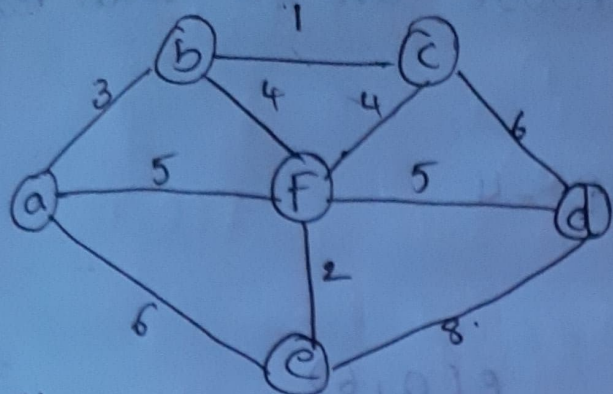


PRIM'S ALGORITHM



vertex $(\downarrow, \rightarrow)$ distance.
any neighbour in min spanning tree.

→ 1st vertex - $a(\text{Null}, \text{Null}) \equiv a(-, -)$

since, it is the 1st vertex it has no adj vertex and no distance.

→ $b(a, 3)$ b has a, c, f as adj. But in min spanning tree only a is present.

$c(-, \infty)$

c does not have a neighbour, so it is considered at ∞ .

$d(-, \infty)$

c, f, e are neighbours.

$e(a, 6)$

a, f, d are neighbours.

only a is in min spanning tree.

$f(a, 5)$

Remaining vertices - Not there in min spanning tree.

$ab - 3$

$ae - 6$

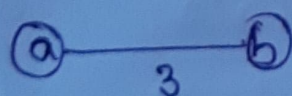
$af - 5$

$b(a, 3) \downarrow$
Min.

$e(a, 6)$

$f(a, 5)$

So, b is added to min spanning tree.



present in min spanning tree.

→ $b(a, 3)$

$c(b, 1)$

b, f, d are neighbours

$d(-, \infty)$

$e(a, 6)$

a, f, d are neighbours

$f(-, \infty)$

a, b, c, d, e are neighbours

When more than 1 neighbours are present in min spanning tree, choose the one with least edge weight.

$$af = 5$$

$$bf = 4$$

$$f(b, 4)$$

$$c(b, 1)$$

$$f(b, 4)$$

$$e(a, 6)$$

Min spanning tree.

$$a(-, -)$$

$$b(a, 3)$$

$$c(b, 1)$$

$$d(c, 6)$$

$$e(a, 6)$$

$$f(\quad) \quad a, b, c \text{ in } \text{min spanning tree.}$$

$$af = 5$$

$$bf = 4$$

$$cf = 4$$

bf (or) cf can be chosen

$$f(b, 4)$$

Least weighted edge - $f(b, 4)$.

Min spanning tree

$$a(-, -)$$

$$b(a, 3)$$

$$c(b, 1)$$

$$f(b, 4)$$

$$e(f, 2)$$

$$ea = 6$$

$$ef = 2$$

$$d(f, 5)$$

$$cd = 6$$

$$df = 5$$

Least weighted edge - $e(f, 2)$

Min spanning tree

$$a(-, -) \quad f(b, 4)$$

$$b(a, 3) \quad e(f, 2)$$

$$c(b, 1)$$

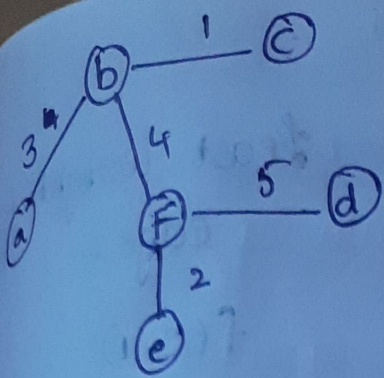
$$d(f, 5)$$

$$df = 5$$

$$de = 8$$

$$dc = 6$$

Add to min spanning tree.



Min Spanning tree -

$a(-, -)$

$d(f, 5)$

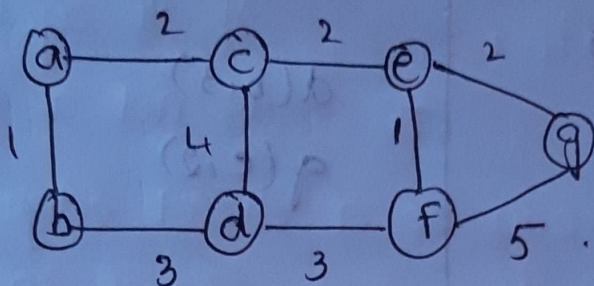
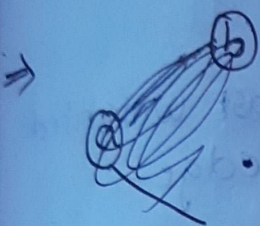
$b(a, 3)$

$e(f, 2)$

$c(b, 1)$

$f(b, 4)$

Total weights = $1 + 3 + 4 + 5 + 2 = 15$



Min spanning tree using Prim's Algorithm.

1st vertex - $a(-, -)$

Min spanning tree - $a(-, -)$

$b(a, 1)$

Least weighted edge - $b(a, 1)$

$c(a, 2)$

Min spanning tree

$a(-, -)$

$b(a, 1)$

$c(a, 2)$

$d(b, 3)$

$e(-, \infty)$

$f(-, \infty)$

$g(-, \infty)$

Least weighted

edge -

$c(a, 2)$

Min spanning tree

Remaining vertices

$a(-, -)$

$b(a, 1)$

$c(a, 2)$

$d(b, 3)$

$e(c, 2)$

$f(-, \infty)$

$g(-, \infty)$

Least weighted

edge -

$e(c, 2)$

Min spanning tree

$a(-, -)$

$b(a, 1)$

$c(a, 2)$

$e(c, 2)$

Min spanning tree

$a(-, -)$

$b(a, 1)$

$c(a, 2)$

$e(c, 2)$

$f(e, 1)$

Min spanning tree

$a(-, -)$

$b(a, 1)$

$c(a, 2)$

$e(c, 2)$

$f(e, 1)$

$g(e, 2)$

Remaining vertices

$d(b, 3)$

$f(e, 1)$

$g(e, 2)$

Least weighted edge -

$f(e, 1)$

Remaining vertices

$d(b, 3)$

$g(e, 2)$

Least weighted edge -

$g(e, 2)$

Remaining vertices

$d(b, 3)$

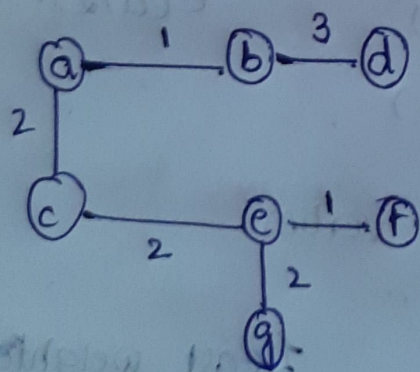
Min spanning tree -

$a(-, -)$, $b(a, 1)$

$c(a, 2)$, $d(b, 3)$

$e(c, 2)$, $f(e, 1)$

$g(e, 2)$



Sum of weights = $1 + 3 + 2 + 1 + 2 + 2 = 11$

ALGORITHM - Prim(G)

Input - A weighted connected graph $G = (V, E)$

V_T - vertices in min spanning tree

E_T - edges in min spanning tree.

Output - E_T , the set of edges composing a min spanning tree.

$V_T \leftarrow \{v_0\}$ the set of tree vertices can be initialized by with any vertex.

$E_T \leftarrow \emptyset$

for $i \leftarrow 1$ to $|V| - 1$ do

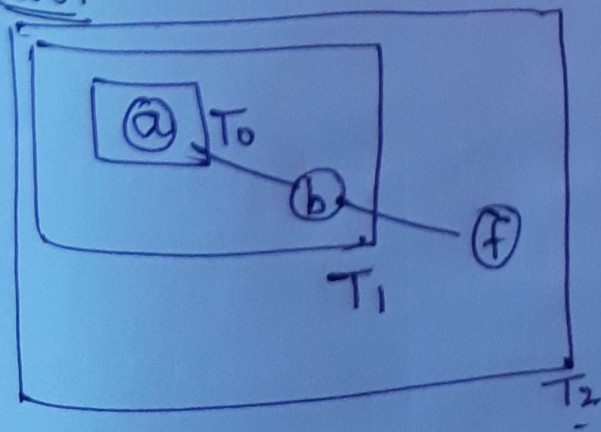
find a minimum-weight edge $e^* = (v^*, u^*)$ among all the edges (v, u) such that v is in V_T and u is in $V - V_T$.

$V_T \leftarrow V_T \cup \{u^*\}$

$E_T \leftarrow E_T \cup \{e^*\}$

return E_T .

PROOF -



~~From~~ To T_{i-1} we add an elt to get T_i .

Prove that T_i gives min spanning tree.