RANDOMIZED_QUICKSORT (A, P, r) if per q = RANDOMIZED_PARTITION (A,p,r); RANDOMIZED_QUICKSORT (A,P,q-1); RANDOMIZED_QUICKSORT (A, q+1, r);

• $(q = RANDOMIZED...) \rightarrow Define X_q = \begin{cases} 1, & \text{if the } q^{th} \text{ smallest element is picked as a pivot} \\ 0, & \text{otherwise} \end{cases}$

$T(n) = X_{1}(T(0) + T(n-1)) + X_{2}(T(1) + T(n-2)) + X_{3}(T(2) + T(n-3)) + \cdots + X_{n}(T(n-1) + T(n))$ $= T(0)[X_1 + X_n] + T(1)[X_2 + X_{n-1}] + \dots + T(n-1)[X_1 + X_n] + n$

$$T(n) = \sum_{q=1}^{n-1} T(q) [X_{q+1} + X_{n-q}] + n$$

$$\begin{array}{c} X: X_1, X_2 ... X_n \\ P_1, P_2 ... P_n \end{array}$$

$$\mathbb{E}(X) = \sum_{i=1}^{n} X_i P_i$$

Expectation:

$$\mathbb{E}(X_q) = \frac{1}{n}$$

$$\mathbb{E}\left(X_{q}\right) = \frac{1}{n} \cdot 1 + 0 \cdots$$

Example:

$$X = \begin{cases} a, P \\ b, q \end{cases} \implies \mathbb{E}(X) = ap + bq$$

$$\begin{split} \mathbb{E}\left(\mathsf{T}(\mathsf{n})\right) &= \mathbb{E}\left(\sum_{q=1}^{\mathsf{n}-1} \mathsf{T}(q) \left[\mathsf{X}_{q+1} + \mathsf{X}_{\mathsf{n}-q}\right]\right) + \mathsf{n} \\ &= \sum_{q=1}^{\mathsf{n}-1} \mathbb{E}\left(\mathsf{T}(q) \left(\mathsf{X}_{q+1} + \mathsf{X}_{\mathsf{n}-q}\right)\right) = \sum_{q=1}^{\mathsf{n}} \mathbb{E}\left(\mathsf{T}(q) \mathsf{X}_{q+1} + \mathsf{T}(q) \mathsf{X}_{\mathsf{n}-q}\right) \end{split}$$

$$= \sum_{q=1}^{n-1} \left(\underbrace{\mathbb{E}\left(T(q) \times_{q+1}\right)}_{n} + \underbrace{\mathbb{E}\left(T(q) \times_{n-q}\right)}_{n} \right)$$

$$= \frac{2}{n} \sum_{q=1}^{n-1} T(q) + n$$

$$\mathbb{E}\left(\mathsf{T}(q)\;\mathsf{X}_q\right) = \frac{\mathbb{E}\left(\mathsf{T}(q)\right)}{\mathsf{n}}$$

$$\mathbb{E}\left(\mathsf{T}(q)\;\mathsf{X}_{\mathsf{n}-q}\right) = \;\;\frac{\mathsf{T}(q)}{\mathsf{n}}$$

$$\mathbb{E}\left(\mathsf{T}(\mathsf{n})\right) = \frac{2}{\mathsf{n}} \sum_{\mathsf{q}=1}^{\mathsf{n}-1} \mathbb{E}\left(\mathsf{T}(\mathsf{q})\right) + \mathsf{n}$$

$$\mathbb{E}(T(n)) = \begin{cases} 1, & \text{if } n=1 \\ \frac{2}{n} \sum_{q=1}^{n-1} \mathbb{E}(T(q)) + n, & n>1 \end{cases}$$

· We need to prove that this solves to $O(n \log n)$ $\exists c : \mathbb{E}(T(n)) \leq c n \log n + 1$

- Induction step:

Assume $\mathbb{E}[T(q)] \leq c q \log q + 1$, for all q < n

$$\mathbb{E}(T(n) = \frac{2}{n} \sum_{q=1}^{n-1} \mathbb{E}(T(q)) + n \leq \frac{2}{n} \sum_{q=1}^{n-1} (c \cdot q \log q + 1) + n$$

$$= \frac{2c}{n} \sum_{q=1}^{n-1} q \log q + 2n - 1$$

Bound:
$$\frac{\sum_{q=1}^{n-1} q \log q}{\sum_{q=1}^{n} q \log q} \leq \int_{1}^{n} x \log x \, dx$$
$$= \frac{1}{\ln 2} \int_{1}^{n} x \ln x \, dx = \frac{1}{\ln 2} \int_{1}^{n} \left(\frac{x^{2}}{2}\right) \ln x \, dx$$

integration by parts:

$$\int f'(x)g(x) dx = \int \frac{1}{\ln 2} \left[\frac{x^2}{2} \ln x - \int_1^n \frac{x^2}{2} \left(\frac{1}{x} \right) dx \right] \\
= \int f(x)g(x) dx = \int \frac{1}{\ln 2} \left[\frac{x^2}{2} \ln x - \frac{x^2}{4} \right] = \int \frac{1}{\ln 2} \left(\frac{n^2}{2} \ln n - \frac{n^2}{4} + \frac{1}{4} \right) \\
= \frac{n^2}{2} \log n - \frac{n^2}{4 \ln 2} + \frac{1}{4 \ln 2}$$

$$\mathbb{E}(T(n) = \frac{2}{n} \sum_{q=1}^{n-1} \mathbb{E}(T(q)) + n \leq \frac{2}{n} \sum_{q=1}^{n-1} (c \cdot q \log q + 1) + n$$

$$= \frac{2c}{n} \sum_{q=1}^{n-1} q \log q + 2n - 1$$

$$\leq \frac{2c}{n} \left(\frac{n^2}{2} \log n - \frac{n^2}{4 \ln 2} + \frac{1}{4 \ln 2} \right) + 2n - 1$$

$$= c \cdot n \log n - \frac{cn}{2 \ln 2} + \frac{c}{2n \ln 2} + 2n - 1$$

$$\leq c \cdot n \log n$$

$$\leq c \cdot n \log n$$

$$\leq c \cdot n \log n + 1$$