

A proposal for a standard process management library

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1 Preamble

This document proposes an extension to the standard library.

The goal of this extension is to provide components and functions that can be used to create, manage and query processes.

1.1 This is a working document

This document outlines the different parts of the proposal but it's not the proposal itself. The proposal is written using LibreOffice and can be found on the Google Code repository as `odt/cpp1y-process-proposal.odt` (and the corresponding `odt/cpp1y-process-proposal.pdf`).

1.2 Rationale

Threads and processes are inherently similar. From an operating system point of view, they may even be implemented using the same mechanisms (as this is the case in the Linux operating system). The main difference between processes and threads lies in the way they manage virtual memory (when available): multiple threads in a process live in the same memory space, while multiple processes on a system live in different memory space. There are advantages and disadvantages to both:

- Little-to-none security mechanism is available to protect the memory space used by a thread from another thread.
- Threads are lighter than processes. This is due to the fact that when scheduling two different processes, the operating system must manage their respective view of the memory and initialize the hardware MMU accordingly. Threads do not share this requirement as they have the same view of the memory.

This is of course a very high-level view since it does not take thread local storage into account.

Process management is an important part of system programming. An implementation would allow C++ to be used to create background programs (such as daemons in POSIX environment or services on Windows) or to start other programs or utilities (as done by a shell or a launcher UI). As of today, programmers have to use specific platform APIs to implement such functionalities in their programs, making them difficult to port to other system architectures.

1.3 Design consideration

Since processes and threads are quite similar, it make sense to proposes a similar interface for both classes. I decided to give this path a try and I believe it worked quite well.

2 Proposal

Please note that both the numbering scheme and the text below are of course subject to changes. They may not match the requirements of the text of the C++ standard. Special Development notes are added to clarify some points or to denote specific issues.

2.1 header <process> synopsis

This section describe components that can be used to create and manage processes. [*Note*: these processes are intended to map one-to-one with operating system processes. *end note*]

```
namespace std {
    class process;

    void swap(process& x, process& y);

    namespace this_process {
        process::id get_id() noexcept;
        template <typename Args...>
            void exec(const string& cmd, Args&&... args);
    }
}
```

2.2 class process

The class process provides a mechanism to create a new operating system process, to join with a process (i.e., wait for a process to complete), and to perform other operations that manage and query the state of a process. A process object uniquely represents a particular operating system process. That representation may be transferred to other process objects in such a way that no two process

objects simultaneously represent the same operating system process. An operating system process is detached when no process object represents that process. Objects of class process can be in a state that does not represent an operating system process. [*Note*: A process object does not represent an operating system process after default construction, after being moved from, or after a successful call to detach or join. *end note*]

```
namespace std {
    class process {
    public:
        typedef implementation-defined
            native_handle_type;

        class id;

        process(process&) = delete;
        process(const process&) = delete;
        process& operator=(const process&) = delete;

        process() = default;
        process(process&& __p) noexcept;

        template <typename F, typename ... Args>
            explicit
            process(F&& f, Args&&... args);

        ~process();

        process& operator=(process&& __p) noexcept;
        void swap(process& __t) noexcept;

        void detach();
        void join();
        bool joinable() const noexcept;

        process::id get_id() const noexcept;
        native_handle_type native_handle() const noexcept;
    };
}
```

2.2.1 class process::id

```
namespace std {
    class process::id
```

```

{
public:
    id() noexcept;
};

bool operator==(process::id x, process::id y);
bool operator!=(process::id x, process::id y);
bool operator<=(process::id x, process::id y);
bool operator>=(process::id x, process::id y);
bool operator<(process::id x, process::id y);
bool operator>(process::id x, process::id y);

template <class Char, class Traits>
    basic_ostream<Char, Traits>&
    operator<< (
        basic_ostream<Char, Traits>& out,
        process::id id);

// hash support
template <class T> struct hash;
template <> struct hash<process::id>;
}

```

An object of type `process::id` provides a unique identifier for each process and a single distinct value for all process objects that do not represent an operating system process. Each operating system process has an associated `process::id` object that is not equal to the `process::id` object of any other operating system process and that is not equal to the `process::id` object of any `std::process` object that does not represent any operating system process.

[*Development note*: at this point, we do not take into account some very specific implementation such as PID namespace in Linux. Under Linux, two processes can share the same process id (PID) if they do not belong to the same PID namespace. Processes in a particular PID namespaces are unaware of the existence of separate processes in a different namespace even if these processes has the same PID. *end note*]

`process::id` shall be a trivially copyable class. The library may reuse the value of a `process::id` of a terminated operating system process that can no longer be joined.

[*Note*: Relational operators allow `process::id` objects to be used as keys in associative containers. *end note*]

`id()` `noexcept`;

Effects: construct an object of type `process::id`.

Postconditions: the object does not represent an operating system process. [

Development note: this doesn't mean that 0 is a good value to represent

the process id. Some systems may assignate PID 0 to a particular process.
end note]

```
bool operator==(process::id x, process::id y);
```

Returns: true only if x and y represent the same operating system process or neither x nor y represents any operating system process.

```
bool operator!=(process::id x, process::id y);
```

Returns: !(x == y)

```
bool operator<(process::id x, process::id y);
```

Returns: A value such that operator< is a total ordering.

```
bool operator<=(process::id x, process::id y);
```

Returns: (x < y) || (x == y)

```
bool operator>(process::id x, process::id y);
```

Returns: !(x <= y)

```
bool operator>=(process::id x, process::id y);
```

Returns: !(x < y)

```
template <class Char, class Traits>  
basic_ostream<Char, Traits>&  
operator<<(basic_ostream<Char, Traits>& out,  
process::id id);
```

Effects: Inserts an unspecified text representation of id into out. For two objects of type process::id x and y, if x == y the process::id objects shall have the same text representation and if x != y the process::id objects shall have distinct text representations.

Returns: out.

```
template <> struct hash<thread::id>;
```

Requires: the template specialization shall meet the requirements of class template hash.

2.2.2 process constructors

`process() noexcept;`

Effects: Construct an object of type process that does not represent an operating system process.

Postcondition: `get_id() != id()`

```
template <class F, class... Args>
explicit process(F&& f, Args&&... args)
```

Requires: F and each Ti in Args shall satisfy the `MoveConstructible` requirements. `INVOKE (DECAY_COPY (std::forward<F>(f)), DECAY_COPY (std::forward<Args>(args)) ...)` shall be a valid expression.

Effects: Construct an object of type process. The new operating system process executes `INVOKE (DECAY_COPY (std::forward<F>(f)), DECAY_COPY (std::forward<Args>(args)) ...)` with the calls to `DECAY_COPY` being evaluated in the constructing process. Any return value from this invocation is ignored. [*Note:* This implies that any exceptions not thrown from the invocation of the copy of f will be thrown in the constructing process, not the new process. *end note*] If the invocation of `INVOKE (DECAY_COPY (std::forward<F>(f)), DECAY_COPY (std::forward<Args>(args)) ...)` terminates with an uncaught exception, `std::terminate` shall be called.

Synchronization: The completion of the invocation of the constructor synchronizes with the beginning of the invocation of the copy of f.

Postcondition: `get_id() == id()`. `*this` represents the newly started process.

Throws: `system_error` if unable to start the new thread.

Error conditions: [*Development note:* not specified yet. *end note*]

```
thread(thread&& x) noexcept;
```

Effects: Constructs an object of type process from x, and sets x to a default constructed state.

Postconditions: `x.get_id() == id()` and `get_id()` returns the value of `x.get_id()` prior to the start of construction.

2.2.3 process destructor

`~process();`

If `joinable()` then `terminate()`, otherwise no effects. [*Note:* Either implicitly detaching or joining a `joinable()` process in its destructor could result in difficult to debug correctness (for `detach`) or performance (for `join`) bugs encountered only when an exception is raised. Thus the programmer must ensure that the destructor is never executed while the process is still joinable. *end note*]

2.2.4 process assignment

`process& operator=(process&& x) noexcept;`

Effects: If `joinable()`, calls `terminate()`. Otherwise, assigns the state of `x` to `*this` and sets `x` to a default constructed state.

Postconditions: `x.get_id() == id()` and `get_id()` returns the value of `x.get_id()` prior to the assignment.

2.2.5 process members

`void swap(process& x) noexcept;`

Effects: Swaps the state of `*this` and `x`.

`bool joinable() const noexcept;`

Returns: `get_id() != id()`

`void join();`

Requires: `joinable()` is true.

Effects: Blocks until the process represented by `*this` has completed.

Synchronization: The completion of the process represented by `*this` synchronizes with the corresponding successful `join()` return. [*Note:* Operations on `*this` are not synchronized. *end note*]

Postconditions: The process represented by `*this` has completed. `get_id() == id()`.

Throws: `system_error` when an exception is required

Error conditions: [*Development note:* not specified yet. *end note*]

```
void detach();
```

Requires: `joinable()` is true.

Effects: The process represented by `*this` continues execution without the calling process blocking. When `detach()` returns, `*this` no longer represents the possibly continuing operating system process. When the process previously represented by `*this` ends execution, the implementation shall release any owned resources.

Postcondition: `get_id() == id()`.

Throws: `system_error` when an exception is required.

Error conditions: [*Development note:* not specified yet. *end note*]

```
id get_id() const noexcept;
```

Returns: A default constructed `id` object if `*this` does not represent an operating system process, otherwise `this_process::get_id()` for the operating system process represented by `*this`.

2.3 namespace `this_process`

```
namespace std {
    namespace this_process {
        process::id get_id() noexcept;

        template <typename Args...>
            void exec(const string& cmd, Args&&... args);
    }
}
```



```
process::id this_process::get_id() noexcept;
```

Returns: an object of type `process::id` that uniquely identifies the current process. No other process shall have this id and this process shall always have this id. The object returned shall not compare equal to a default constructed `process::id`.

```
template <typename Args...>
void exec(const string& cmd, Args&&... args);
```

Effects: replaces the current process image by a new process image. `cmd` specify the command to execute, and `args...` is the command argument list. If the underlying system call executes, the function never returns and the new process image shall have the same id as the calling process. [*Development note:* the choice of `std::string` to represent the command is due to the fact that not all operating systems allows unicode strings to be used as commands. An open question on this subject is to be found below. *end note*]

Throws: `system_error` when the underlying system call fails to execute properly.

Error conditions: [*Development note:* not specified yet. *end note*]

3 Technical considerations

3.1 About `this_process::exec()`

The `execve()` system call which is used to implement this template function has the following prototype:

```
int execve(const char *filename ,
           char *const argv [],
           char *const envp []);
```

This proposal eludes the `envp` parameter for now (it may appear in a subsequent proposal, see below).

The `argv` argument is an array of `char *const`. In order to derive this array from an `args...` list, we need a specialized mechanism. In the test implementation, I use `std::stringstream` (to convert any `Args` type into its string representation) and I allocate the individual C strings before I copy their content. The relevant (tentative) code is:

```

void __unpack_to_strings(vector<char*>& __l)
{ __l.push_back(NULL); }

template <typename _Arg0, typename... _Args>
void __unpack_to_strings(
    vector<char*>& __l,
    _Arg0&& __arg0,
    _Args&&... __args)
{
    stringstream __stream;
    __stream << __arg0;
    string __out = __stream.str();
    char *__s = new char[__out.length()+1];
    copy(__out.begin(), __out.end(), __s);
    __s[__out.length()] = 0;
    __l.push_back(__s);
    __unpack_to_strings(__l, __args...);
}

```

Such kind of code is likely to be pervasive in C++11 codebases: transforming an argument pack into a container of things (whatever the thing is) is very likely to be an algorithm which may be used quite often by C++ programmers. There might be room here for a standard utility function, although I understand that a generic algorithm might be difficult to devise. I may replace the code above by a better version (see below, not definitive) but ideally, a standard algorithm would be far better.

```

template <class _Adder>
void __unpack_to(_Adder&& __adder)
{ }

template <
    class _Adder,
    class _Arg0,
    class... _Args>
void __unpack_to(
    _Adder&& __adder,
    _Arg0&& __arg0,
    _Args&&... __args)
{
    __adder(__arg0);
    __unpack_to(__adder, __args...);
}

```

Another implementation would be based on a conversion from `std::tuple<Args...>` (and thus would first require the creation of a tuple) to a container of things - using either an implicit conversion or a user-supplied conversion.

3.2 Is it possible to implement fork() on a Windows system?

The Windows process management API does not propose any strict equivalent to POSIX fork(). This is outlined by Microsoft in. In the same document, Microsoft endorse CreateProcess as a rough equivalent to the fork/exec use case. While this use case is often said to be prevalent, it's not the only use case we want to address.

However, our research found that the implementation of fork() is still feasible:

- The open source cygwin environment proposes a fork() function that works like its POSIX counterpart.
- Scilab for Windows (another open source project) has implemented its own version of fork() to help its conversion to the Windows platform.
- The Subsystem for Unix-based Applications (SUA) for Windows (also known as Interix) proposes a fully compliant POSIX subsystem on top of the Windows API, including a working fork() function.

Both implementation leverage the public, low-level NT API to create a child process that the relevant characteristics with its parent process.

3.3 Further implementation notes

The full interface has already been implemented for both Linux and Windows (a BSD implementation shall be quite similar to the Linux one), and the full implementation has been open-sourced on Google Code:

<https://code.google.com/p/edt-process-cpp1y/>

As it's a PoC implementation it may differ from this proposal by several points.

4 Open questions

4.1 Should the proposal implement environment variable management?

I have the feeling that this would be an improvement over the current proposal. Such extension would contain:

- A way to retrieve the value of an environment variable (i.e. a "getenv()" function).
- A way to setup or kill a particular environment variable (i.e. "setenv()" and "unsetenv()" functions).
- A way to list all environment variables (i.e. an "environ()" function that would fetch the content of the POSIX C variable "environ").

If such extension is implemented in this proposal then another overload of `this_process::exec()` shall be needed as well.

4.2 Should `exec()` take a basic `_string<>` for its cmd parameter ?

Windows `_execv()` function has a wide string counterpart named `_wexecv()`. Most posix systems don't have any equivalent to this function.

I believe that UTF8 strings can be used to represent all possible command names, but support for UTF16 command names might be of interest as well on systems that support them. It's possible to convert UTF16 to UTF8 but the conversion comes with a performance penalty.