# Matroska File Format (under construction!)

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# 1 Introduction

This document is intended to be used by developers who want to implement support for the MATROSKA file format in their applications, but who want to build this support from scratch rather than using existing implementations, or people who just want to understand the MATROSKA file format in detail. Thus, the file format itself is described, the usage of existing libraries isn't.

This document does not replace the official documentation<sup>1</sup>. It is less condensed, but not necessarily complete. Especially, in the case that MATROSKA supports Digital Restrictions Management one day, I will expressively not document that part. Also, typos in element IDs are never impossible.

When speaking about element occurence, elements can be mandatory or not, elements may be present several times inside a parent element or not etc. Occurence restrictions will be indicated using expressions like = 1 or  $\ge 1$  etc. Those restrictions will exclude cases which do not technically render a file unusable or ambigous, but which are unreasonable, like a file with no **Segmentuld**, see section 5.2. The same way it would be weird (but not make a file unusable) to have a **Chapters** element (which is supposed to describes chapters) which is empty. An element that must occur at least once is a reasonable file is called "mandatory". When an element is really mandatory, i.e the file or a part of it is useless when it's missing, it will be labeled as  $\ge 1$  (!) or = 1 (!). An example would be the codec ID of a track, without which a track cannot be decoded at all.

The official Matroska specification pages use the following interpretation of "mandatory" and "default": When an element has a default value that is used if the element itself is not present, the value cannot be missing, thus the element is inherently mandatory. This interpretation of "mandatory" being weird, this document considers an element mandatory when it must be physically present in the file. Also, default values can only be valid values. Consequently, a mandatory element cannot have a default value because if it had one, it couldn't be mandatory anymore.

In this document, element names are always printed like **THIS**, element values are printed like **\$THIS**, as in "if **\$THISFLAG**=1, ...".

If you have any questions concerning this document, if you have comments, additions, if you have found an error, or if you want to contact me for whatever reason,

<sup>&</sup>lt;sup>1</sup>http://www.matroska.org/technical/specs/index.html

#### MATROSKA file format

please send me an e-mail (include 'matroska' in the topic!). You can contact me in german, english or french, whatever you prefer. Just don't ask me if you can ask something or if I could document some Digital Restrictions Management.

This document is powered by LaTeX, so changing the order of certain tables or the style of those tables etc. is, with certain limits, possible within a few seconds.

Screenshots of real life file structures are used to illustrate the file structure. All of them have been made using the EBML Tree Viewer in AVI-Mux GUI.

## 2 EBML - basics

EBML files use integers of variable size. This way, the file format doesn't waste space with storing 32 or even 64 bit integers in placed where they *might sometimes* occur. The way the size is coded is inspired by the UTF-8 encoding format.

# 2.1 Unsigned Integer Values of Variable Length ("vint")

The length of an integer is equivalent to  $length = 1 + [number\_of\_leading\_zero\_bits]$ . All integers use big endian. You could use more than 7 leading zeros, then the first byte would be 0x00, however, this would only be needed if integers longer than 56 bits are required. This is forbidden in MATROSKA files.

#### Example: 3A 41 FE:

The first byte 3A (0011 1010) has 2 leading zeros, resulting in a total length of 3 bytes. The first '1' in the byte (0011 1010) is just needed to finish the sequence of leading zeros and can't be used to store the value either. Thus, it is reset to obtain the value this byte sequence represents. The result is then 0x1A41FE. As you can see, you lose one bit per byte to know how long a number is, and you can use 7 bits per byte to store the integer's value itself.

Of course, the value 0x1A41FE could also be written as 10 1A 41 FE or 08 00 1A 41 FE (do the decoding on a piece of paper if it's not clear), however, when writing EBML files, the shortest possible encoding should be used to avoid wasting space, which is the very point of this coding scheme.

## **Unknown Length**

All bits after the leading zeros being set to one, such as FF or 7F FF, indicates an *unknown length*. Muxers shall avoid writing unknown length values whenever possible. The only exception is the last Level 0 element of a file. If encoding a number as described above results in such a sequence, it must be encoded again with a greater destination length. Example: When encoding 16383 as described above, the result is 7F FF. In 7F FF, all bits after the leading zero are set, which would indicate an unknown length. That means, the length is increased to 3, and the number is encoded again to 20 3F FF.

#### Note

It is possible to use a lookup table to determine the total length from the first byte.

The Matroska file format does not allow integer lengths greater than 8, meaning that the number of leading zeros is not higher than 7 and that the total length can always be retrieved from the first byte.

#### 2.2 EBML elements

One piece of information is stored the following way:

The length of ID shall be called s\_ID, the length of size shall be called s\_size. Elements that contain other EBML Elements are called *EBML Master elements*.

Generally, the order of EBML elements inside a parent element is not fixed. In some cases, a certain order is recommended, but it is never mandatory. Especially, no element order should be assumed inside small parent elements.

# 2.3 Signed Integer Values of Variable Length (svint)

Signed integers have the following value: Read the integer as Unsigned Integer and then subtract

# 2.4 Data Types

Whereas vints are used in the header section of EBML elements, the data types describes in this section occur in the data section.

## 2.4.1 Signed and Unsigned Integers (int and uint)

Integers, signed as well as unsigned, are stored in big endian byte order, with leading 0x00 (in case of positive values) and 0xFF (in case of negative values) being cut off (example for int: -257 is 0xFE 0xFF). An int/uint may not be larger than 8 bytes.

#### 2.4.2 Float

A Float value is a 32 or 64 bit real number, as defined in IEEE. 80 Bit values have been in the specification, but have been removed and should not be used. The bytes are stored in big endian order.

## 2.4.3 Types of Strings

**String** refers to an ASCII string. **UTF-8** refers to a string that is encoded as UTF-8

# 3 MATROSKA files - Top-Level elements

MATROSKA files only have two different top level elements:

```
EBML (header: 5 bytes, data: 35 bytes, pos.: 0): master
Segment (header: 12 bytes, data: 4,696,195,032 bytes, pos.: 40): master
```

## 3.1 EBML

This header describes the contents of an EBML file. There should be only one EBML header in one file. Any further EBML headers do not render a file invalid, but shall be ignored by any application reading the file. Files with more than one EBML header could be created for instance if two or more files are appended by using the copy /b command.

## 3.2 Segment

A **SEGMENT** contains multimedia data, as well as any header data necessary for replay. There can be several **SEGMENT**s in one MATROSKA file, but this is not encouraged to be done, as not many tools are able to handle multisegment MATROSKA files correctly. If you want to replay multisegment MATROSKA files on Windows, please use Haali Media splitter<sup>2</sup>

<sup>&</sup>lt;sup>2</sup>http://haali.cs.msu.ru/mkv/

# 4 EBML - The EBML file header

The **EBML** top level element contains a description of the file type, such as EBML version, file type name, file type version etc.

Obviously, this header being missing makes it necessary to guess the file type.

Table 1: The EBML element (Top-Level)		
Element	Description	
uint, $\# \le 1$	indicates the version of the EBML Writer that has been used	
EBMLVersion	to create a file	
ID: 42 86		
def: 1		
uint, $\# \le 1$	indicates the minimum version an EBML parser needs to be	
EBMLREADVERSION	compliant with to be able to read the file	
ID: 42 F7		
def: 1		
uint, $\# \le 1$	indicates the length of the longest EBML-ID the file contains.	
EBMLMAXIDLENGTH	In case of matroska, this value is 4. Any EBML-ID which is	
ID: 42 F2	longer than the value of this element shall be considered	
def: 4	invalid.	
uint, $\# \le 1$	indicates the maximum s_size value the file contains. Any	
EBMLMAXSIZELENGTH	EBML element having an s_size value greater than EBML-	
ID: 42 F3	MaxSizeLength shouldl be considered invalid.	
def: 8		

**EBML** continued on next page

Element	Description
string, $\# \le 1$	describes the contents of the file. In the case of a MATROSKA
DOCTYPE	file, its value is 'matroska'
ID: 42 82	
def: matroska	
<b>uint,</b> # $\leq 1$	indicates the version of the <b>\$DocType</b> writer used to create
DOCTYPEVERSION	the file
ID: 42 87	
def: 1	
uint, $\# \le 1$	indicates the minimum version number a \$DocType parser
DOCTYPEREADVERSION	must be compliant with to read the file.
ID: 42 85	
def: 1	

 $Index \rightarrow page 2 \qquad \qquad \textit{end of } \textbf{EBML}$ 

As you can see, in the case of Matroska files all child elements of the **EBML** element have a default value. Thus, an empty **EBML** element would technically introduce a Matroska file (with file type version 1, maximum ID length 4, maximum size length 8 etc.) correctly. However, I don't recommend to push the specifications like this.

It is not recommended to use either IDs or size values greater than 8 bytes. While it's clear that 8 bytes are enough to represent any size of anything on any hard disc, one might think about using IDs larger than 8 bytes. However, since the ID is considered an integer, treating IDs larger than 8 bytes is difficult on current CPUs, which are limited to 64 bit for simple integer operations.

# 5 Level 1 - Elements inside Segments

## 5.1 Overview

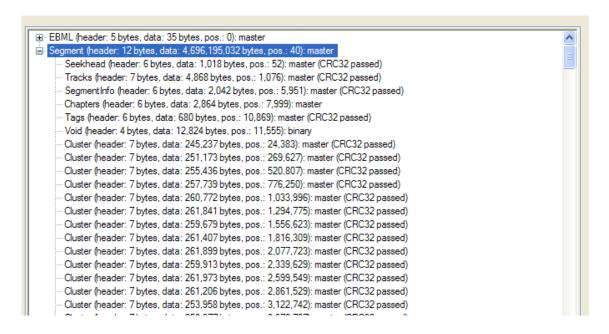


Table 2: The SEGMENT element (Top-Level)		
Element	Description	
<b>Master,</b> # = 1	SEGMENTINFO contains general information about a seg-	
SEGMENTINFO (→3)	ment, like an UID, a title etc. This information is not really	
ID: 15 49 A9 66	required for playback, but should be there ( $\rightarrow$ section 5.2).	
Master, $\# \ge 0$	A SEEKHEAD is an index of elements that are children of	
SEEKHEAD $(\rightarrow 4)$	<b>SEGMENT.</b> It can point to other <b>SEEKHEADS</b> , but not to itself.	
ID: 11 4D 9B 74	If all non-Cluster precede all Clusters ( $\rightarrow$ section 5.5),	
	a SEEKHEAD is not really necessary, otherwise, a missing	
	SEEKHEAD leads to long file loading times or the inability to	
	access certain data.	
Master, $\# \ge 0$	A CLUSTER contains video, audio and subtitle data. Note	
Cluster ( $\rightarrow$ 16)	that a MATROSKA file could contain chapter data or attach-	
ID: 1F 43 B6 75	ments, but no multimedia data, so <b>CLUSTER</b> is not a manda-	
	tory element.	

**SEGMENT** continued on next page

Element	Description
Master, $\# \ge 0$	A TRACKS element contains the description of some or all
Tracks $(\rightarrow 6)$	tracks (preferably all). This element can be repeated once in
ID: 16 54 AE 6B	a while for backup purposes. A file containing only chapters
	and attachments does not have a TRACKS element, thus it's
	not mandatory.
Master, $\# \leq 1$	The Cues element contains a timestamp-wise index to
<b>C</b> UES (→18)	CLUSTERS, thus it's helpful for easy and quick seeking.
ID: 1C 53 BB 6B	
Master, $\# \le 1$	The ATTACHMENTS element contains all files attached to this
Attachments (→26)	SEGMENT.
ID: 19 41 A4 69	
<b>Master,</b> # = 1	The CHAPTERS elements contains the definition of all chap-
Chapters $(\rightarrow 21)$	ters and editions of this SEGMENT
ID: 10 43 A7 70	
Master, $\# \le 1$	The TAGS element contains further information about the
<b>TAGS</b> (→28)	SEGMENT or elements inside the SEGMENT that is not really
ID: 12 54 C3 67	required for playback.

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end of **SEGMENT** 

# 5.2 SegmentInfo

The **SegmentInfo** element contains general information about the **Segment**, such as its duration, the application used for writing the file, date of creation, a unique 128 bit ID, to name a few only. Information included in the **SegmentInfo** element is not required for playback, but should be written by any Matroska muxer.

```
EBML (5: 35 bytes at 0): master

Segment (12: 1,621,732,571 bytes at 40): master

Seekhead (6: 1,018 bytes at 52): master

Tracks (7: 1,144 bytes at 1,076): master

CRC32 (2: 4 bytes at 2,233): C5 DB 1F 65

SegmentUID (3: 16 bytes at 2,239): 2E E8 2D 73 11 D5 DF 34 8C 6C 3D E7 43 11 32 A2

Duration (3: 4 bytes at 2,258): 57168000.00

WritingApp (3: 41 bytes at 2,265): AVI-Mux GUI 1.16.4, Mar 15 2004 17:19:57

MuxingApp (3: 55 bytes at 2,309): matroska muxer by Alexander Noe, build date Mar 14 2004

Title (3: 5 bytes at 2,367): Virus

TimecodeScale (4: 3 bytes at 2,375): 100,000

Void (9: 1,884 bytes at 2,382): binary

Cluster (7: 428,891 bytes at 4,275): master
```

(read: <element name> (<s\_size + s\_ID>: <size> bytes at <position in file>: value)

Table 3: The SegmentInfo element, child of Segment $(\rightarrow 2)$	
Element	Description
<b>char[16]</b> , # = 1	a unique 128 bit number identifying a <b>SEGMENT</b> . Obviously,
SEGMENTUID	a file can only be referred to by another file if a SEGMEN-
ID: 73 A4	TUID is present, however, playback is possible without that
	UID.
<b>utf-8,</b> # ≤ 1	contains the name of the file the <b>Segment</b> is stored in. Since
SEGMENTFILENAME	renaming files is easy, the reliability of this element's value
ID: 73 84	should not be overrated.
<b>char[16]</b> , # ≤ 1	contains the unique 128 bit ID of the SEGMENT that is re-
PrevUID	played before the currently active SEGMENT, i.e. the ID of
ID: 3C B9 23	the SEGMENT that should be loaded if the user tries to seek
	to a timecode earlier than the earliest timecode of the ac-
	tive <b>Segment</b> . That <b>segment</b> should, of course, be easy to
	locate, for instance in a file in the same directory.
<b>utf-8,</b> # ≤ 1	contains the name of the file in which the <b>Segment</b> having
PREVFILENAME	the ID \$PREVUID is stored. PREVFILENAME should not be
ID: 3C 83 AB	considered reliable for the same reason as SEGMENTFILE-
	NAME, however, it could be the first filename the player is
	looking for when the SEGMENT described in PREVUID is
	needed

**SEGMENTINFO** continued on next page

Element	Description
<b>char[16],</b> # ≤ 1	contains the unique 128 bit ID of the SEGMENT that is re-
NEXTUID	played after the currently active <b>SEGMENT</b> , i.e. the ID of the
ID: 3E B9 23	SEGMENT that should be loaded if the user tries to seek to
	a timecode after the end of the active SEGMENT. Like PRE-
	vUID, the corresponding SEGMENT should be easy to locate.
<b>utf-8,</b> # ≤ 1	contains the name of the file in which the SEGMENT having
NEXTFILENAME	the ID <b>\$NEXTUID</b> is stored. <b>NEXTFILENAME</b> shall not be
ID: 3E 83 BB	considered reliable for the same reason as SEGMENTFILE-
	NAME.
uint, $\# \le 1$	Each scaled timecode in a MATROSKA file is multiplied by
TIMECODESCALE	TIMECODESCALE to obtain a timecode in nanoseconds. Note
ID: 2A D7 B1	that not all timecodes are scaled!
float, $\# \leq 1$	The <b>DURATION</b> indicates the duration of the <b>SEGMENT</b> . The
DURATION	duration measured in nanoseconds is scaled and is thus
ID: 44 89	equal to <b>\$Duration</b> * <b>\$TimecodeScale</b> . This element
	should be written.
<b>utf-8,</b> # ≤ 1	Contains a general name of the SEGMENT, like "Lord of
TITLE	the Rings - The Two Towers". No language can be at-
ID: 7B A9	tached to the title, however, Tags ( $\rightarrow$ section 5.9) could be
	used to define several titles for a segment. This is not yet
	commonly done, though.
string, $\# = 1$	contains the name of the library that has been used to create
MUXINGAPP	the file (like "libmatroska 0.7.0"). This element should be
ID: 4D 80	written by any muxer! Especially if non-compliant files are
	encountered, this help to know who must be blamed for that
	file.
<b>utf-8,</b> # = 1	contains the name of the application used to create the file
WRITINGAPP	(like "mkvmerge 0.8.1"). This element should be written for
ID: 57 41	the same reason as MUXINGAPP.
int, # $\leq 1$	contains the production date, measured in nanoseconds rel-
DATEUTC	atively to Jan 01, 2001, 0:00:00 GMT+0h
ID: 44 61	

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end of SEGMENTINFO

## 5.3 SeekHead

The **Seekhead** element contains a list of positions of Level 1 elements in the **Segment**. Each pair (element id, position) is stored in one **Seek** element:

Table 4: The SEEKHEAD element, child of SEGMENT ( $ ightarrow$ 2)		
Element	Description	
Master, $\# \geq 1$	One SEEK element contains an EBML-ID and the position	
<b>S</b> eek (→5)	within the SEGMENT at which an element with this ID can	
ID: 4D BB	be found.	

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end of SEEKHEAD

Table 5: The SEEK element, child of SEEKHEAD ( $ ightarrow 4$ )		
Element	Description	
<b>uint</b> , # = 1	The SEEKID element contains the EBML-ID of the element	
SEEKID	found at the given position	
ID: 53 AB		
<b>uint,</b> # = 1	The SEEKPOSITION element contains the position relatively	
SEEKPOSITION	to the <b>S</b> EGMENT's data at which an element with the ID	
ID: 53 AC	\$SEEKID can be found.	

Index  $\rightarrow$ page 2 end of Seek

Not all Level 1 elements need to be included. Typical **SeekHeads** either include a list of all Level 1 elements, or a list of all Level 1 elements except for **Clusters** ( $\rightarrow$  section 5.5). **SeekHeads** can also include references to other **SeekHeads** if there is, for example, a small **SeekHead** at the beginning of the file and a larger one at its end.

The following picture illustrates the **SeekHead** element in a real file. Note that the EBML Tree Viewer replaced Level 1 IDs in **SeekID** with their human-readable name:



## 5.4 Tracks

The Tracks element contains information about the tracks that are stored in the Segment, like track type (audio, video, subtitles), the used codec, resolution and sample rate. All tracks shall be described in one (or more, but preferably only one) Tracks element.

Each track is described in one **TRACKENTRY**. Theoretically, using the **TRACKUID**, information about one track could be spread over different **TRACKENTRY**s, the UID would allow to know which track the information applies to, however, it is highly discouraged to stretch the specification like this.

Also, an empty TRACKS element would be rather useless, but should not lead to a parser error since the file can be played if all tracks are defined *somewhere*. Especially pure chapter files might have an empty TRACKS element if the muxer doesn't catch the case that no tracks are present and consequently creates an empty TRACKS element.

An example of a **TrackEntry** element can be found on ( $\rightarrow$  page 25)

Table 6: The TRACKS element, child of SEGMENT ( $ ightarrow$ 2)		
Element	Description	
Master, $\# \geq 1$	One TRACKENTRY element describes one track of the SEG-	
TrackEntry $(\rightarrow 7)$	MENT	
ID: AE		

 $Index \rightarrow page \ 2 \qquad \qquad \textit{end of Tracks}$ 

Table 7: The TRACKENTRY element, child of TRACKS ( $ ightarrow 6$ )	
Element	Description
<b>uint,</b> # = 1 (!)	defines an identification number of the track. This number
TRACKNUMBER	cannot be equal to 0. This number is used by the <b>BLOCK</b> and
ID: D7	SIMPLEBLOCK structures.
<b>uint,</b> # = 1	is a unique identificator of the track within the file. It cannot
TRACKUID	be equal to 0
ID: 73 C5	

TRACKENTRY continued on next page

Element	Description
<b>uint,</b> # = 1 (!)	defines the type of a track, i.e. video, audio, subtitle etc.
TrackType $(\rightarrow 13)$	
ID: 83	
bool, $\# \leq 1$	When FLAGENABLED is 1, track is used
FLAGENABLED	
ID: B9	
def: 1	
bool, $\# \leq 1$	When FLAGDEFAULT is 1, the track should be selected by
FLAGDEFAULT	the player by default. Obviously, if no video track and/or no
ID: 88	audio track has a default flag, one video track and one audio
def: 1	track should be chosen by the player, whereas no subtitle
	should be enabled if no subtitle has a default flag.
<b>bool,</b> $\# \le 1$	When FLAGFORCED is 1, the track must be played. When
FLAGFORCED	several subtitle tracks are forced, the one matching the au-
ID: 55 AA	dio language should be chose. An example would be a sub-
def: 0	title track that cannot be disabled, like the one you find on
	the german DVD "Eiskalte Engel" when you select english
	audio. Since this flag can only be used to apply a restriction
	on digital content, it must be qualified as Digital Restrictions
	Management.
<b>bool,</b> $\# \le 1$	When <b>FLAGLACING</b> is 1, the track may contain laced blocks.
FLAGLACING	A parser that supports all types of lacing ( $\rightarrow$ section 6.2) can
ID: 9C	safely ignore this flag.
def: 0	
<b>uint,</b> # $\leq 1$	indicates the number of frames a player must be able to
MINCACHE	cache during playback. This is for instance interesting if a
ID: 6D E7	native MPEG4 file with frames in coding order is played.
def: 0	
uint, $\# \le 1$	indicates the maximum cache size a player needs to cache
MAXCACHE	frames. A value of NULL means that no cache is required.
ID: 6D F8	

TRACKENTRY continued on next page

Element	Description
uint, $\# \le 1$	This value indicates the number of nanoseconds a frame
DEFAULTDURATION	lasts. This value is applied if no \$DURATION value is in-
ID: 23 E3 83	dicated for a frame or if lacing ( $\rightarrow$ section 6.1) is used. A
	value of 0 means that the duration of frames of the track
	is not necessarily constant (e.g. variable framerate video,
	or Vorbis audio). <b>DEFAULT DURATION</b> should be written for
	each track with a constant frame rate since it makes seeking
	easier.
float, $\# \leq 1$	Every timecode of a block (cluster timecode + block
TRACKTIMECODESCALE	timecode) is multiplied by this value to obtain the real time-
ID: 23 31 4F	code of a block.
<b>utf-8,</b> $\# \le 1$	A NAME element contains a human-readable name for the
NAME	track. Note that you can't define which language this track
ID: 53 6E	name is in. You have to use Tags ( $\rightarrow$ section 5.9)) if you
	want to use several titles in different languages for the same
	track.
string, $\# \leq 1$	specifies the language of a track, using ISO639-2 <sup>3</sup> . This
LANGUAGE	is NOT necessarily the language of \$NAME, for example a
ID: 22 B5 9C	german AC3 track could be called "German - AC3 5.1" or
def: eng	"Deutsch - AC3 5.1" or "Allemand AC3 5.1" etc.
<b>string,</b> # = 1 (!)	The CODECID specifies the Codec <sup>4</sup> which is used to decode
CODECID	the track.
ID: 86	
binary, $\# \leq 1$	CODECPRIVATE contains information the codec needs before
CODECPRIVATE	decoding can be started. An example is the Vorbis initializa-
ID: 63 A2	tion packets for Vorbis audio.
<b>utf-8,</b> # ≤ 1	CODECNAME is a human-readable name of the Codec
CODECNAME	
ID: 25 86 88	
uint, $\# \ge 0$	An ATTACHMENTLINK contains the UID of an attachment
ATTACHMENTLINK	that is used by this track.
ID: 74 46	

TRACKENTRY continued on next page

<sup>&</sup>lt;sup>3</sup>http://lcweb.loc.gov/standards/iso639-2/englangn.html <sup>4</sup>http://matroska.org/technical/specs/codecid/index.html

Element	Description
Master, $\# \le 1$	VIDEO contains information that is specific for video tracks
VIDEO (→8)	
ID: E0	
Master, $\# \le 1$	AUDIO contains information that is specific for audio tracks
<b>A</b> UDIO (→9)	
ID: E1	
Master, $\# \le 1$	CONTENTENCODINGS contains information about (lossless)
CONTENTENCODINGS	compression or encryption of the track
(→10)	
ID: 6D 80	

 $Index \rightarrow page \ 2$ 

end of TRACKENTRY

Obviously, the VIDEO element must be present for video tracks, whereas the AUDIO element must be present for audio tracks. Although it doesn't make sense to have both elements in one TRACKENTRY element, it wouldn't make a file unplayable.

Table 8: The VIDEO element, child of TRACKENTRY ( $\rightarrow$ 7)	
Element	Description
<b>uint,</b> # = 1	Width of the encoded video track in pixels
PIXELWIDTH	
ID: B0	
<b>uint,</b> # $\leq 1$	Height of the encoded video in pixels
PIXELHEIGHT	
ID: BA	
uint, $\# \le 1$	Number of Pixels to be cropped from the bottom
PIXELCROPBOTTOM	
ID: 54 AA	
def: 0	
uint, $\# \le 1$	Number of Pixels to be cropped from the top
PIXELCROPTOP	
ID: 54 BB	
def: 0	

**VIDEO** continued on next page

Element	Description
uint, $\# \le 1$	Number of Pixels to be cropped from the left
PIXELCROPLEFT	
ID: 54 CC	
def: 0	
uint, $\# \le 1$	Number of Pixels to be cropped from the right
PIXELCROPRIGHT	
ID: 54 DD	
def: 0	
uint, $\# \le 1$	Width of the video during playback
DISPLAYWIDTH	
ID: 54 B0	
def: <b>\$P</b> IXELWIDTH	
uint, $\# \le 1$	Height of the video during playback
DISPLAYHEIGHT	
ID: 54 BA	
def: <b>\$P</b> IXEL <b>H</b> EIGHT	
uint, $\# \le 1$	Unit \$DISPLAYWIDTH and \$DISPLAYHEIGHT is measured
DISPLAYUNIT	in. This can be $0\rightarrow$ pixels, $1\rightarrow$ centimeters, $2\rightarrow$ inches
ID: 54 B2	
def: 0	

Index  $\rightarrow$ page 2 end of VIDEO

**\$PIXELCROPXXXX** is applied on **\$PIXELXXX**, so the output is cropped after decoding, but before stretching it to the dimensions indicated with **\$DISPLAYXXXX**.

Table 9: The Audio element, child of TrackEntry ( $ ightarrow 7$ )	
Element	Description
<b>uint,</b> # $\leq 1$	Indicates the sample rate the track is encoded at in Hz
SAMPLINGFREQUENCY	
ID: B5	
def: 8 kHz	

**AUDIO** continued on next page

Element	Description
uint, $\# \le 1$	Indicates the sample rate the track must be played at in
Оитрит-	Hz. The default value of this element is equal to \$SAM-
SAMPLINGFREQUENCY	PLINGFREQUENCY.
ID: 78 B5	
uint, $\# \leq 1$	Number of channels of the audio track
CHANNELS	
ID: 9F	
def: 1	
uint, $\# \le 1$	Bits per sample, this is usually used with PCM-Audio.
Віт Дертн	
ID: 62 64	

Index →page 2 end of Audio

Table 10: The ContentEncodings element, child of TrackEntry ( $\rightarrow$ 7)	
Element	Description
Master, $\# \geq 1$	A CONTENTENCODING-element describes one compression
CONTENTENCODING	or encryption that has been used on this track.
(→11)	
ID: 62 40	

 $Index \rightarrow page \ 2$ 

end of ContentEncodings

Table 11: The ContentEncoding element, child of		
	CONTENTENCODINGS $(\rightarrow 10)$	
Element	Description	
uint, $\# \le 1$	Tells when to decode according to this pattern. The de-	
CONTENTENCODING-	coder starts with the CONTENTENCODING that has the high-	
ORDER	est ContentEncodingOrder.	
ID: 50 31		
def: 0		

CONTENTENCODING continued on next page

Element	Description
uint, $\# \le 1$	Defines which parts of the track are compressed or en-
CONTENTENCODING-	crypted this way
<b>S</b> COPE (→14)	
ID: 50 32	
def: 1	
<b>uint,</b> # $\leq 1$	Describes which type of encoding is described. $0 \rightarrow com-$
CONTENTENCODING-	pression, $1 \rightarrow$ encryption
ТүрЕ	
ID: 50 33	
def: 0	
Master, $\# \le 1$	If <b>CONTENTENCODINGTYPE=0</b> , this element describes how
CONTENTCOMPRESSION	it is compressed
(→12)	
ID: 50 34	
Master, $\# \le 1$	If CONTENTENCRYPTION = 1, this element describes how it
CONTENTENCRYPTION	is encrypted
(→??)	
ID: 50 35	

Index  $\rightarrow$ page 2

end of CONTENTENCODING

The **CONTENTENCODING** element allows to apply not only encryption, but also lossless compression to a track. This can be used to compress text subtitles, but also to remove sync headers from audio packets. For example, each AC3 frame starts with 0B 77, and there is no real point in saving those two bytes for each frame in a MATROSKA file. For a simple AC3 file, this does make sense because there it can be used to find a new frame start if data is damaged.

Table 12: The CONTENTCOMPRESSION element, child of		
	CONTENTENCODING $(\rightarrow 11)$	
Element	Description	
uint, $\# \le 1$	The CONTENTCOMPALGO element says which algorithm was	
CONTENTCOMPALGO	used for this compression.	
(→15)		
ID: 42 54		
def: 0		

**CONTENTCOMPRESSION** continued on next page

Element	Description
binary, $\# \leq 1$	Contains settings that are required for decompression.
CONTENTCOMPSETTINGS	These settings are specific for each compression algorithm.
ID: 42 55	For example, it contains the striped header bytes when
	\$CONTENTCOMPALGO=3 (→ page 25).

Index  $\rightarrow$ page 2

end of ContentCompression

Table 13: Values of TRACKTYPE, child of TRACKENTRY ( $ ightarrow 7$ )	
Value	Description
0x01	track is a video track
0x02	track is an audio track
0x03	track is a complex track, i.e. a combined video and audio track
0x10	track is a logo track
0x11	track is a subtitle track
0x12	track is a button track
0x20	track is a control track

end of TRACKTYPE

Table 14: Bits in ContentEncodingScope, child of	
CONTENTENCODING $(\rightarrow 11)$	
Value	Description
1	all frames
2	the track's CODECPRIVATE
4	the CONTENTCOMPRESSION in the next CONTENTENCODING (next as in next in decoding order)

end of CONTENTENCODINGSCOPE

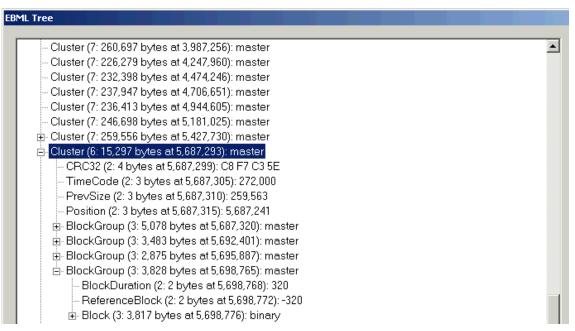
Here is one example of a possible **TrackEntry** element: A DTS-audio track that is using header striping. The **ContentCompSettings** element contains the four bytes each DTS frame starts with.

Table 15: Values of CONTENTCOMPALGO, child of	
CONTENT COMPRESSION $(\rightarrow 12)$	
Value	Description
0	zlib
1	bzlib
2	lzo1x
3	header striping

end of CONTENTCOMPALGO

## 5.5 Cluster

A **CLUSTER** contains multimedia data and usually spans over a range of a few seconds. The following picture shows a typical cluster:



Although sticking to this order of the elements is not mandatory, it is recommended not to have any non-BlockGroup/SimpleBlock after the first Block-Group/SimpleBlock, because it's bad if the entire cluster must be read before it can be used just because the timecode is stored at the end.

Table 16: The Cluster element, child of Segment $(\rightarrow 2)$	
Element	Description
<b>uint,</b> # $\leq 1$	The Cluster timecode is the timecode all block timecodes are
TIMECODE	indicated relatively to.
ID: E7	
def: 0	
<b>uint,</b> # $\leq 1$	The POSITION element indicates the position of the begin-
Position	ning of its parent element inside its grand parent element.
ID: A7	This can help to resync in case of damaged data, but is of no
	use if no data is damaged.

**CLUSTER** continued on next page

Element	Description
uint, $\# \le 1$	Indicates the size of the preceding cluster in bytes. This
PREVSIZE	helps to seek backwards, and to find the preceding cluster,
ID: AB	without having to look at METASEEK or CUE data. This is
	also helpful to resync, e.g. if the EBML-ID of the preceding
	CLUSTER is damaged.
Master, $\# \ge 0$	Contains a <b>BLOCK</b> along with some attached information
BlockGroup (→17)	like references
ID: AO	
binary, $\# \ge 0$	This is a <b>BLOCK</b> ( $\rightarrow$ page 39) without additional attached in-
SIMPLEBLOCK	formation. Since a SIMPLEBLOCK does not require a BLOCK-
ID: A3	GROUP around it, it causes less overhead. SIMPLEBLOCK is
	Matroska v2.

 $Index \rightarrow page \ 2 \qquad \qquad \textit{end of Cluster}$ 

Table 17: The BLOCKGROUP element, child of CLUSTER $(\rightarrow 16)$	
Element	Description
<b>binary,</b> # = 1 (!)	contains data to be replayed. See page 39 for details.
Вгоск	
ID: A1	
int, # $\geq 0$	Timecode of a frame, relative to the BLOCK's timecode, of
REFERENCEBLOCK	a frame that needs to be decoded before this <b>BLOCK</b> can be
ID: FB	decoded.
int, # $\leq 1$	Indicates the scaled duration of the BLOCK. If this
BLOCKDURATION	value is not written, it is assumed to be (1) the differ-
ID: 9B	ence <timecode block="" next="" of="" same="" stream="" the=""> -</timecode>
	<pre><timecode> (2) equal to DefaultDuration (for the last</timecode></pre>
	block of each stream).
	As a consequence, the <b>DURATION</b> element is mandatory for
	every BLOCK of subtitle tracks, unless a subtitle is indeed
	supposed to disappear only directly before the next one ap-
	pears. But even then it is recommended to write <b>DURATION</b> .

Index →page 2

end of BLOCKGROUP

## **5.6** Cues

The Cues element contains information helpful (but not necessary) for seeking. Each piece of information, called a CuePoint, contains a timestamp, and a list of pairs (track number, (cluster position[, block number within cluster])). Generally, a CuePoint should only point to keyframes.

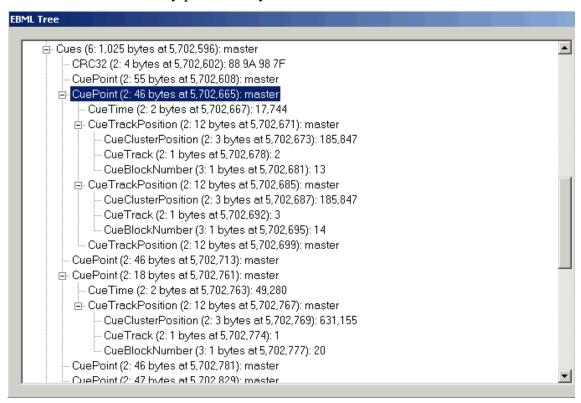


Table 18: The CUES element, child of SEGMENT $(\rightarrow 2)$	
Element	Description
Master, $\# \geq 1$	One CUEPOINT contains one entry point (or a list of entry
CuePoint $(\rightarrow 19)$	points with one point for one track) for one timecode.
ID: BB	

Index  $\rightarrow$ page 2 end of Cues

Table 19: The CUEPOINT element, child of CUES ( $ ightarrow$ 18)	
Element	Description
<b>uint,</b> # = 1 (!)	The timecode of the Clusters or Blocks that are referred
CUETIME	to by this CuePoint
ID: B3	
Master, $\# \geq 1$	A position where a Cluster or Block can be found with
CUETRACKPOSITIONS	the timecode \$CUETIME.
(→20)	
ID: B7	

 $Index \rightarrow page \ 2$ 

end of CUEPOINT

Table 20: The CueTrackPositions element, child of CuePoint ( $\rightarrow$ 19)	
Element	Description
<b>uint,</b> $\# \ge 1$ (!)	Track for which a position is given. This track number is the
CUETRACK	same as TrackEntry ( $\rightarrow$ Table 7)::TrackNumber.
ID: F7	
<b>uint,</b> $\# \ge 1$ (!)	The position of the cluster the referred block is found in.
CUECLUSTERPOSITION	This position is relative to the <b>Segment</b> 's ( $\rightarrow$ Table 2) data
ID: F1	section.
uint, $\# \leq 1$	The block with timecode \$CUETIME is the \$CUEBLOCK-
CUEBLOCKNUMBER	NUMBER-th BLOCK/SIMPLEBLOCK inside the CLUSTER at
ID: 53 78	position \$CUECLUSTERPOSITION.

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end of CUETRACKPOSITIONS

# 5.7 Chapters - Editions and ChapterAtoms

The CHAPTERS element contains a list of all editions and chapters found in this SEGMENT. Chapters in MATROSKA files are more powerful than chapters on DVDs, their handling is, however, way more complex.

Table 21: The CHAPTERS element, child of SEGMENT ( $ ightarrow$ 2)	
Element	Description
Master, $\# \geq 1$	One EDITIONENTRY describes one Edition. Just like with
EditionEntry ( $\rightarrow$ 22)	<b>TRACKENTRY</b> ( $\rightarrow$ Table 7), theoretically you could spread
ID: 45 B9	information about one Edition over different EDITIONEN-
	TRYS and use \$EDITIONUID to find out which edition the
	EDITIONENTRY is referring to, but it's highly discouraged.

Index  $\rightarrow$ page 2 end of Chapters

An edition contains one set of chapter definitions, so having several editions means having several sets of chapter definitions. This case is used when using this as a playlist - playing one chapter after the other while having gaps between the chapters.

Table 22: The EDITIONENTRY element, child of CHAPTERS $(\rightarrow 21)$	
Element	Description
uint, $\# \le 1$	<b>\$EDITIONUID</b> is the UID of the edition. This element is
EDITIONUID	mandatory if you want to apply one or more titles to an
ID: 45 BC	edition
bool, $\# \leq 1$	When \$EDITIONFLAGHIDDEN is 1, this edition should not
EDITIONFLAGHIDDEN	be available via the user interface
ID: 45 BD	
def: 0	
bool, $\# \leq 1$	When \$EDITIONFLAGDEFAULT is 1, this edition should be
EDITIONFLAGDEFAULT	selected by the player as default
ID: 45 DB	
def: 0	

**EDITIONENTRY** continued on next page

Element	Description
<b>bool,</b> $\# \le 1$	When <b>\$EDITIONFLAGORDERED</b> is 1, this edition contains a
EDITIONFLAGORDERED	playlist. When <b>\$EDITIONFLAGORDERED</b> is 0, it contains a
ID: 45 DD	simple DVD like chapter definition.
def: 0	
Master, $\# \geq 1$	One CHAPTERATOM contains the definition of one chapter.
ChapterAtom ( $\rightarrow$ 23)	This element is the only one in MATROSKA files that can con-
ID: B6	tain itself recursively - in this case to define subchapters.

Index  $\rightarrow$ page 2

end of **EditionEntry** 

## The following picture shows an ordered edition:

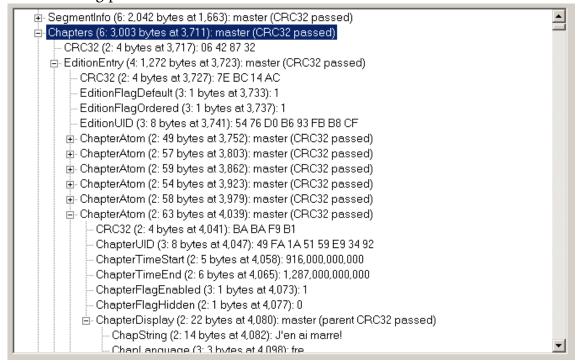


Table 23: The ChapterAtom element, child of EditionEntry ( $ ightarrow$ 22),	
child of ChapterAtom ( $ ightarrow 23$ )	
Element	Description
<b>uint,</b> # = 1	The UID of this chapter. It must be unique within the file.
CHAPTERUID	
ID: 73 C4	

**CHAPTERATOM** continued on next page

Element	Description
uint, $\# \le 1$	The unscaled timecode the chapter starts at. As the value
CHAPTERTIMESTART	is unsigned, a chapter cannot start earlier than at timecode
ID: 91	0, even whereas timecodes up to -30.000 are possible for
def: 0	multimedia data.
uint, $\# \le 1$	The unscaled timecode the chapter ends at. The default
CHAPTERTIMEEND	value is the start of the next chapter or the end of the parent
ID: 92	chapter or the end of the segment, whatever exists, in that
	order.
<b>bool,</b> $\# \leq 1$	When \$CHAPTERFLAGHIDDEN is 1, the chapter should not
CHAPTERFLAGHIDDEN	be visible in the user interface, but should be played back
ID: 98	normally.
def: 0	
<b>bool,</b> # ≤ 1	When \$CHAPTERFLAGENABLED is 0, the chapter should be
CHAPTERFLAGENABLED	skipped by the player
ID: 45 98	
def: 1	
<b>char[16]</b> , $\# \le 1$	This element can only occur if <b>\$EDITIONFLAGORDERED</b> =1.
CHAPTERSEGMENTUID	The SEGMENT of which the UID is \$CHAPTERSEGMENTUID
ID: 6E 67	is used instead of the current SEGMENT. Obviously, this
	SEGMENT should be easy to find, like when it is the first
	segment of a file in the same directory.
uint, $\# \leq 1$	The edition to use inside the SEGMENT selected via CHAP-
CHAPTERSEGMENT	TERSEGMENTUID. The timecodes \$CHAPTERTIMESTART
-EditionUID	and \$CHAPTERTIMEEND refer to playback timecodes of that
ID: 6E BC	edition, i.e. the timecodes are relative to that playlist. This
	is called "nested Editions" and is NOT SUPPORTED by Haali
	Media Splitter.
Master, $\# \leq 1$	Contains a list of tracks the chapter applies to.
CHAPTERTRACKS	
(→24)	
ID: 8F	
Master, $\# \geq 0$	Contains all chapter titles
CHAPTERDISPLAY	
(→25)	
ID: 80	

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end of CHAPTERATOM

A useful application for the **CHAPTERFLAGHIDDEN** element in connection with ordered editions is the following: You have a couple of episodes of a series, but want to save space by only saving the intro and outtro once. You create one playlist (ordered edition) per episode, and another playlist playing all episodes in a row. Whereas in the first case you might want to play intro and outtro for each episode, you might not want to do that in the second case.

If you don't want to make the three parts intro - movie - outtro selectable via the user interface when playing single episodes, you call the intro-chapter "Episode - blah" and hide the movie- and the outtro chapter using CHAPTERFLAGHIDDEN=1. Then, the playlist playing all episodes would be *intro - episode 1 - episode 2 - ... - last episode - outtro*, whereas the other playlists would be *intro - episode* N - *outtro*. The name of the intro chapter would be set to "Episode n".

Table 24: The ChapterTracks element, child of ChapterAtom ( $ ightarrow$ 23)	
Element	Description
uint, $\# \geq 1$	One number of a track a chapter is used with.
CHAPTERTRACKNUMBER	
ID: 89	

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end of CHAPTERTRACKS

Table 25: The ChapterDisplay element, child of ChapterAtom ( $ ightarrow$ 23)	
Element	Description
<b>utf-8,</b> # ≤ 1	A title of a chapter
CHAPSTRING	
ID: 85	
string, $\# \ge 0$	The language of <b>\$CHAPSTRING</b> as defined in ISO639-2 <sup>5</sup>
CHAPLANGUAGE	
ID: 43 7C	
def: eng	
<b>utf-8,</b> # ≥ 0	A country the title is used in. For example, a german title in
CHAPCOUNTRY	Germany might be different than the title used in Austria.
ID: 43 7E	

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end of ChapterDisplay

<sup>&</sup>lt;sup>5</sup>http://lcweb.loc.gov/standards/iso639-2/englangn.html#two

# 5.8 Attachments

Theoretically, any file type can be attached to a MATROSKA file, however, this possibility is usually used to attach pictures like CD covers or fonts required to display a subtitle track correctly. Obviously, attaching executable files would allow for MATROSKA files to contain viruses - a scenario that is not exactly the indended application of attachments or anything else MATROSKA is capable of.

Table 26: The Attachments element, child of Segment ( $ ightarrow$ 2)	
Element	Description
Master, $\# \geq 1$	Describes and contains one attached file
AttachedFile ( $\rightarrow$ 27)	
ID: 61 A7	

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end of ATTACHMENTS

Table 27: The AttachedFile element, child of Attachments ( $ ightarrow$ 26)	
Element	Description
<b>utf8,</b> # ≤ 1	A human-readable description of the file
FILEDESCRIPTION	
ID: 46 7E	
<b>utf8,</b> # ≤ 1	The name that should be proposed by a demuxer when ex-
FILENAME	tracting the file
ID: 46 6E	
string, $\# \le 1$	MIME type of the file, like
FILEMIMETYPE	
ID: 46 60	
binary, $\# \le 1$	The file itself
FILEDATA	
ID: 46 5C	
<b>uint,</b> # = 1	The UID of that file, just like TRACKUID, CHAPTERUID etc.
FILEUID	The UID is required if a <b>TRACKENTRY</b> ( $\rightarrow$ Table 7) wants to
ID: 46 AE	refer to this Attachment.

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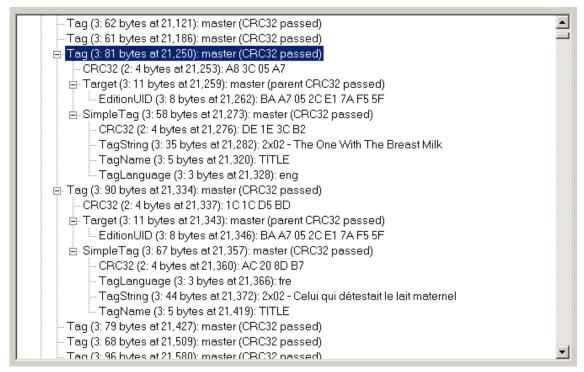
end of ATTACHEDFILE

# **5.9** Tags

Table 28: The TAGS element, child of SEGMENT $(\rightarrow 2)$	
Element	Description
Master, $\# \geq 1$	One TAG element describes one Tag
<b>T</b> AG (→29)	
ID: 73 73	

Index  $\rightarrow$ page 2 end of TAGS

TAGS provide additional information<sup>6</sup> not important for replay. A TAGS element contains a number of TAG elements. Each TAG element contains a list of UIDs (usually TRACKUIDs or EDITIONUIDs), and a list of SIMPLETAGS, each one containing a name and a value:



If no TARGETS are specified, then the TAG is a global TAG refering to the entire SEGMENT. Of course, two different TAG elements can contain identical TARGETS.

<sup>6</sup>http://www.matroska.org/technical/specs/tagging/index.html

Table 29: The TAG element, child of TAGS ( $\rightarrow$ 28)							
Element Description							
Master, $\# \leq 1$	Describes which elements a Tag applies to						
Targets $(\rightarrow 30)$							
ID: 63 CO							
Master, $\# \geq 1$	Each SIMPLETAG contains one tag that applies to each target						
SimpleTag ( $\rightarrow$ 31)	in TARGETS						
ID: 67 C8							

Index →page 2 end of TAG

Note that there is nothing like a TAGUID.

Table 30: The TARGETS element, child of TAG ( $ ightarrow$ 29)							
Element	Description						
uint, $\# \le 1$	This number describes the logical level of the object the Tag						
TARGETTYPEVALUE	refers to						
(→??)							
ID: 68 CA							
def: 50							
<b>utf-8,</b> # ≤ 1	A string describing the logical level of the object the Tag is						
TARGETTYPE	refering to						
ID: 63 CA							
uint, $\# \ge 0$	The UID of a track the tag is referring to						
TRACKUID							
ID: 63 C5							
uint, $\# \ge 0$	The UID of an edition the tag is referring to. Note that this						
EDITIONUID	is the only way to apply titles to an edition						
ID: 63 C9							
uint, $\# \ge 0$	The UID of a chapter the tag is referring to						
CHAPTERUID							
ID: 63 C4							
uint, $\# \ge 0$	The UID of an attachment the tag is referring to						
ATTACHMENTUID							
ID: 63 C6							

 $Index \rightarrow page \ 2 \qquad \qquad \textit{end of Targets}$ 

Table 31: The SIMPLETAG element, child of TAG ( $ ightarrow$ 29)							
Element	Description						
utf-8, $\# \ge 1$ (!)	Name of the tag.						
TAGNAME							
ID: 45 A3							
string, $\# \le 1$	\$TAGLANGUAGE is the language of \$TAGNAME. Note that						
TAGLANGUAGE	the default here is 'und', whereas the default track / chapter						
ID: 44 7A	title language is 'eng'.						
def: und							
bool, $\# \leq 1$	When 1, this title and language is the original title given to						
TAGORIGINAL	the item						
ID: 44 84							
def: 1							
<b>utf-8,</b> # ≤ 1	The value of the tag when it is a string						
TAGSTRING							
ID: 44 87							
binary, $\# \leq 1$	The 'value' of the tag when it's a binary tag						
TAGBINARY							
ID: 44 85							

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end of **SIMPLETAG** 

## 5.9.1 A few common Tags

- **TITLE**, Target: EditionUID: used to define names for Editions. This is exactly what you can see in the screenshot above.
- BPS, Target: TrackUID: used to define the bitrate of a track
- FPS, Target: TrackUID: used to define the framerate of a track

# 6 MATROSKA block Layout and Lacing

## 6.1 Basic layout of a Block

is slow

```
A MATROSKA block has the following format:
BLOCK {
  vint TrackNumber
  sint16 Timecode // relative to Cluster timecode
  int8 Flags // lacing, keyframe, discardable
  if (lacing) {
    int8 frame_count-1
    if (lacing == EBML lacing) {
        vint size[0]
        svint size [1..frame count-2]
    } else
    if (lacing == Xiph lacing) {
      int8 size[size of <leading (frame_count-1) frames> / 255 + 1]
    }
  int8[] data
}
The following bits are defined for FLAGS:
 Bit 0x80: keyframe:
             No frame after this frame can reference any frame before
             this frame and vice versa (in AVC-words: this frame is an
             IDR frame). The frame itself doesn't reference any other
             frames.
 Bits 0x06: lace type
             00 - no lacing
             01 - Xiph lacing
             11 - EBML lacing
             10 - fixed-size lacing
 Bit 0x08 : invisible: duration of this block is 0
 Bit 0x01 : discardable: this frame can be discarded if the decoder
```

The following flags are only defined for Matroska v2 and can thus only be used in a **SIMPLEBLOCK**: *keyframe*, *invisible*, *discardable*. The type of lacing in use defines how the **SIZE** values are to be read.

## 6.2 Lacing

Lacing is a technique that allows to store more than one atom of data (like one audio frame) in one block, with the goal to decrease overhead, without losing the ability to separate the frames in a lace later again.

Generally, the size of the last frame in a Lace is not stored, as it can be derived from the total block size, the size of the block header and the sum of the sizes of all other frames.

Frame duration values are not preserved! That means, it is highly recommended **not** to use lacing if the frame duration is not constant, like Vorbis audio.

#### 6.2.1 Xiph Lacing

The size of each frame is coded as a sum of int8. A value smaller than 255 indicates that the next value refers to the next frame.

#### Example

size = { 187, 255, 255, 120, 255, 0, 60 } means that there are 4 frames with 187, 630, 255, 60 bytes.

### 6.2.2 EBML Lacing

```
Size of first frame ("frame 0") of a lace = size[0]
Size of frame i of a lace: size[i] - size[i-1]
```

### 6.2.3 Fixed Lacing

Fixed Lacing is used if all frames in a lace have the same size. Examples are AC3 or DTS audio. In this case, knowing the number of frames is enough to calculate the size of one frame. Consequently, there are no size values.

## 7 Overhead of Matroska files

The scope of this section is explaining how to predict the overhead of a MATROSKA file before muxing, and without analysing any of the source files excessively. This section assumes that **BLOCKGROUPS** and **BLOCKS** are used, and that no **SIMPLE-BLOCKS** are used. If you want to estimate overhead of files that use **SIMPLE-BLOCKS**, you get about the same overhead as with **BLOCKS** without **BLOCKDURA-TION**, **REFERENCEBLOCK** or **BLOCKGROUP**.

### 7.1 Overhead of BLOCKGROUPS

First, here again the layout of a typical BLOCKGROUP

```
BlockGroup <size>
  Block <size> <number, flag, timecode>
  [ Reference <size> <val> ]
```

The EBML identication for **BLOCKS** and **BLOCKGROUPS** are 1 byte each, so that the structure above, not counting **REFERENCES**, takes:

• BlockGroup < 128 bytes: **8 bytes** 

• BlockGroup < 16kbytes: 10 bytes

• BlockGroup < 2MBytes: **12 bytes** 

BLOCKGROUPS larger than 2MBytes are extremely unlike, and even BLOCKGROUPS larger than 16kBytes won't occur often, compared to BLOCKGROUPS between 128 bytes and 16 kBytes. That means, assuming an overhead of 10 bytes for BLOCKGROUPS without REFERENCES usually results in a good approximation.

#### 7.1.1 video

In a typical video stream, there are a lot of frames with 1 REFERENCE (P-Frames, Delta-Frames), and a few keyframes. Typical rations are 100:1. There might also be frames with 2 REFERENCES (B-Frames), e.g. native MPEG4 streams. Assuming a ratio of 66:33:1 for B:P:K, and assuming a bitrate far below 3,2 MBit/s (meaning that typical B- and P-frames are smaller than 16 kB), that causes about 15 bytes of

overhead per frame. If there are no B-Frames, there are about 13 bytes per frame.

Example: 2 hours, 25 fps.

The video stream will cause around 2,3 MB of overhead.

### 7.1.2 audio - without lacing

As audio does usually not have any **REFERENCES** (all audio frames are keyframes), one audio frame will take 8 or 10 bytes of overhead. For MP3, AC3, DTS and AAC, frames causing 8 bytes of overhead are unlikely. They are more likely for Vorbis.

Example: MP3 audio, 24ms per frame, duration: 2h

This stream will cause 3MB of overhead.

## 7.1.3 audio - with lacing

#### 1. CBR+CFR: fixed lacing

In this case, *fixed lacing* (see section 6.2.3) is used. With fixed lacing, the overhead is the normal **BLOCKGROUP** overhead, plus 1 byte for the lace header. Assuming that **BLOCKGROUPS** are not larger than 16k, that means that the overhead per frame is equal to 11 / frame\_count

Example: AC3 audio, 448 kbps, 1792 bytes per frame, 32ms per frame

1.) 8 frames per lace.

overhead for one frame = 11/8 = 1,375 bytes = 1 byte / 23,3 ms.

2.) 9 frames per lace.

overhead for one frame = 11/9 = 1,222 bytes = 1 byte / 26,2 ms.

3.) 10 frames per lace.

overhead for one frame = 13/10 = 1,3 bytes = 1 byte / 24,6 ms.

An AC3 stream of 2 hours with 9 frames per lace will cause 270kB of overhead.

#### 2. no CBR, but almost all frames smaller than 255 bytes: XIPH lacing

In this case, XIPH lacing (see section 6.2.1) is used, meaning that the overhead of a **BLOCKGROUP** is equal to normal BlockGroup overhead + frame\_count, meaning that the overhead per frame is about (11+frame\_count)/frame\_count, if there are frame\_count frames in each lace. Again, if the **BLOCKGROUP**s are larger than 16kBytes, then the overhead is (13+frame\_count)/frame\_count.

In other words, the ratio in bytes / frame will always be between about 1,2 and

2,5 for audio streams with mainly small frames.

Although XIPH lacing is also defined for larger frames, EBML lacing is usually more effective then.

**3. otherwise: EBML lacing** Assuming that the difference in size between 2 consecutive frames is smaller than 8191, 1 or 2 bytes are needed to code the size of each frame, additionally to the normal **BLOCKGROUP** overhead.

As a result, we get 3 possible estimations:

a) worst case That means, a lace with frame\_count frames using EBML lacing will cause not more than ((11 or 13)+2\*frame\_count)/frame\_count bytes of overhead per frame.

```
Example 1: 16 frames per lace, BLOCKGROUP > 16kB, worst case: overhead <= (13 + 2*16)/16 = 2,8 bytes / frame.

Example 2: 8 frames per lace, BLOCKGROUP < 16kB, worst case: overhead <= (11 + 2*8)/8 = 3,4 bytes / frame.
```

**b) best case** The best case is obviously that 2 consecutive frames differ by not more than 62 bytes. In that case, one byte is needed to code the size of one frame. However, the first frame might need to bytes, if it is larger than 126 bytes.

```
Example 1: 16 frames per lace, BLOCKGROUP > 16kB, best case: overhead <= (13 + 1*16)/16 = 1,8 bytes / frame.

Example 2: 8 frames per lace, BLOCKGROUP < 16kB, best case: overhead <= (11 + 1*8)/8 = 2,4 bytes / frame.
```

**c) average case** This is the case you need for optimal overhead prediction. Unfortunately, the average case depends on the compression format of the corresponding audio track, its bitrate, maybe even the encoder that has been used. The easiest way to gather data on the average case of EBML lace header overhead is to simulate the lace results of different files that are likely to be used. Candidates are MPEG 1/2/4 audio and Vorbis, but not AC3 or DTS.

I have run a simulation with the following file types:

MPEG 1 Layer 3 (128 and 192 kbps, 48 kHz), HE-AAC (224 kbps and 96 kbps, 44,1 kHz), LC-AAC (268 kbps, 44,1 kHz)

The results obtained from those files are discussed on the following pages. The lace behaviour simulation has been run using mls<sup>7</sup> ( short for 'matroska lace sim-

<sup>7</sup>http://www-user.tu-chemnitz.de/~noe/Video-Zeug/mls/

ulator'). Note that it would be required to run the simulation and to evaluate the results as follows for each audio format, in each bitrate, maybe even with each encoder, for which results as accurate as possible shall be predicted.

The results for the lace header size are as follows:

		Lace	header	overhe	ead per	frame	@ <x< th=""><th>&gt; Fran</th><th>nes per</th><th>lace</th></x<>	> Fran	nes per	lace
Audio Format		4	8	12	16	24	32	48	64	96
MP3	@ 128 kbps	1,39	1,29	1,26	1,24	1,22	1,22	1,21	1,20	1,20
MP3	@ 192 kbps	1,50	1,41	1,38	1,37	1,36	1,35	1,34	1,34	1,33
HE-AAC	C @ 224 kbps	1,39	1,29	1,25	1,24	1,22	1,21	1,20	1,20	1,20
HE-AAC	C @ 64 kbps	1,34	1,23	1,19	1,18	1,16	1,15	1,14	1,14	1,13
LC-AAC	268 kbps	1,31	1,19	1,16	1,14	1,12	1,11	1,10	1,09	1,09

Applications using libmatroska for MATROSKA file creation are using 8 frames per lace. As a consequence, the overhead for a track using EBML lacing can be predicted to an acceptable accuracy if the audio format is known.

As you can also see, larger laces hardly affect the overhead caused by the lace headers of **Blocks** from a certain size on.

However, larger laces mean fewer **Blocks** and thus fewer **BlockGroups**, so the total overhead per frame, including the overhead caused by overhead outside of the **Blocks**, is worth a look. Here are the results with the same test files as above

		Overhead per frame @ <x> Frames per lace</x>								
Audio Format		4	8	12	16	24	32	48	64	96
MP3	@ 128 kbps	4,14	2,67	2,17	1,93	1,68	1,56	1,48	1,41	1,33
MP3	@ 192 kbps	4,25	2,79	2,30	2,06	1,81	1,75	1,61	1,54	1,47
HE-AAC	2 @ 224 kbps	4,14	2,66	2,23	2,05	1,76	1,62	1,48	1,40	1,33
HE-AAC	C @ 64 kbps	4,09	2,61	2,11	1,86	1,62	1,49	1,40	1,34	1,27
LC-AAC	@ 268 kbps	4,06	2,57	2,07	1,82	1,66	1,51	1,37	1,30	1,22

Now lets take the 2nd table and find out how much overhead that means in a real movie of 2 hours.

In the case of the mp3 files used in that example, one frame lasts 24ms. In the case of our LC-AAC file, one frame lasts 23,22 ms, and for the HE-AAC file we get 46,44ms.

Thus a file of 2 hours will have the following number of frames:

MP3 - 300,000

LC-AAC - 310,000

HE-AAC - 155,000.

First, lets use the default setting of libmatroska (8 frames per lace) and calculate the overhead a muxing app using libmatroska would cause when muxing those files into a movie:

- MP3 @ 128: overhead = 300,000 \* 2,67 = 801,000 bytes
- MP3 @ 192: overhead = 300,000 \* 2,79 = 837,000 bytes
- **HE-AAC** @ **224**: overhead = 155,000 \* 2,66 = 412,300 bytes
- **LC-AAC** @ **268**: overhead = 310,000 \* 2,57 = 796,700 bytes

With 24 frames per lace, an MP3 block would have a duration of 576ms, an HE-AAC block even about 1 second. That means, when seeking in a file, an awkward impression of the audio being missing for a moment could occur. Thus, larger laces than 1 second are highly discouraged. Nevertheless, let's analyze the overhead in our file for laces of 24 and 96 frames each, and compare the overhead to the one caused by libmatroska. Here is the corresponding table:

			Frames per lace				
Audio Format			8	24	96		
MP3	@	128 kbps	782kB	492kB	389kB		
MP3	@	192 kbps	1	530kB			
HE-AAC	@	224 kbps	402kB	266kB	201kB		
HE-AAC	@	64 kbps	395kB	245kB	192kB		
LC-AAC	@	268 kbps	778kB	502kB	369kB		

As you can see, putting 24 frames in one block, compared to 8 frames, saves some overhead. However, putting 96 frames in one **Block** instead of 24 saves less overhead than 24 compared to 8. As 96 frames per lace will usually cause uncomfortable seeking, it is recommended not to put more than about 24 frames in one **Block**.

### 7.2 Overhead of Clusters

Although most of the overhead is caused by **BlockGroups**, the amount of overhead caused by **Clusters** themselves is noticeable as well.

Here again the basic layout of a **CLUSTER**:

```
Cluster <size>
  [ CRC32 ]
  TimeCode <size> <timecode>
  [ PrevClusterSize <size> <prevsize> ]
  [ Position <size> <position> ]
  { BlockGroup }
```

First, some conventions:

- each Cluster has a size between 16kB and 2MB
- each Cluster may begin between 16MB and 4GB

As typical movie files are designed to fit on 1 or 2 CDs, or 2 or 3 of them fill one DVD, point 2 will be true for most of the clusters in typical files.

With the abovementioned restrictions on **CLUSTERS**, the overhead inside one Cluster will be:

• Cluster ID + <size>: 7 bytes

• **CRC32**: 6 bytes

• TIMECODE: 5 bytes

• PrevClusterSize: 5 bytes

• **Position**: 5 bytes

• **Seekhead** entry for **Cluster**: 17 bytes

Depending on the muxing settings, the overhead caused by one **CLUSTER** will be between 12 and 45 bytes.

**Example**: Assuming a size of 1 MB per **CLUSTER**, that means an overhead rate of 0,001% - 0,005%, or up to 100 kB in a file of 2GB.

## 7.3 Overhead caused by Cues

Here again the layout of a **CUEPOINT**:

Assuming that a CUEPOINT only points into one certain track, the overhead is:

• CuePoint: 2 bytes

• CueTime: 5 bytes

• CueTrackPosition: 2 bytes

• CueClusterPosition: 6 bytes

• CueTrack: 3 bytes

• CueBlockNumber: 4 bytes

Total: 22 bytes.

**Example:** Assuming that there is a **CUEPOINT** each 4 seconds (1 keyframe in 100 frames), this adds on overhead of 0,22 bytes / frame

There can also be **CUEPOINTs** for audio tracks. In that case, as every frame will be a keyframe, the number of **CUEPOINTs** only depends on the muxing application. Predicting the overhead requires to know its behaviour.

# 8 Links

## Matroska pages / software:

```
http://www.matroska.org
http://haali.cs.msu.ru/mkv/
http://www.alexander-noe.com/
http://de.wikipedia.org/wiki/Matroska
http://www.matroska.info/
http://ld-anime.faireal.net/guide/jargon.matroska-en
```