

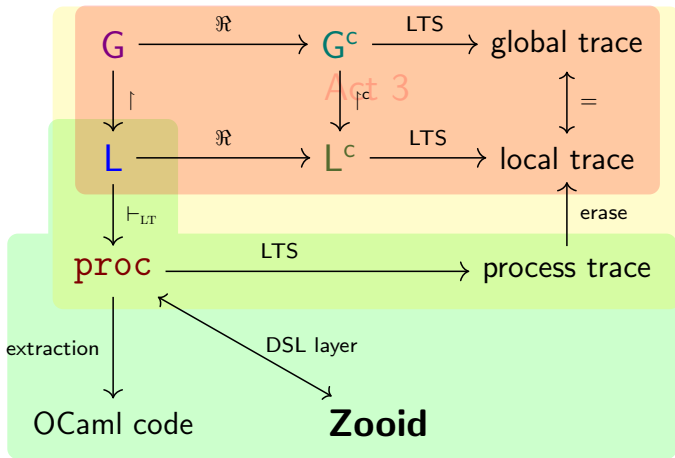
## **Act II**

**Smol-Zooid: multiparty with shallower embedding**

# Goals

1. Certifying **individual** processes of a **distributed system**
2. Extracting runnable code
3. Avoiding complex formalisations of binders, whenever possible

# Overview



# Smol Zooid

- We combine **shallow/deep embeddings** of binders
    - Processes are defined inductively
    - Values are standard Gallina values
    - We use DeBruijn indices for the deeply embedded binders
  - SZooid constructs are **well-typed by construction**
- 
- We leverage **Coq code extraction** mechanism
  - For simplicity, SZooid does not cover choices

## Core Processes

In: <http://github.com/emtst/gentleAdventure>

```
Inductive proc :=  
| Inact | Rec (e : proc) | Jump (X : nat)  
| Send (p : participant) {T : type}  
    (x : interp_type T) (k : proc)  
| Recv (p : participant) {T : type}  
    (k : interp_type T -> proc).
```

# Payload Types

We need to define a type for payload types:

- We need a decidable equality on payload types
- We need a decidable equality on payload values

**Inductive** type := Nat | Bool | ...

**Definition** interp\_type : type -> Type := ...

# Semantics: overview

$$P \xrightarrow{E} P'$$

- What is an event?
- How to manage recursion?

## Semantics: events

The semantics is an LTS:

- the labels are the **communication events**
- it is parameterised by a **payload interpretation function**
- traces are obtained as the greatest fixpoint of the LTS step

**Inductive** action := a\_send | a\_recv.

**Record** event interp\_payload :=

```
{ action_type   : action;  
  subj          : participant;  
  party         : participant;  
  payload_type  : type;  
  payload       : interp_payload payload_type }.
```



## Semantics: Recursion Variables

`p_unroll` exposes the first communication action in a process (unfold recursion):

```
Definition p_unroll e :=  
match e with  
| Rec e' => p_subst 0 e e'  
| e' => e'  
end.
```

## Semantics: Recursion Variables

```
Fixpoint p_subst d e' e :=  
match e with  
| Rec e => Rec (p_subst d.+1 e' e)  
| Jump X => if X == d then p_shift d 0 e' else e  
| ...  
end.
```

```
Example ex_p_subst:  
p_subst 0 (ping_Alice) (Rec (Jump 1)) = Rec ping_Alice.
```

## Semantics: recursion unrolling

*(\* unroll uT. Alice!0. T to Alice!0. uT. Alice!0. T \*)*

Example ex\_p\_unroll:

```
p_unroll (Rec (@Send Alice Nat 0 (Jump 0)))  
= @Send Alice Nat 0 (Rec (@Send Alice Nat 0 (Jump 0))).
```

## Semantics: step

The step of the LTS is defined as a **function**:

```
Definition step' e E :=  
  match e with  
  | Send p T x k =>  
    if (action_type E == a_send) && (party E == p) &&  
      (eq_payload (payload E) x)  
    then Some k else None  
  | Recv p T k => ... | _ => None  
end.  
Definition step e := step' (p_unroll e).
```

## Semantics: step

**Definition** event\_alice: event interp\_type :=  
{ | action\_type := a\_send;  
 from := Bob;  
 to := Alice;  
 payload\_type := Nat;  
 payload := 0 | }.

**Example** ex\_step: step infinite\_ping\_Alice event\_alice  
= Some infinite\_ping\_Alice.

# Local Types

- Local types for processes
- Notion of “Being well-typed”
- Simultaneous construction of processes & well-typeness proof

## Local Types

```
Inductive lty :=  
  | l_end  
  | l_jump (X : nat)  
  | l_rec (k : lty)  
  | l_send (p : participant) (T : type) (l : lty)  
  | l_recv (p : participant) (T : type) (l : lty).
```

# Type System

```
Inductive of_lty : proc -> lty -> Prop :=  
| lt_Send      p T k L x :  
    of_lty k L -> of_lty (@Send p T x k) (l_send p T L)  
| ...  
.
```

```
Example ex_of_lty:  
of_lty infinite_ping_Alice  
(l_rec (l_send Alice Nat (l_jump 0))).
```



## Smol Zooid: Smart Constructors

- It would be tedious to type up both a local type and a process
- Users would need to provide a proof that processes are well-typed

We define **SZooid** (Smol Zooid), to write well-typed processes by construction, avoiding repetition.

## Smol Zooid: Smart Constructors

**Definition**  $\text{SZooid } L := \{ p \mid \text{of\_lty } p \ L \}.$

**Definition**  $\text{z\_Send } p \ T \ x \ L \ (k : \text{SZooid } L)$   
   $: \text{SZooid } (\text{l\_send } p \ T \ L)$   
   $:= \text{exist } \_ \_ (\text{lt\_Send } p \ x \ (\text{proj2\_sig } k)).$

...

## Smol Zoid: Smart Constructors

```
Example Alice_smart_constructor :=  
  z_Rec (z_Send Bob Nat 42 (z_Jump 0)).  
  
(*  
  Alice_smart_constructor:  
    SZoid (l_rec (l_send Bob Nat (l_jump 0)))  
  *)
```

# Conclusion

1. Define processes and local types
2. Semantics of processes
3. Automatic construction of local types
  - We also have *code extraction*
  - and *subject reduction*