Student Information

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Answer 1

a) Write a context-free grammar for the language $L_1 = \{w | w \in \{a, b\}^* \wedge \text{ w has twice as many b's as a's}\}$.

Let G be the grammar (V, Σ, R, S) for the language L_1 where

$$\begin{aligned}
V &= \{S, a, b\}, \\
\Sigma &= \{a, b\}, \\
R &= \{S \rightarrow aSbSb, \\
&\quad S \rightarrow bSaSb, \\
&\quad S \rightarrow bSbSa, \\
&\quad S \rightarrow e\}.
\end{aligned}$$

b) Write a context-free grammar for the language $L_2 = \{a^n b^m | m, n \in \mathbb{N} \wedge m \leq n \leq 2m\}$.

Let G be the grammar (V, Σ, R, S) for the language L_2 where

$$\begin{aligned}
V &= \{S, a, b\}, \\
\Sigma &= \{a, b\}, \\
R &= \{S \rightarrow aSbSa, \\
&\quad S \rightarrow aSaSb, \\
&\quad S \rightarrow bSaSa, \\
&\quad S \rightarrow bSa, \\
&\quad S \rightarrow aSb, \\
&\quad S \rightarrow e\}.
\end{aligned}$$

c) Formally define and draw a PDA that accepts L_1.

Let $(K, \Sigma, V, \Delta, p, \{q\})$ be the sextuple for the language L_1 where K, Σ, V, R are defined on part a. Δ contains the following transitions:

1) $((p, e, e), (q, S))$
2) $((p, e, A), (q, x))$ for each rule $A \rightarrow x$ in R
3) $((q, a, a), (q, e))$ for each $a \in \Sigma$

Then,

$$\begin{aligned}\Delta = \{&((p, e, e), (q, S)),\\ &((p, e, S), (q, aSbSa)),\\ &((p, e, S), (q, bSaSa)),\\ &((p, e, S), (q, aSaSb)),\\ &((p, e, S), (q, e)),\\ &((p, a, a), (q, e)),\\ &((p, b, b), (q, e))\}.\end{aligned}$$

d) Write a context-free grammar for the language $L_3 = L_1 \cup L_2$.

Let $G_1 = (V_1, \Sigma_1, R_1, S_1)$ for the language L_1 and $G_2 = (V2, \Sigma_2, R_2, S_2)$ for the language L_2.

Then $G_3 = (V_1 \cup V2 \cup \{S\}, \Sigma_1 \cup \Sigma_2, R_1 \cup R_2 \cup \{S \rightarrow S_1, S \rightarrow S_2\})$ for the language L_3.

$$\begin{aligned}V_3 &= \{S, S_1, S_2, a, b\}\\ \Sigma &= \{a, b\}\\ R &= \{S \rightarrow S_1,\\ &\quad S \rightarrow S_2,\\ &\quad S_1 \rightarrow aS_1bS_1b\\ &\quad S_1 \rightarrow bS_1aS_1b\\ &\quad S_1 \rightarrow bS_1bS_1a\\ &\quad S_1 \rightarrow e\\ &\quad S_2 \rightarrow aS_2bS_2a\\ &\quad S_2 \rightarrow aS_2aS_2b\\ &\quad S_2 \rightarrow bS_2aS_2a\\ &\quad S_2 \rightarrow bS_2a\\ &\quad S_2 \rightarrow aS_2b\\ &\quad S_2 \rightarrow e\\ &\}.\end{aligned}$$

Answer 2

Given $G1 = V, \Sigma, R, S$ where $V = 0, 1, S, A$, $\Sigma = 0, 1$, and $R = S \to AS|e, A \to A1|0A1|01$

a) Show that $G1$ is ambiguous.

For the string 00111, there are two possible leftmost derivations.

b)

Answer 3

a)

b)