Agenda

1) Revision quizzes
2) unique element (rest all coming thrice)
3) every element will occur twice except two ele

} very important

2.5 hrs
class notes revision → 25 min step
solving assign

If the last bit(0^{th} bit) of a number is 1 (rest can be anything), then the number is?

6 users have participated

A	Even	33%
B ✓	Odd	67%
C	Cannot be determined	0%
D	1	0%

25 → 11001

24 → 11000

What is the time complexity of the operation to check whether the i-th bit of a number with N bits is set or not?

no. A

→ check ith bit : A & (1<<i) → 0 (off), non-zero (on)

→ set ith bit : A | (1<<i)

→ toggle ith bit : A^ (1<<i)

What is the time complexity to count the total number of set bits of an integer with N bits ?

TC: O (no. of bits)

↓

O(N)

(Based on last class ques logic)

↓

no. bits = 32

32 itr ⇒ O(1)

Q. Given A[], all elements occurs thrice except one element. Find the unique element.

A = [4 5 5 4 1 6 6 4 5 6]
    0 1 2 3 4 5 6 7 8 9

ans = 1

A = [5 7 5 9 7 11 11 7 5 11]
    0 1 2 3 4 5 6 7 8 9

ans = 9

→ Brute-force : two loops

for (go to every element)

TC: $O(n^2)$

for (find freq of ele by travelling complete)

Demand : TC → O(n)

SC → O(1)

A = [5 7 5 9 7 11 11 7 5 11]

0 1 2 3 4 5 6 7 8 9

ans = 1001 = 9

3 2 1 0

5 → 0 1 0 1

7 → 0 1 1 1

5 → 0 1 0 1

9 → 1 0 0 1

7 → 0 1 1 1

11 → 1 0 1 1

11 → 1 0 1 1

7 → 0 1 1 1

5 → 0 1 0 1

11 → 1 0 1 1

4 6 6 10

1 0 0 1 → 9

0 to 31

0th bit ⇒ 10 (10%3 == 1)

1th bit ⇒ 6 (6%3 == 0)

2nd bit → 6 (6%3 == 0)

3rd bit → 4 (4%3 == 1)

```
                                  ↷N
function uniqueNum ( A[] )

    ans = 0;

    for ( i → 0 to 31 ) {

        count = 0;

        // in how many elements of A[], ith bit is "on"

        for (j → 0 to N-1) {

            if ( (A[j] & (1<<i)) != 0) {

                count++;

            }

        }

        if ( count % 3 == 1 ) {

            // set ith bit in ans

            ans = ans | (1<<i);

        }

    }

    return ans;

}
```

TC: O(N)

SC: O(1)

```
function uniqueNum(A[], N)
    ans = 0;
    for (i → 0 to 31) {
        count = 0;
        // in how many elements of A[], ith bit is "on"
        for (j → 0 to N-1) {
            if ((A[j] & (1<<i)) != 0) {
                count++;
            }
        }
        if (count % 3 == 1) {
            // set ith bit in ans
            ans = ans | (1<<i);
        }
    }
    return ans;
}
```

A = [8 5 8 8 6 5 5]
0 1 2 3 4 5 6

	3	2	1	0
8 →	1	0	0	0
5 →	0	1	0	1
8 →	1	0	0	0
8 →	1	0	0	0
6 →	0	1	1	0
5 →	0	1	0	1
5 →	0	1	0	1
count :	3	4	1	3

ans = 0 1 1 0 0
= 6

to do : dry run this eg. using code

A = [5 7 5 9 7 11 11 7 5 11]
0 1 2 3 4 5 6 7 8 9

Q. Given an A[] in which all elements will occur twice except two elements. Find those two elements.

A = [4 5 4 1 6 6 5 3] ⇒ 1 3
    0 1 2 3 4 5 6 7

A = [9 5 6 7 6 2 7 5] ⇒ 2 9
    0 1 2 3 4 5 6 7

Brute force : $O(n^2)$ {two loops}

↓ better

$O(n \log n)$ {sorting}

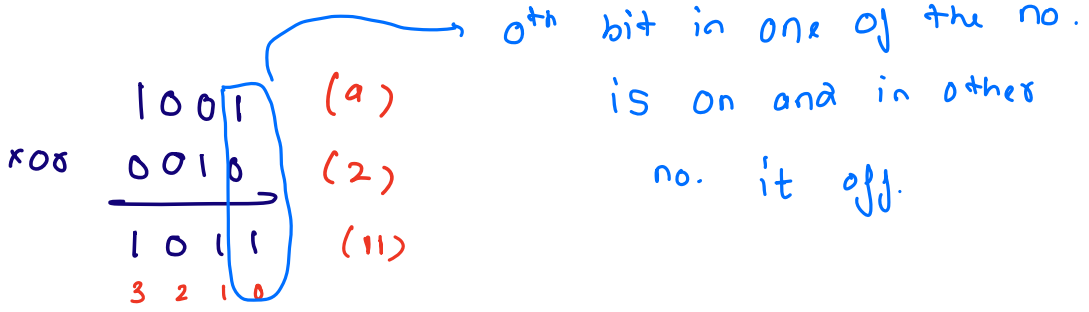
A = [9 5 6 7 6 2 7 5]
    0 1 2 3 4 5 6 7

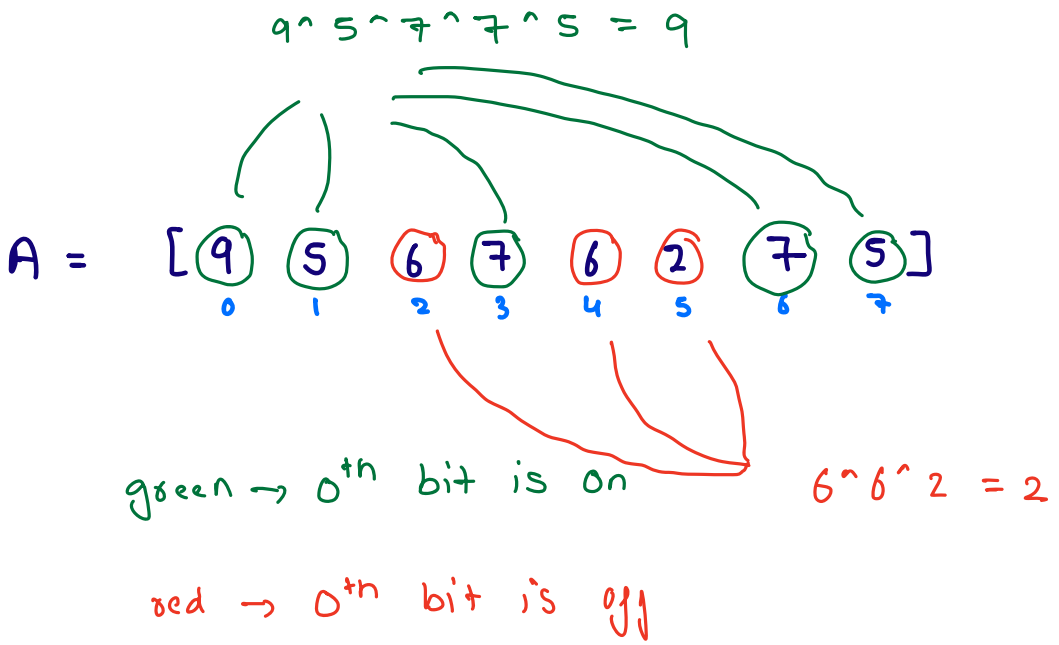
xorAll = 9 ^ 5 ^ 6 ^ 7 ^ 6 ^ 2 ^ 7 ^ 5

= 9 ^ 2

= 11

xor of two numbers, find those numbers.





→ A = [9 5 6 7 6 11 7 5]
0 1 2 3 4 5 6 7

i) xorAll = xor of entire array = 9^5^6^7^6^11^7^5

= 9^11

= 2

```
  1001
xor 1011
-------
  0010
  3210
```

ii) In variable "xorAll" find the first bit which is on

0 to 31 ⇒ ans : 1

iii) categorising the array in two parts

→ ele in which 1st bit is on

ele in which 1st bit is off

A = [9 5 6 7 6 11 7 5]
0 1 2 3 4 5 6 7

. . . . — — — —
3 2 1 0

num1 = 6 ^ 7 ^ 6 ^ 11 ^ 7 = 11
(1st bit is on)

num2 = 9 ^ 5 ^ 5 = 9
(1st bit is off)

```
                    ↱ N
function solve ( A[] ) {

    //find xor of all elements
    xorAll = 0;
    for (i -> 0 to N-1) {
        xorAll = xorAll ^ A[i];
    }

    // in xorAll, find the pos^n of rightmost set bit
    pos = -1;
    for(i -> 0 to 31) {
        if ( ( xorAll & (1<<i) ) != 0) {
            pos = i;
            break;
        }
    }

    // categorise array into 2 parts based of bit at pos
    num1 = 0; // contribution of ele in which bit at pos is on
    num2 = 0; // contribution of ele in which bit at pos is off
    for (i -> 0 to N-1) {
        if ( (A[i] & (1 << pos)) != 0) {
            num1 = num1 ^ A[i];
        }
        else {
            num2 = num2 ^ A[i];
        }
    }

    print(num1);
    print(num2);
}
```

TC : O(N)

SC : O(1)

```
function solve (A[], N) {

    //find xor of all elements
    xorAll = 0;
    for (i -> 0 to N-1) {
        xorAll = xorAll ^ A[i];
    }

    // in xorAll, find the pos^n of rightmost set bit
    pos = -1;
    for (i -> 0 to 31) {
        if ((xorAll & (1<<i)) != 0) {
            pos = i;
            break;
        }
    }

    // categorise array into 2 parts based of bit at pos
    num1 = 0; // contribution of ele in which bit at pos is on
    num2 = 0; // contribution of ele in which bit at pos is off

    for (i -> 0 to N-1) {
        if ((A[i] & (1 << pos)) != 0) {
            num1 = num1 ^ A[i];
        }
        else {
            num2 = num2 ^ A[i];
        }
    }

    print(num1);
    print(num2);
}
```

A = [5 9 5 6 6 2 13 2]
0 1 2 3 4 5 6 7

1) xorAll = 5^9^5^6^6^2^13^2
= 9^13 = 4

9 → 1 0 0 1
13 → 1 1 0 1
4 → 0 1 0 0
3 2 2 0

2) pos = 2

3) A = [5 9 5 6 6 2 13 2]
0 1 2 3 4 5 6 7

num1 = 5^5^6^6^13 = 13

num2 = 9^2^2 = 9

3 2 1 0

to do: dry run on one eg.