

Security: Empirical measures and theoretical bases

Andrew Ruef

EVERYTHING IS HACKED

RISK ASSESSMENT / SECURITY & HACKTIVISM

Schneier on Security

All four major browsers take
Pwn2Own hacking competition

Security defenses keep getting better, but so too do hac

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About Me

[← "Unbreakable" Encryption Almost Certainly Isn't](#)

[Police Disabling Their Own Voice Recorders →](#)

Heartbleed

[Heartbleed](#) is a catastrophic bug in OpenSSL:

Maryland Hacked Again:
4 Weeks

March 20, 2014 2:41 PM

Security on NBC NEWS.com

Heartland Payment Systems hacked

Payments processor said breach did not involve merchant, cardholder data

Target credit card hack: What you need to know

What can science do

- Measure the world
- Develop theories

MEASURE THE WORLD

Hackers are not tornados



Spam

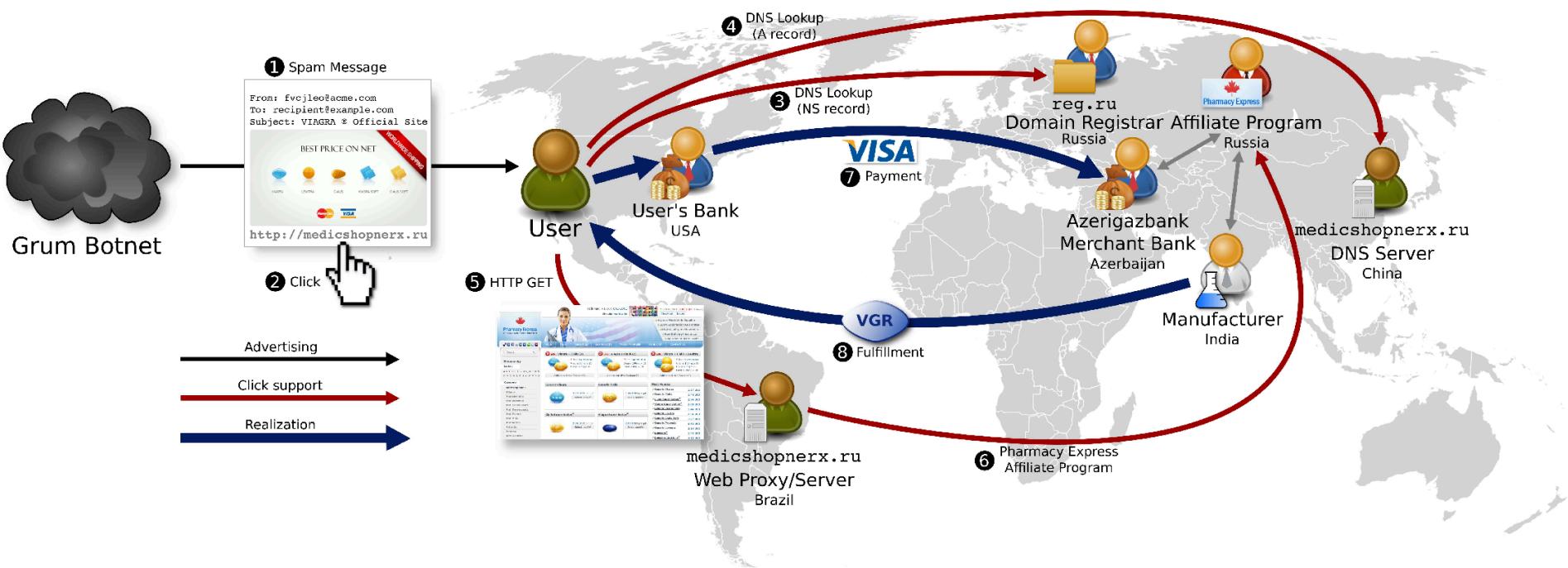


Figure 1: Infrastructure involved in a single URL's value chain, including advertisement, click support and realization steps.

Zeus

**UNITED STATES DISTRICT COURT
EASTERN DISTRICT OF NEW YORK**

MICROSOFT CORP., FS-ISAC, INC., and NATIONAL
AUTOMATED CLEARING HOUSE ASSOCIATION,

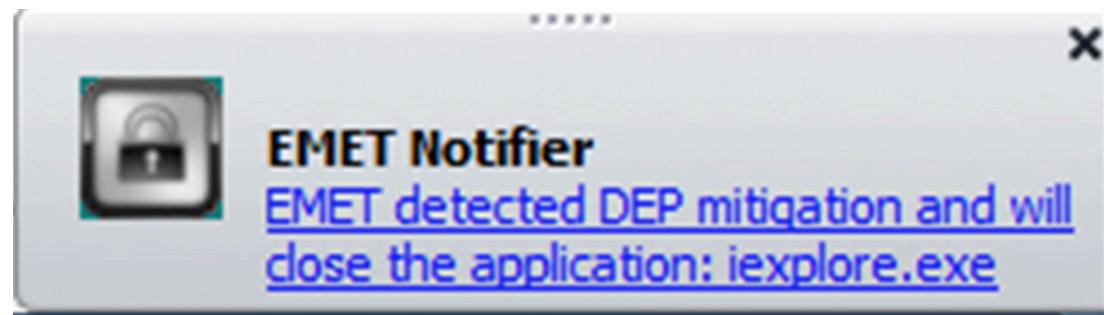
Plaintiffs

v.

JOHN DOES 1-39 D/B/A Slavik, Monstr, IOO, Nu11,
nvidiag, zebra7753, lexa_Mef, gss, iceIX, Harderman,
Gribodemon, Aqua, aquaSecond, it, percent, cp01, hct,
xman, Pepsi, miami, miamibc, petr0vich, Mr. ICQ, Tank,
tankist, Kusunagi, Noname, Lucky, Bashorg, Indep, Mask,
Enx, Benny, Bentley, Denis Lubimov, MaDaGaSka,
Vkontakte, rfcid, parik, reronic, Daniel, bx1, Daniel
Hamza, Danielbx1, jah, Jonni, jtk, D frank, duo,
Admin2010, h4x0rdz, Donsft, mary.J555, susanneon,
kainehabe, virus_e_2003, spanishp, sere.bro, muddem,
mechan1zm, vlad.dimitrov, jheto2002, sector.exploits
AND JabberZeus Crew AND YEVHEN KULIBABA AND
YURIY KONOVALENKO, CONTROLLING COMPUTER
BOTNETS THEREBY INJURING PLAINTIFFS, AND THEIR
CUSTOMERS AND MEMBERS,

Defendants.

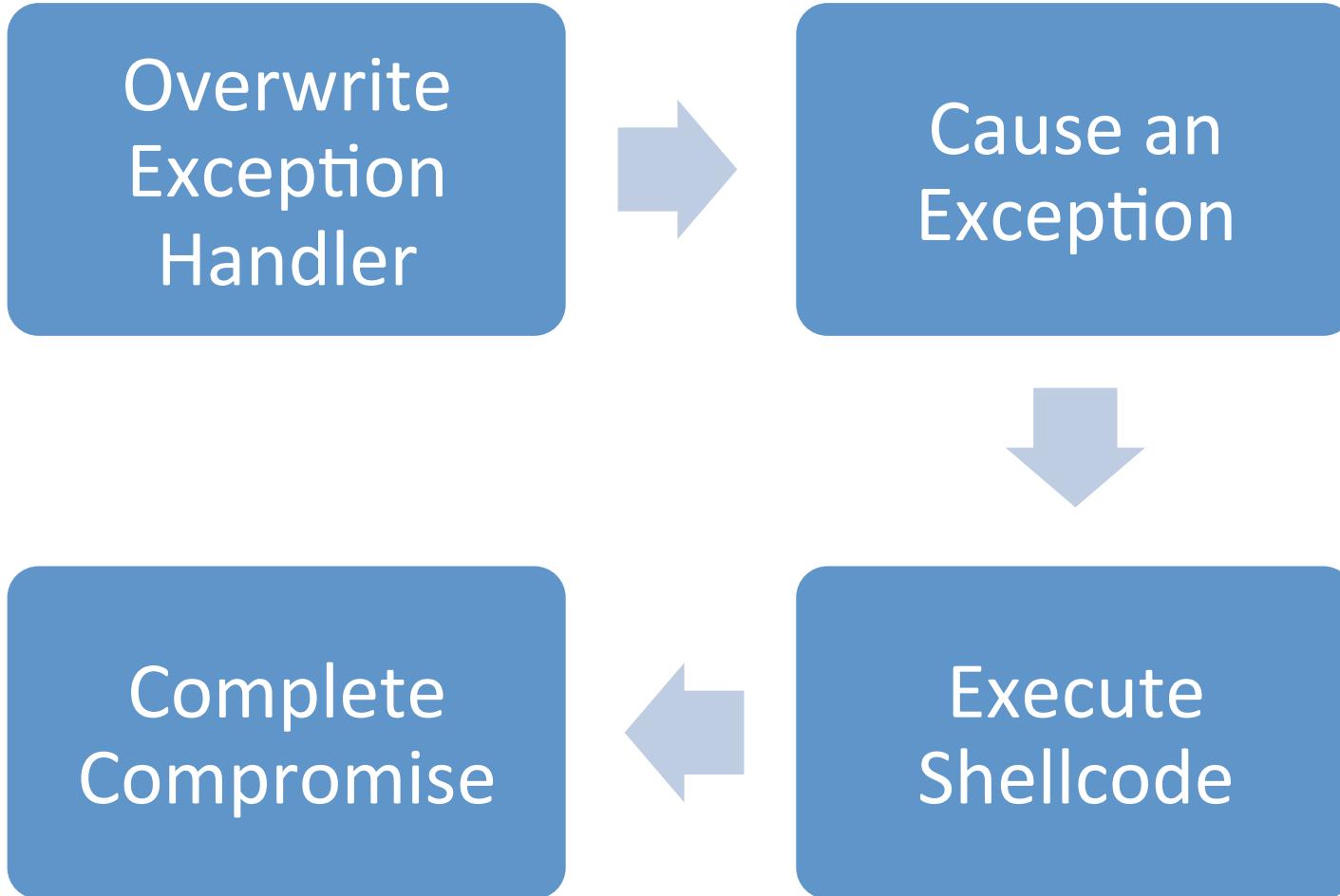
Develop mitigations



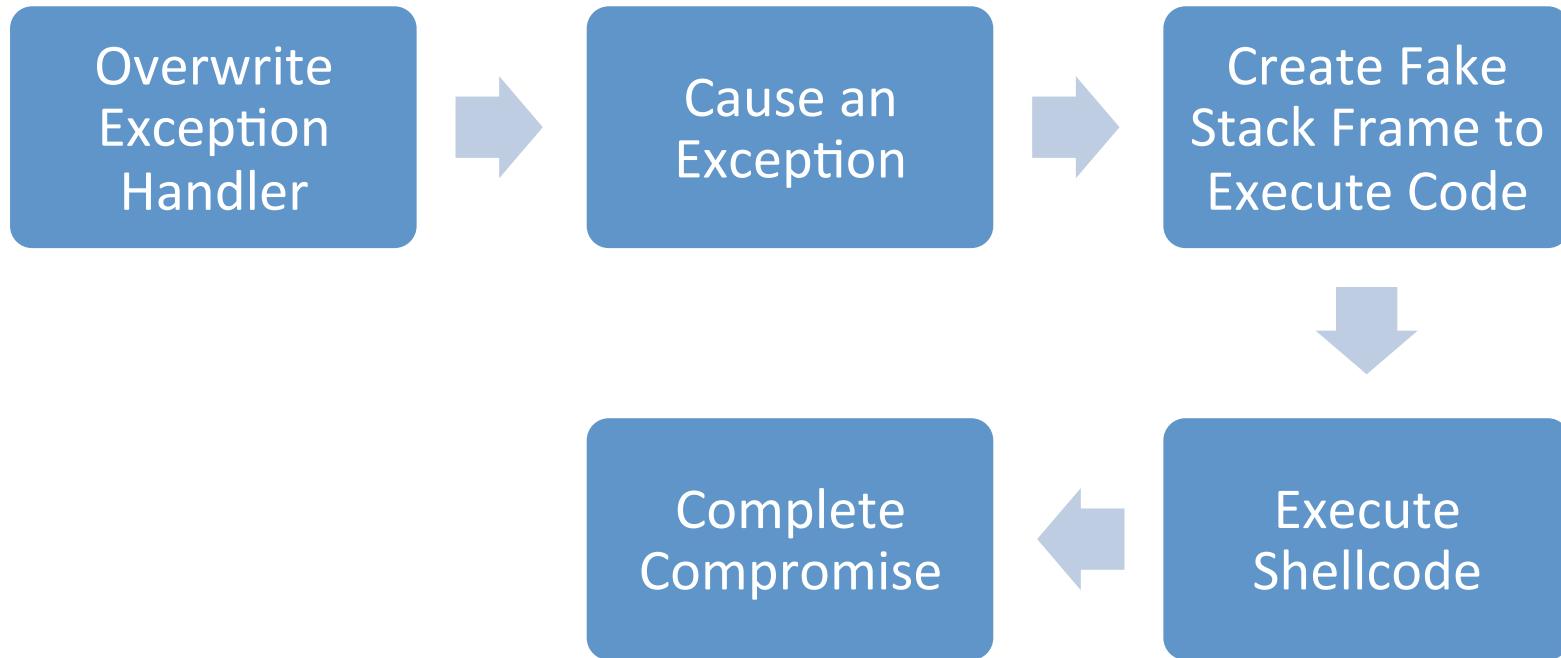
Problem (1996)



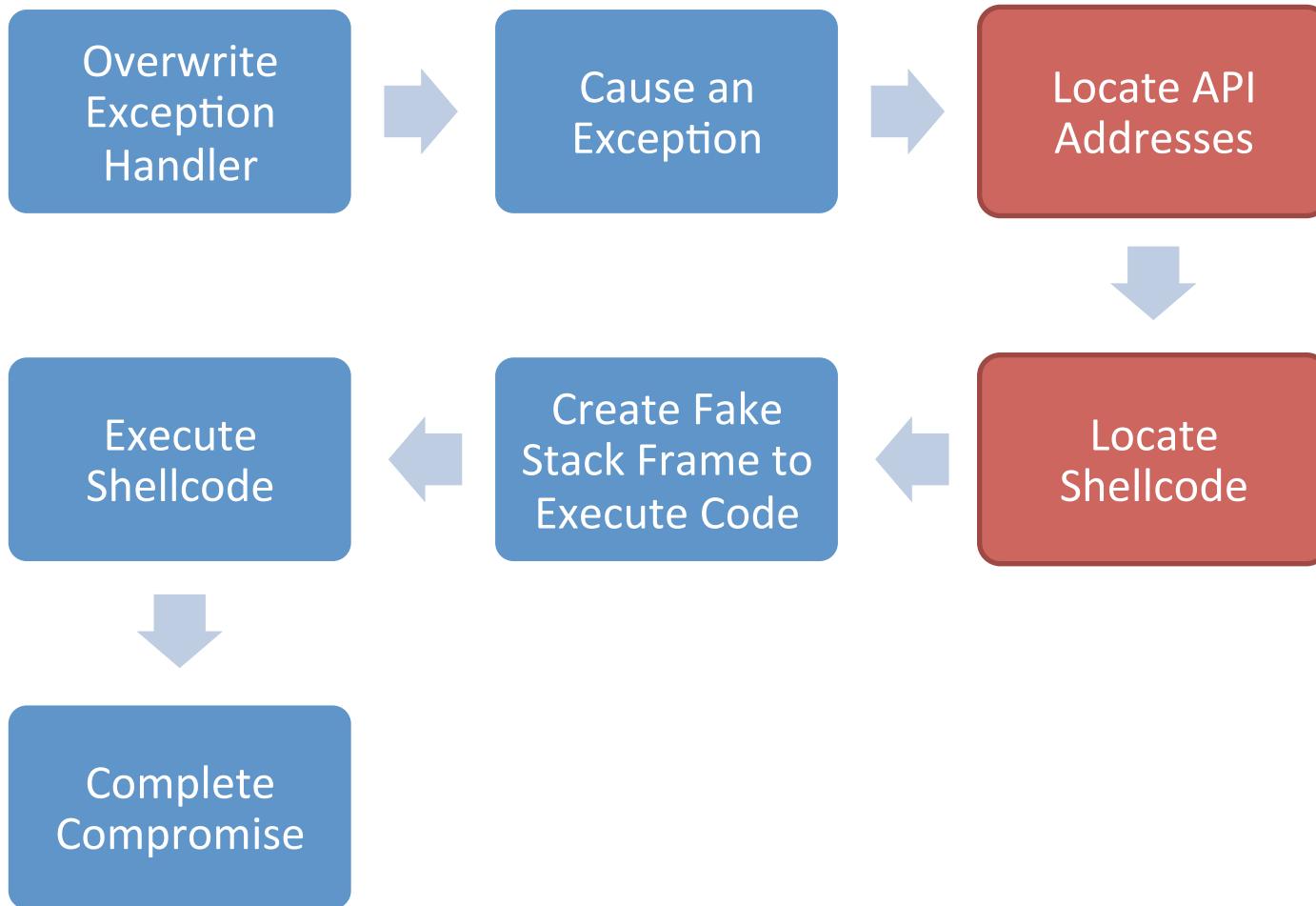
Problem (2003)



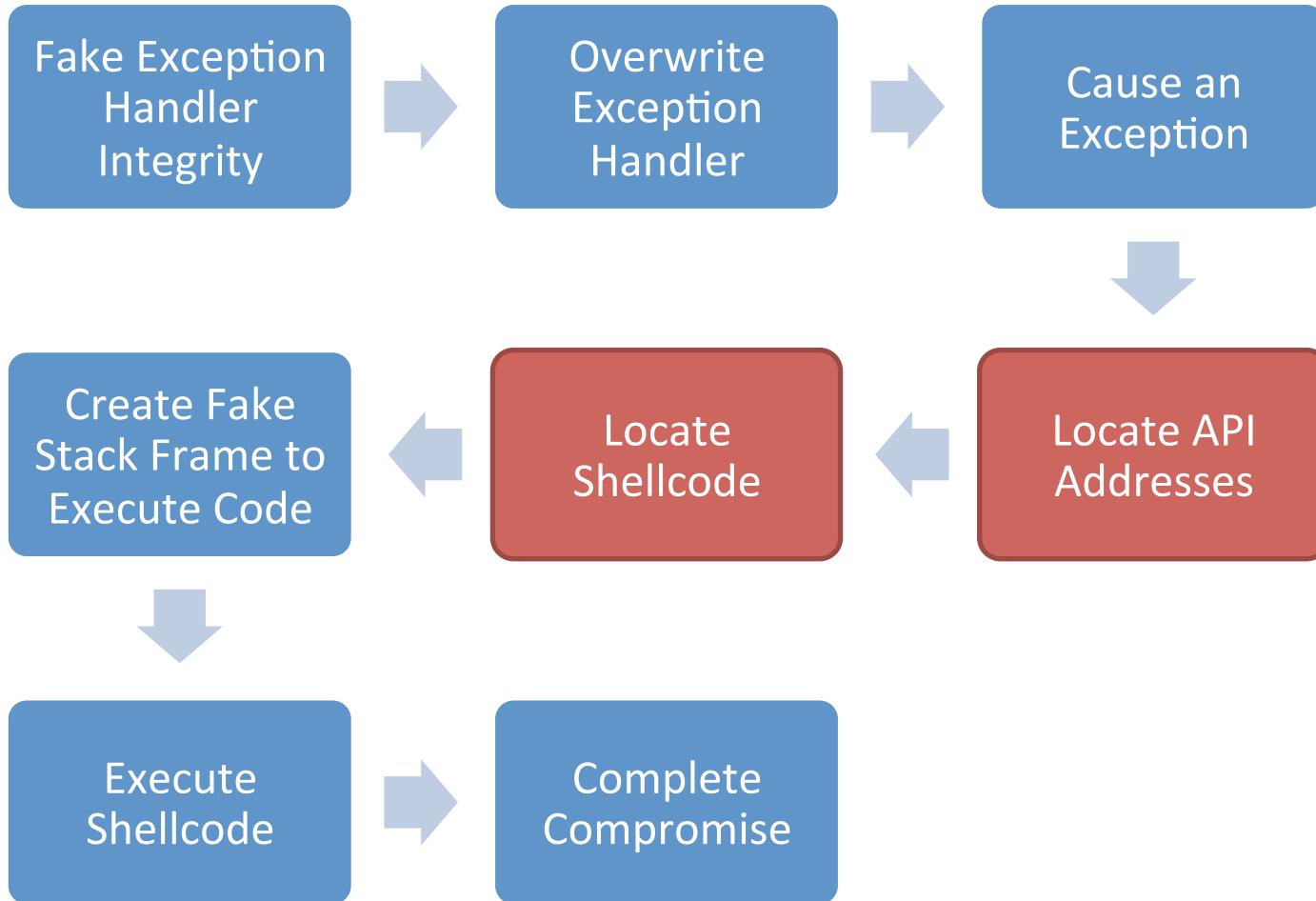
Problem (2004)



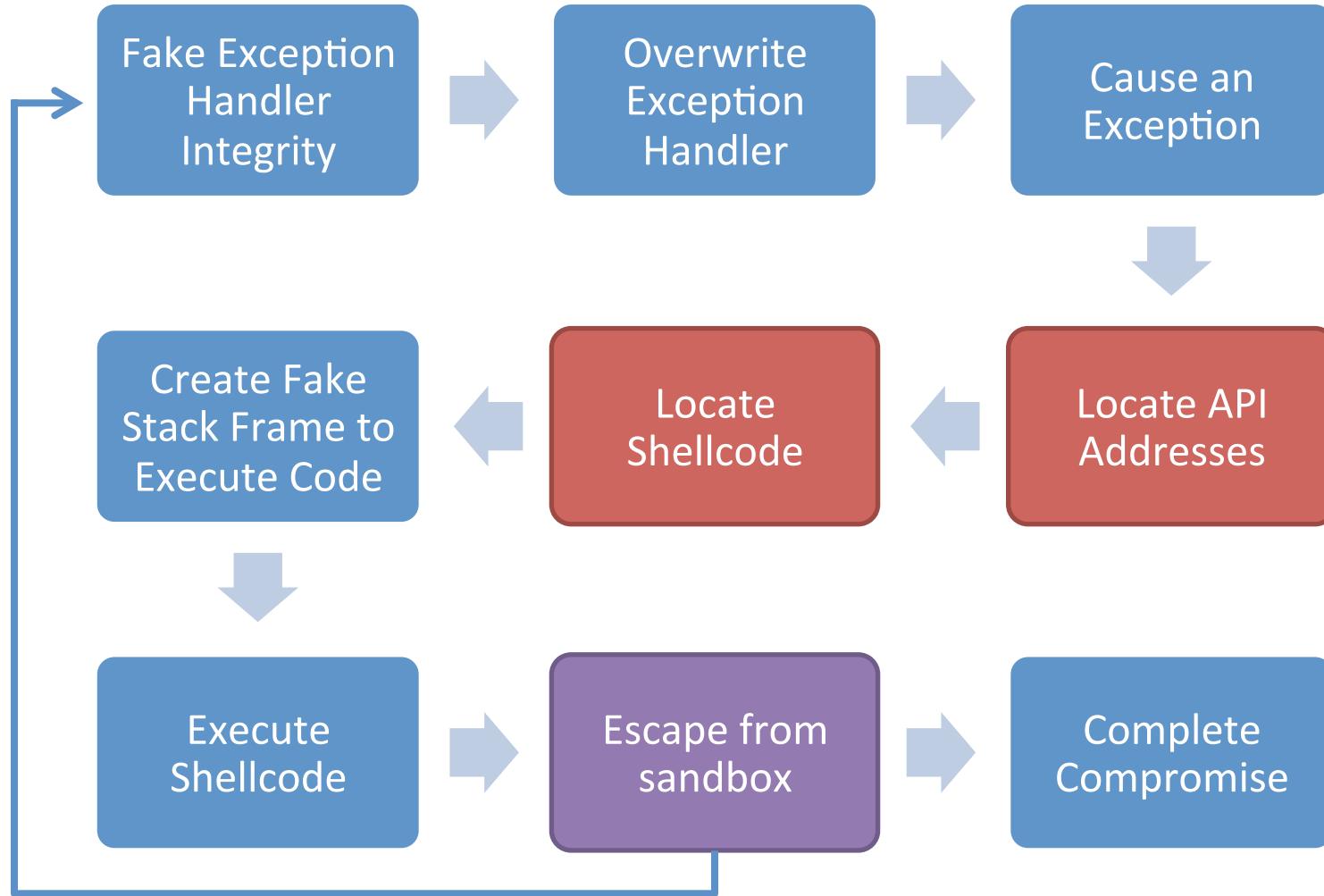
Problem (2007)



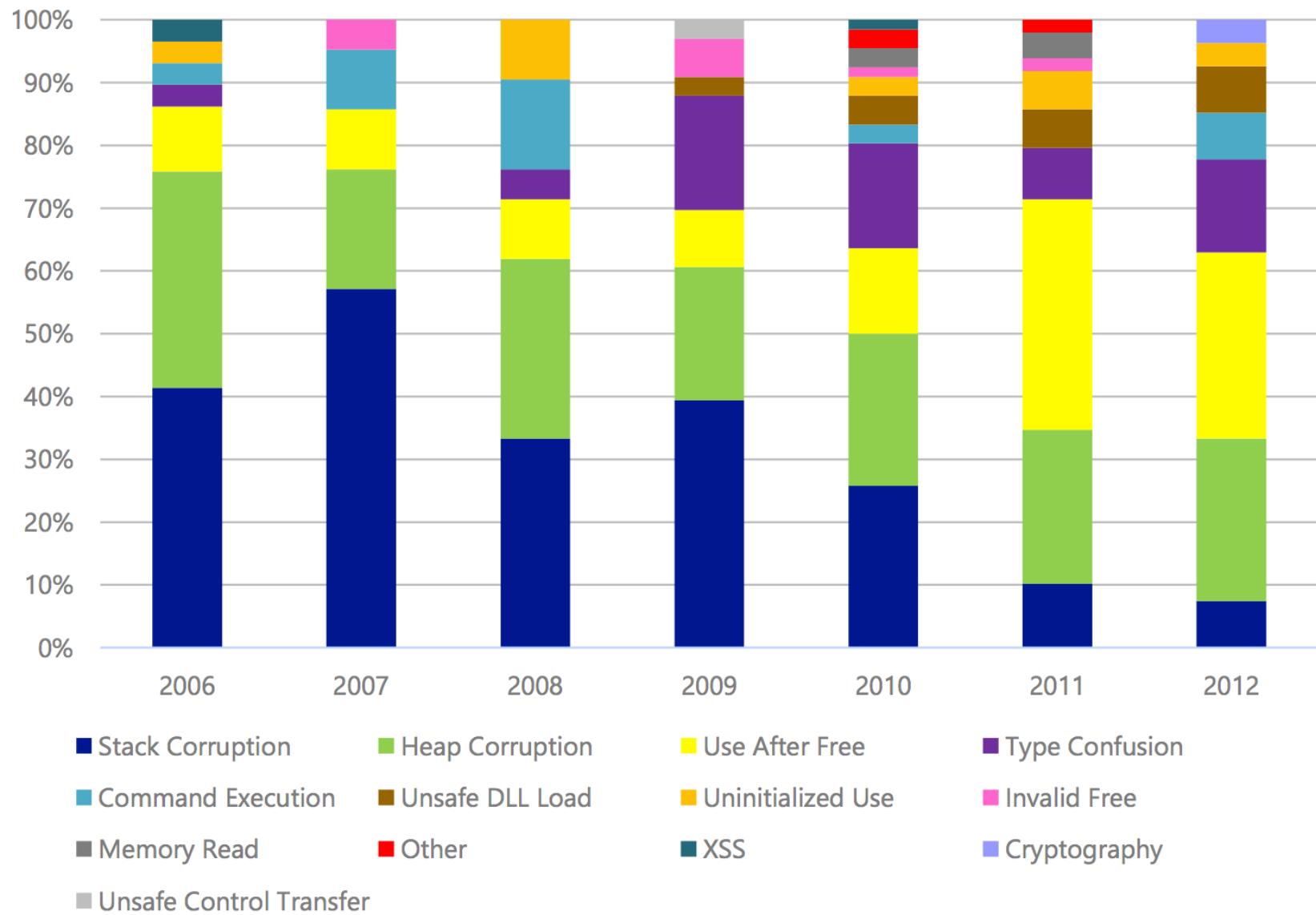
Problem (2008)



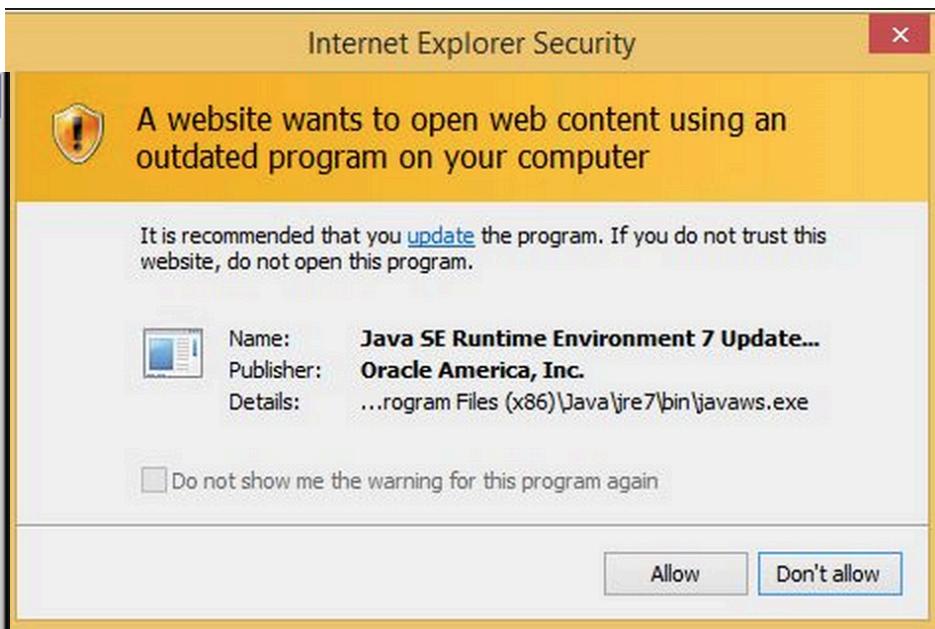
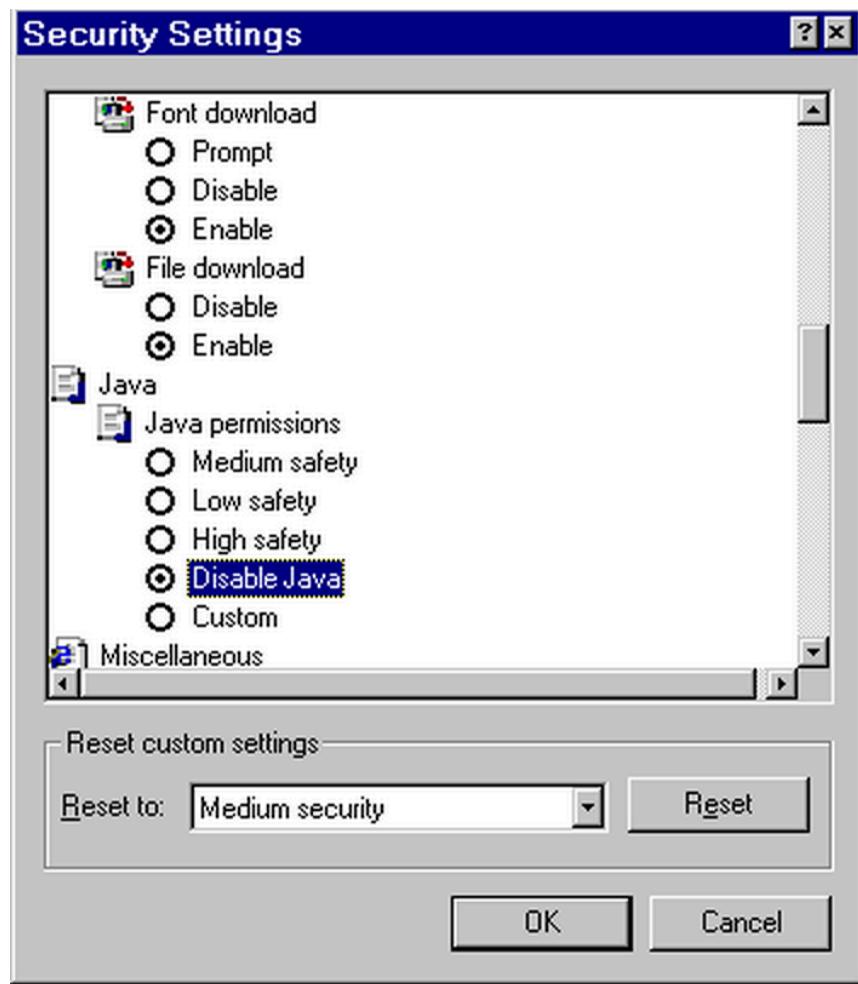
Problem (2013)



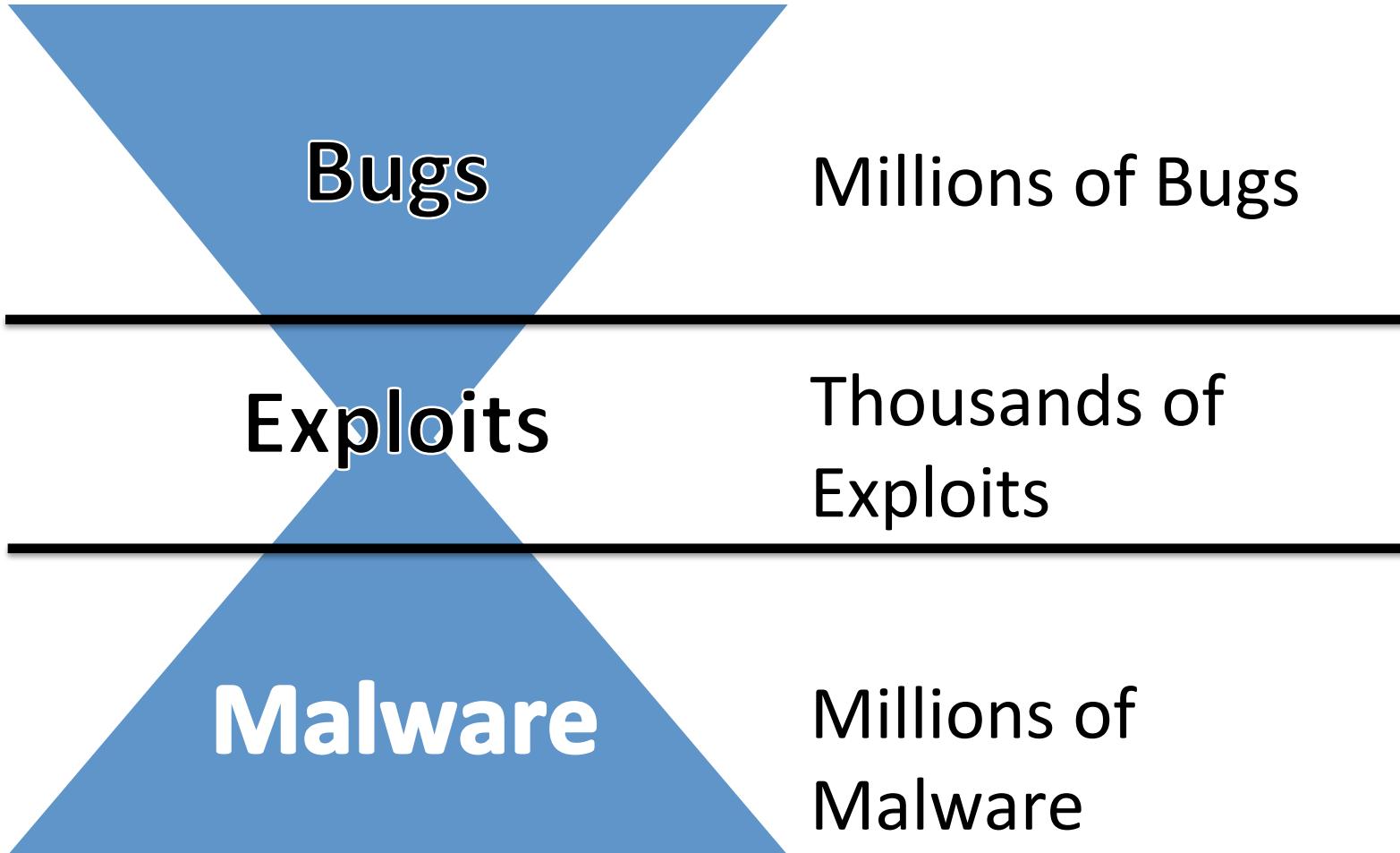
IT WORKS



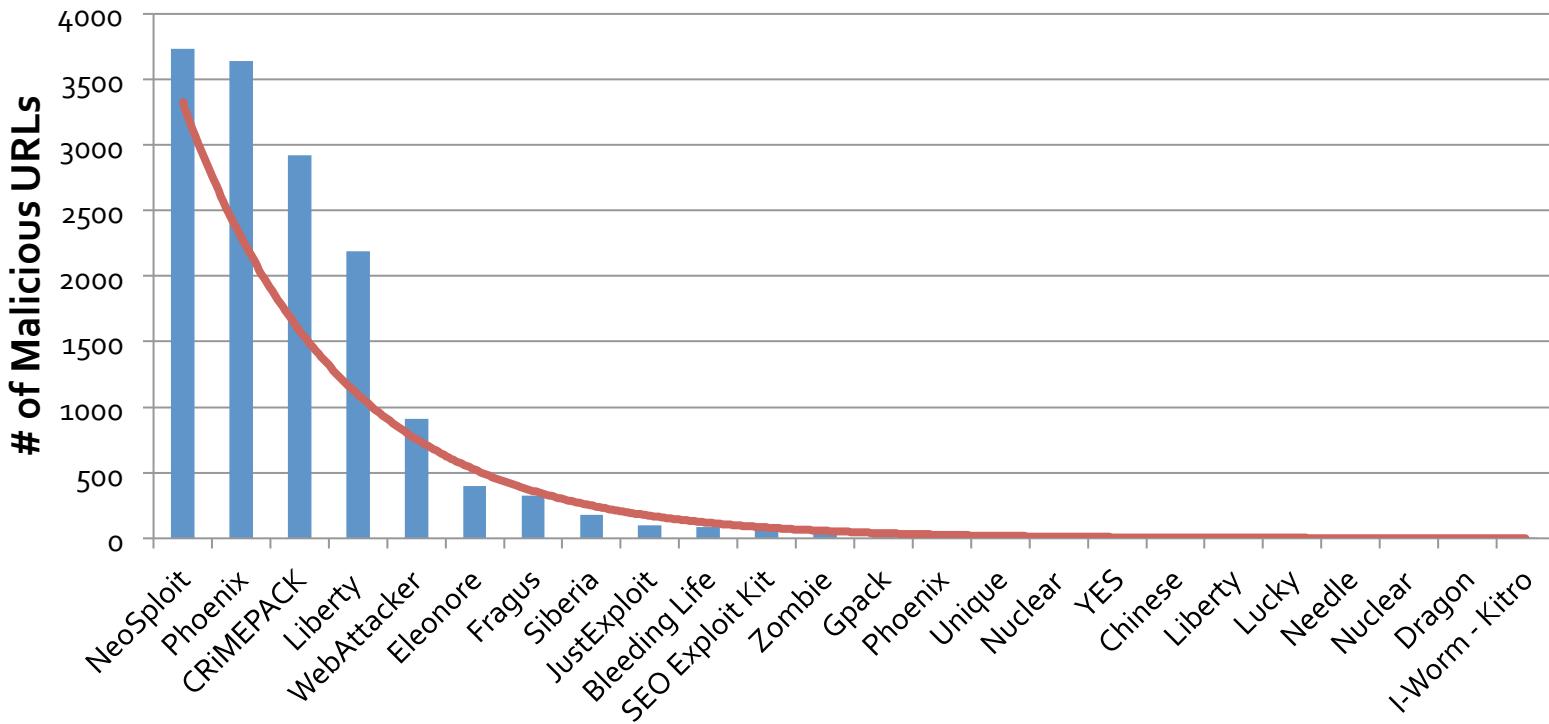
Enterprise configuration



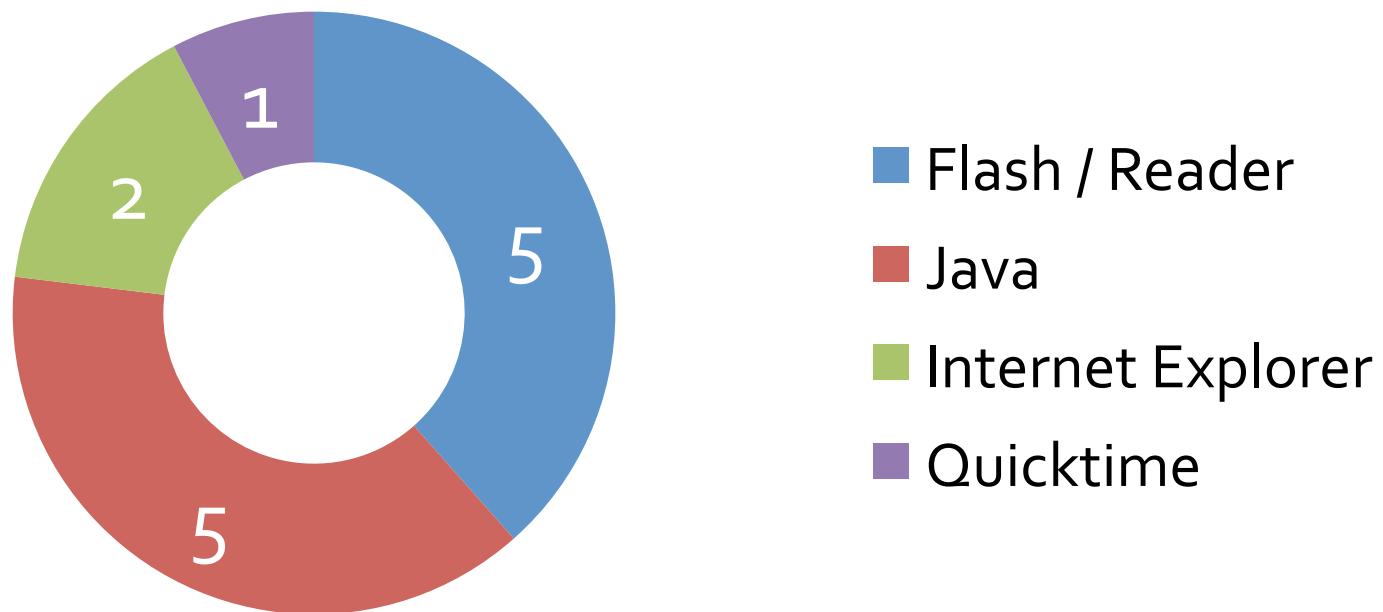
Why Study Exploitation?



Clear Market Leaders



Limited Target Support



Low Quality Exploits

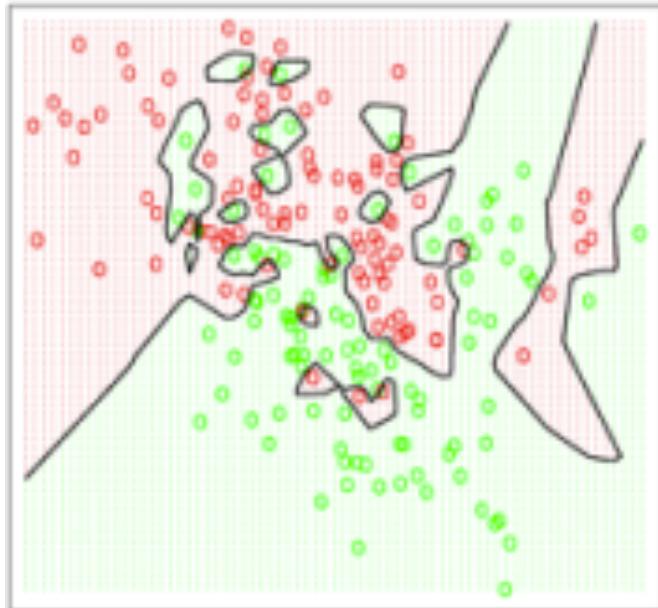
Memory Corruption (19)

| | |
|------------------|----|
| Defeated by DEP | 14 |
| Defeated by ASLR | 17 |
| Defeated by EMET | 19 |

Logic Flaws (8)

| | |
|----------------------------|---|
| No Java in Internet Zone | 4 |
| No EXEs in PDFs | 1 |
| No Firefox or FoxIt Reader | 2 |

Empiricism can be dangerous



still teaching computer janitors of tomorrow
to lose the battles of yesterday
drive.google.com/file/d/0B_mjsP...

◀ ▶ ⚡ ...

Google Docs



DEVELOP THEORIES

Full abstraction

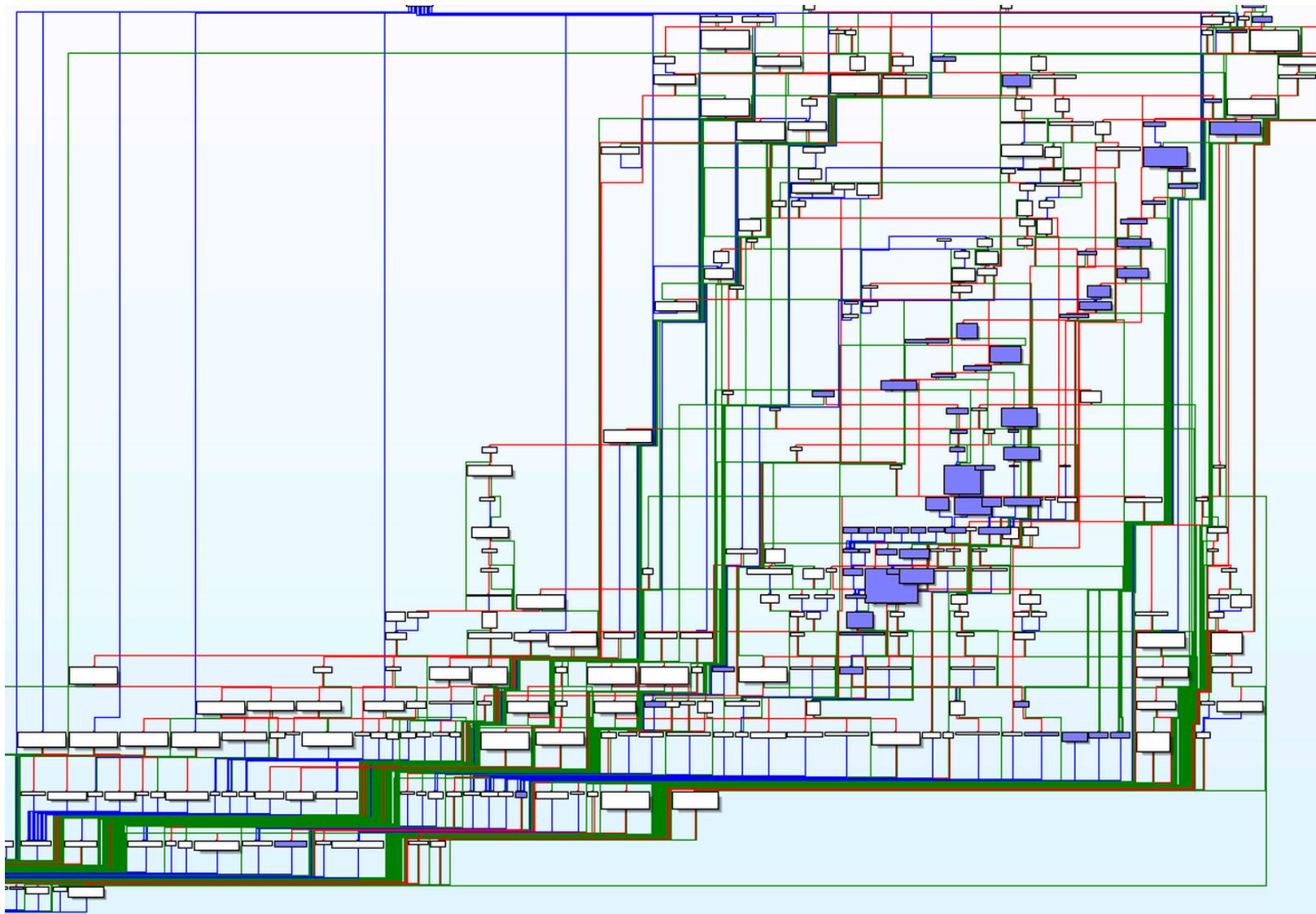
```
#include <stdio.h>
#include <stdlib.h>
#include <sys/types.h>
#include <arpa/inet.h>

void serveur1(portServ ports)
{
    int sockServ1, sockServ2, sockClient;
    struct sockaddr_in monAddr, addrClient, addrServ2;
    socklen_t lenAddrClient;

    if ((sockServ1 = socket(AF_INET, SOCK_STREAM, 0)) == -1) {
        perror("Erreur socket");
        exit(1);
    }
    if ((sockServ2 = socket(AF_INET, SOCK_STREAM, 0)) == -1) {
        perror("Erreur socket");
        exit(1);
    }

    bzero(&monAddr, sizeof(monAddr));
    monAddr.sin_family = AF_INET;
    monAddr.sin_port = htons(ports.port1);
    monAddr.sin_addr.s_addr = INADDR_ANY;
    bzero(&addrServ2, sizeof(addrServ2));
```

Full abstraction



Control Flow Integrity

```
bool lt(int x, int y) {  
    return x < y;  
}  
  
bool gt(int x, int y) {  
    return x > y;  
}  
  
sort2(int a[], int b[], int len)  
{  
    sort( a, len, lt );  
    sort( b, len, gt );  
}
```

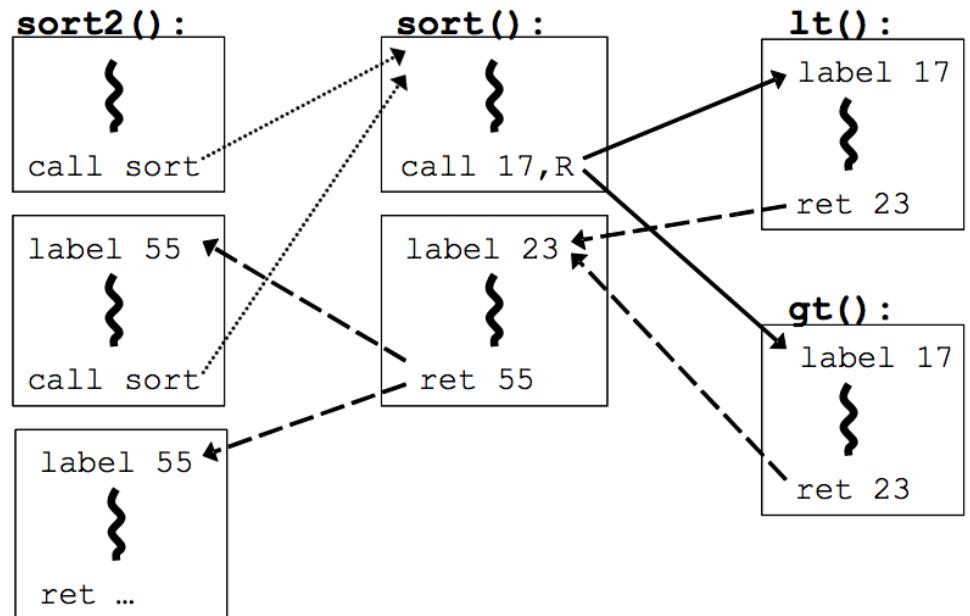


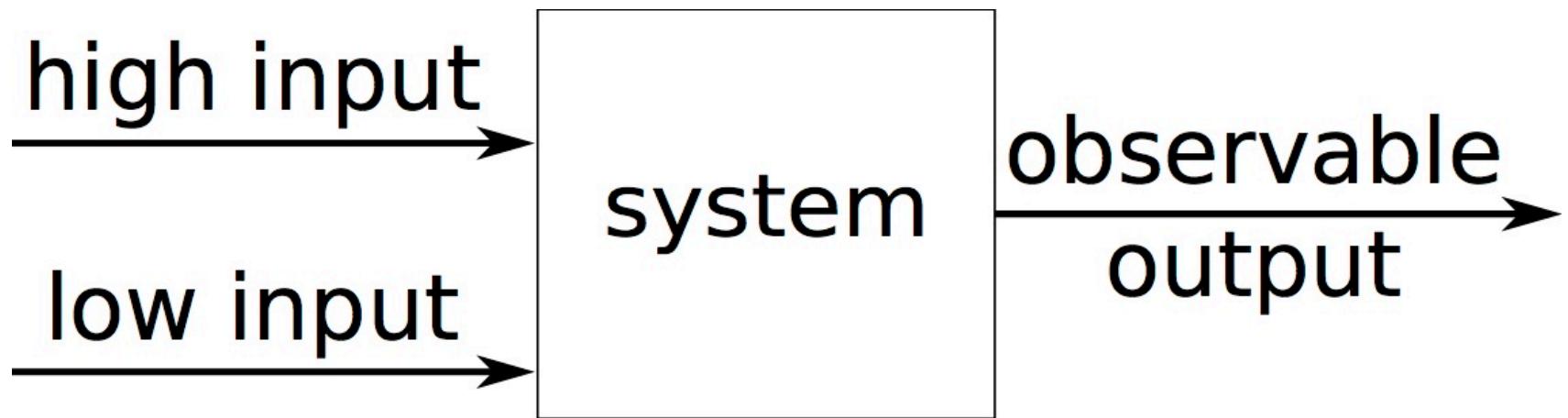
Figure 1: Example program fragment and an outline of its CFG and CFI instrumentation.

Control Flow Integrity (CFI)

The screenshot shows a debugger window displaying assembly code. A red arrow points from the top of the slide down to the instruction at address 100B832C. A green bracket is positioned below the instruction at address 100B8344, spanning from the start of the instruction to the end of the label 'loc_100B8348'. The assembly code is as follows:

```
100B8325 mov      eax, [ebx]
100B8327 mov      esi, [eax+50h]
100B832A mov      ecx, esi
100B832C call    ds:_guard_check_icall_fptr
100B8332 mov      ecx, ebx
100B8334 call    esi
100B8336 mov      edi, [ebp+var_9AC]
100B833C xor      edx, edx
100B833E mov      ecx, [edi+30h]
100B8341 cmp      ecx, 1
100B8344 jbe    short loc_100B8348
```

Noninterference



Enforcing noninterference

$$[\text{E1-2}] \quad \vdash \exp : \text{high} \qquad \frac{h \notin \text{Vars}(\exp)}{\vdash \exp : \text{low}}$$

$$[\text{C1-3}] \quad [pc] \vdash \text{skip} \quad [pc] \vdash h := \exp \qquad \frac{\vdash \exp : \text{low}}{[low] \vdash l := \exp}$$

$$[\text{C4-5}] \quad \frac{[pc] \vdash C_1 \quad [pc] \vdash C_2}{[pc] \vdash C_1; C_2} \qquad \frac{\vdash \exp : pc \quad [pc] \vdash C}{[pc] \vdash \text{while } \exp \text{ do } C}$$

$$[\text{C6-7}] \quad \frac{\vdash \exp : pc \quad [pc] \vdash C_1 \quad [pc] \vdash C_2}{[pc] \vdash \text{if } \exp \text{ then } C_1 \text{ else } C_2} \qquad \frac{[\text{high}] \vdash C}{[low] \vdash C}$$

Verification



SLAM



Software Analyzers

IT WORKS

development version of CompCert is the only compiler we have tested for which Csmith cannot find wrong-code errors. This is not for lack of trying: we have devoted about six CPU-years to the task. The apparent unbreakability of CompCert supports a strong argument that developing compiler optimizations within a proof framework, where safety checks are explicit and machine-checked, has tangible benefits for compiler users. (PLDI 2011)

Better testing (Heartbleed)

```
2561     /* Read type and payload length first */
2562     hbtype = *p++;
2563     n2s(p, payload);
2564     p += 2;
2565     pl = p;
2566
2567     if (s->msg_callback)
2568         s->msg_callback(0, s->version, TLS1_RT_HEARTBEAT,
2569                         &s->s3->rrec.data[0], s->s3->rrec.length,
2570                         s, s->msg_callback_arg);
2571
2572     if (hbtype == TLS1_HB_REQUEST)
```

1 Taking false branch →

```
          s->msg_callback(0, s->version, TLS1_RT_HEARTBEAT,
                          &s->s3->rrec.data[0], s->s3->rrec.length,
                          s, s->msg_callback_arg);
```

2 ← Assuming 'hbtype' is equal to 1 →

3 ← Taking true branch →

Better testing (Heartbleed)

```
2581     buffer = OPENSSL_malloc(1 + 2 + payload + padding);  
2582     bp = buffer;  
2583  
2584     /* Enter response type, length and copy payload */  
2585     *bp++ = TLS1_HB_RESPONSE;  
2586     s2n(payload, bp);  
2587     memcpy(bp, pl, payload);
```

4 ← Tainted, unconstrained value used in memcpy size

CONCLUSION

Measure things

- It can tell you what you should focus on
- It can tell you when what you are doing is working

Develop theories

- Understand and explain phenomenon
- Build systems right the first time