y) = f_2(x, y)$, we have either $f(x, y) = (0, 0)$ or $f(x, y) = (1, 1)$. Let k be the number of examples in the training sample such that $f(x, y) = (1, 1)$. Then $n - k$ is the number of examples in the training sample such that $f(x, y) = (0, 0)$. Consider the case where $v' = (v, v)$. Then $L(v') = \sum_i \log P(y_i | x_i, v) - c \sum_k v_k^2 = L(v) = \sum_i \log \frac{e^{v' \cdot f(x, y)}}{\sum_{y' \in V} e^{v' \cdot f(x, y')}} = k \log \frac{2^{2v}}{ke^{2v} + n - k} + (n - k) \log \frac{1}{ke^{2v} + n - k} - 2cv^2 = 2kv - k \log(ke^{2v} + n - k) + (n - k) \log(1) + (k - n) \log(ke^{2v} + n - k) - 2cv^2 = 2kv - n \log(ke^{2v} + n - k) - 2cv^2$. Now consider the case where $v'' = (v_1, v_2)$, $v_1 \neq v_2$. Following the same simplification steps as before, we have $L(v'') = k(v_1 + v_2) - n \log(ke^{v_1 + v_2} + n - k) - c(v_1^2 + v_2^2)$. Now to show v^* satisfies $v_1 = v_2$, we need only fix $v_1 + v_2 = 2v$ and show $L(v') > L(v'')$. Since $v_1 + v_2 = 2v$, $2kv - n \log(ke^{2v} + n - k) = k(v_1 + v_2) - n \log(ke^{v_1 + v_2} + n - k)$, so we only need to show $2v^2 < v_1^2 + v_2^2$. Let $v_1 = v + \alpha$ so $v_2 = v - \alpha$. Then $v_1^2 + v_2^2 = (v + \alpha)^2 + (v - \alpha)^2 = 2v^2 - 2v\alpha + 2v\alpha + 2\alpha^2 = 2v^2 + 2\alpha^2$. Since $|\alpha| > 0$, $\alpha^2 > 0$, so $2v^2 < 2v^2 + 2\alpha^2$ and therefore $L(v') > L(v'')$ so v^* satisfies $v_1^* = v_2^*$. \square

b

Now we define $L(v) = \sum_i \log P(y_i | x_i, v) - c \sum_k |v_k|$. As before, we let k be the number of examples in the training sample such that $f(x, y) = (1, 1)$, so $n - k$ is the number of examples in the training sample such that $f(x, y) = (0, 0)$. Again, we evaluate $L(v')$ and $L(v'')$ where $v' = (v, v)$ and $v'' = (v_1, v_2)$, $v_1 \neq v_2$. Following the same simplification steps as in the previous part, we have $L(v') = 2kv - n \log(ke^{2v} + n - k) - 2c|v|$ and $L(v'') = k(v_1 + v_2) - n \log(ke^{v_1 + v_2} + n - k) - c(|v_1| + |v_2|)$. Again, fixing $v_1 + v_2 = 2v$, $2kv - n \log(ke^{2v} + n - k) = k(v_1 + v_2) - n \log(ke^{v_1 + v_2} + n - k)$, so to find $\max\{L(v'), L(v'')\}$, we need only find $\min\{2|v|, |v_1| + |v_2|\}$. As before, let $v_1 = v + \alpha$ so $v_2 = v - \alpha$. For $-\alpha < v < \alpha$, $2|v| < |v + \alpha| + |v - \alpha| = |v_1| + |v_2|$, and for $-v \leq \alpha \leq v$, $2|v| = |v + \alpha| + |v - \alpha| = |v_1| + |v_2|$. Thus, v^* satisfies $-v_1^* - v_2^* \leq v_1^* - v_2^* \leq v_1^* + v_2^*$. \square

2

Since $v^* = \operatorname{argmax}_v L(v)$, $L'(v^*) = 0$. We know $\frac{dL(v)}{dv_k} = \sum_{i=1}^n f_k(x_i, y_i) - \sum_{i=1}^n \sum_{y \in Y} P(y|x_i, v^*) f_k(x_i, y)$, so we set $\sum_{i=1}^n f_k(x_i, y_i) - \sum_{i=1}^n \sum_{y \in Y} P(y|x_i, v^*) f_k(x_i, y) = 0$. Since $f_k(x_i, y_i) = 1$ if and only if $x_i = x_k$ and $y_i = y_k$ (and $f_k(x_i, y_i) = 0$ otherwise), $\sum_{i=1}^n f_k(x_i, y_i) = \operatorname{Count}(x_k, y_k)$. Since $\sum_{y \in Y} P(y|x_i, v^*) f_k(x_i, y) = 0$ if $x_i \neq x_k$ and $\sum_{y \in Y} P(y|x_i, v^*) f_k(x_i, y) = \sum_{y \in Y} P(y|x_k, v^*) f_k(x_k, y)$ if $x_i = x_k$, $\sum_{i=1}^n \sum_{y \in Y} P(y|x_i, v^*) f_k(x_i, y) = \operatorname{Count}(x_k) \sum_{y \in Y} P(y|x_k, v^*) f_k(x_k, y)$. If we take $P(x|y, v) = \frac{\operatorname{Count}(x, y)}{\operatorname{Count}(x)}$, we have $\operatorname{Count}(x_k) \sum_{y \in Y} P(y|x_k, v^*) f_k(x_k, y) = \operatorname{Count}(x_k) \sum_{y \in Y} \frac{\operatorname{Count}(x_k, y)}{\operatorname{Count}(x_k)}$. $\operatorname{Count}(x_k) \frac{1}{\operatorname{Count}(x_k)} \sum_{y \in Y} \operatorname{Count}(x_k, y) f_k(x_k, y) = \sum_{y \in Y} \operatorname{Count}(x_k, y) f_k(x_k, y)$. Since $f_k(x_k, y) = 1$ if $y = y_k$ and $f_k(x_k, y) = 0$ if $y \neq y_k$, $\sum_{y \in Y} \operatorname{Count}(x_k, y) f_k(x_k, y) = \operatorname{Count}(x_k, y_k)$. Then $\frac{dL(v)}{dv_k} = \operatorname{Count}(x_k, y_k) - \operatorname{Count}(x_k, y_k) = 0$. Since $L(v)$ is concave, it has one maximum and no minimum, so we have found the global maximum, v^* , and therefore it must satisfy $P(x|y, v^*) = \frac{\operatorname{Count}(x, y)}{\operatorname{Count}(x)}$. \square

3

a

We choose the log-linear model with inputs $X = V$, labels $Y = V \cup V'$, $d = 2$, function $f : X \times Y \rightarrow \mathbb{R}^2$ where $f(x, y) = (f_1(x, y), f_2(x, y))$, $f_1(x, y) = 1$ if $x = y$ and $f_1(x, y) = 0$ if $x \neq y$, and $f_2(x, y) = 1$ if $x = y'$ and $f_2(x, y) = 0$ if $x \neq y'$.

b

Using model defined above, $f(x, x) = (1, 0)$, $f(x, x') = (0, 1)$, $f(x, y) = (0, 0)$ where $y \notin \{x, x'\}$. Then we also have $\sum_{y \in Y} e^{(v_1, v_2) \cdot f(x, y)} = (v_1, v_2) \cdot (0, 0) + e^{(v_1, v_2) \cdot (1, 0)} + e^{(v_1, v_2) \cdot (0, 1)}$, $\forall x \in X$ where the parameter vector $v = (v_1, v_2)$. Using these substitutions and the probabilities given in the problem, we have the following equations:

$$P(\text{the}|\text{the}) = \frac{e^{(v_1, v_2) \cdot (1, 0)}}{(|V'| - 2)e^{(v_1, v_2) \cdot (0, 0)} + e^{(v_1, v_2) \cdot (1, 0)} + e^{(v_1, v_2) \cdot (0, 1)}} = 0.4$$

$$P(\text{eht}|\text{the}) = \frac{e^{(v_1, v_2) \cdot (0, 1)}}{(|V'| - 2)e^{(v_1, v_2) \cdot (0, 0)} + e^{(v_1, v_2) \cdot (1, 0)} + e^{(v_1, v_2) \cdot (0, 1)}} = 0.3$$

$$P(\text{dog}|\text{the}) = \frac{e^{(v_1, v_2) \cdot (0, 0)}}{(|V'| - 2)e^{(v_1, v_2) \cdot (0, 0)} + e^{(v_1, v_2) \cdot (1, 0)} + e^{(v_1, v_2) \cdot (0, 1)}} = \frac{0.3}{|V'| - 2}$$

c

Using the equations from the previous part, we evaluate $e^{(v_1, v_2) \cdot (0, 0)} = e^0 = 1$, $e^{(v_1, v_2) \cdot (1, 0)} = e^{v_1}$, $e^{(v_1, v_2) \cdot (0, 1)} = e^{v_2}$ and substitute, giving us the following equations:

$$\frac{1}{|V'| - 2 + e^{v_1} + e^{v_2}} = \frac{3}{10(|V'| - 2)}$$

$$\frac{e^{v_2}}{|V'| - 2 + e^{v_1} + e^{v_2}} = \frac{3}{10}$$

$$\frac{e^{v_2}}{|V'| - 2 + e^{v_1} + e^{v_2}} = \frac{4}{10}$$

Then $|V'| - 2 + e^{v_1} + e^{v_2} = \frac{10(|V'| - 2)}{3}$, so solving for v_1 and v_2 and substituting this in yields: $v_2 = \ln(|V'| - 2)$ and $v_1 = \ln(\frac{4}{3}) + \ln(|V'| - 2)$.

Then our parameter vector $v = (v_1, v_2) = (\ln(\frac{4}{3}) + \ln(|V'| - 2), \ln(|V'| - 2))$.