

Message-passing concurrency





Erlang rationale?

Designed by Ericsson for telecoms systems.

These must be ...

... concurrent,

... high-availability, and

... fault-tolerant.





Erlang rationale?

Used today by WhatsApp!, Cisco, Bet365, ...

Because it is ...

... concurrent,

... high-availability, and

... fault-tolerant.



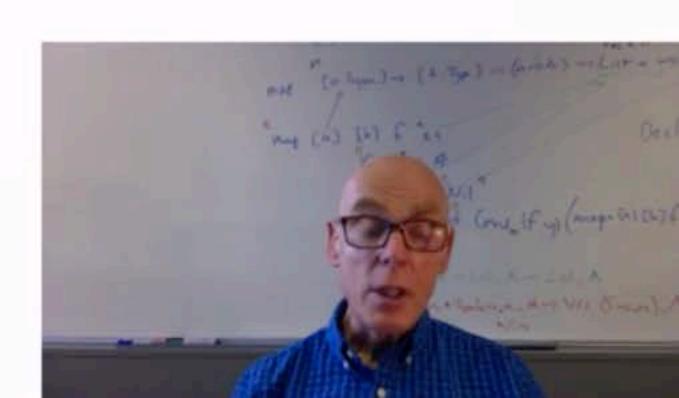


Based on functional programming

Functional is simpler: values and not references, pointers etc.

Functional supports high-level patterns e.g. map/reduce.

With no side-effects ("immutable data") algorithms clearer.



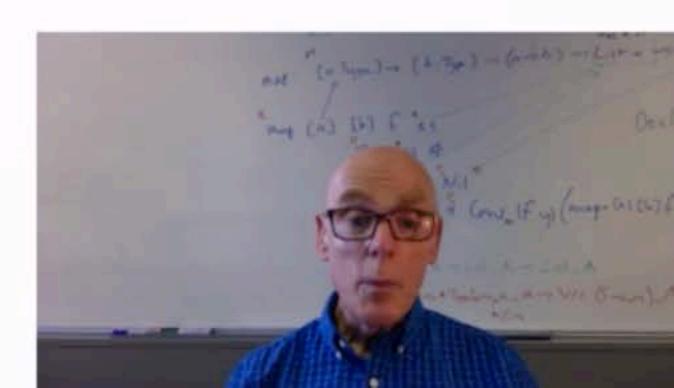


Immutability as a design pattern

Structures whose state doesn't change ...

... if you want a different structure, create a new one.

Thread-safe programming ... safe caching and sharing ... consistency



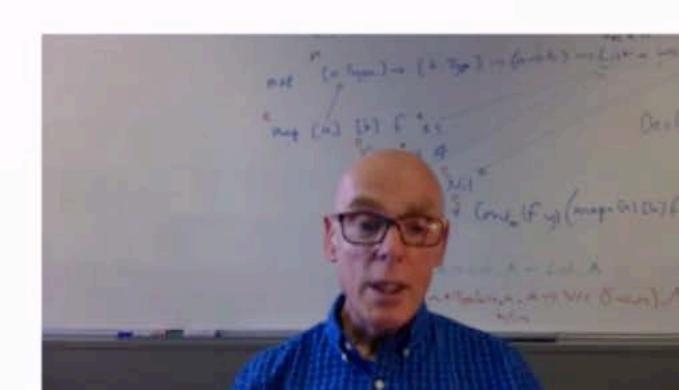


A pragmatic approach to side-effects

At its core, Erlang is functional ...

...it does allow some side-effects (e.g. communication) ...

... but not others (no updatable Java-style variables).





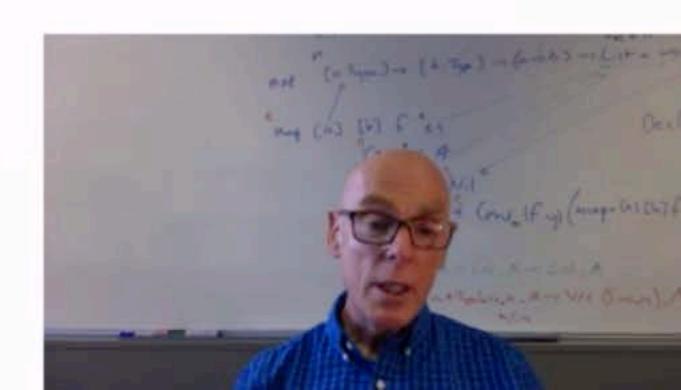
Concurrency

Modularity: different "processes" in different processes.

Robustness: "let it fail" and others processes sort it out.

Concurrency for design (independence) ...

... versus parallelism for efficiency (happen at same time).



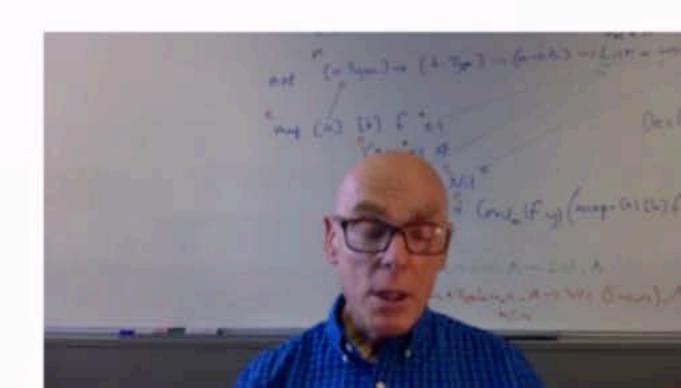


Message-passing concurrency

Different concurrent processes "share nothing" ...

...and only communicate by passing messages.

No worries about shared state, thread safety, ...





A beautiful design

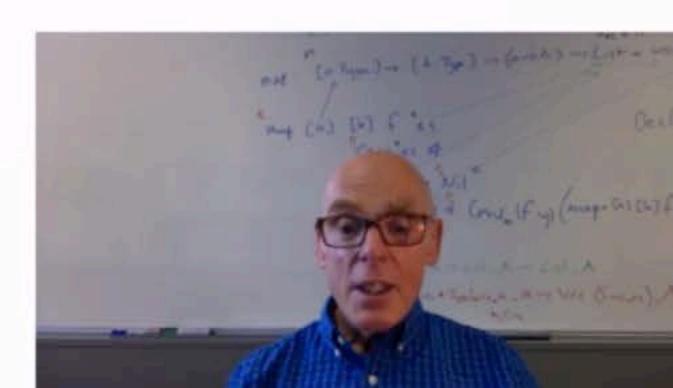
Mailboxes and message handling

Process errors and trapping exits

OTP generics

Concurrency and distribution

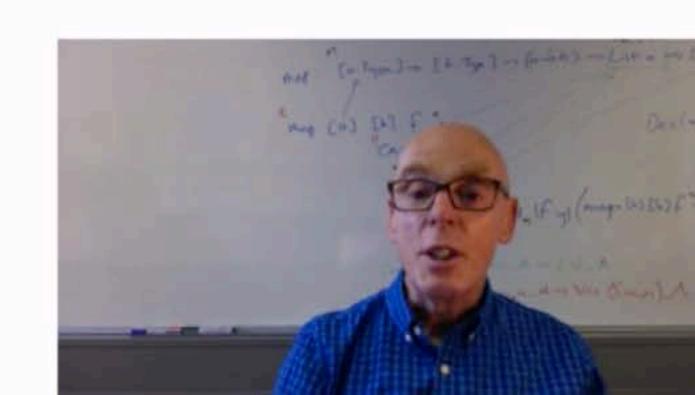




University of Kernt



Processes and messages



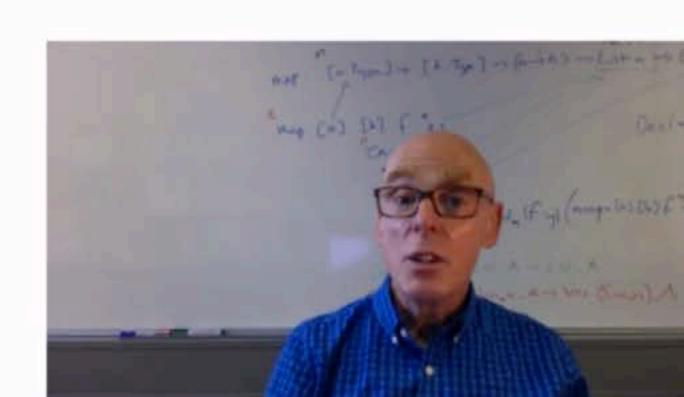


Message-passing concurrency

Different concurrent processes "share nothing" ...

...and only communicate by passing messages.

No worries about shared state, thread safety, ...





The "actor model"

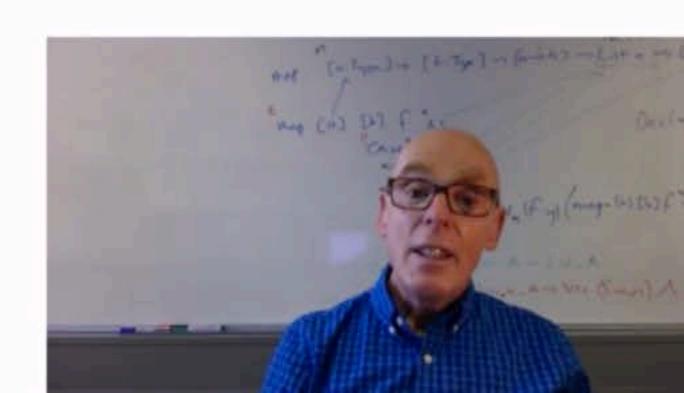
Carl Hewitt's model for the foundations of computation ...

... everything is an actor

... actors communicate by messages

... actors can create other actors

... and change behaviour according to messages.





The design of Erlang's concurrency

Erlang designed to build solutions in a particular problem space.

Choices to support that ... the roads not chosen.





Processes versus threads

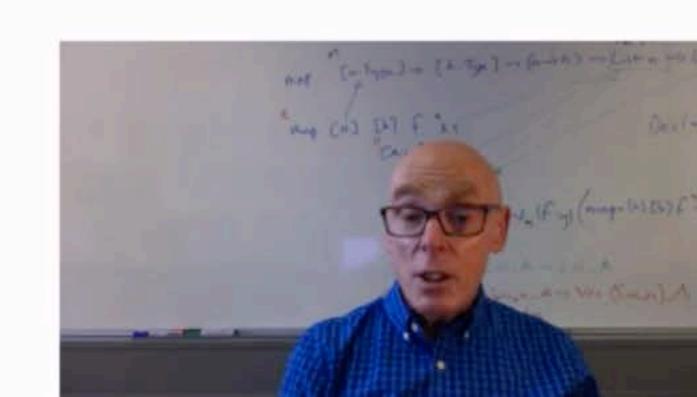
Processes share nothing with each other.

In the virtual machine, not the operating system ... like "green threads" but ...

...more lightweight than threads

...more numerous

... more control: stack, heap and garbage collection per process.





Synchronous versus asynchronous

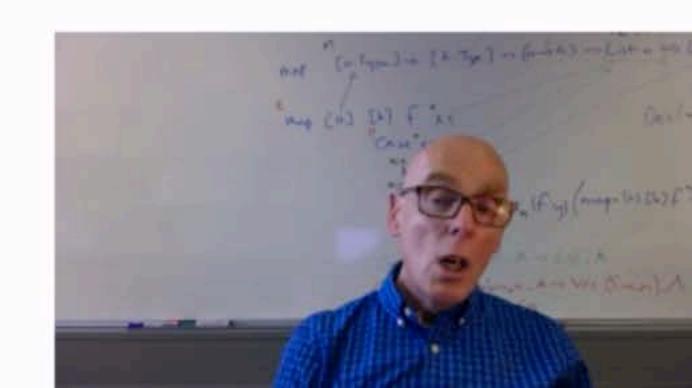
Message send and message receive are decoupled.

... avoids a common source of deadlock

...directly supports processes distributed across a network

...supports more complex message handling

...can simulate synchrony through send/receive protocol





Identifiers versus channels

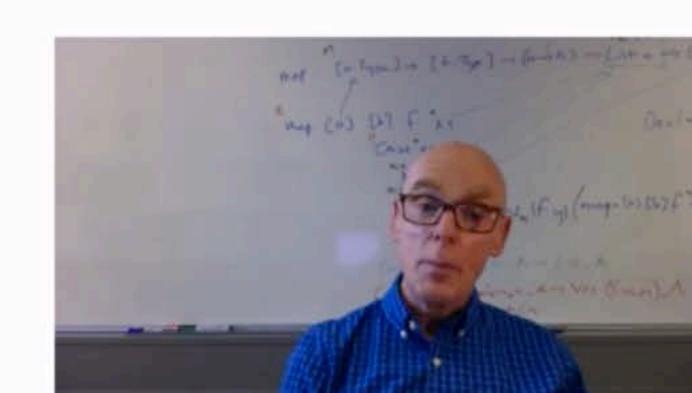
Messages are sent to a particular process id (Pid) ...

... rather than along a channel between processes.

The communication structure is less evident in the program ...

... but supports dynamic systems more readily.

Named processes give more visible long-lived structures.



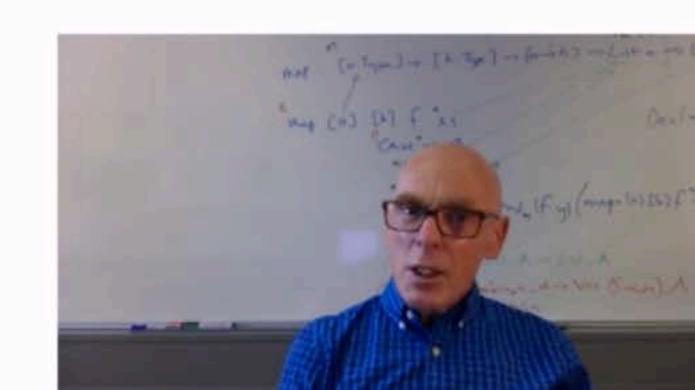


Concurrency versus parallelism

Concurrency means the possibility of running independently, perhaps at the same time.

On a single processing element, concurrent processes timeshare, mediated by a scheduler.

On a multicore processor, there is – potentially at least – concurrent activity on each core, with each scheduled separately.





A beautiful design

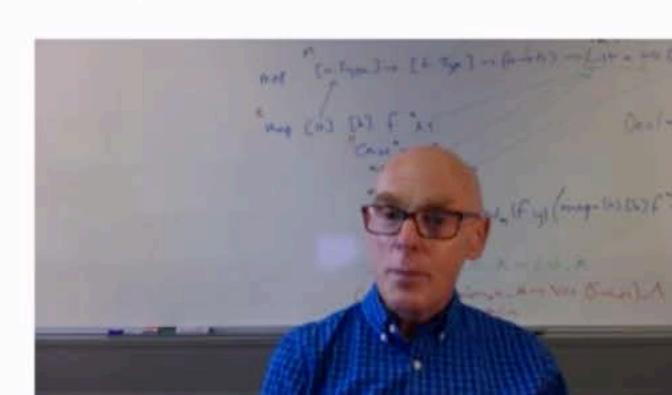
Mailboxes and message handling

Process errors and trapping exits

OTP generics

Concurrency and distribution





University of Kernt



Processes and messages in practice





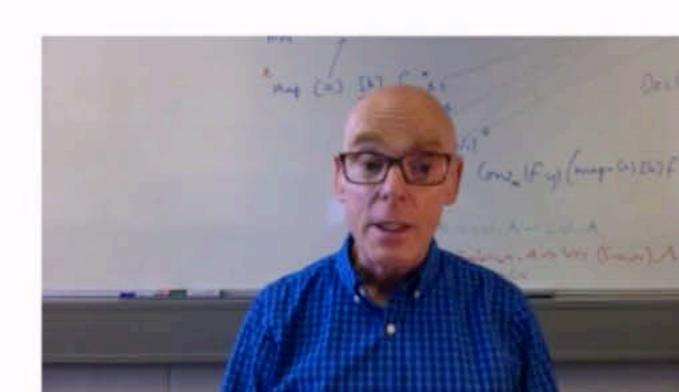
Erlang concurrency in a nutshell

spawn – create a process

send a message

self() – give the Process Identifier (Pid) of a process

receive - handle a message





What is a process?

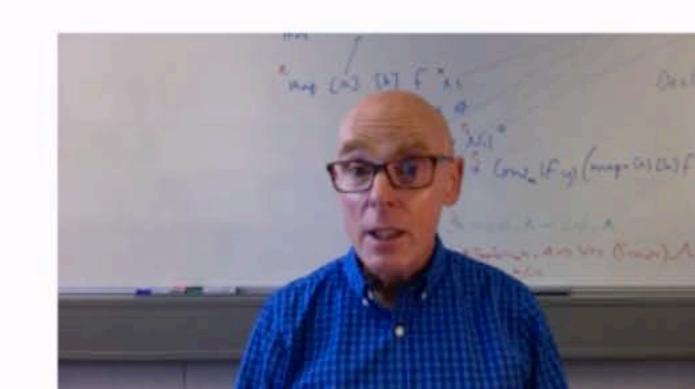
A process is a separate computation

... running in its own space

... time-sharing with other processes

A process runs an Erlang function

... and terminates when / if that function terminates





spawn - create a process

spawn (Mod, Fun, Args)

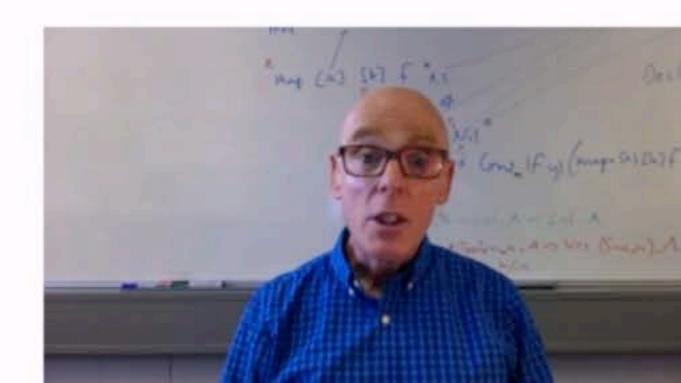
Mod is a module,

Fun is a function in Mod,

Args is a list of arguments to
pass to Fun at startup.

For example:

spawn(foo,bar,[])





spawn returns a process id

```
A call like

spawn (Mod, Fun, Args)

returns the process id (Pid)

of the process created.
```

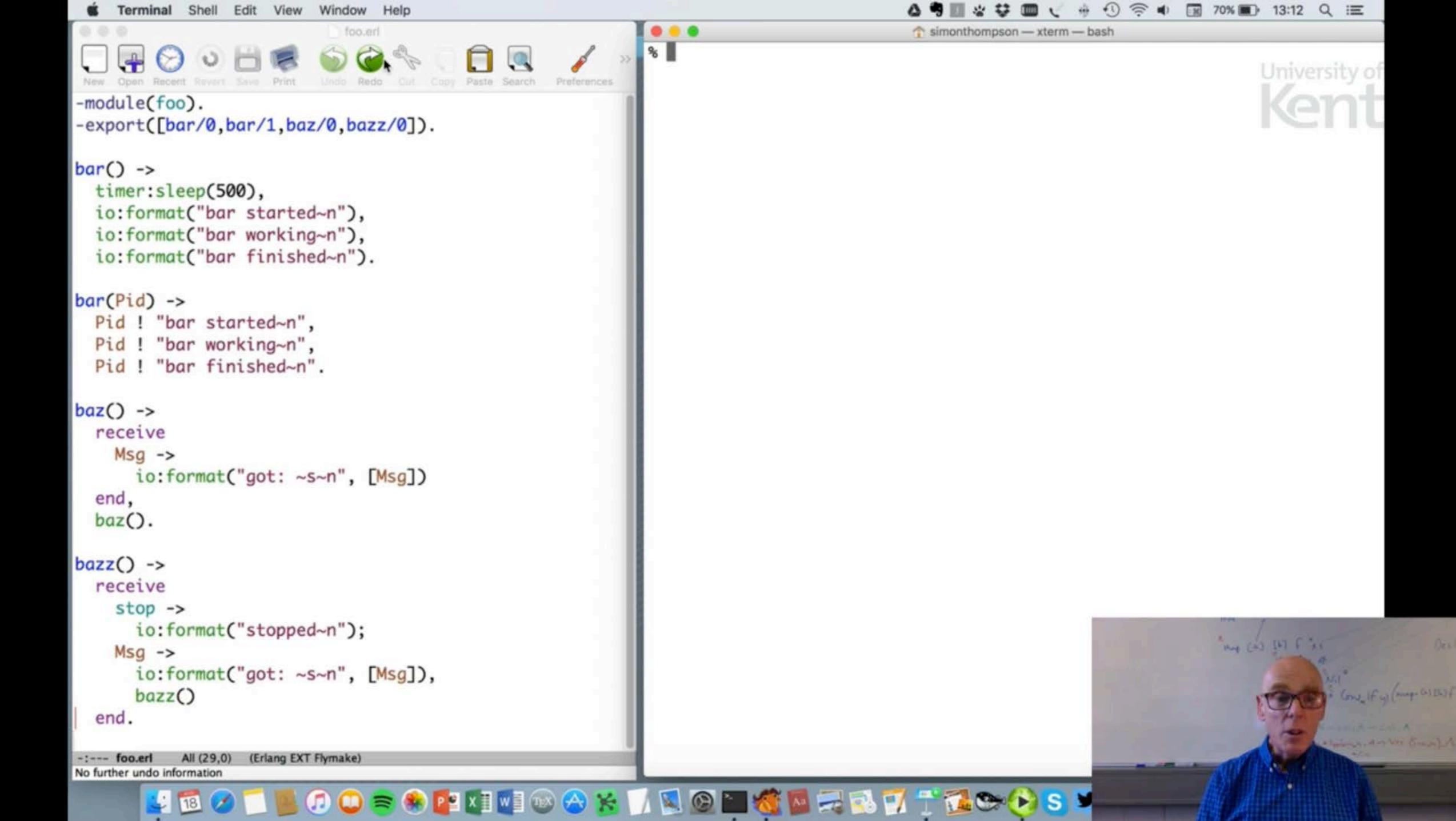
It is typically captured like this:

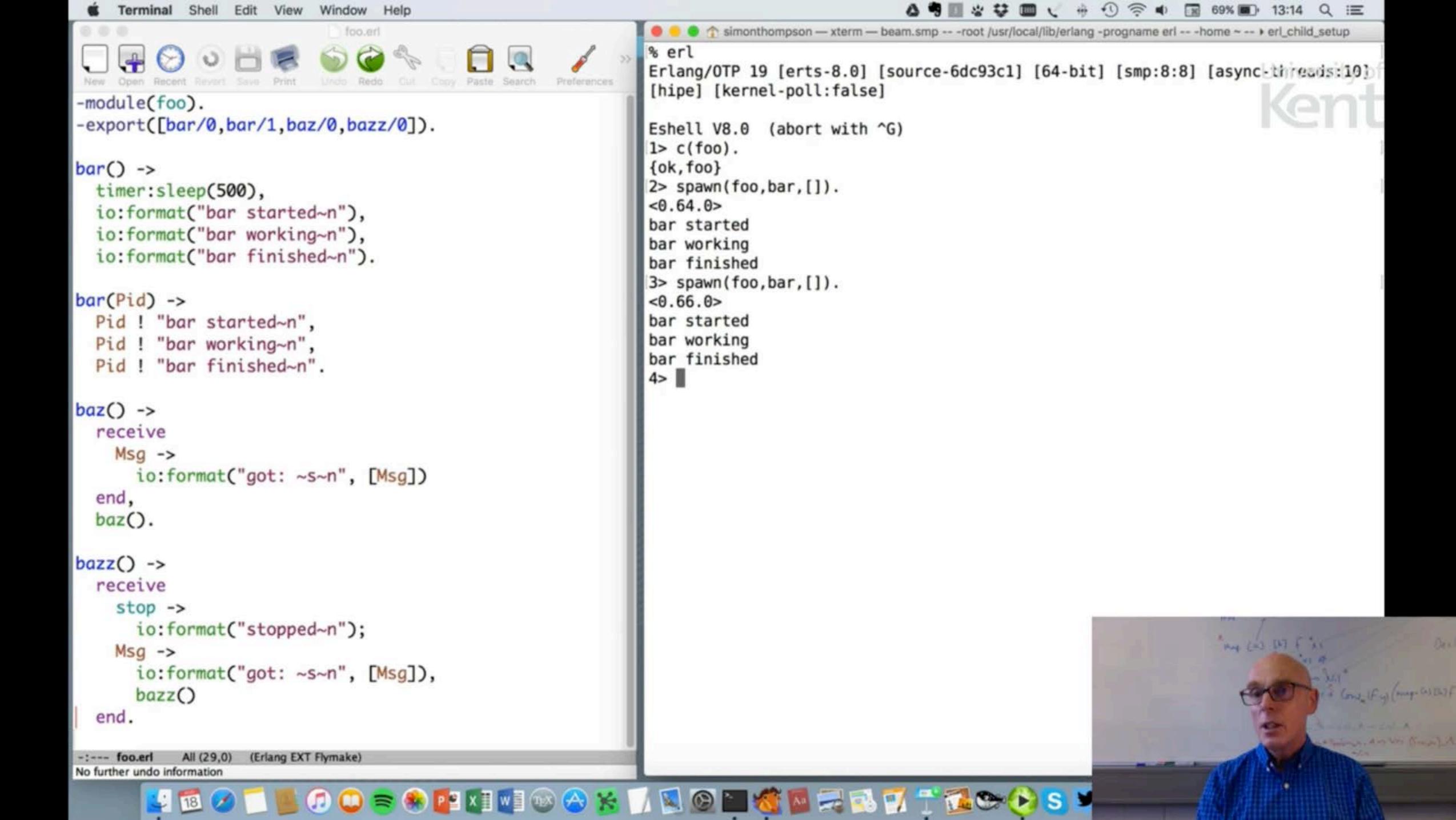
```
Proc = spawn(foo,bar,[])
```

```
-module(foo).
-export([bar/0]).

bar() ->
  timer:wait(500),
  io:format("bar started~n"),
  io:format("bar working~n"),
  io:format("bar finished~n").
```









! - send a message

Pid! Msg

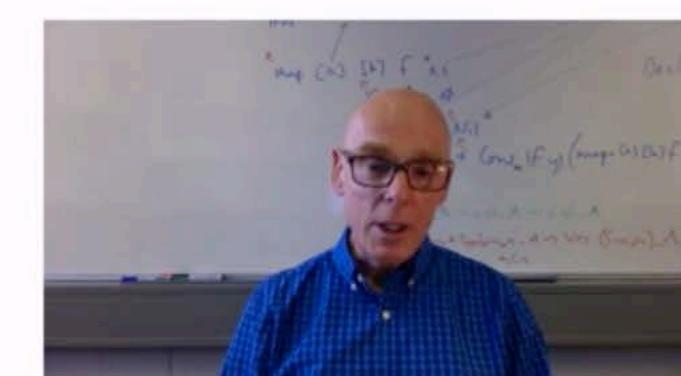
sends the message Msg to the process with process id Pid ...

We say "sends Msg to Pid".

The effect is to put Msg into the mailbox of process Pid.

```
-module(foo).
-export([bar/1]).

bar(Pid) ->
  Pid ! "bar started~n",
  Pid ! "bar working~n",
  Pid ! "bar finished~n".
```

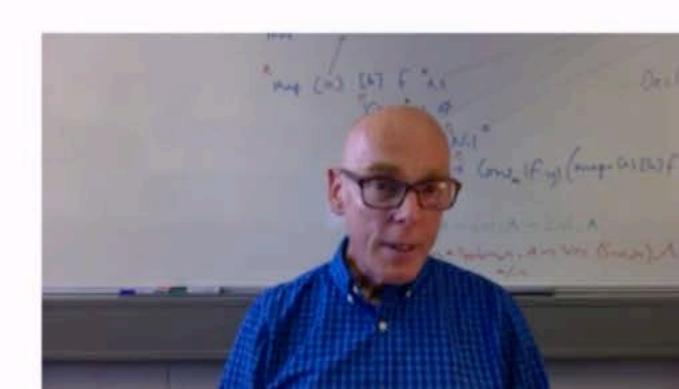


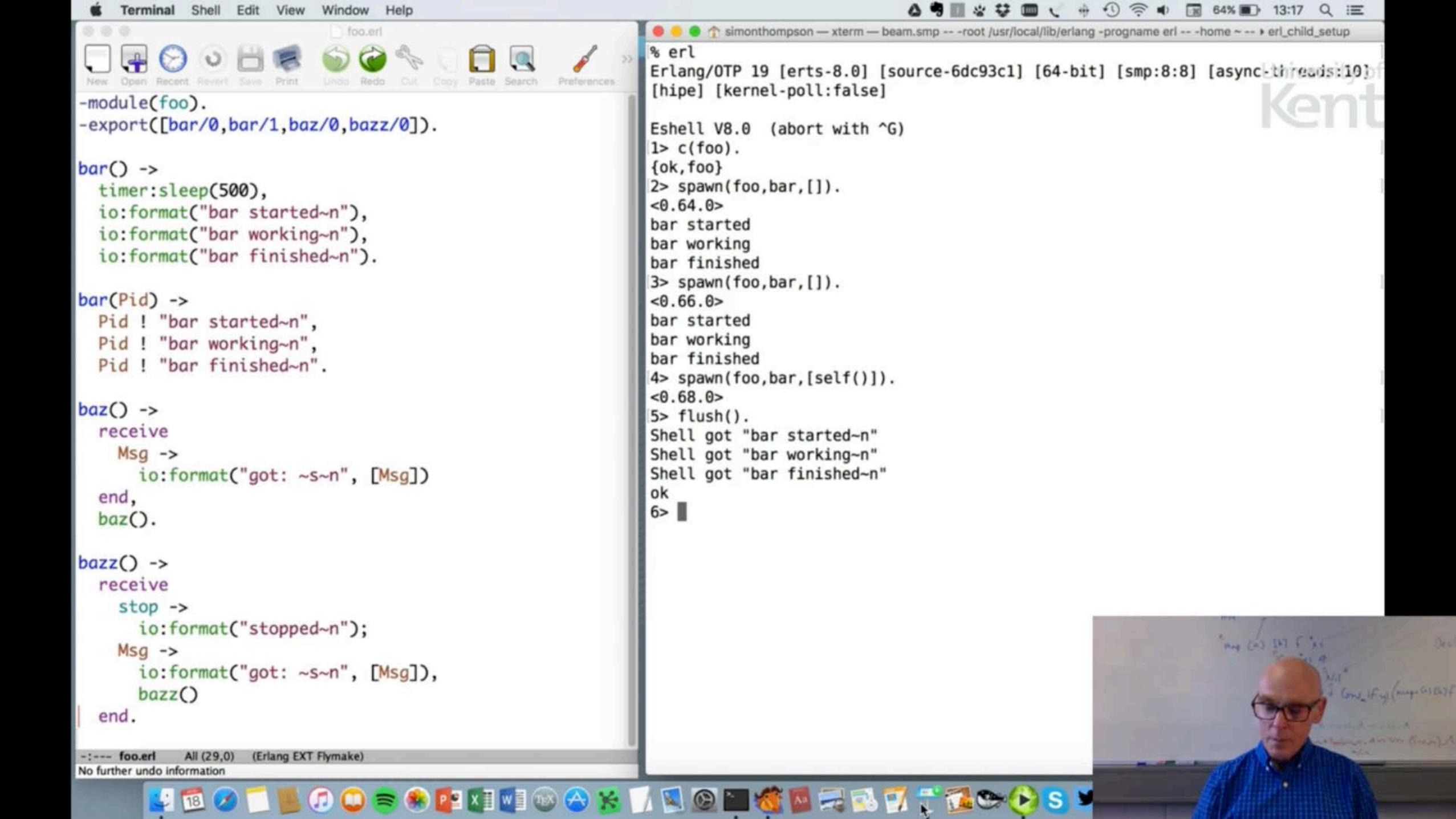


self() - what's my Pid?

self() gives the Pid of the process in which it's called.

In particular, we can get the Pid of the Erlang shell process by calling self() there.

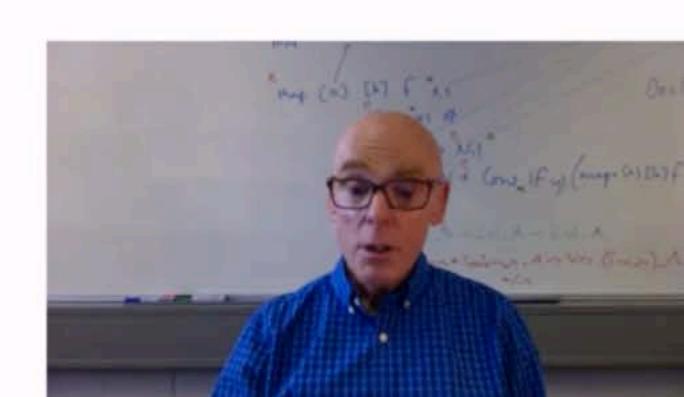






```
receive
Msg -> ... handle Msg ...
end
```

takes the message Msg out of the mailbox and handles it.





```
receive
Msg -> ... handle Msg ...
end
```

takes the message Msg out of the mailbox and handles it.

```
-module(foo).
-export([baz/0]).

baz() ->
  receive
    Msg ->
    io:format("got: ~s~n",[Msg])
  end,
  baz().
```





The general receive construct lets us pattern match on messages and choose between alternatives.

Just like a case statement.



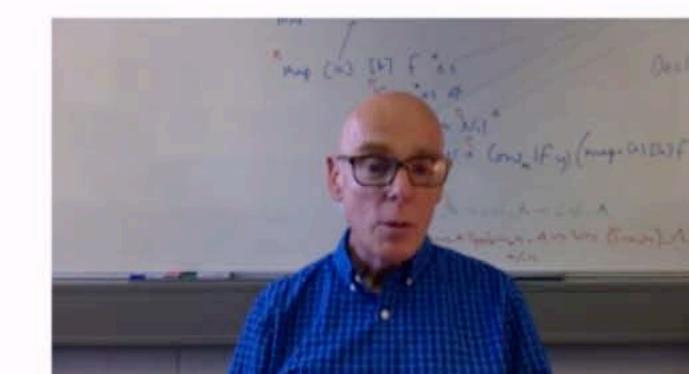


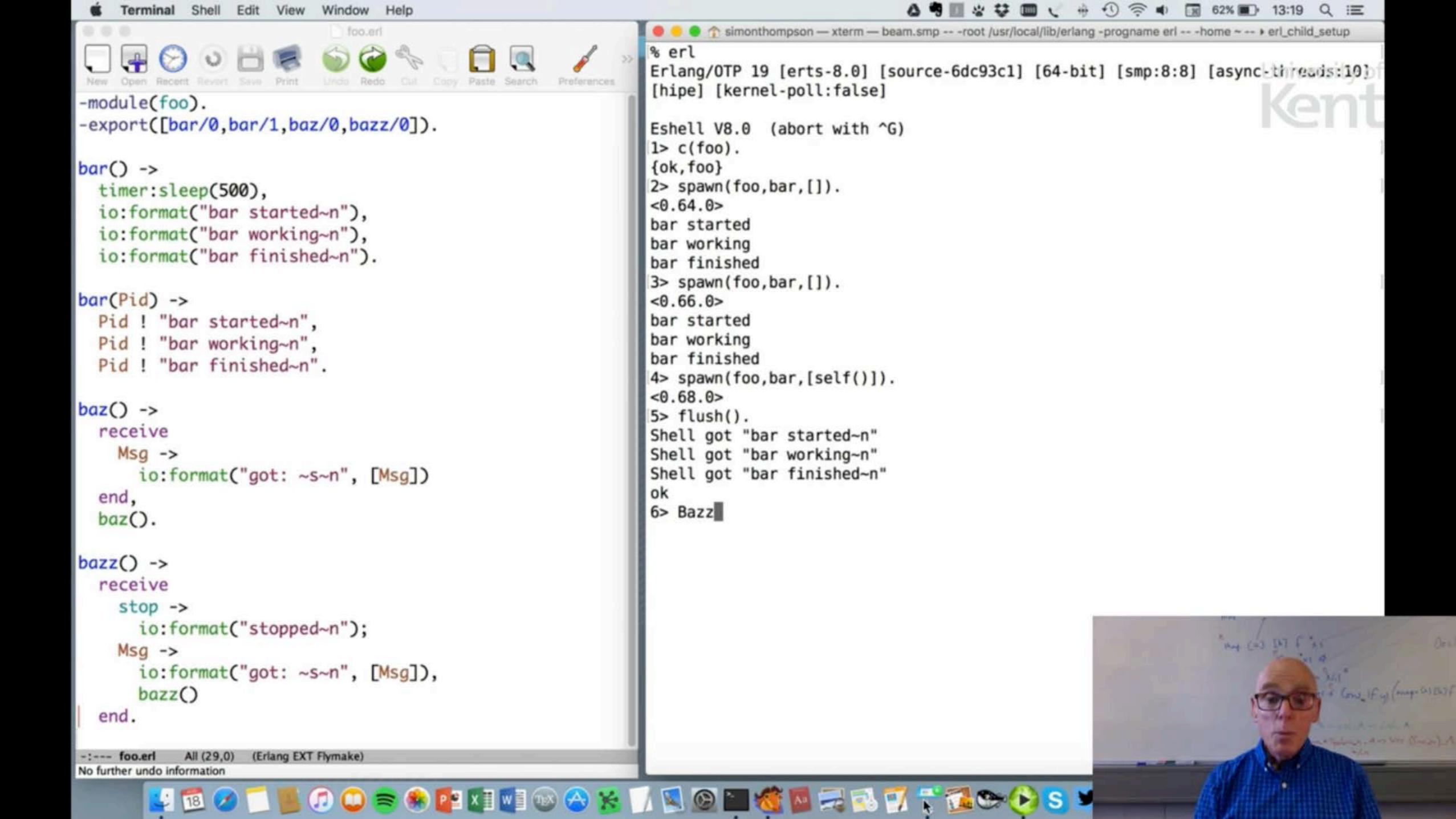
The general receive construct lets us pattern match on messages and choose between alternatives.

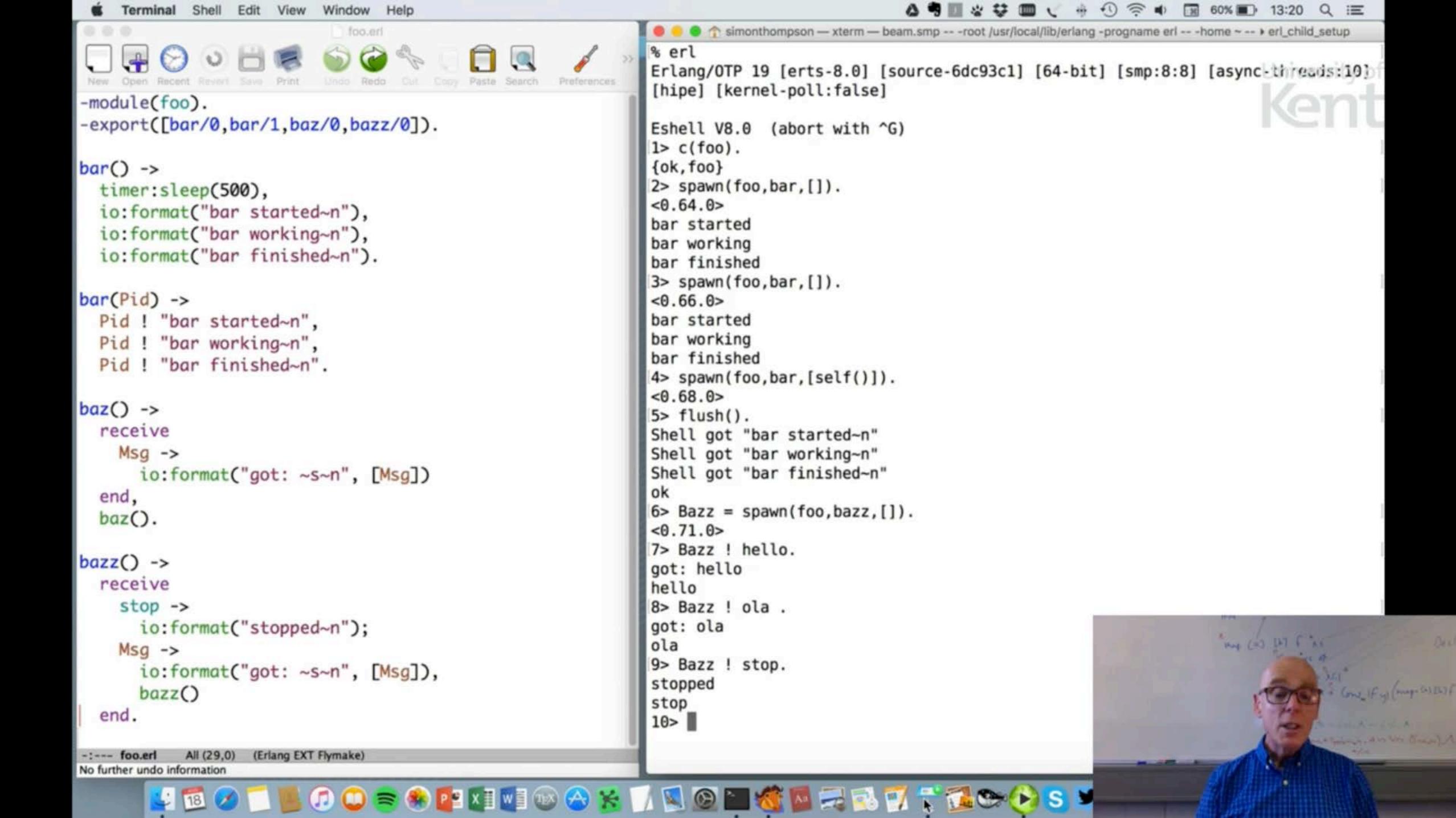
Just like a case statement.

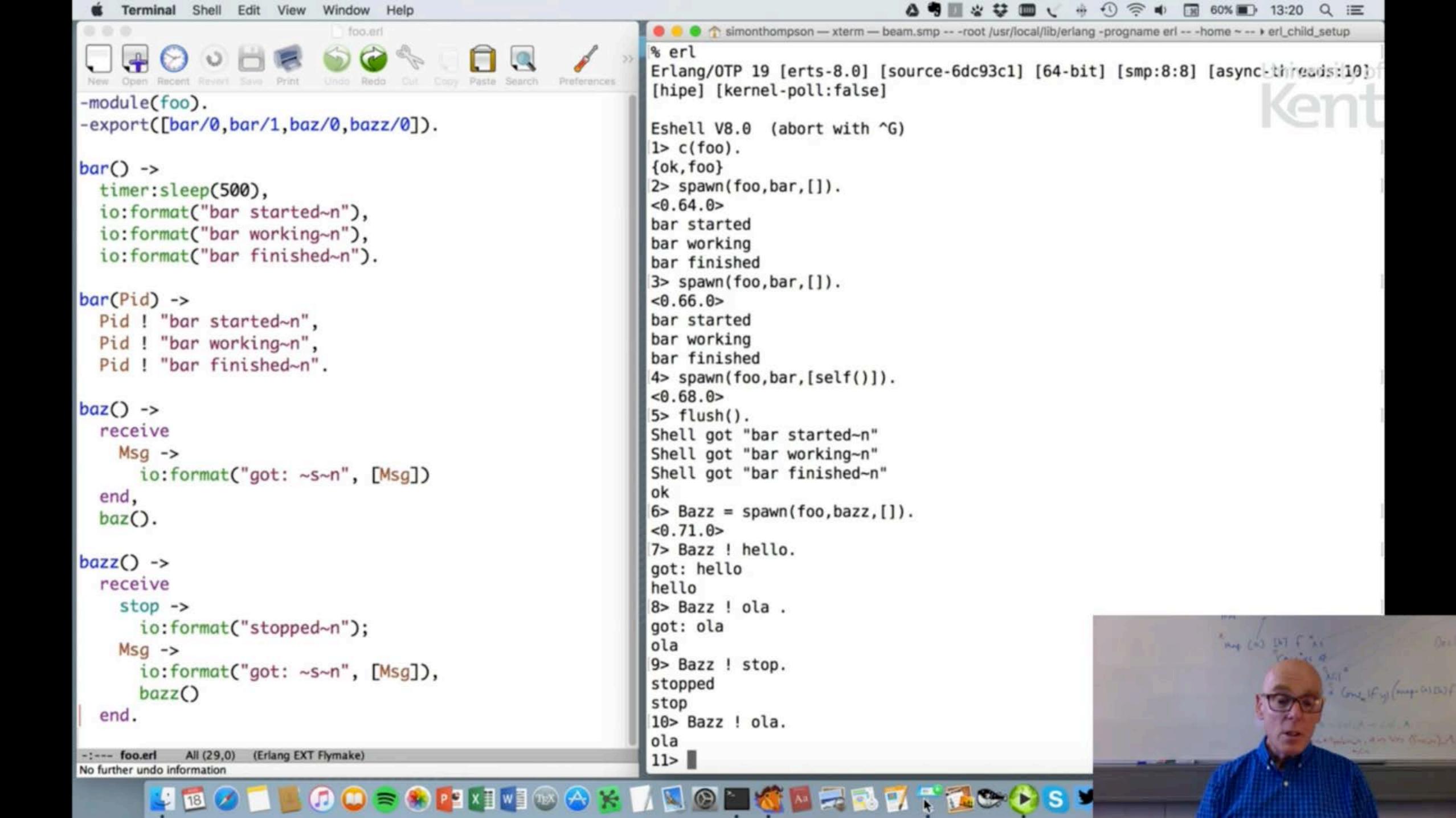
```
-module(foo).
-export([bazz/0]).

bazz() ->
   receive
    stop ->
        io:format("stopped~n");
    Msg ->
        io:format("got: ~s~n",[Msg]),
        bazz()
   end.
```











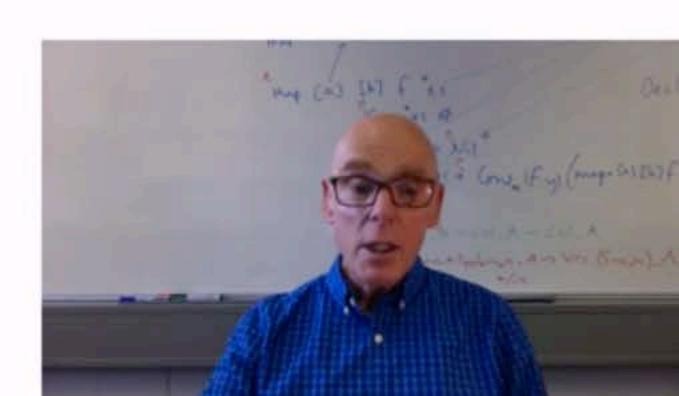
Erlang concurrency in a nutshell

spawn – create a process

send a message

self() – give the Process Identifier (Pid) of a process

receive - handle a message



University of Kernt



Joe Armstrong

EXPERT SYSTEM DEVELOPER, AND ONE OF THE CREATORS OF ERLANG AT ERICSSON





```
area({square,X}) ->
    X*X;
area({rectangle,X,Y}) ->
    X*Y.
```



```
{square,X} ->
   X*X;
{rectangle,X,Y} ->
   X*Y.
```



```
area({square,X}) ->
    X*X;
area({rectangle,X,Y}) ->
    X*Y.
```



```
Pid = spawn(demo1, area, []),
Pid! {square, 10}
                         -module(demo1).
                         -export([area/0]).
                         area() ->
                             receive
                               {square,X} ->
                                X*X;
                               {rectangle,X,Y} ->
                                X*Y
                            end,
                            area().
```



```
area() ->
  receive
     {square,X} ->
       print(X*X);
     {rectangle,X,Y} ->
       print(X*Y)
  end,
  area().
```



Return to sender

```
Pid = spawn(Mod, area, []),
Pid ! {self(), {square, 10}},
receive
   Reply ->
                                   -module(demo2).
      Reply
                                   -export([area/0]).
end.
                                   area() ->
                                      receive
                                        {From, {square,X}} ->
                                          From ! X*X;
                                        {From, {rectangle,X,Y}} ->
                                          From ! X*Y
                                      end,
                                      area().
```



Return to sender

```
Pid = spawn(Mod, area, []),
Pid ! {self(), {square, 10}},
receive
   {Pid, Reply} ->
                                   -module(demo2).
      Reply
                                   -export([area/0]).
end.
                                   area() ->
                                      receive
                                        {From, {square,X}} ->
                                          From ! {self(), X*X};
                                        {From, {rectangle,X,Y}} ->
                                          From ! {self(), X*Y}
                                      end,
                                      area().
```



Four concurrency primitives

- Pid = spawn(...)
- Pid ! Message
- receive Pattern -> Actions; ... end
- self()



```
receive
   Pattern1 ->
     Actions1;
   ...
   PatternN ->
     ActionsN
after Time ->
   TimeoutActions
end
```



Registered processes

```
Only the parent knows its child

Pid = spawn(...)

Pid = spawn()

register(Name, Pid)
```

Pid = whereis(Name)