

Decentralized Systems

Optimistic Rollups

Smart Contract Scalability

- Ethereum is “the slowest computer” (and most expensive) in the world
- **Every verifier executes everything**
 - Worst redundancy possible
- We can extend payment channels to **state channels**
 - Limited to two-player interactions (e.g., chess)
- Can we do better?

Layer-2 Blockchains

- **Blockchains on top of blockchains**
- Goal:
 - Layer-2 (L2) transactions are run on **a few machines**
 - The underlying L1 blockchain **guarantees correctness**
- We'll see a technique called **optimistic rollup**
 - Reference implementation: [Arbitrum](#)

References

- Kalodner et al. [Arbitrum: Scalable, private smart contracts](#). Usenix Security 2018.
- Finematics. ROLLUPS - The Ultimate Ethereum Scaling Strategy? Arbitrum & Optimism Explained – YouTube [video](#), 2021.
- Arbitrum. [Inside Arbitrum Nitro](#).

Optimistic Rollup

Rollup Transaction

- A transaction that “rolls up” several blocks of the L2 blockchain and bundles it all in one transaction
- Posted by a **Manager** that created the blocks
- Blockchain state as a Merkle tree, input messages as attachments to the L1 block
 - We’ll see more on this later
- Meaning: “The state with hash A and this input leads to state with hash B”

Optimistic Rollup Idea

- There is a **happy path** followed **most of the time**
 - 1) A **manager** publishes its assertion “from state A and this input we get to state B” (and send these messages/currency from this chain to these L1 wallets)
 - 2) Some other entities (confusingly also called managers) verify that the computation is correct
 - 3) If after some time (e.g., blocks) nobody disputes that, the new state is **confirmed**

Optimistic Rollup: Challenge

- Managers need to **stake** some currency
- If an assertion is **challenged**:
 - The L2 chain's progress is paused
 - The conflicting managers play a **game** on the L1 chain
 - The L1 smart contract of the L2 chain will make sure it will be won by whoever is correct
 - The loser pays part of its stake to the winner, the rest goes somewhere else (e.g., the verifier)
- It *should* be irrational to willingly post **wrong states**

Challenge: Bisection Game

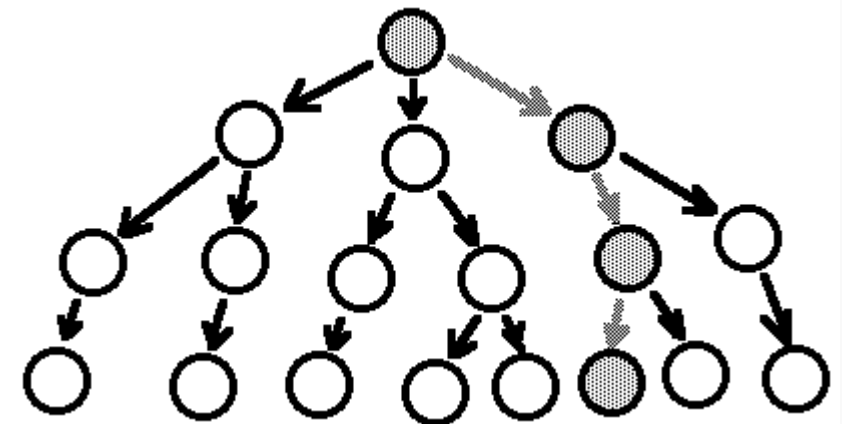
- A: "State A leads to state B in N steps"
- B: "I challenge it!"
- A: "State A leads to state C in $N/2$ steps, and C to B in $N/2$ steps"
- B: "Prove you get to from A to C (or C to B) in $N/2$ steps"
- A: "You get from A to D in $N/4$ steps, and from D to C in $N/4$ steps"
- ...

Bisection Game: Result

- A and B get in $\log(n)$ steps to **a single step**
 - n is the number of execution steps in the whole “rollup block”
 - If either doesn't answer within a deadline, **they lose**
- That step gets **replayed on the main chain** to decide the winner
- How to do this?

Running a Single Step

- Remember that a state is a **Merkle tree**
- A reveals **only the parts of the state needed to run the instruction**
 - **Merkle proof:** you expand the nodes containing the parts of the state touched by the instruction
 - The L1 chain **verifies the hashes**
 - It **runs the instruction**
 - It **verifies the result**



Sequencing Input

- A **sequencer** takes all the messages, compresses them, and posts them to the blockchain as a binary “blob”
 - Such blobs are **erased after some time** (~18 days in Ethereum)
 - Enough time to **challenge** the state, then they're lost
 - Blocks have been introduced as **Proto-Danksharding** in March '24 (input was stored in CALLDATA before)
- Blobs are way cheaper than storing blockchain data
 - The security comes from the fact they're signed, as usual

Avoiding Censorship

- To make sure the sequencer can't censor messages, you can also **post** the ignored ones **on the blockchain**
- They will **have to be included** in the next rollup block
- No punishment for the sequencer: they may not have seen it in good faith

Putting It Together

- An L2 chain that is secure if the L1 chain is secure
- ...and if there's at least a honest manager verifying the rollup transactions
- The L2 chain can be faster and cheaper, because it's verified by much fewer nodes
 - Also, gas limits can be way larger
- Users have ways to bring currency from the L2 to the L1 chain

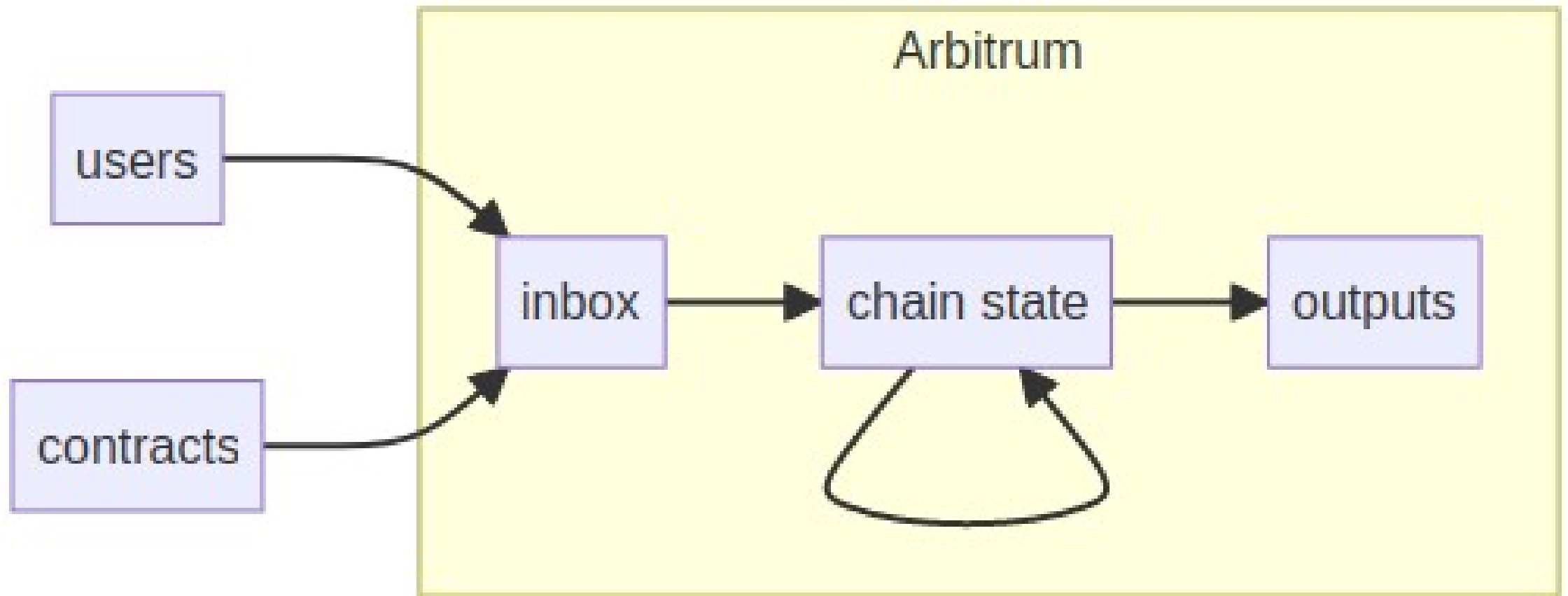
Arbitrum

About Arbitrum

- The most popular L2 network at the moment
- Based on the protocol described in a 2018 USENIX Security [paper](#)
- Currently used by many users, while still under development

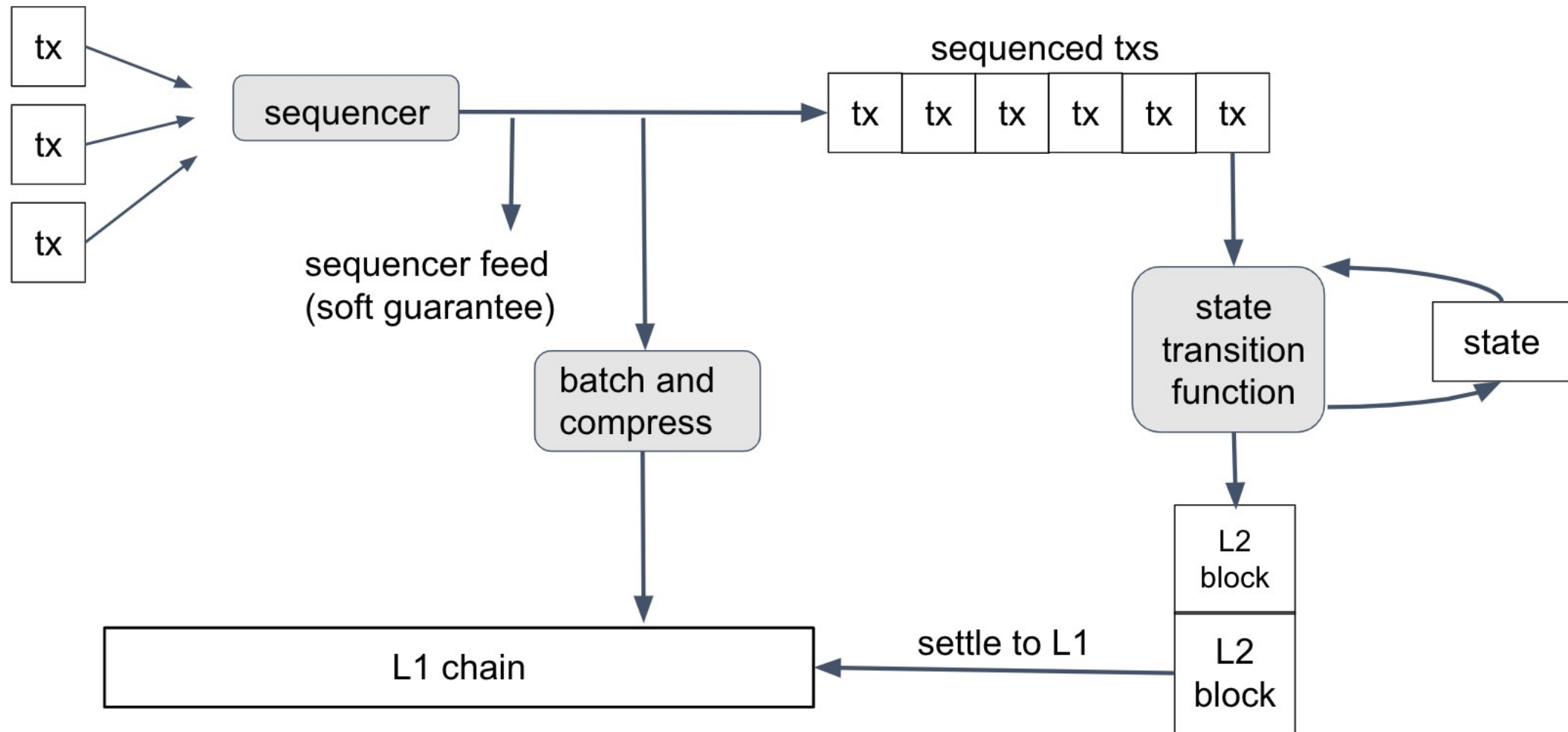


Architecture (1)



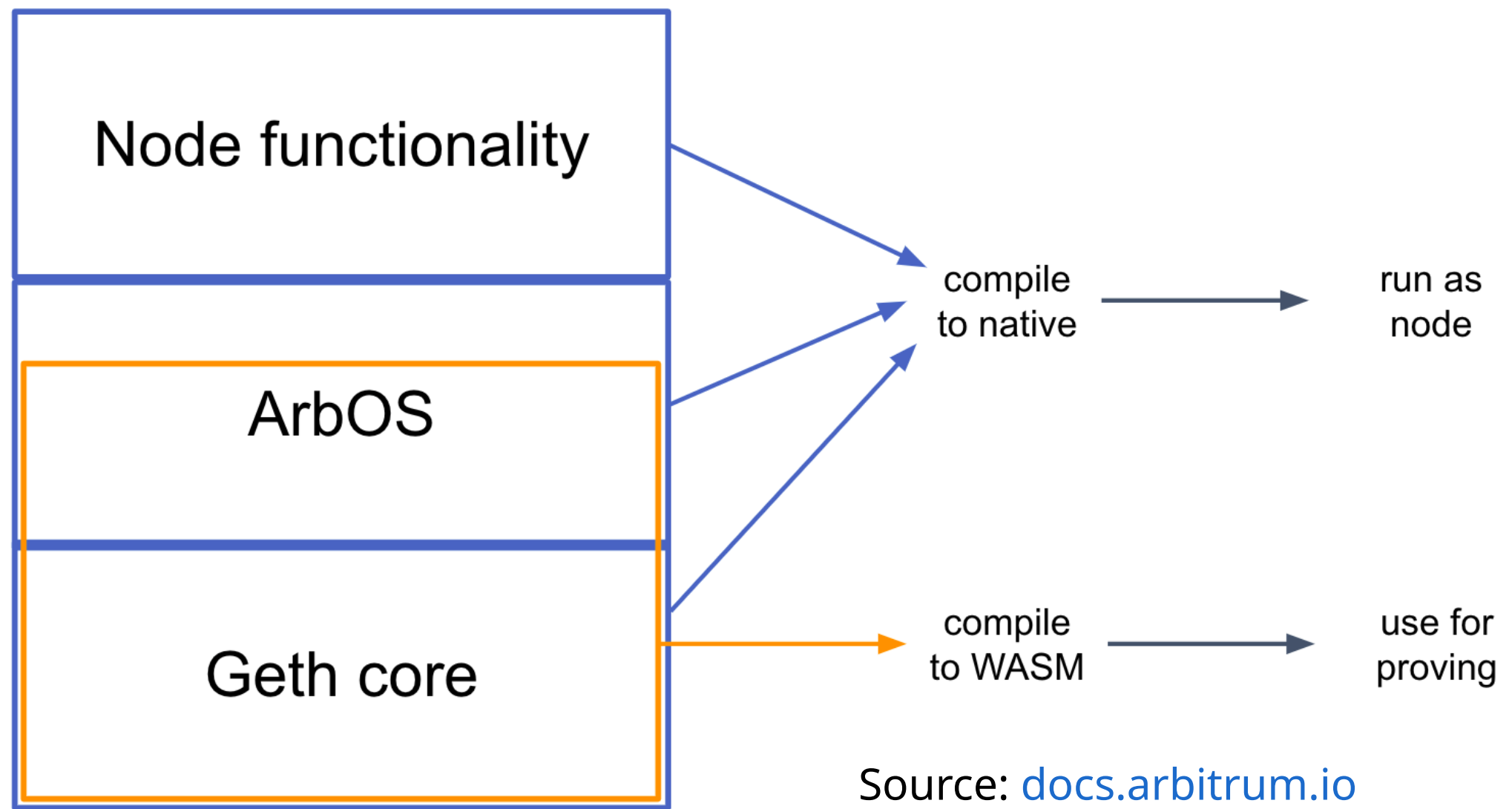
Source: docs.arbitrum.io

Architecture (2)

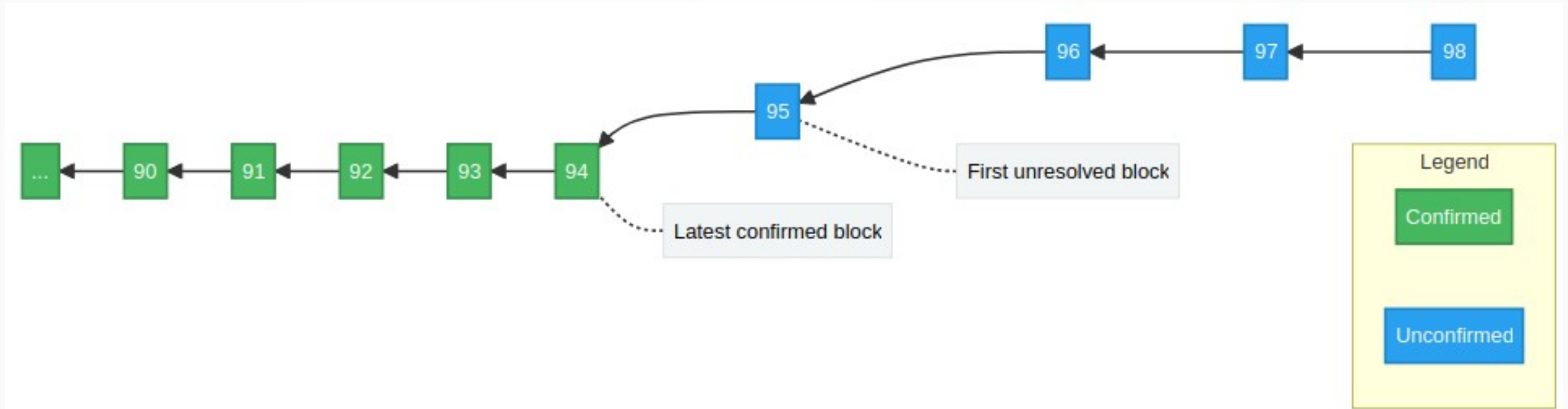


Source: docs.arbitrum.io

Implementation

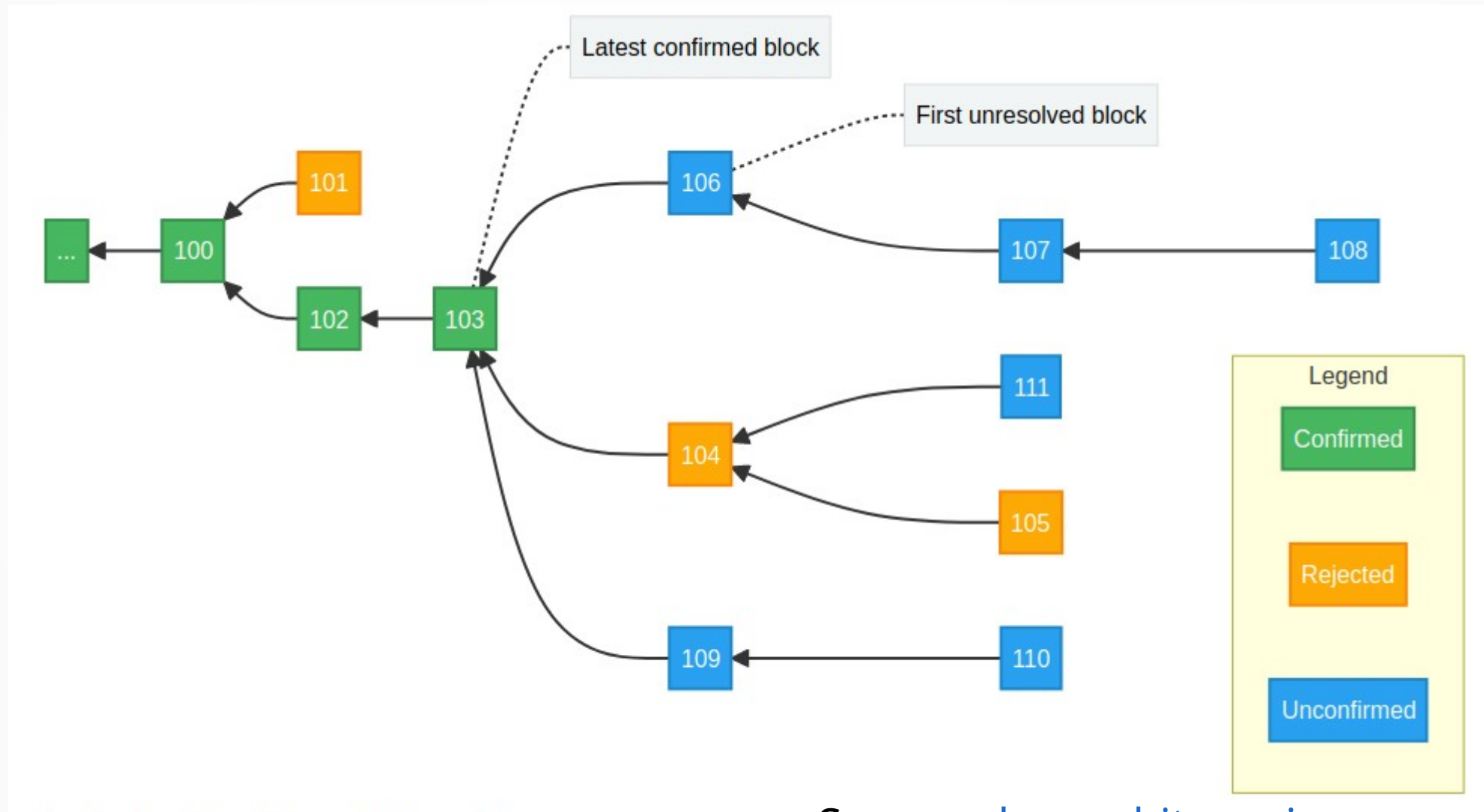


Happy Path



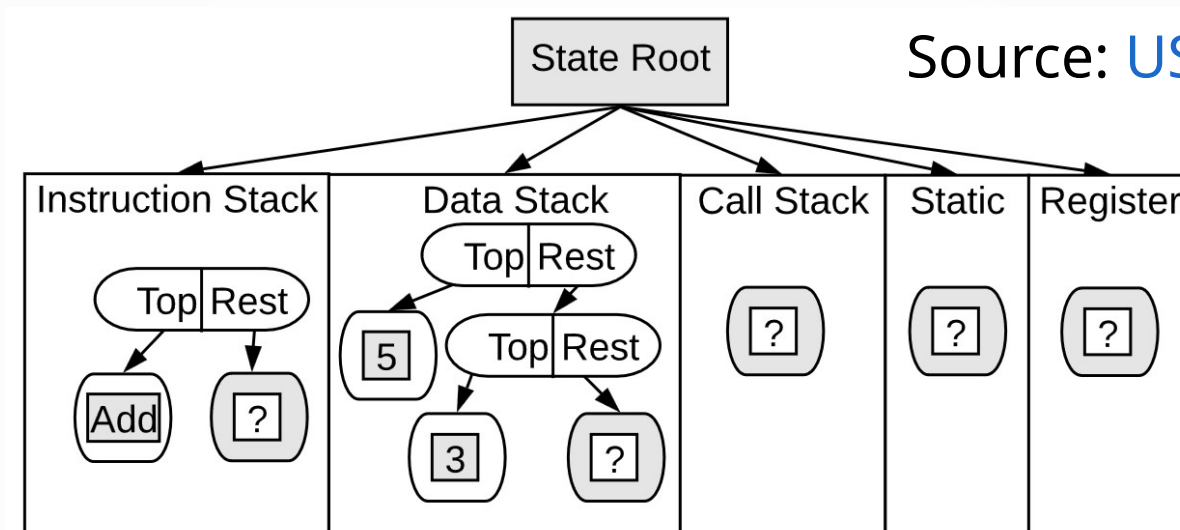
Source: docs.arbitrum.io

(Very) Unhappy Path



Source: docs.arbitrum.io

Virtual Machine for Proofs



- The state as a Merkle tree is visible to the VM
- Micro-instructions that only touch elements close to the top of the stack
- EVM instructions are library calls
- $O(1)$ (rather than $O(\log n)$) to run on the L1 blockchain

K-Way Dissection

- Alice: "State X leads to Y in N steps"
- Bob: "No. State X leads to:
 - $X_{1,B}$ in N/K steps
 - $X_{2,B}$ in N/K steps from $X_{1,B}$
 - ...
 - $X_{K,B}$ in N/K steps from $X_{K-1,B}$ "
- A: "No. State $X_{2,B}$ leads to
 - $X_{(2,1),A}$ in N/K^2 steps..."
- Less rounds and messages involved in the dissections

Status

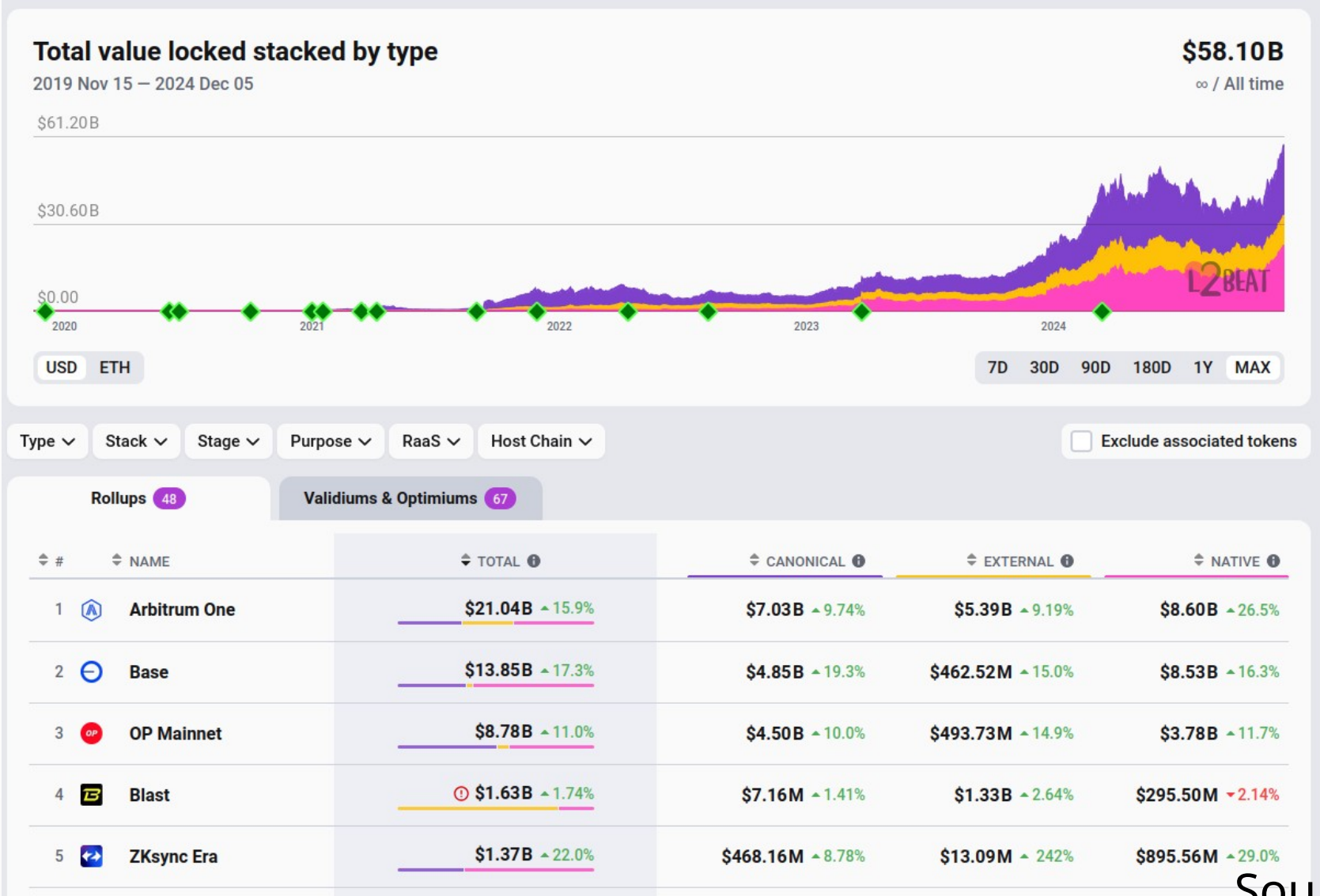
- Currently, **partly centralized**
- Ownership through a DAO governed by tokens
- Validators through an allow list
- The sequencer is centralized
 - It can only delay transactions, not prevent them since they can be posted on the L1 blockchain

L2 Panorama

Zero-Knowledge Rollups

- A promising alternative is based on **crypto proofs** of correctness
- The crypto machinery belongs to the category of **non-interactive zero-knowledge proofs**
 - Proof of correctness that do not reveal any specific information beyond the validity of the statement
 - **zk-SNARK**: zero-knowledge succinct non-interactive argument of knowledge
- Computation-intensive, difficult to prove generic computation

Existing L2 Chains



Source: l2beat.com