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## *Micro : Assignment 2*

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**Code :**

<https://drive.google.com/drive/folders/1KpUcmQ6zx6vvuzRWXDNyOYjgNSWSKXmF?usp=sharing>

The assignment consists of 3 classes :

- Montgomery\_M\_M
- Part1
- Part2

### ★ **Montgomery\_M\_M :**

- **private boolean enter\_to\_reduction :**  
Variable to know if it enters the addition reduction step or not .
- **boolean check () :**  
To return the enter\_to\_reduction variable when called .

- **Calculate\_unchanged\_values(BigInteger N) :**

It calculates the unchanging values which depend on the modulus only .

$$R = 2^{(\text{bit length of } N)} .$$

$$N' = -1 * N.\text{modInverse}(R).$$

$$R\_Inverse = R.\text{modInverse}(N).$$

It returns them as an array of BigInteger of size =3.

- **mog\_mod\_mul(BigInteger a, BigInteger b, BigInteger N, BigInteger R, BigInteger n\_bar, BigInteger R\_Inverse, int part)**

```
public BigInteger mog_mod_mul(BigInteger a, BigInteger b, BigInteger N, BigInteger R, BigInteger n_bar, BigInteger R_Inverse, int part) {
    // set enter_to_reduction to false
    enter_to_reduction = false;
    // calculate a' = a R mod N      b' = b R mod N
    BigInteger a_bar = (a.multiply(R)).mod(N);
    BigInteger b_bar = (b.multiply(R)).mod(N);
    // t= a' b'      m= t N' mod R      t= (t + mN) /R
    BigInteger t = a_bar.multiply(b_bar);
    BigInteger m = (t.multiply(n_bar)).mod(R);
    t = (t.add(m.multiply(N))).divide(R);
    BigInteger c_bar = BigInteger.valueOf(0);
    // if t < N  c'=t      else      c' = t-N
    if (t.compareTo(N) < 0)
        c_bar = t;
    else {
        if (part == 1)
            c_bar = (t.subtract(N));
        else if (part == 2) {
            enter_to_reduction = true ;
            for (int i = 0; i < 10; i++)
                c_bar = (t.subtract(N));
        }
    }
    // c = c' R^(-1) mod N
    return ((c_bar.mod(N)).multiply(R_Inverse)).mod(N);
}
```

- First it sets the enter\_to\_reduction variable equals false.

- Calculate

$$\rightarrow a' = a * R \bmod N$$

$$\rightarrow b' = b * R \bmod N$$

$$\rightarrow T = a' * b'$$

$$\rightarrow M = t * N' \bmod R$$

$$\rightarrow T = (t + m * N) / R$$

- Initialize  $c'$  equals zero .
- Check if  $t < N$  then  $C' = t$
- Other wise check if part 1 is calculated just put  $C' = t - N$
- If part 2 then set `enter_to_reduction` equals true and make the subtraction for 10 times without changing the result .
- Finally calculate  $c = c' * R^{(-1)} \bmod N$  and returns it .

## ★ Part1:

The Part1 class contains 4 functions “ the past implementation “ with modification :

- Function to calculate the time in msec of decryption in the **main** function .
- Instead of multiplication in the **modExp** function , an object from the `Montgomery_M_M` class is taken to call the function **mog\_mod\_mul** which calculates the result instead of multiplication .

Challenges i faced was to decrease the time as much as possible so i observed that  $R$  ,  $N'$  and  $(R_{\text{inverse\_mod } N})$  depend only on the  $N$  so they can be calculated once using **Calculate\_unchanged\_values** function and sent as a parameter to function **mog\_mod\_mul** .

## Important functions :

```
private static BigInteger modExp(BigInteger a, BigInteger exponent, BigInteger N) {
    Montgomery_M_M obj = new Montgomery_M_M();
    BigInteger [] temp = obj.Calculate_unchanged_values(N);
    // Note: This code assumes the most significant bit of the exponent is 1, i.e., the exponent is not zero.
    BigInteger result = a;
    int expBitlength = exponent.bitLength();
    for (int i = expBitlength - 2; i >= 0; i--) {
        result = obj.mog_mod_mul(result, result, N, temp[0], temp[1], temp[2], part: 1);
        if (exponent.testBit(i)) {
            result = obj.mog_mod_mul(result, a, N, temp[0], temp[1], temp[2], part: 1);
        }
    }
    return result;
}
```

- First object from Montgomery is taken .
- “Calculate \_unchanged\_values “ is called to calculate  $R, N'$  and  $R_{\text{inverse\_mod } N}$  and store them in array temp to be calculated once .
- Initialize the result equals to a.
- Loop on the bit length of the exponent from **n-2** to **0** .
- Call the “ mog\_mod\_mul “ to calculate  $\text{result} * \text{result} \bmod N$ .
- If the bit equals one , calculate  $\text{result} * a \bmod N$  using mog .
- Finally returns the result .

This function is called twice , one from the encrypt function to calculate  $m^e \bmod N$  and the other time from the decrypt method to calculate  $c^d \bmod N$ .

## Sample Runs :

```
"C:\Program Files\Java\jdk-11.0.2\bin\java.exe" "-javaa  
Original message = 4fde46bdb80fea15b33cb386609556a80ca7  
Ciphertext = 4267a27937f4965b788331c3149d7e53d6328db8ea  
Decrypted message = 4fde46bdb80fea15b33cb386609556a80ca  
Time = 194 msec
```

```
Run: Montgomery_M_M x  
"C:\Program Files\Java\jdk-11.0.2\bin\jav  
Original message = 171f4fe326bc0d8c2f3b1c  
Ciphertext = 179566bcb32fab4bbd7a2cf6207c  
Decrypted message = 171f4fe326bc0d8c2f3b1  
Time = 200 msec  
Process finished with exit code 0
```

## ★ Part 2 :

The part2 contains 4 functions :

- **The main function :**
  - Declare the lengths in an array of size = 6 .
  - Loop on this array and for each length call the `Calculate_for_each_length` function .
  - Get the result and calculate how many ones are over 20 times it succeeded and print it .
- **`Calculate_for_each_length(BigInteger d) :`**
  - Initialize result array of size = 20 .
  - Loop on this array and each time call `Calculate_experiment` and store the result either 0 or 1 in the array .
- **`boolean modExp(BigInteger a, BigInteger exponent, BigInteger N , boolean bit) :`**

Looks like the one in part 1 but with little difference , it takes a **boolean bit** as a parameter which equals true if the target bit is assumed to be one and false otherwise .

  - An object from `Montgomery _M_M` class is taken .
  - Calculate the unchanged values .
  - Initialize the result equals a .
  - Either the target bit is assumed the square step should be calculated so set  $result = result * result \bmod N$  using `mod` .

- If the target bit is assumed to be one set  
 $\text{result} = \text{result} * m \bmod N$ .
  - Start looping on the exponent bit length from **n-3** to **0**.
  - Only for the 3rd most bit check if it enters the reduction  
 save the boolean res to be returned when finished.
  - Continue looping as part 1.
  - Return res.
- **BigInteger average (LinkedList<BigInteger> list) :**
    - Loop on the list to get the sum of its element.
    - Calculate average = (sum)/ size of list.
    - Return average.
- **int Calculate\_experiment(BigInteger d) :**
    - Initialize 4 linked lists to store the time .  
 If the target bit assumed to equal 1  
 T 11 -> it enters the reduction.  
 T 10 -> it doesn't enter .  
 If the target bit assumed to equal 0  
 T 01 -> it enters the reduction.  
 T 00 -> it doesn't enter .
    - Call the mod Exp with bit = true and if it returns true, store  
 the time in T11. otherwise , store it in T 10 .
    - Call the mod Exp with bit = false and if it returns true,  
 store the time in T01. otherwise , store it in T 00 .
    - Calculate the average of each linked list .

★ By printing the average values i noticed that the function of time isn't so accurate so i tried to fix this with this check .

- Check if  $\text{avg } 11 > \text{avg } 10$  and the  $\text{abs}(\text{avg } 00 - \text{avg } 01)$  is less than the  $\text{abs}(\text{avg } 11 - \text{avg } 10)$  , return 1 as succeeded
- Otherwise return 0 .

## Sample Runs :

Length	No of succeeded
3	8
5	9
10	11
20	9
50	13
100	14

I observed that the output is not accurate but i think this because of my laptop as i hear his sound cannot handle all these samples .