

## Lecture 20:-

## Kleene theorem III.

'For each regex there is a corresponding FA'.

$$\underbrace{(Y_1 + Y_2)^+}_{PA_1} \underbrace{Y_1 Y_2^+}_{PA_2} \underbrace{(Y_1 + Y_3)^*}_{PA_3} \dots$$

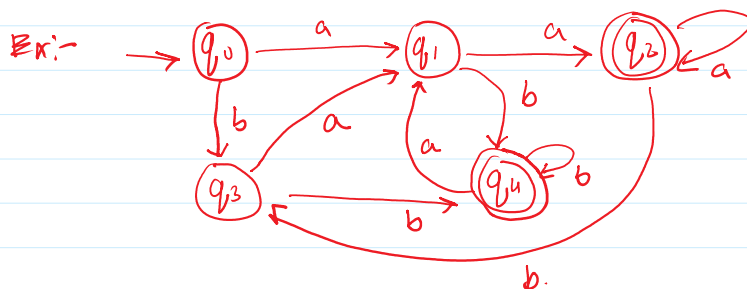
Three issues in combining.

- 1- Sum, Union, +
- 2- Concatenation.
- 3- Closure. \*.

closure:-  $\rightarrow$  During processing Any state encountered final the resultant state will also be final.

$\rightarrow$  The final state will always be written together with initial.

$\rightarrow$  The initial for the first time will also be marked as final.



$$(Y_1 + Y_2 + Y_4 Y_5)^*$$

old state

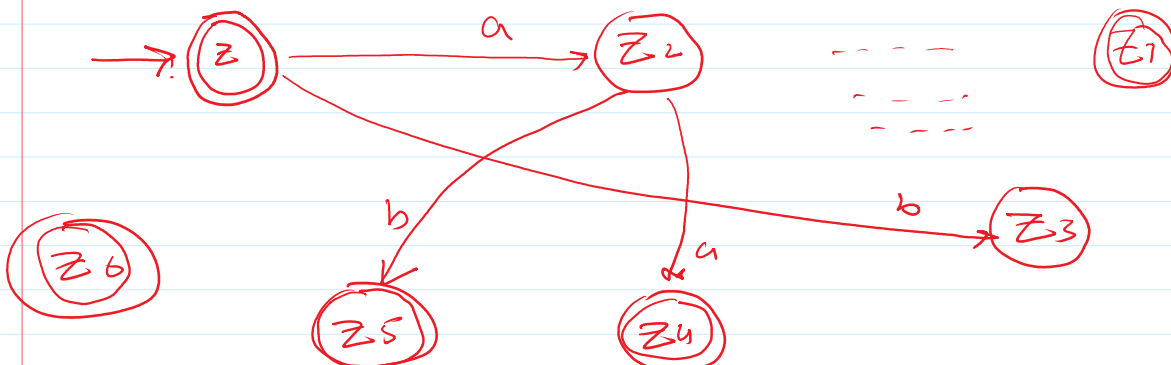
$$\begin{aligned} Z_1^+ &\equiv q_0 \\ Z_2 &\equiv q_1 \\ Z_3 &\equiv q_3 \\ Z_4^+ &\equiv (q_2, q_0) \\ Z_5^+ &\equiv (q_4, q_0) \\ Z_6^+ &\equiv (q_2, q_0, q_1) \\ Z_7^+ &\equiv (q_4, q_0, q_3) \end{aligned}$$

Transition at 'a'

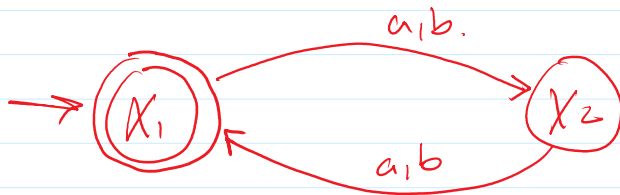
$$\begin{aligned} Z_2 &\equiv q_1 \\ Z_4^+ &\equiv (q_2, q_0) \\ Z_2 &\equiv q_1 \\ Z_6^+ &\equiv (q_2, q_0, q_1) \\ Z_2 &\equiv (q_1, q_1) \\ Z_6^+ &\equiv (q_2, q_0, q_1, q_2, q_0) \\ Z_2 &\equiv (q_1, q_1, q_1) \end{aligned}$$

Transition at 'b'.

$$\begin{aligned} Z_3 &\equiv q_3 \\ Z_5^+ &\equiv (q_4, q_0) \\ Z_5^+ &\equiv (q_4, q_0) \\ Z_3 &\equiv (q_3, q_3) \\ Z_7^+ &\equiv (q_4, q_0, q_3) \\ Z_7^+ &\equiv (q_3, q_3, q_4, q_0) \\ Z_7^+ &\equiv (q_4, q_0, q_3, q_4, q_0) \\ &\equiv (q_3, q_4, q_0) \\ &\equiv (q_4, q_0, q_3) \\ &\equiv Z_7 \end{aligned}$$



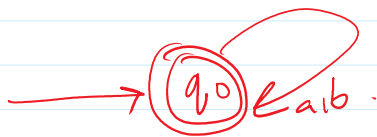
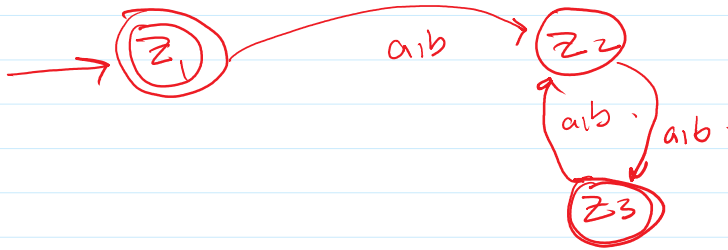
HW.



old state  
 $z_1^- \equiv X_1$   
 $z_2 \equiv X_2$   
 $z_3^+ \equiv X_1$

transition at 'a'  
 $z_2 \equiv X_2$   
 $z_3^+ \equiv (X_1, X_1)$   
 $z_2 \equiv X_2$

transition at 'b'  
 $z_2 \equiv X_2$   
 $z_3^+ \equiv X_1$   
 $z_2 \equiv X_2$



old state  
 $z_1^+ \equiv q_0$   
 $z_2^+ \equiv q_0$

Transition at 'a'  
 $z_2^+ \equiv (q_0, q_0)$   
 $z_2^+ \equiv (q_0, q_0)$

Transition at 'b'  
 $z_2^+ \equiv (q_0, q_0)$   
 $z_2^+ \equiv (q_0, q_0)$

