

# Propositional Equivalences

Course Code: CSC 1204

Course Title: Discrete Mathematics



**Dept. of Computer Science**  
**Faculty of Science and Technology**

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# Lecture Outline



## 1.2 Propositional Equivalences

- Tautology
- Contradiction
- Contingency
- Logical Equivalences

# Objectives and Outcomes



- **Objectives:** To understand the terms Tautology, Contradiction, Contingence with examples, to understand the standard logical equivalences, to determine whether a compound proposition is a Tautology or Contradiction, to determine whether two compound propositions are logically equivalent.
- **Outcomes:** Students are expected to be able to write the definitions of Tautology, Contradiction and Contingency with examples, be able to determine whether a compound proposition is a Tautology or Contradiction using a Truth Table and standard logical equivalences, be able to determine whether two compound propositions are logically equivalent using a Truth Table and logical equivalences.

# Tautology



***Tautology***: A compound proposition that is always true is called a tautology.

**Examples:**

a)  $p \vee \neg p$

b) The professor is either a woman or a man

c) People either like watching TV or they don't

# Contradiction



***Contradiction***: A compound proposition that is always false is called a contradiction.

**Examples:**

a)  $p \wedge \neg p$

b)  $x$  is prime and  $x$  is an even integer greater than 8

c) All men are good and all men are bad

# Examples of *Tautology* and *Contradiction*



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**TABLE 1** Examples of a Tautology and a Contradiction.

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
T	F	T	F
F	T	T	F

# Contingency



**Contingency:** A compound proposition that is neither a tautology nor a contradiction is called a contingency. In other words, a compound proposition whose truth value is not constant is called a contingency.

Examples:

a)  $p \rightarrow \neg p$

b)  $p$

c)  $\neg p$

# How to determine whether a compound proposition is a Tautology or Contradiction?



- We can determine whether a compound proposition is a **Tautology** or **Contradiction** in **two ways**:
  - 1) Using a **truth table** – The **easiest** way to see if a compound proposition is a tautology or contradiction is to **use a truth table**. Show that the compound proposition is always true
  - 2) Using (**Laws of**) **Logical Equivalences**



# Tautology : Example



Show that  $[\neg p \wedge (p \vee q)] \rightarrow q$  is a *tautology* using a *Truth Table*

# Solution



$p$	$q$	$\neg p$	$p \vee q$	$\neg p \wedge (p \vee q)$	$[\neg p \wedge (p \vee q)] \rightarrow q$
T	T				
T	F				
F	T				
F	F				

# Solution



$p$	$q$	$\neg p$	$p \vee q$	$\neg p \wedge (p \vee q)$	$[\neg p \wedge (p \vee q)] \rightarrow q$
T	T	F			
T	F	F			
F	T	T			
F	F	T			

# Solution



$p$	$q$	$\neg p$	$p \vee q$	$\neg p \wedge (p \vee q)$	$[\neg p \wedge (p \vee q)] \rightarrow q$
T	T	F	T		
T	F	F	T		
F	T	T	T		
F	F	T	F		

# Solution



$p$	$q$	$\neg p$	$p \vee q$	$\neg p \wedge (p \vee q)$	$[\neg p \wedge (p \vee q)] \rightarrow q$
T	T	F	T	F	
T	F	F	T	F	
F	T	T	T	T	
F	F	T	F	F	

# Solution



$p$	$q$	$\neg p$	$p \vee q$	$\neg p \wedge (p \vee q)$	$[\neg p \wedge (p \vee q)] \rightarrow q$
T	T	F	T	F	T
T	F	F	T	F	T
F	T	T	T	T	T
F	F	T	F	F	T

Since the truth table shows all the true values of compound proposition  $[\neg p \wedge (p \vee q)] \rightarrow q$  are true(T), so it is a tautology.

# Class Work



- 1) Determine whether  $\neg (p \wedge q) \vee p$  is a tautology or contradiction.
- 2) Determine whether  $p \wedge (q \wedge \neg p)$  is a tautology or contradiction.

# Logical Equivalences



- Compound propositions that have the same truth values in all possible cases are called **logically equivalent**.
- **Definition** : Compound propositions  $p$  and  $q$  are *logically equivalent* if  $p \leftrightarrow q$  is a tautology (denoted by  $p \equiv q$  **or**  $p \Leftrightarrow q$  )

NOTE: We will use the notation:  $p \equiv q$



# How to determine whether two compound propositions are logically equivalent?



- We can determine whether two compound propositions are logically equivalent in two ways:
  - 1) Using a **Truth Table**
  - 2) Using (laws of ) **Logical Equivalences**



# Using a Truth Table to determine whether two compound propositions are logically equivalent

- **Two compound propositions are *logically equivalent*** if they always have the same truth values in the corresponding rows.
- Construct a truth table for the given two compound propositions **[in one table]**
- **If the truth values of both of the compound propositions are same in the corresponding rows, then they are logically equivalent.**
- **If the true values of both of the compound propositions are different in one or more rows, then they are NOT logically equivalent.**

# Example 1



Show that  $p \leftrightarrow q$  is logically equivalent to  $(p \rightarrow q) \wedge (q \rightarrow p)$

$p$	$q$	$p \rightarrow q$	$q \rightarrow p$	$p \leftrightarrow q$	$(p \rightarrow q) \wedge (q \rightarrow p)$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	T	F	F	F
F	F	T	T	T	T

Since the truth values of both of the compound propositions are same in the corresponding rows, they are logically equivalent.

# Class Work



Show that  $p \vee (q \wedge r)$  and  $(p \vee q) \wedge (p \vee r)$  are logically equivalent

# Solution



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**TABLE 5** A Demonstration That  $p \vee (q \wedge r)$  and  $(p \vee q) \wedge (p \vee r)$  Are Logically Equivalent.

$p$	$q$	$r$	$q \wedge r$	$p \vee (q \wedge r)$	$p \vee q$	$p \vee r$	$(p \vee q) \wedge (p \vee r)$
T	T	T	T	T	T	T	T
T	T	F	F	T	T	T	T
T	F	T	F	T	T	T	T
T	F	F	F	T	T	T	T
F	T	T	T	T	T	T	T
F	T	F	F	F	T	F	F
F	F	T	F	F	F	T	F
F	F	F	F	F	F	F	F

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Since the truth values of both of the compound propositions are same in the corresponding rows, they are logically equivalent.

# Logical Equivalences

Table 6 ( page 24 ) → Rosen, 7<sup>th</sup> edition



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TABLE 6 Logical Equivalences.	
<i>Equivalence</i>	<i>Name</i>
$p \wedge \mathbf{T} = p$ $p \vee \mathbf{F} = p$	Identity laws
$p \vee \mathbf{T} = \mathbf{T}$ $p \wedge \mathbf{F} = \mathbf{F}$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r = p \vee (q \vee r)$ $(p \wedge q) \wedge r = p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) = (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) = p$	Absorption laws
$p \vee \neg p = \mathbf{T}$ $p \wedge \neg p = \mathbf{F}$	Negation laws



# A very Useful Logical Equivalence(ULE)

$$p \rightarrow q \equiv \neg p \vee q$$



# Example 1

Show that  $\neg(p \rightarrow q)$  and  $p \wedge \neg q$  are logically equivalent.

## Solution:

$$\begin{aligned}\neg(p \rightarrow q) &\equiv \neg(\neg p \vee q) && \text{by ULE} \\ &\equiv \neg(\neg p) \wedge \neg q && \text{by the second De Morgan law} \\ &\equiv p \wedge \neg q && \text{by the double negation law}\end{aligned}$$



# Example 7 (page 26)



Show that  $\neg(p \vee (\neg p \wedge q))$  and  $\neg p \wedge \neg q$  are logically equivalent by developing a series of logical equivalences.

## Solution:

$\neg(p \vee (\neg p \wedge q))$	$\equiv \neg p \wedge \neg(\neg p \wedge q)$	by the second De Morgan law
$\equiv \neg p \wedge [\neg(\neg p) \vee \neg q]$		by the first De Morgan law
$\equiv \neg p \wedge (p \vee \neg q)$		by the double negation law
$\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q)$		by the second distributive law
$\equiv \mathbf{F} \vee (\neg p \wedge \neg q)$		because $\neg p \wedge p \equiv \mathbf{F}$
$\equiv (\neg p \wedge \neg q) \vee \mathbf{F}$		by the commutative law for disjunction
$\equiv \neg p \wedge \neg q$		by the identity law for $\mathbf{F}$

Consequently  $\neg(p \vee (\neg p \wedge q))$  and  $\neg p \wedge \neg q$  are logically equivalent.

# Exercise



Show that  $(\neg p \wedge (p \vee q)) \rightarrow q$  is a **tautology** using a series of logical equivalences.

# Solution



$$(\neg p \wedge (p \vee q)) \rightarrow q$$

$$\equiv ((\neg p \wedge p) \vee (\neg p \wedge q)) \rightarrow q$$

Distributive Law

$$\equiv (F \vee (\neg p \wedge q)) \rightarrow q$$

Negation Law

$$\equiv (\neg p \wedge q) \rightarrow q$$

Identity Law

$$\equiv \neg (\neg p \wedge q) \vee q$$

ULE

$$\equiv (\neg(\neg p) \vee \neg q) \vee q$$

De Morgan's Law

$$\equiv (p \vee \neg q) \vee q$$

Double Negation Law

$$\equiv p \vee (\neg q \vee q)$$

Associative Law

$$\equiv p \vee T$$

Domination Law

$$\equiv T \quad \text{So, } (\neg p \wedge (p \vee q)) \rightarrow q \text{ is a tautology.}$$

# Summary



- What is Tautology and Contradiction? What is Contingency?
- How to show/determine whether two compound propositions are logically equivalent?
  - Using a truth table
  - Using logical equivalences
- How to show whether a compound proposition is a tautology?
  - Using a truth table
  - Using logical equivalences
- **Note:** Make sure you learn the important Logical Equivalences in Table 6 (page 24) & ULE (  $p \rightarrow q \equiv \neg p \vee q$  )
- **Practice @ Home:** Relevant Odd-numbered Exercises (e.g. 1, 3, 7, 9, 11, 15, 17 )

# Practice @ Home



## \* Practice questions 1-4 **without using a Truth Table**

1. Determine whether  $(\neg p \wedge (q \rightarrow p)) \rightarrow \neg q$  is tautology.
2. Determine whether  $(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$  is tautology.
3. Show that  $(p \wedge (p \rightarrow q)) \rightarrow q$  is a tautology.
4. Show that  $((p \rightarrow q) \wedge (p \rightarrow r))$  and  $(p \rightarrow (q \wedge r))$  are logically equivalent.

## \*\* Practice relevant Odd-Numbered Exercises

**=> 1, 3, 7, 9, 11, 15, 17**



# Answer 1

$$(\neg p \wedge (q \rightarrow p)) \rightarrow \neg q$$

$$\equiv (\neg p \wedge (\neg q \vee p)) \rightarrow \neg q$$

$$\equiv ((\neg p \wedge \neg q) \vee (\neg p \wedge p)) \rightarrow \neg q$$

$$\equiv ((\neg p \wedge \neg q) \vee F) \rightarrow \neg q$$

$$\equiv (\neg p \wedge \neg q) \rightarrow \neg q$$

$$\equiv \neg(\neg p \wedge \neg q) \vee \neg q$$

$$\equiv (p \vee q) \vee \neg q$$

$$\equiv p \vee (q \vee \neg q)$$

$$\equiv p \vee T$$

$$\equiv T$$

So,  $(\neg p \wedge (q \rightarrow p)) \rightarrow \neg q$  is a tautology.

[ULE]

[Distributive Law]

[Negation Law]

[Identity Law]

[ULE]

[De Morgan's & Double Negation Law]

[Associative Law]

[Negation Law]

[Domination Law]

# Answer 2



$$(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$$

$$\equiv (\neg q \wedge (\neg p \vee q)) \rightarrow \neg p \text{ [You MUST write the names of the laws]}$$

$$\equiv (\neg q \wedge \neg p) \vee (\neg q \wedge q) \rightarrow \neg p$$

$$\equiv (\neg q \wedge \neg p) \vee F \rightarrow \neg p$$

$$\equiv (\neg q \wedge \neg p) \rightarrow \neg p$$

$$\equiv \neg (\neg q \wedge \neg p) \vee \neg p$$

$$\equiv q \vee p \vee \neg p$$

$$\equiv q \vee T$$

$$\equiv T$$

So,  $(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$  is a tautology.



# Answer 3



$$(p \wedge (p \rightarrow q)) \rightarrow q$$

$$\equiv (p \wedge (\neg p \vee q)) \rightarrow q$$

$$\equiv ((p \wedge \neg p) \vee (p \wedge q)) \rightarrow q$$

$$\equiv (F \vee (p \wedge q)) \rightarrow q$$

$$\equiv (p \wedge q) \rightarrow q$$

$$\equiv \neg(p \wedge q) \vee q$$

$$\equiv (\neg p \vee \neg q) \vee q$$

$$\equiv \neg p \vee (\neg q \vee q)$$

$$\equiv \neg p \vee T$$

$$\equiv T$$

**ULE** (Substitution for  $\rightarrow$  )

Distributive Law

Negation Law

Identity Law

**ULE** (Substitution for  $\rightarrow$  )

First De Morgan's Law

Associative Law

Negation Law

Domination Law





# Books

- *Discrete Mathematics and its applications with combinatorics and graph theory (7<sup>th</sup> edition)* by Kenneth H. Rosen [Indian Adaptation by KAMALA KRITHIVASAN], published by McGraw-Hill



# References

1. Discrete Mathematics, *Richard Johnsonbaugh*, Pearson education, Inc.
2. Discrete Mathematical Structures, *Bernard Kolman, Robert C. Busby, Sharon Ross*, Prentice-Hall, Inc.
3. *SCHAUM'S outlines Discrete Mathematics*(2<sup>nd</sup> edition), by *Seymour Lipschutz, Marc Lipson*