

Reduction of an E-R Schema to Tables

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Steps in Database Design

- **Requirements Analysis**
 - user needs; what must database do?
- **Conceptual Design**
 - high level descr (often done with ER model)
- **Logical Design**
 - translate ER into DBMS data model
- **Schema Refinement**
 - consistency, normalization
- **Physical Design**
 - indexes, disk layout
- **Security Design**
 - who accesses what, and how

Reduction of an E-R Schema to Tables

- Converting an E-R diagram to a relational database:
 - Each **entity set** is converted to its' own table.
 - Each **relationship** *can be* (**but may not be**) converted to its' own table.
- Each table has a **number of columns**, which generally corresponding to the **attributes** in the corresponding entity or relationship set.
- The resulting tables can be modified in a variety of ways to support performance, space, or other requirements.

ER Diagram



Entity



Attribute



Relationship



Weak
Entity

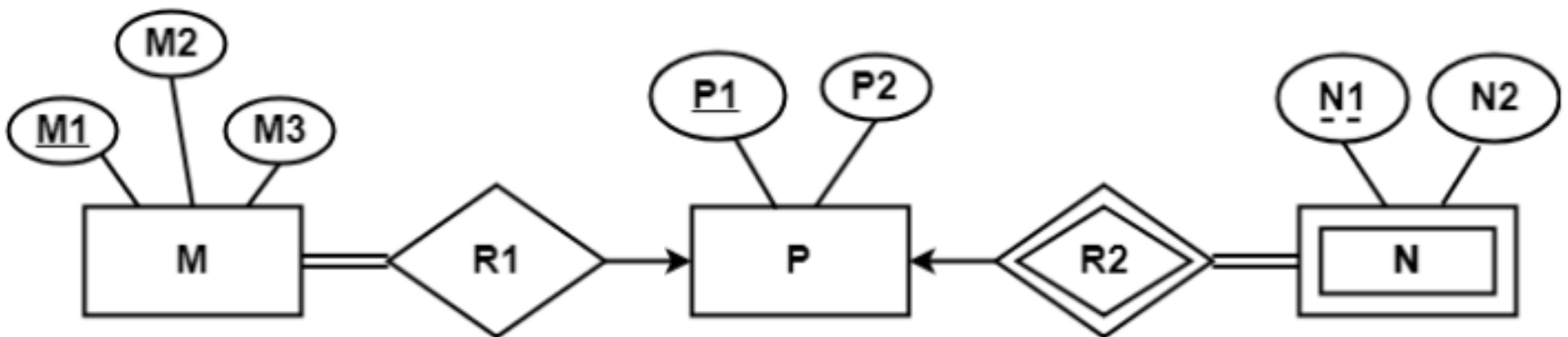


Multivalued
Attribute



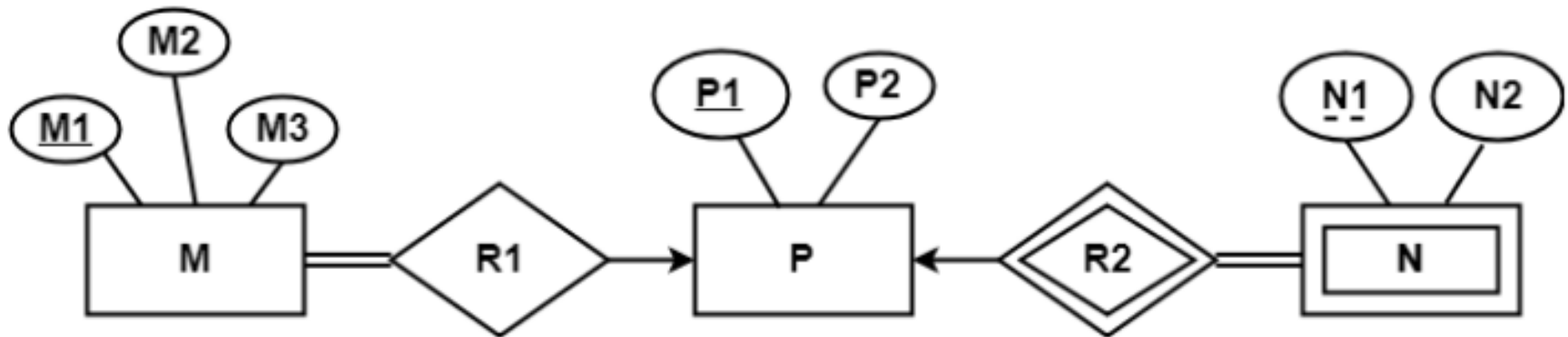
Weak
Relationship

Elements in ER diagrams



1. Representing Strong Entity Sets

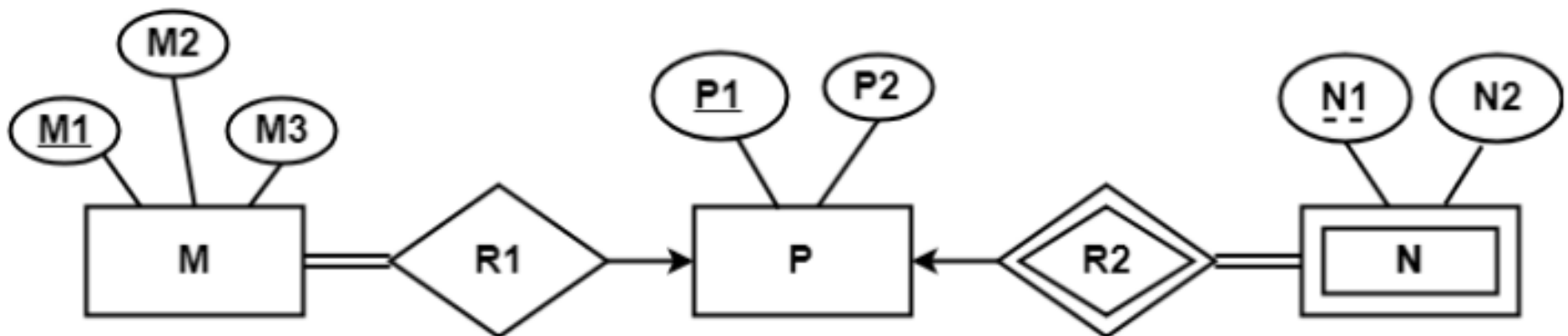
- A **strong entity set** reduces to a table with the **same attributes**.



- M and P are strong entity sets**
- M (M1, M2, M3)
- P (P1, P2)

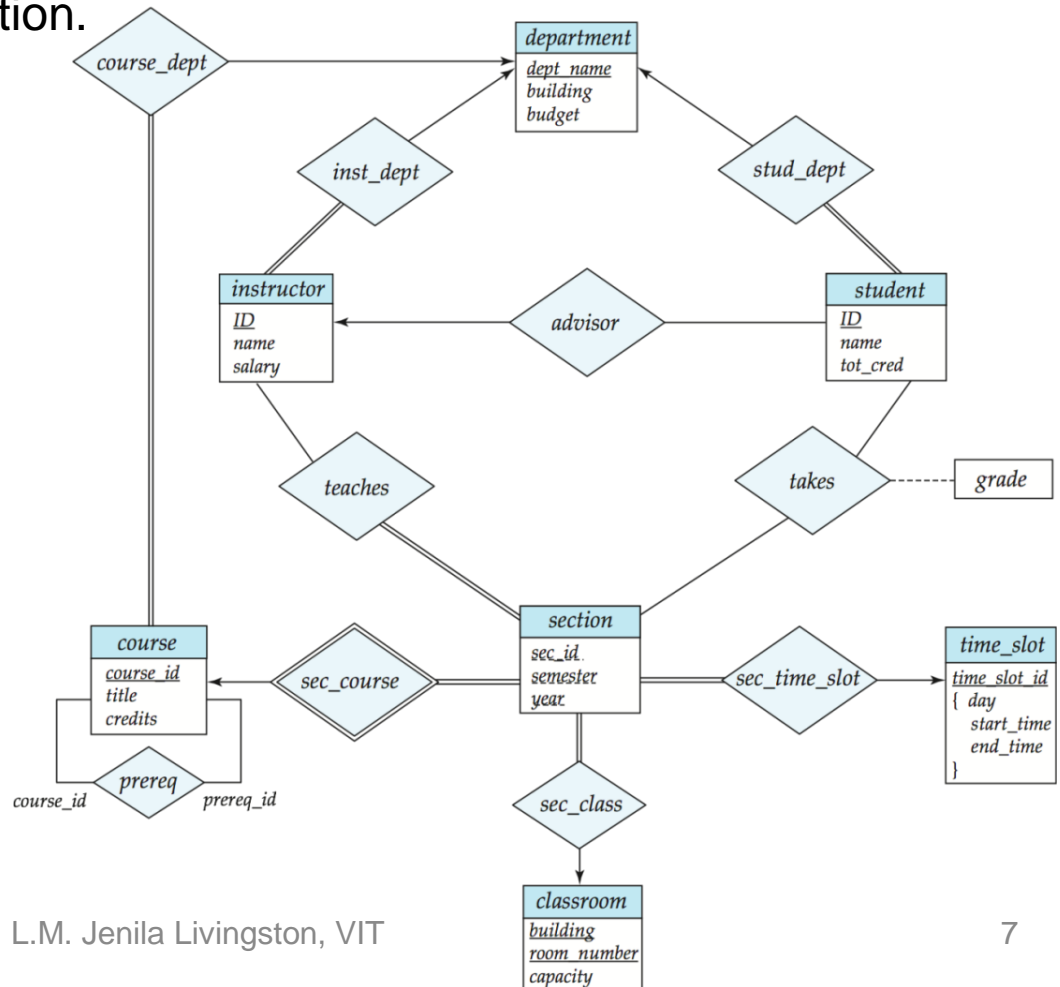
2. Representing Weak Entity Sets

- A weak entity set becomes a table that **includes a column** for the primary key of the identifying strong entity set.
- Weak entity identified by the double rectangle and corresponding relationship identified by the diamond
- N (N1, N2, P1)



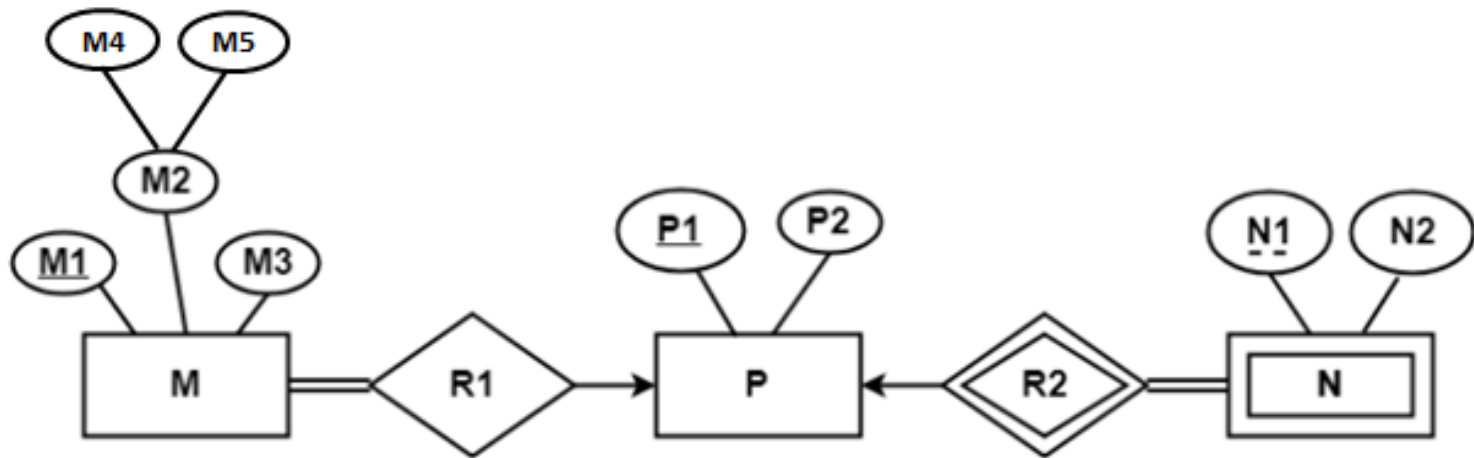
Representing Weak Entity Sets

- In the above ER diagram Section is a weak entity set.
- It is identified by the strong entity set course id. So course_id to be included in the section relation.
- (identified by the double diamond relationship)



3. Composite Attributes

- **Composite attributes** are broken up:
 - Example: attribute *name* with components *first-name* and *last-name* becomes two attributes in the corresponding table – *name_first-name* and *name_last-name*.
- M (M1, M2_M4, M2_M5, M3)



4. Multi-valued Attributes

- A **multi-valued attribute** M of entity E is represented by a **new table** with the following attributes:
 - The primary key of E
 - An attribute corresponding to multi-valued attribute M

- Example:

Entity Set:

employee with attributes *id#*, *name*, *phone#*

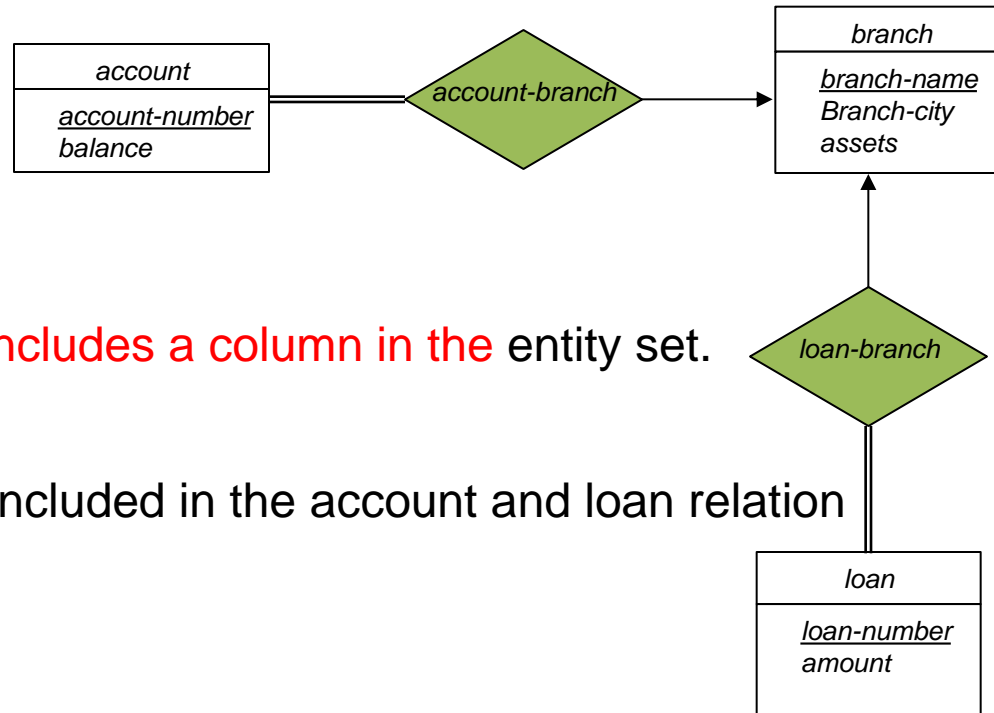
Tables:

employee (*id#*, *name*)

dependent (*id#*, *phone*)

5. Representing Total Participation

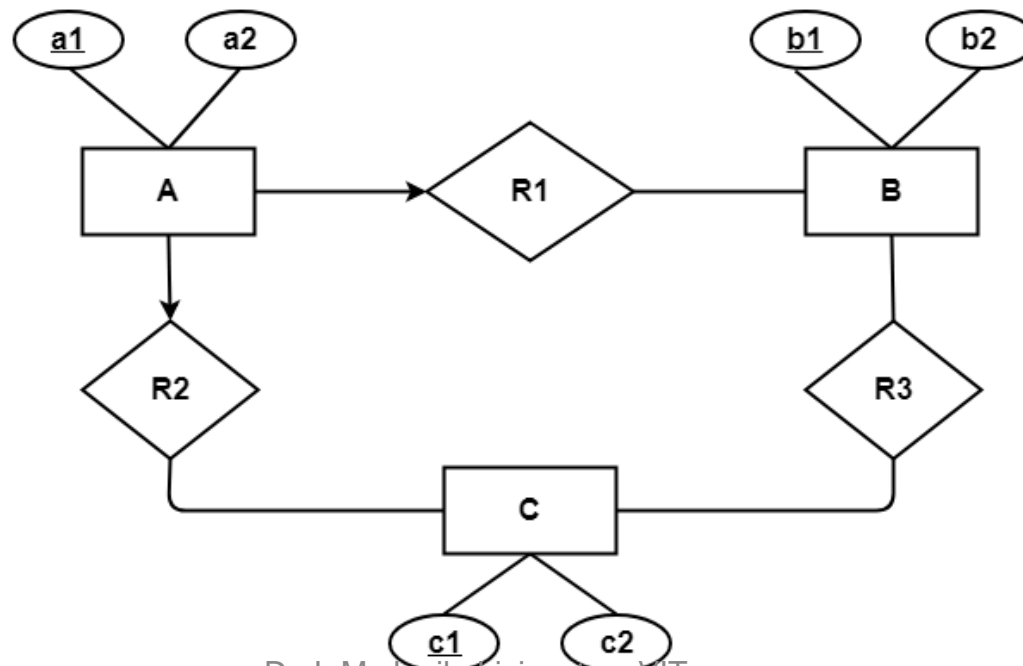
- If every entity in an entity set must participate in a relationship set, then that entity set is said to have total participation in the relationship; indicated by a double-line.



- A total participation **includes a column in the** entity set.
- **branch_name** to be included in the account and loan relation

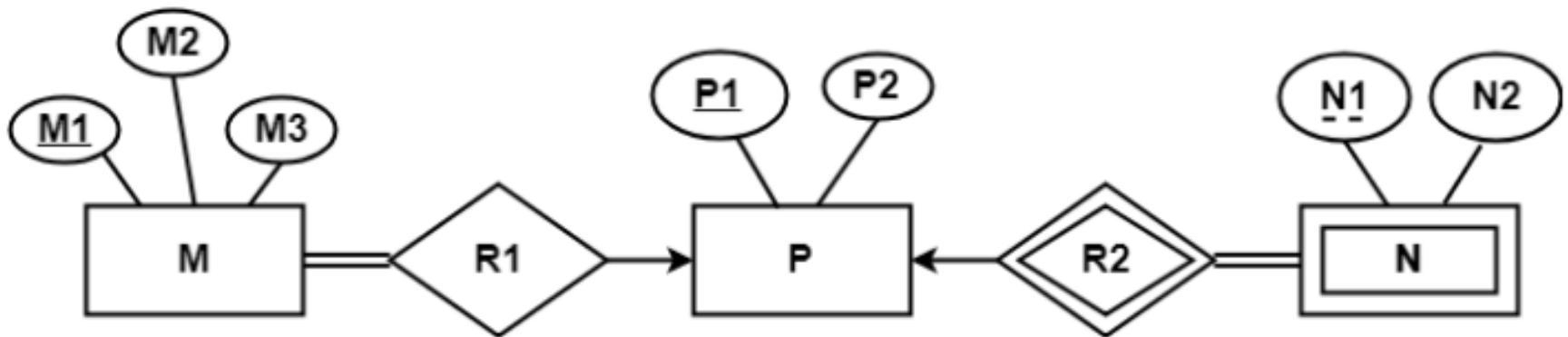
6. Representing Relationship Sets

- A **many-to-many relationship** set is represented as a table with **columns for the primary keys of the two participating entity sets**, and any descriptive attributes of the relationship set.
- Example: R3 (b1, c1)



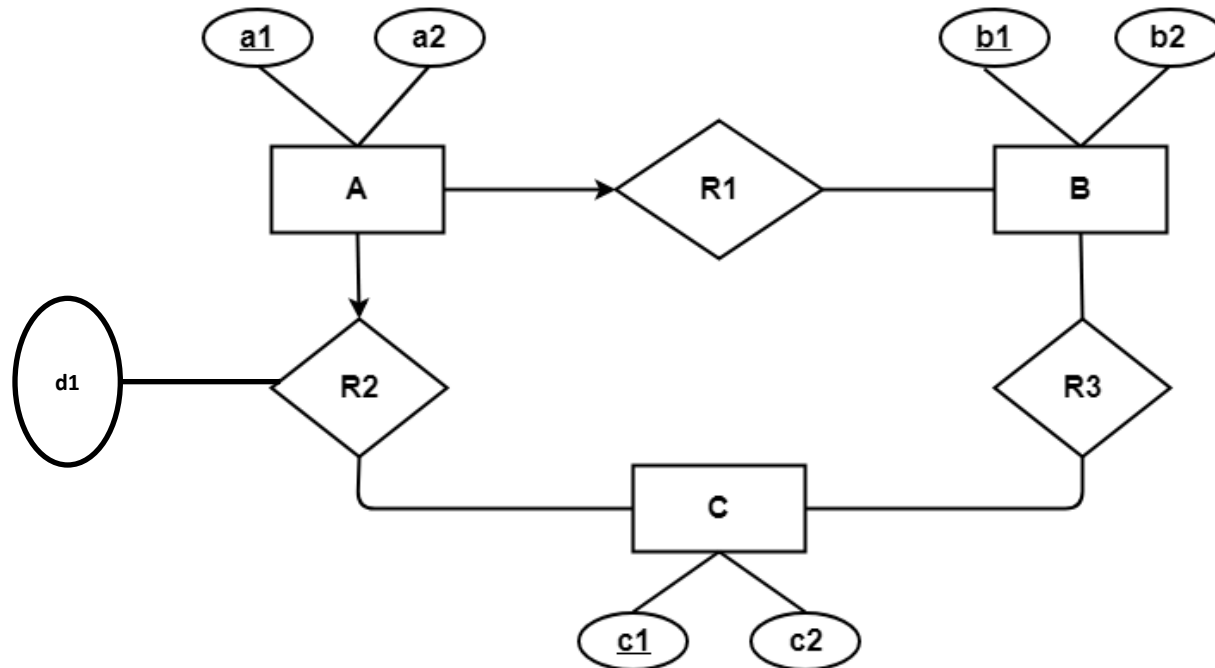
Representing Relationship Sets cont.

- A **many-to-one relationship** set can be represented just like a many-to-many relationship.
- Technically this is not necessary, and in some cases it does not result in a good design.
- If a relationship contains its own attribute then form a relation



Representing Relationship Sets

- In many-to-one relationship, if a relationship contains its own attribute then form a relation
- $R2(d1, a1, c1)$



Representing Relationship Sets, cont.

- For **one-to-one relationship** sets, the **extra attribute can be added to either of the tables** corresponding to the two entity sets.
- If no attribute, no need to form a table

7. Representing Specialization

- Note: This discussion assumes a 2-level inheritance hierarchy.
 - Exercise: Generalize it to an arbitrarily deep hierarchy.
- Method 1:
 - Form a table for the higher level entity set.
 - Form a table for each lower level entity set, including the primary key of the higher level entity set and local attributes.

<i>table</i>	<i>attributes</i>
<i>person</i>	<i>name, street, city</i>
<i>customer</i>	<i>name, credit-rating</i>
<i>employee</i>	<i>name, salary</i>

- One Drawback: getting information about specific entities requires accessing two tables

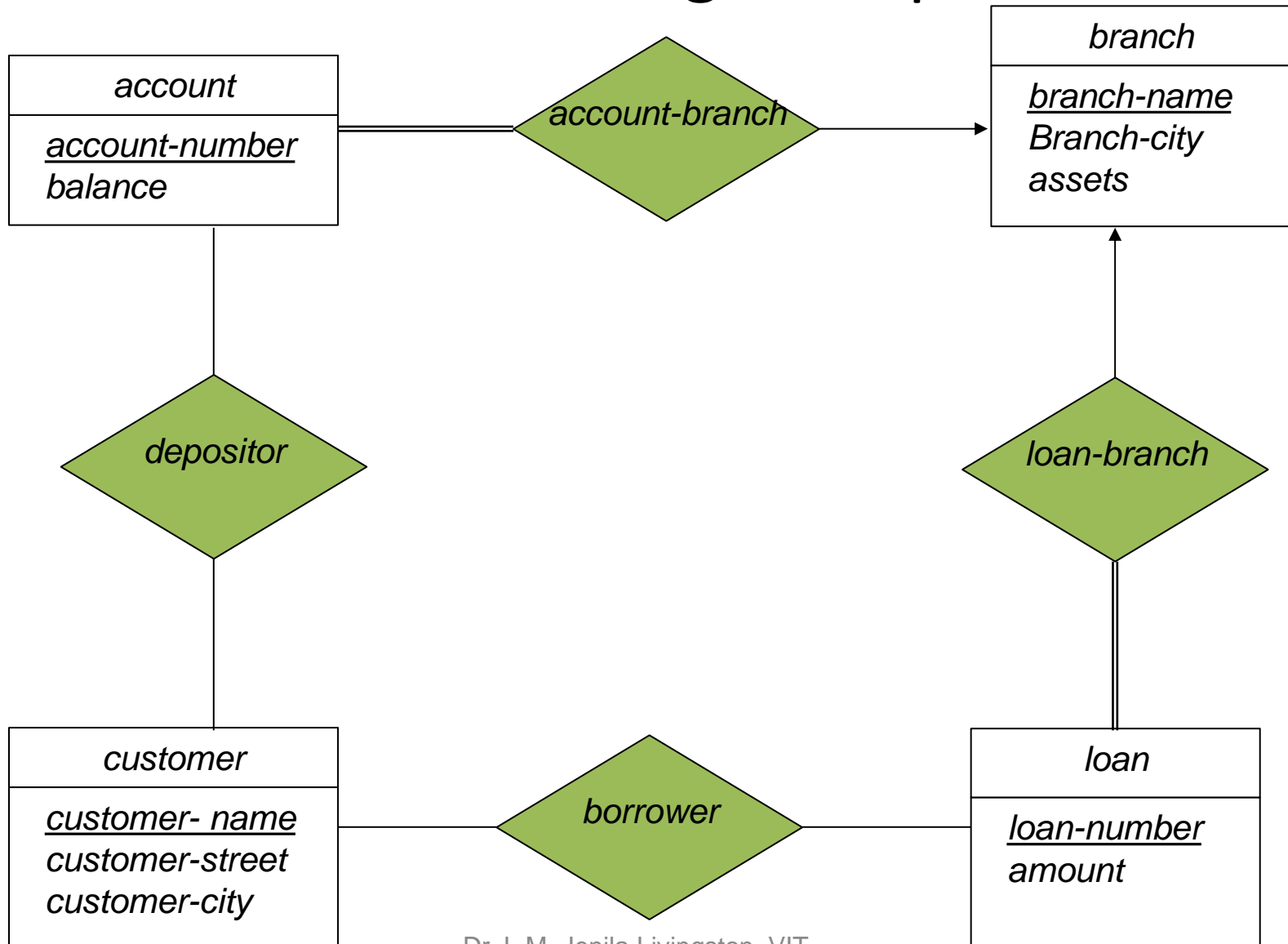
Representing Specialization

- Method 2:
 - Form a table for each entity set with all local and inherited attributes

table	table attributes
<i>person</i>	<i>name, street, city</i>
<i>customer</i>	<i>name, street, city, credit-rating</i>
<i>employee</i>	<i>name, street, city, salary</i>

- This method has obvious redundancies.
 - Particularly bad for persons who are both customers and employees.
- If specialization is total, the table for the generalized entity is redundant.
 - Temptation is to delete the *person* table; still might be needed for foreign key constraints.

E-R Diagram for the Banking Enterprise



Relational Schemes for the Banking Enterprise

- The following relational schemes result:

branch (*branch-name*, *branch-city*, *assets*)

customer (*customer-name*, *customer-street*, *customer-city*)

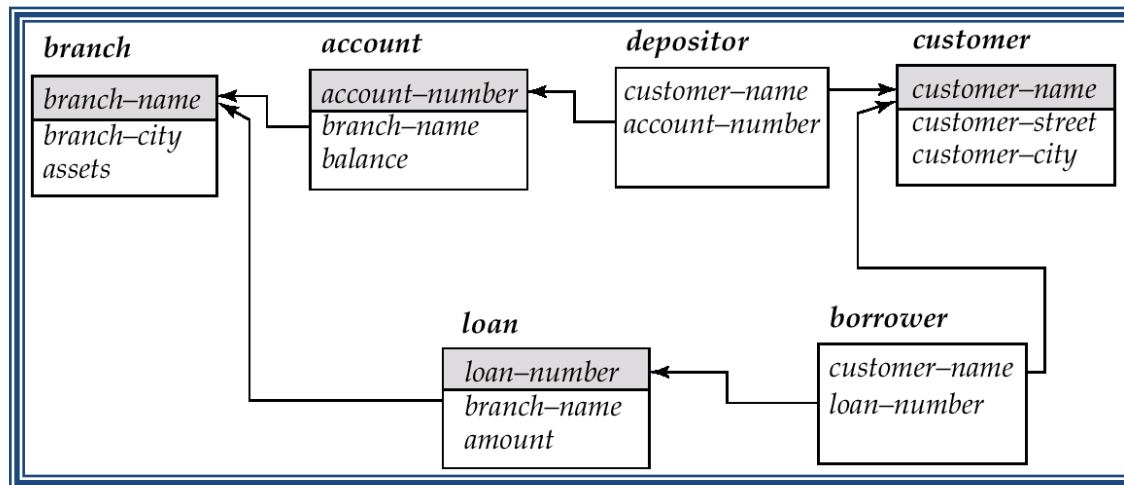
account (*account-number*, *branch-name*, *balance*)

loan (*loan-number*, *branch-name*, *amount*)

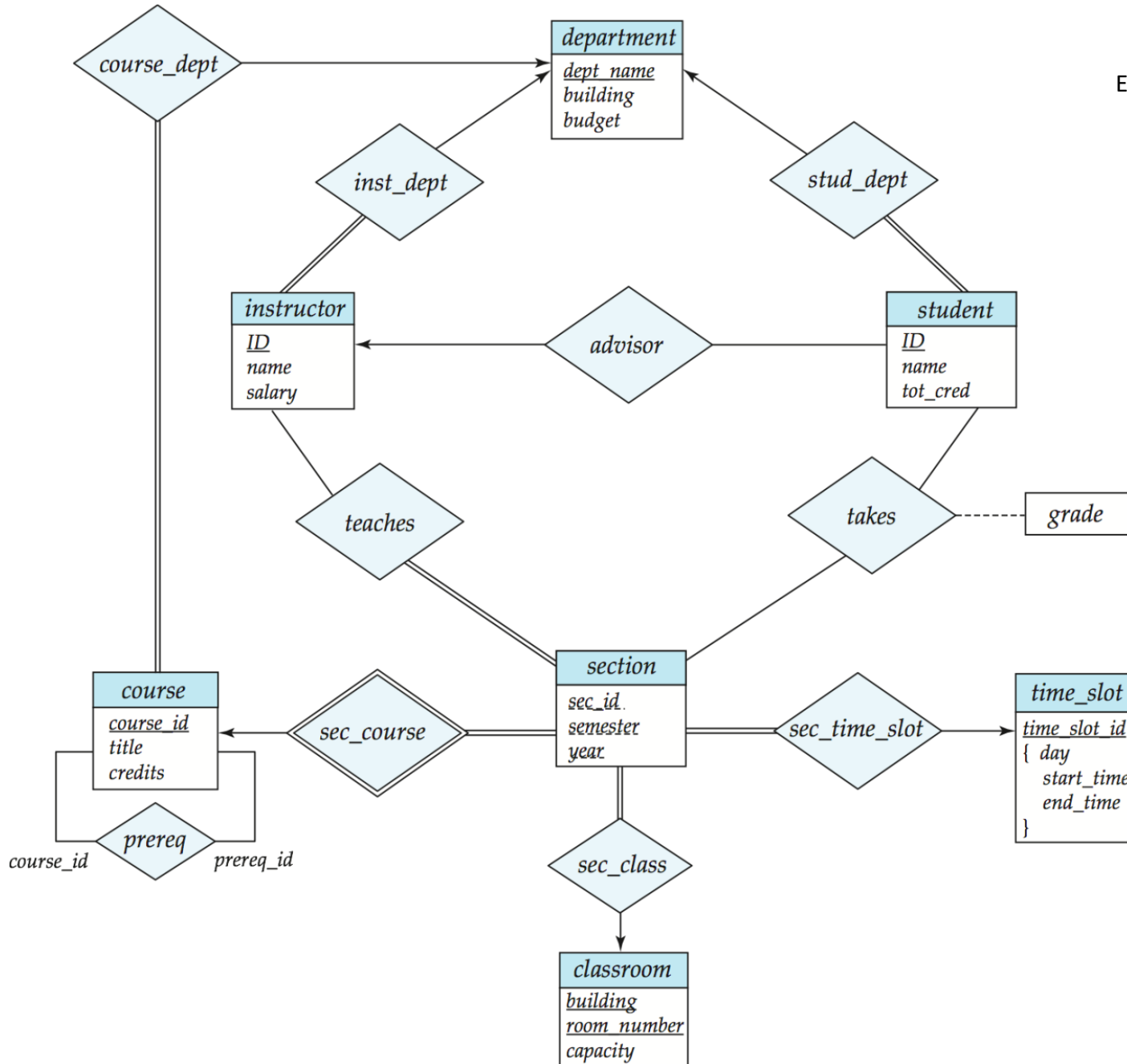
depositor (*customer-name*, *account-number*)

borrower (*customer-name*, *loan-number*)

Schema Diagram for the Banking Enterprise



E-R Diagram for a University



Relational Schema for a University

Classroom (building, room-number, capacity)

Department (dept-name, building, budget)

Course (course-id, title, dept-name, credits)

Instructor (ID, name, depart-name, salary)

Section (course-id, sec-id, semester, year, building, room-number, time-slot-id)

Teaches (ID, course-id, sec-id, semester, year)

Student (ID, name, dept-name, tot-cred)

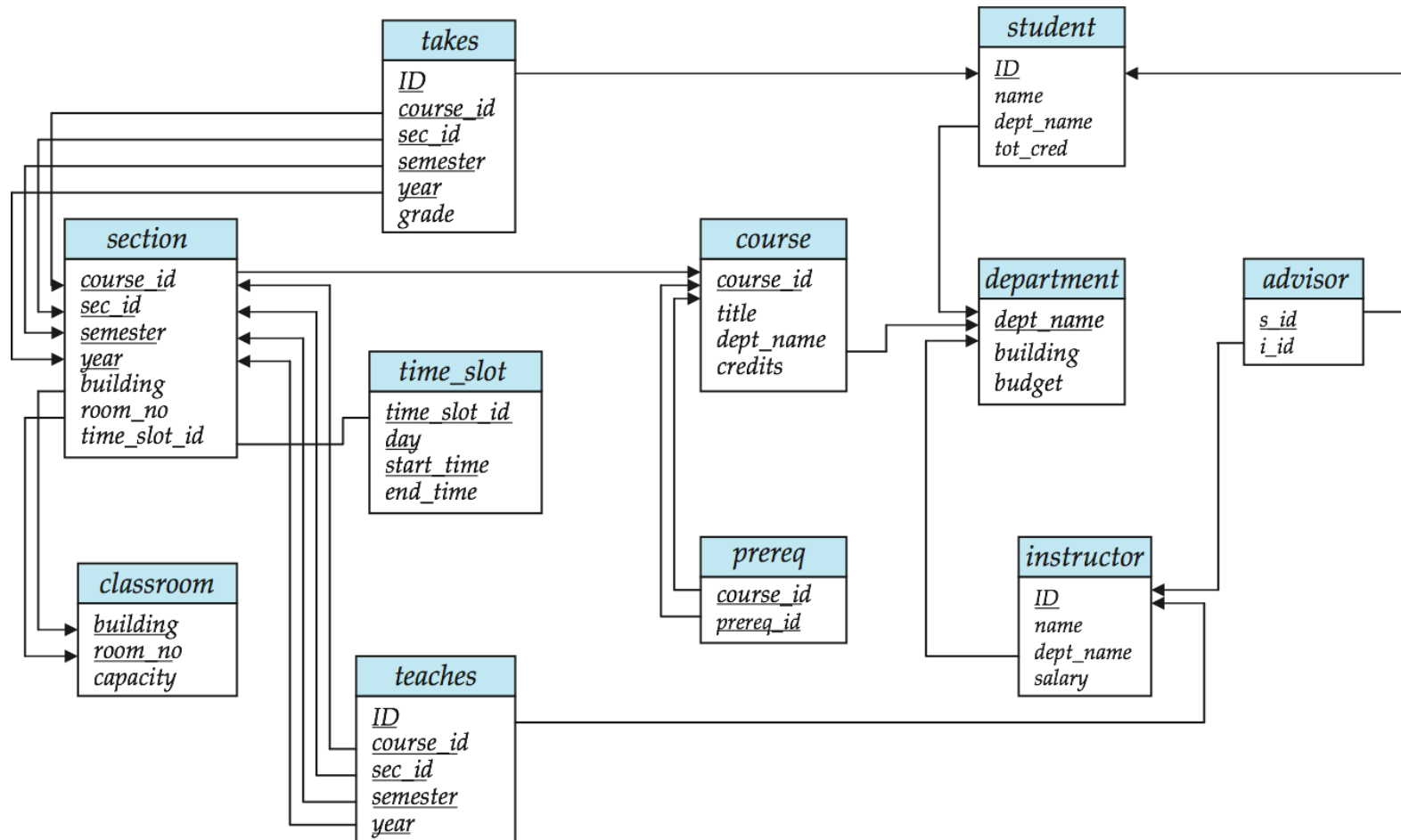
Takes (ID, course-id, sec-id, semester, year, grade)

Advisor (s-ID, i-ID)

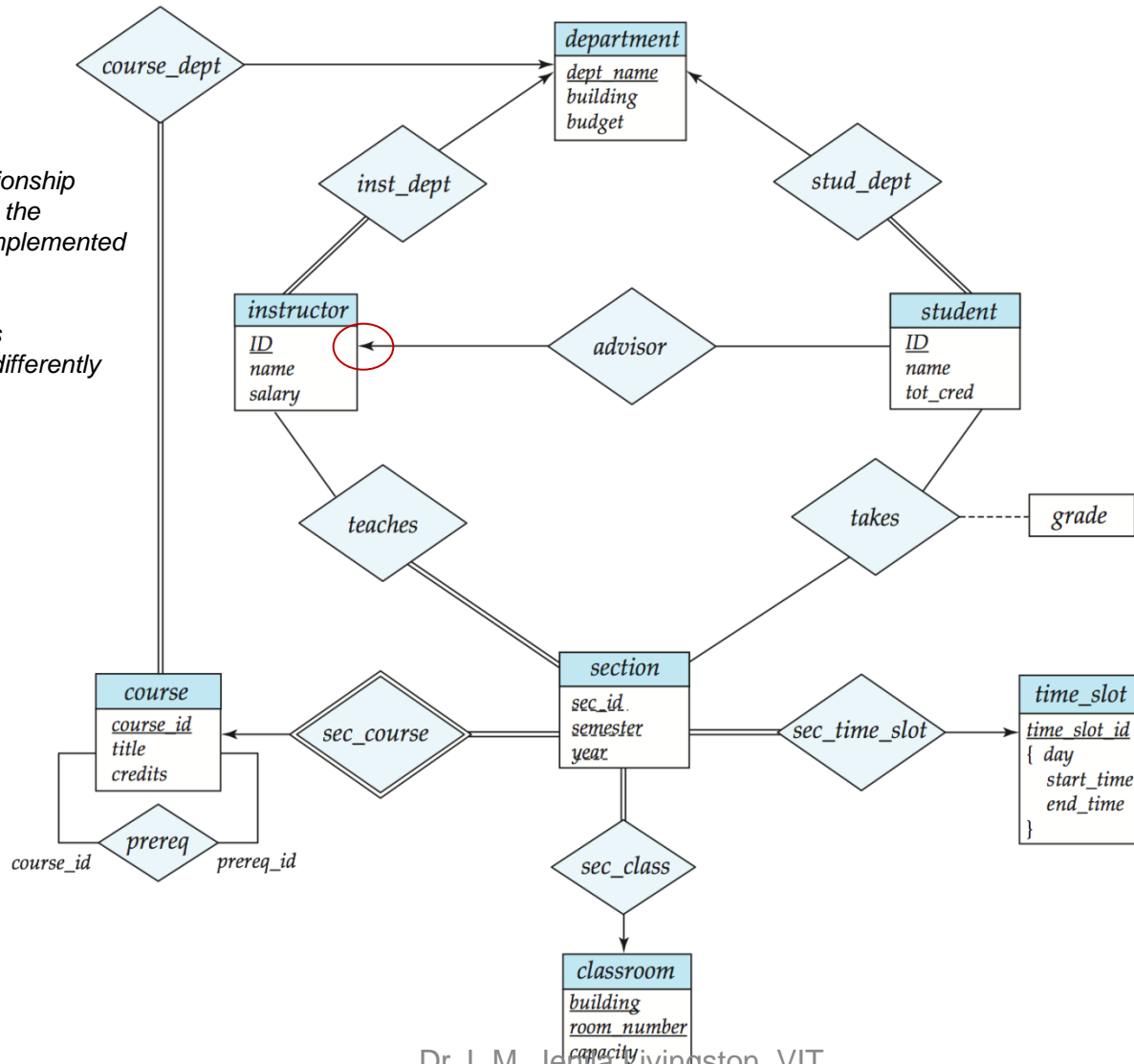
Time-slot (time-slot-id, day, start-time, end-time)

Prereq (course-id, prereq-id)

Schema Diagram for a University



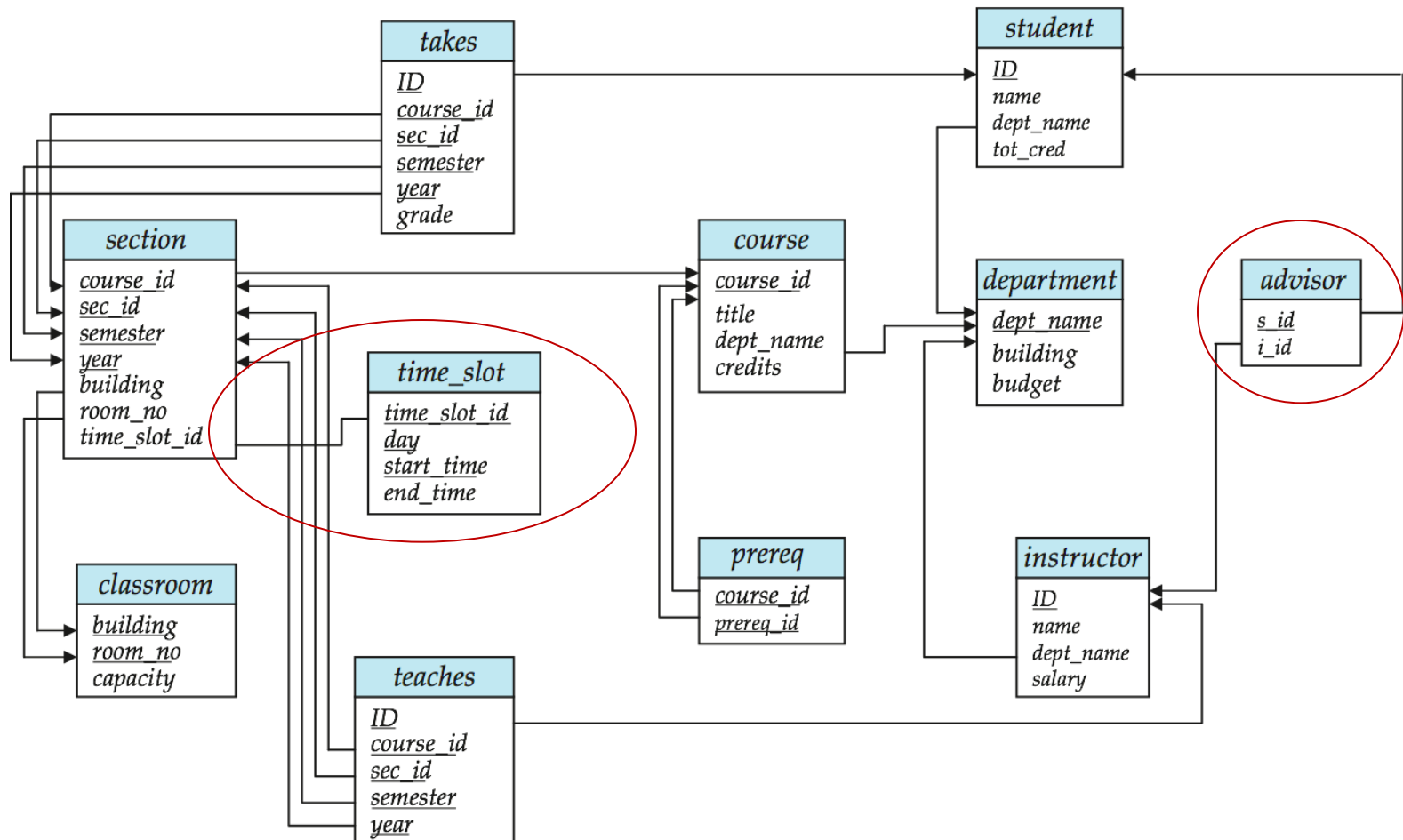
E-R Diagram for a University



Note the advisor relationship
Is one-to-many, but in the
schema it has been implemented
as its' own table..

Similarly, time—slot is
implemented slightly differently
than described...

Schema Diagram for a University



Thank You