

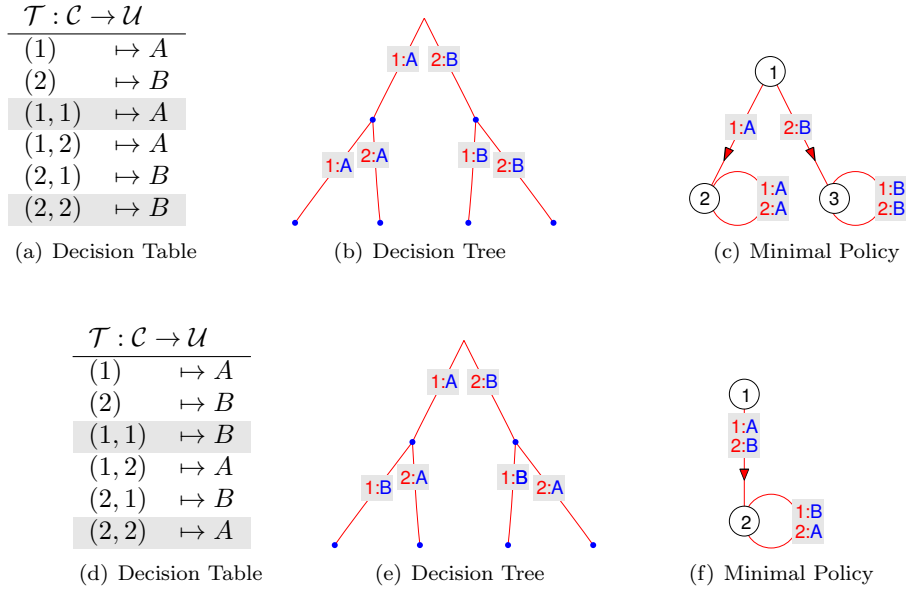
1 Formalization

Let \mathcal{Y} be a set of observations and \mathcal{U} a set of control actions. The original formalization defines a decision table as a map $\mathcal{T} : \mathcal{C} \rightarrow \mathcal{U}$, where $\mathcal{C} = \bigcup_{i=1}^n \mathcal{Y}^i$ for some $n \in \mathbb{N}$. This definition is often stronger than necessary, as there may be contexts c in which certain observations $y \in \mathcal{Y}$ are never encountered. In such a context, any assignment $\mathcal{T}(c) \in \mathcal{U}$, produces the same policy, although the choice of $\mathcal{T}(c)$ can affect a policy's reducibility.

Example 1 (Equivalent Decision Tables) Suppose $\mathcal{Y} = \{1, 2\}$, $\mathcal{U} = \{A, B\}$, $\mathcal{C} = \bigcup_{i=1}^2 \mathcal{Y}^i$ and

$$\mathcal{T}(c) = \begin{cases} A & c \in \{(1), (1, 2)\} \\ B & c \in \{(2), (2, 1)\} \\ \mathcal{T}'(c) & \text{otherwise.} \end{cases}$$

is an optimal policy, for arbitrary $\mathcal{T}' : \mathcal{C} \rightarrow \mathcal{U}$. However, depending on the choice of \mathcal{T}' (highlighted in the tables below), the completed policy can have differently-sized minimal representations:



Instead, we propose an “incompletely-determined” formalization:

Definition 1 (Policies) Given a set \mathcal{Y} of observations, recursively construct

$$\mathcal{C}_0 = \{\emptyset\} \quad \text{and} \quad \mathcal{C}_{i+1} = \{(c, y_{i+1}) : c \in \mathcal{C}_i, y_{i+1} \in \mathcal{Y}_c\}, \quad (1)$$

where $\mathcal{Y}_c \subseteq \mathcal{Y}$ are the observations that may be seen in context c . Let $\mathcal{C} = \bigcup_{i=0}^{\infty} \mathcal{C}_i$. A **policy** P is then a tuple $\langle \mathcal{C}, \mathcal{U}, \mathcal{T}, \mathcal{Y} \rangle$, where \mathcal{U} is some decision set and $\mathcal{T} : \mathcal{C} \setminus \{\emptyset\} \rightarrow \mathcal{U}$.

Definition 2 (Completely-Determined Policies) If $\mathcal{Y} = \bigcup_c \mathcal{Y}_c$ and $\mathcal{C} = \mathcal{Y}^{\leq n}$ for some $n \in \mathbb{N}$, then $P = \langle \mathcal{C}, \mathcal{U}, \mathcal{T}, \mathcal{Y} \rangle$ is **completely determined**. A **completion** of P is a policy $P' = \langle \mathcal{C}', \mathcal{U}', \mathcal{T}', \mathcal{Y} \rangle$ such that

$$\mathcal{Y} \subseteq \mathcal{Y}', \quad \mathcal{C}' = \bigcup_{i=0}^{\infty} (\mathcal{Y}')^i, \quad \mathcal{U} \subseteq \mathcal{U}', \quad \text{and} \quad \mathcal{T}'|_{\mathcal{C}} = \mathcal{T}. \quad (2)$$

Let $\text{Comp}(P)$ be the set of completions of the policy P .

Definition 3 (FSM Representations) An **FSM representation** (or just **representation**) is a tuple $\langle \mathcal{C}, \mathcal{R}, \mathcal{U}, \mathcal{S}, \mathcal{T}, \mathcal{Y} \rangle$ (abbreviated to $\langle \mathcal{R}, \mathcal{S} \rangle$ when $P = \langle \mathcal{C}, \mathcal{U}, \mathcal{T}, \mathcal{Y} \rangle$ is given), with “states” $\mathcal{S} \subseteq \mathbb{N}$ and state assignments $\mathcal{R} : \mathcal{C} \rightarrow \mathcal{S}$, such that

$$\mathcal{R}(c) = \mathcal{R}(c') \quad \text{and} \quad y \in \mathcal{Y}_c \cap \mathcal{Y}_{c'} \implies \mathcal{T}(c, y) = \mathcal{T}(c', y). \quad (3)$$

Let $\text{Rep}(P)$ be the set of representations of the policy P .

Definition 4 (Minimal Representations) The *size* of an FSM representation is the cardinality of its state set. A representation $\langle \mathcal{R}, \mathcal{S} \rangle$ of P is **minimal** if $|\mathcal{S}| = \min\{|\mathcal{S}'| : \langle \mathcal{R}', \mathcal{S}' \rangle \in \text{Rep}(P)\}$. A representation $\langle \mathcal{R}', \mathcal{S}' \rangle$ is a **reduction** of the representation $\langle \mathcal{R}, \mathcal{S} \rangle$ if there is a surjection $\phi : \mathcal{S} \rightarrow \mathcal{S}'$ such that $\mathcal{R}' = \phi(\mathcal{R})$.

Example 2 If $\mathcal{C} = \{c_1, c_2, \dots\}$, then we have a canonical representation $\langle \mathcal{R}, \mathcal{S} \rangle$, where

$$\mathcal{S} = \{1, \dots, |\mathcal{C}|\} \quad \text{and} \quad \mathcal{R} : c_k \mapsto k. \quad (4)$$

It can be shown that the size of a minimal representation of a policy P is equal to the minimum size of the minimal representations of its completions, i.e.

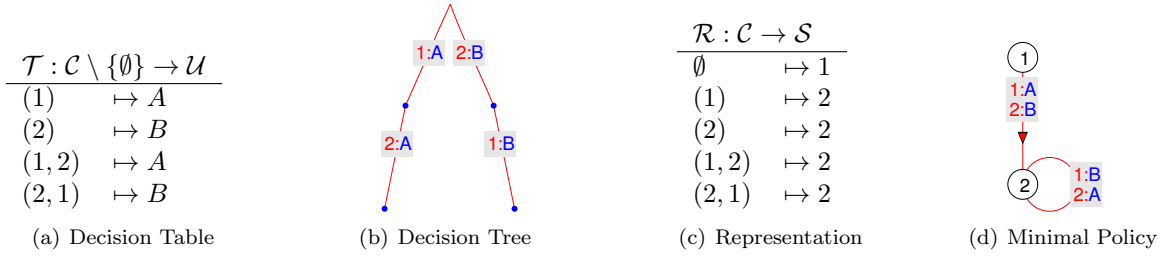
$$\min\{|\mathcal{S}'| : \langle \mathcal{R}', \mathcal{S}' \rangle \in \text{Rep}(P)\} = \min\{|\mathcal{S}'| : \langle \mathcal{R}', \mathcal{S}' \rangle \in \text{Rep}(P'), P' \in \text{Comp}(P)\} \quad (5)$$

Incompletely-determined policies allow more freedom in representation reduction, as shown in the next example.

Example 3 (Incompletely-Determined Policies) Let $\mathcal{C} = \{\emptyset, (1), (2), (1, 2), (2, 1)\}$, $\mathcal{U} = \{A, B\}$, and

$$\mathcal{T}(c, y) = \begin{cases} A & c \in \{(1), (1, 2)\} \\ B & c \in \{(2), (2, 1)\} \end{cases}.$$

Observe that the minimal policy is the same as that of the completely-determined policy in Example 1(f).



2 Algorithm

To find a minimum representation of a given policy, we first compute a graph of reducibility relations, then compute a minimal clique-covering. For practical computation of reducibility, we'll start with the weaker condition of compatibility.

2.1 Reducibility Relations

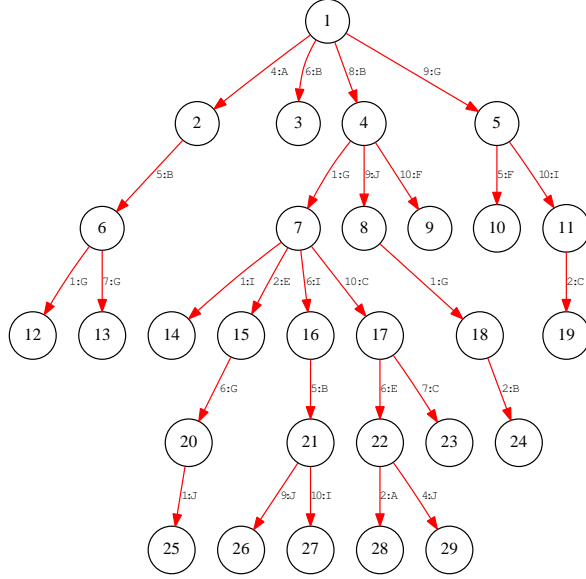
Definition 5 (Reducibility) For a given policy $P = \langle \mathcal{C}, \mathcal{U}, \mathcal{T}, \mathcal{Y} \rangle$, two contexts $c_1, c_2 \in \mathcal{C}$ are **reducible** (write $c_1 \sim c_2$) if there exists a representation $\langle \mathcal{R}, \mathcal{S} \rangle$ of P such that $\mathcal{R}(c_1) = \mathcal{R}(c_2)$. Likewise, for a given representation $R = \langle \mathcal{R}, \mathcal{S} \rangle$, two states $s_1, s_2 \in \mathcal{S}$ are **reducible** if there exists a reduction $(\phi, \langle \mathcal{R}', \mathcal{S}' \rangle)$ of R such that $\phi(s_1) = \phi(s_2)$.

Observe that for any representation $\langle \mathcal{C}, \mathcal{R}, \mathcal{U}, \mathcal{S}, \mathcal{T}, \mathcal{Y} \rangle$, the contexts $c_1, c_2 \in \mathcal{C}$ are reducible if and only if the states $\mathcal{R}(c_1)$ and $\mathcal{R}(c_2)$ are reducible. Observe also that for incompletely-determined policies, reducibility is a symmetric but not-necessarily-transitive relation

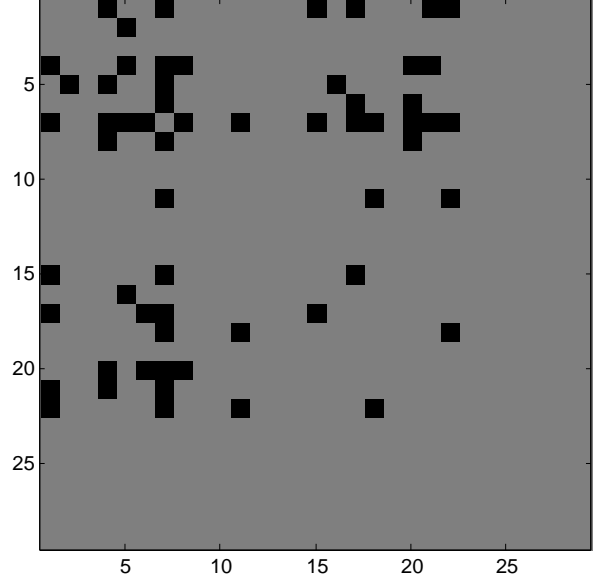
Example 4 (Non-Transitive Reducibility) Suppose $\mathcal{Y} = \{1, 2, 3\}$, $\mathcal{C} = \{\emptyset, (1), (2), (1, 3), (2, 3)\}$, $\mathcal{U} = \{A, B\}$, and

$$\mathcal{T}(c) = \begin{cases} A & c \in \{(1), (1, 3)\} \\ B & c \in \{(2), (2, 3)\} \end{cases}. \quad (6)$$

Observe that, under this policy, $\emptyset \sim (1)$ and $\emptyset \sim (2)$, but $(1) \not\sim (2)$, since $\mathcal{T}(1, 3) \neq \mathcal{T}(2, 3)$.



(a) Canonical Representation



(b) Compatibility Matrix

However, it can be shown that, under a completely-determined policy, reducibility induces an equivalence relation. In either case, we compute reducibility using the following criterion:

Lemma 1 *Two contexts $c_1, c_2 \in \mathcal{C}$ are reducible iff*

$$\mathcal{T}(c_1, s) = \mathcal{T}(c_2, s) \quad \text{for all } s \in \mathcal{Y}^* \quad \text{such that } (c_1, s), (c_2, s) \in \mathcal{C} \quad (7)$$

This informs the following algorithm

Algorithm 1: Compute Reducibility Relations

Input: A representation $\langle \mathcal{C}, \mathcal{R}, \mathcal{U}, \mathcal{S}, \mathcal{T}, \mathcal{Y} \rangle$

Output: A reducibility matrix $A : \mathcal{S} \times \mathcal{S} \rightarrow \{\text{true}, \text{false}\}$.

```

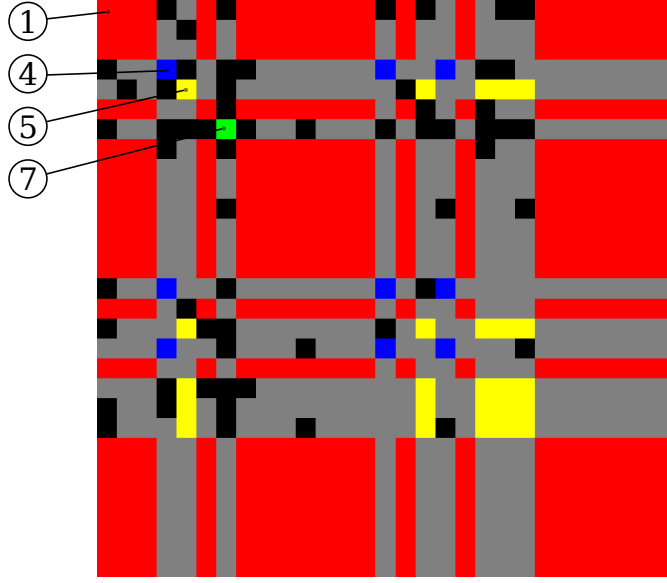
 $A(s_1, s_2) \leftarrow \text{true}$  for all  $s_1, s_2 \in \mathcal{S}$ .
repeat
   $isChanged \leftarrow \text{false}$ 
  for  $s_1 < s_2 \in \mathcal{S}$  do
    if  $A(s_1, s_2) = \text{true}$  then
      for  $c_1 \in \mathcal{R}^{-1}(s_1), c_2 \in \mathcal{R}^{-1}(s_2)$  do
        for  $y \in \mathcal{Y}_{c_1} \cap \mathcal{Y}_{c_2}$  do
          if  $\mathcal{T}(c_1, y) \neq \mathcal{T}(c_2, y)$  or  $\sim A(\mathcal{R}(c_1, y), \mathcal{R}(c_2, y))$  then
             $A(s_1, s_2) \leftarrow \text{false}$ .
             $isChanged \leftarrow \text{true}$ .
          end if
        end for
      end for
    end for
  end for
until  $\sim isChanged$ 

```

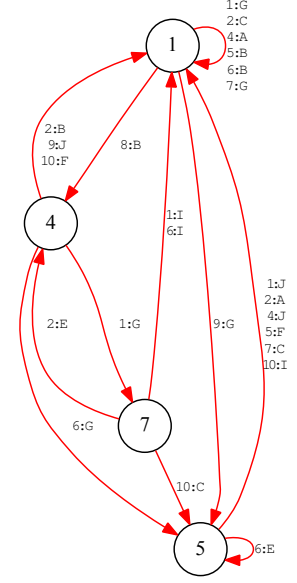
Now, although reducibility is not an equivalence relation, any reduction $\phi : \mathcal{S} \rightarrow \mathcal{S}'$ induces an equivalence relation, partitioning \mathcal{S} into cliques of mutually-reducible states, i.e.

$$\mathcal{S} = \bigsqcup_{s' \in \mathcal{S}'} \phi^{-1}(s'), \quad \text{where } \phi(s_1) = \phi(s_2) \implies A(s_1, s_2) \quad (8)$$

Thus, a minimal reduction induces a minimal clique partition of the reducible states of a representation.



(c) Greedy Clique Covering of 4(b)



(d) Reduced Representation

2.2 Assembling Cliques

Notation 1 (Arrow notation) For a policy $P = \langle \mathcal{C}, \mathcal{U}, \mathcal{T}, \mathcal{Y} \rangle$, write $c \rightarrow c'$ if $c = (c_1, \dots, c_i) \in \mathcal{C}_i \subseteq \mathcal{C}$ and $c' = (c_1, \dots, c_i, y) \in \mathcal{C}_{i+1} \subseteq \mathcal{C}$, for some i . For a representation $\langle \mathcal{R}, \mathcal{S} \rangle$ of P , write $s_1 \rightarrow s_2$ if there are $c_1 \in f^{-1}(s_1)$ and $c_2 \in f^{-1}(s_2)$ such that $c_1 \rightarrow c_2$.

We propose the following, greedy, approximate algorithm. Although we have no proof that this algorithm

Algorithm 2: Compute Clique Covering

Input: A representation $\langle \mathcal{C}, \mathcal{R}, \mathcal{U}, \mathcal{S}, \mathcal{T}, \mathcal{Y} \rangle$ with $s_1 < s_2$ only if $s_2 \not\rightarrow s_1$.

Input: A reducibility matrix $A : \mathcal{S} \times \mathcal{S} \rightarrow \{\text{true}, \text{false}\}$ as computed by Algorithm 1.

Output: A partition function $\phi : \mathcal{S} \rightarrow \mathcal{S}'$ with $\phi(s_1) = \phi(s_2)$ only if $A(s_1, s_2)$.

```

 $\mathcal{S}' \leftarrow \mathcal{S}$ 
 $\phi \leftarrow id_{\mathcal{S}}$ 
 $unused \leftarrow \mathcal{S}$ 
while  $|unused| > 0$  do
   $s_1 \leftarrow \min(unused)$ 
   $unused \leftarrow unused \setminus \{s_1\}$ 
  for  $s_2 \in unused$  do
    if  $A(s_1, s_2)$  then
       $\phi(s_2) \leftarrow s_1$ 
       $unused \leftarrow unused \setminus \{s_2\}$ 
    end if
  end for
end while

```

produces minimal representations of a given policy, it is not inconceivable that this or another greedy policy could work. In general, the Minimal Clique Covering problem is NP-Complete, but the tree structure of the decision policy is an a constraint that may simplify the problem.

3 Code

The main MATLAB routine is in `decision_script.m`. By default it generates a random decision tree, whose branching is governed by a Poisson distribution.

3.1 Console Output

Initial FSM:

```
-----
01: 08-->(05, 02),
02: 02-->(05, 03), 06-->(03, 04), 10-->(01, 05),
03: 07-->(04, 06), 08-->(07, 07),
04:
05: 05-->(10, 08), 07-->(06, 09),
06: 04-->(09, 10), 05-->(05, 11), 07-->(09, 12),
07:
08: 09-->(07, 13),
09: 03-->(08, 14),
10:
11:
12: 04-->(01, 15), 06-->(06, 16),
13: 01-->(03, 17),
14: 07-->(01, 18),
15:
16:
17:
18:
```

Final clique covering: {1, 2, 4, 5, 7, 8, 9, 10, 11, 13, 15, 16, 17, 18} {3, 12} {6} {14}

3.2 Visualizations

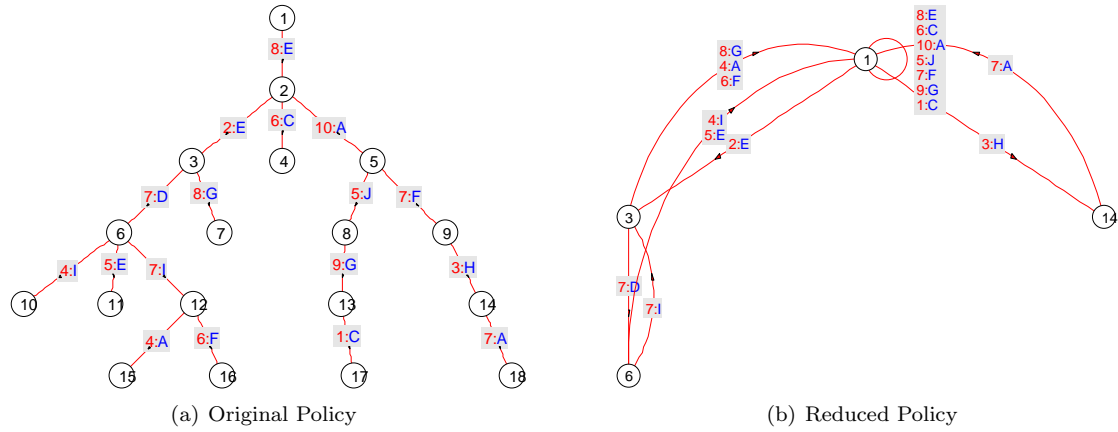


Figure 4: MATLAB Visualizations

An optional section of this script uses Graphviz to produce high-quality FSM visualizations.

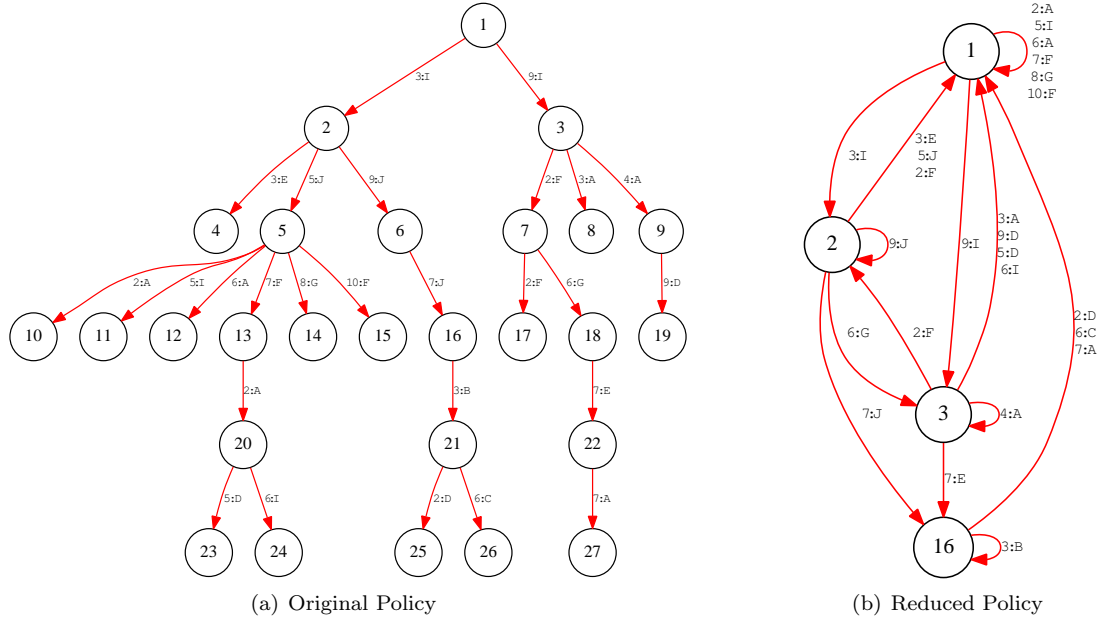


Figure 5: Graphviz Visualizations

This makes a system call to the script `drawgraph.sh`, which runs in the bash shell, and uses a few non-standard graphics utilities. In Ubuntu, these dependencies should be resolved by the command

```
$ sudo apt-get install ghostscript texlive-extra-utils graphviz evince
```