## $\Sigma_n$ and $\Pi_n$

Recall that a set  $A \subset \mathbb{N}_0$  is enumerable if and only if there exists a decidable set  $B \subset \mathbb{N}_0 \times \mathbb{N}_0$  such that A is the projection of B. Let's identify sets and predicates so that we can say that a property A(x) of natural numbers is enumerable if and only if it can be represented in the form

$$A(x) \Leftrightarrow \exists y B(x,y)$$

where B(x, y) is some decidable property.

**Q:** What can be said about other combinations of quantifiers?

**Q:** What properties are representable in the form

$$A(x) \Leftrightarrow \exists y \exists z C(x, y, z)$$

where C is some decidable property?

We can replace consecutive quantifiers by using one with computable numbering of pairs:

$$C''(x, \sigma(y, z)) \Leftrightarrow C(x, y, z) \text{ and } A(x) \Leftrightarrow \exists w C''(x, w).$$

**Q:** What properties can be represented in the form

$$A(x) \Leftrightarrow \forall y B(x, y)$$

where B(x, y) is a decidable property.

**A:** The properties with enumerable negations:

$$\neg A(x) \Leftrightarrow \neg \forall y B(x, y)$$
  
 $\Leftrightarrow \exists y (\neg B(x, y))$ 

We are implicitly using the fact that decidability is preserved under negation.

We now progress to a general definition:

**Definition 1.** A property A belongs to the class  $\Sigma_n$  if it can be represented in the form

$$A(x) \Rightarrow \exists y_1 \forall y_2 \exists y_3 \cdots B(x_1, y_1, y_2, \dots, y_n)$$

with n alternating quantifiers and B is a decidable property.

If the n alternating quantifiers starts with  $\forall$ , then we obtain the definition of the class  $\Pi_n$ .

**Theorem 1.** (a) The class  $\Sigma_n$  and  $\Pi_n$  does not change if we allow groups of quantifiers of the same type instead of a single quantifier.

(b) If a predicate belongs to  $\Sigma_n$  then its negation belongs to  $\Pi_n$  and vice versa.

**Theorem 2.** (a) The intersection and union of two sets of the class  $\Sigma_n$  belongs to  $\Sigma_n$ .

(b) The intersection and union of two sets of the class  $\Pi_n$  belong to  $\Pi_n$ .

## Example 1.

$$A(x) \Leftrightarrow \exists y \forall z B(x, y, z)$$

$$C(x) \Leftrightarrow \exists u \forall v D(x, u, v)$$

Then

$$A(x) \wedge C(x) \Leftrightarrow \exists y \exists u \forall z \forall v [B(x, y, z) \wedge D(x, u, v)]$$

The property is decidable and we can complie the quantifiers to get it to belong to  $\Sigma_n$ .

**Theorem 3.** The classes  $\Sigma_n$  and  $\Pi_n$  are "hereditary downward" with respect to m-reducibility in the following way:

If  $A \leq_m B$  and  $B \in \Sigma_n$  then  $A \in \Sigma_n$  or  $B \in \Pi_n$  then  $A \in \Pi_n$ .

## Unviersal Sets in $\Sigma_n$ and $\Pi_n$

**Goal:** Show that the classes  $\Sigma_n$  and  $\Pi_n$  are distinct for different n.

**How?** We find a universal set for  $\Sigma_n$  and show that it cannot belong to  $\Sigma_k$  for any k < n.

**Theorem 4.** For any n, the class  $\Sigma_n$  contains a set universal for this class. The complement of this set will be universal for the class  $\Pi_n$ .

**Theorem 5.** Universal  $\Sigma_n$ -sets do not belong to the class  $\Pi_n$ . Similarly universal  $\Pi_n$ -sets do not belong to  $\Sigma_n$ .

Remark 1. The hierarchy is strict.