

SOFTENG370 Notes 2017

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1 Lecture 1

1.1 Generics

Operating System The software that makes the computer usable. Using modern computers without an OS is "impossible"

Examples: Windows, OSX, Linux, Unix, iOS, Android, etc...

1.2 Approaches to Understanding

Minimalist

- mostly going to be using this one
- OS contains minimum amount of software to function
- archlike

Maximalist

- All software comes with standard OS release.
- Contains many utilities and programs.
- ubuntuish

1.3 Usable vs Efficient

- make sure you make OS suited for needs
- either specialised or more general purpose
- Think of who you expect to use the system
- If creating a realtime system with potentially thousands of operations in a short amount of time, have to consider efficiency
- Same with battery life if you expect the system to be used in a mobile setting.

1.4 OS themes

Manager Model

- OS is collection of managers, ensuring proper use of devices.
- Managers are independent.
- look out for everything associated with computer
- tie in with hardware. Current state of HW lets OS do more/less things

Onion Model

- Onions have layers (Abstractions)
- resources contained in lower layers.
- Lower layers can't access higher level layers but other way around possible
- Very difficult to get these layers 'right'
- can use in terms of security. Very good idea

Resource Allocator Model

- similar to manager model
- emphasis on fairness and providing services

Dustbin Model

- contains middleware that not considered part of OS
- Sees OS as bits no-one wants to do

Getting Work Done Model

- Idea of it is we use computers to do something else.
- Goal for OS is to help be able to get it all done.

1.5 OS design

1.5.1 Themes

All in one

- All OS components freely interact with each other
- MS-DOS and Early Linux

Separate Layers (Onion Model)

- Simplify verification and debugging
- Correct design difficult to get

Modules

- All in one with modules for some features
- Linux and Windows.

Microkernels

- Client/Server model
- make OS as small as possible
- **Exokernel** puts kernel outside. OS's job only need to authenticate people to use hardware.

VMs

- Java is an example of this

1.5.2 MS-DOS

- Written to provide the most functionality in the least amount of space
- not divided into modules
- Something exokernels trying to do. Make application program access hardware directly.

1.5.3 Early Unix

- UNIX OS in 2 parts. **Kernel** and **System Programs**
- Provides:
 - File System
 - CPU scheduling
 - Memory management
 - Other OS functions

- Ken Thompson and Dennis Ritchie
- Make OS as simple as possible.
- Simple 2 letter commands.
- Ideas of pipelining and process communication

1.5.4 THE Multiprogramming System

- THE was the first to use the layered system
- Contains 6 layers:
 - 5 User programs
 - 4 Input/Output buffering
 - 3 Operator-Console device driver
 - 2 Memory Management
 - 1 CPU scheduling
 - 0 Hardware

1.5.5 WinNT and Client/Server

- WinNT still being still run
 - Win10 now has Windows Subsystem for Linux
- NT provide env subsystem to run code written for differnt OS
- NT and successors are hybrid systems. Parts are layered but some merged to improve performance.

2 Lecture 2: History of OS

- Started at mainframes.
 - Early PDAs were similar to mainframes. Had no memory protection.
- Then go to Minicomputers
- And then desktop
- And how handheld computers

Each of these stages go through cycle of:

1. No software
2. Compulers
3. Multiuser
4. Networked
5. Clustered
6. Distributed Systems
7. Multiprocessor & Fault tolerant.

2.1 Total Control

- Computers expensive in 50s. Data and programs were saved on paper tape.
- Programmers knew how the computer worked. They were very knowledgable about computers.
 - Prepared program and data cards
 - do setup
 - control computer
 - debug
- Computers did 10,000s instructions per second, but were idle a lot of the time.

2.2 Properties of old OS

- **IO polling**, since no other programs running in background, therefore just waiting on input and able to just poll.
- No file system
- No memory management or security
- OS defined by decisions made by user.
- Single program at a time

2.3 Progression: Operators & Offlining

Operators

- Goal is to reduce the time CPU was doing nothing.
- Operators now just "use" the computer. No need for programmer.
 - If something crashes, then just start the next program.
 - Batch similar jobs together, maximise usage of computer.

Offlining

- Form of parallelism in early computing.
- With Big Expensive Computer BEC, but they are just waiting for IO a lot of time. Therefore want to make IO as fast as possible.
- Use smaller computers to convert slower paper to faster magnetic tape. Then that magnetic tape is used as IO for the BEC
- This is the same for output. Have another smaller cheaper computer offload the output magnetic tape from BEC to a printer.

Resident Monitor

- Keep some code in memory.
- It did the work that some operators were doing.
 - clearing memory
 - reading start of new program that needs to be loaded.
 - Can also do some of the IO routines.

Control Programs Standardise the language to communicate with the Resident Monitor. Had tags for things such as \$JOB (for signifying jobs), \$FTN (When fortran compiler needed), and \$END (signifying end of program)

Conclusions from this

- Memory management and file system still not present. Therefore still need to reset if anything bad happens.
- Security patchy at best
- Still need IO polling
- Standard IO routines for programmers
- 2 programs in memory, but one executed
- User interface was JCL (Job Control Language)
- Output of program can be input of another.

2.4 Changes in Hardware

- **Disk drives** provide faster IO.
- Processors that you can **interrupt** also means that there is no more reliance on polling.
- IO devices and CPU concurrent execution, and use local buffer.

SPOOLING (Simultaneous Peripheral Operation On-Line) Meaning that when interrupt, contents of cards read to disk. Therefore current program interrupted. Multiprogramming Putting multiple programs on at once. Need more memory to do this. Now also need for scheduler to manage multiple users' program needs. Need to figure out how to manage stuff. Priority of jobs, how much time to allocate for these jobs, etc... No memory protection, so programs could overwrite other program's chunk of memory. Java is an example of something that doesn't give you direct access to memory in JVM. Memory Protection better done by hardware than having software impose limits.

2 Modes for memory protection User/Restricted Mode Kernel Mode (SU)

Systems works differently in the 2 modes.

Mode bit can be used to signify which mode its in. If do something in User mode, then go to Kernel mode and show exception before going back to User mode.

Memory Protection Process gets area of memory that it can use Base and Limit register set for each process and how much memory it can have.

Batch systems REREAD SLIDES