

lecture 3:- (i) prove that $(S^+)^* = (S^*)^*$.

$S^+ = \{ \text{All possible concatenations of alphabets in } S \text{ excluding } \epsilon \}$.

$\Sigma = \{x\}$.

$\Sigma^+ = \{x, xx, xxx, \dots\}$.

$\Sigma^* = \{ \epsilon, \Sigma^+ \}$.

LHS.

$(S^+)^* = \{ \epsilon \in \text{all possible concatenations of } S^+ \}$.

$= \{ \epsilon \in \text{all possible concatenations of } \{ \text{All possible concatenations of } S \} \}$.

$= \{ \epsilon \in \text{all possible concatenations of } S \}$.

$= S^+.$

RHS:- $S^+ = \{ \epsilon \in \text{all possible concatenations of } S \}$.

$(S^+)^* = \{ \epsilon \in \text{all possible " " " } S^+ \}$.

$= \{ \epsilon \in \text{all " " " " " } \{ \epsilon \in \text{all possible concatenations of } S \} \}$.

$= \{ \epsilon \in \text{all possible concatenations of } S \}$.

$= S^+.$

(i) $(S^+)^+ = S^+ \quad (\text{HW})$.

(ii) $(S^*)^+ = (S^+)^* \quad (\text{HW})$.

Example:- let $S = \{a, bb, bab, abaab\}$.

$bbab$
 $bbbab$.

(i) $\underline{abba} \underline{baab} \overset{x}{ab} \in S^+ \quad \text{false.}$

(ii) $\overset{x}{\underline{abaab}} \overset{x}{\underline{bbab}} \overset{x}{\underline{baab}} \in S^+.$

(iii) Does S^* contains any word with odd b's?
No.

Example:- $L = \{ \text{A Concatenation of two words, will exist iff. two words are Not the Same} \}$.
 Can Such a language exist.

Let $w_1 \in L$ & $w_2 \in L$ $w_1 \neq w_2$.
 $\rightarrow w_1 w_2 \in L$.

Let $w_1 w_2 \in L$ & $w_1 \in L$ $w_1 w_2 \neq w_1$.
 $\rightarrow w_1 w_2 w_1 \in L$

$w_1 w_2 w_1 \in L$ & $w_2 \in L$ $w_1 w_2 w_1 \neq w_2$.
 $\rightarrow \underline{w_1 w_2 w_1} w_2 \in L$.

$(w_1 w_2)(w_1 w_2)$ $w_1 w_2 \neq w_1 w_2$.

therefore such a language does Not exist.

Four ways for defining a language.

1- Descriptive

2- Recursive

3- Regular Expression. (Week six). online Course.

4- Finite Automata.

Certificate Print.
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Descriptive:-
 - Semi formal way.
 - plain English.

Ex:- $L = \{ \text{strings of even length} \}$ $\Sigma = \{b\}$.

$L = \{ bb, bbbb, bbbbbb, \dots \}$.

$L = \{ \text{strings not starting with a} \}$ $\Sigma = \{a, b, c\}$.

$L = \{ b, c, ba, bb, bc, ca, cb, cc, baa, \dots \}$.

$L = \{ \text{strings of length 3} \}$ $\Sigma = \{0, 1, 2\}$.

$$= \{ 000, 001, 002, \dots \}$$

$$L = \{ \text{strings ending in 1} \}. \quad \Sigma = \{0, 1\}.$$

$$= \{ 1, 01, 11, 001, 011, 101, 111, \dots \}.$$

$$\text{EQUAL} = \{ \# \text{ of } a's = \# \text{ of } b's \} \quad \Sigma = \{a, b\}.$$

$$= \{ \Lambda, ab, ba, aabb, bbba, abab, baba, abba, baab, \dots \}.$$

$$\text{EVEN-EVEN} = \{ \text{Even number of } a's \text{ \& \text{ even number of } b's} \}. \\ \Sigma = \{a, b\}.$$

$$= \{ \Lambda, aa, bb, aaaa, bbbb, aabb, bbba, abab, baba, \dots \}.$$

$$\text{INTEGER} = \{ \text{All concatenations of } \Sigma, - \text{ will not be alone and} \\ \text{it always at the beginning} \}.$$

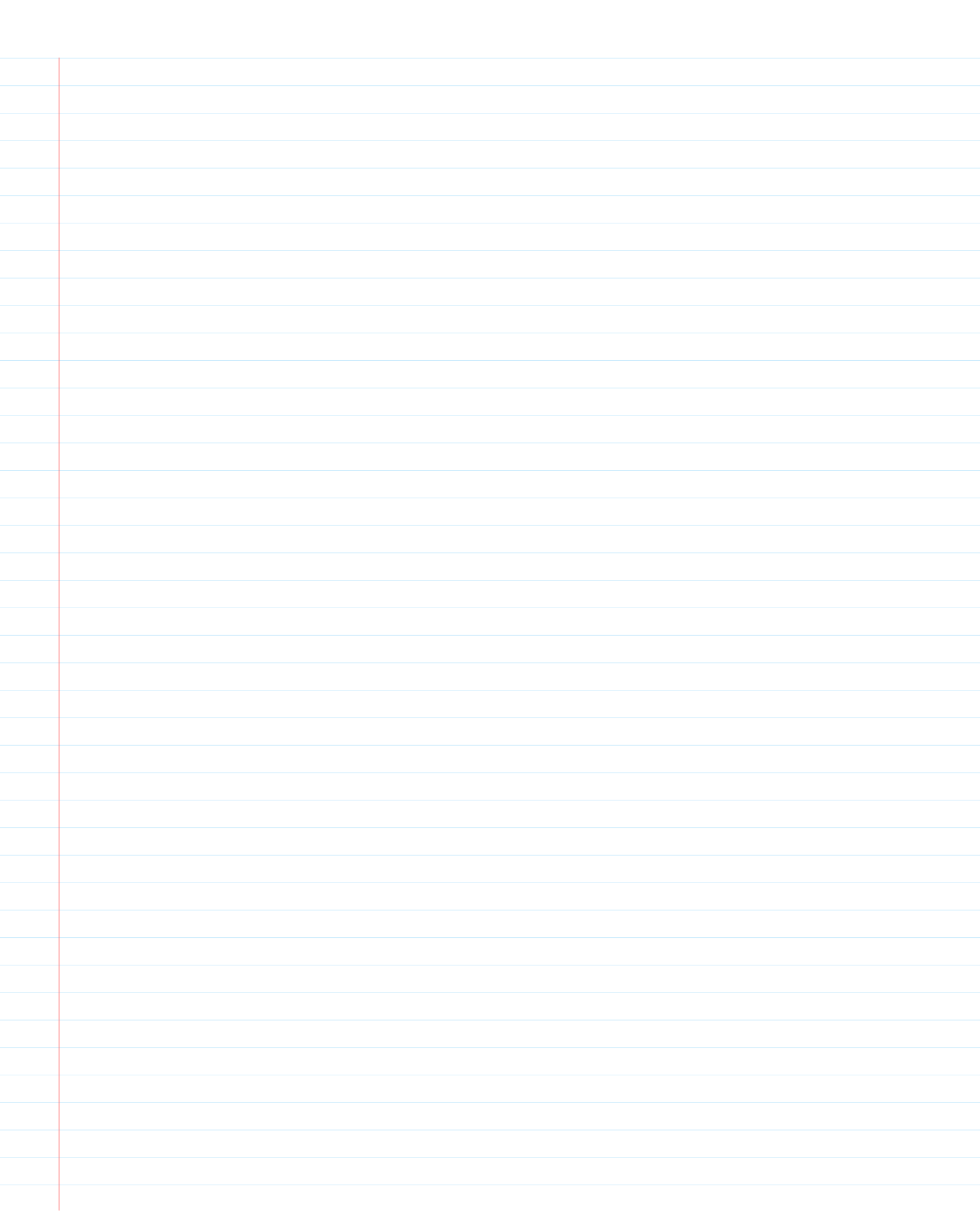
$$\Sigma = \{ -, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9 \}.$$

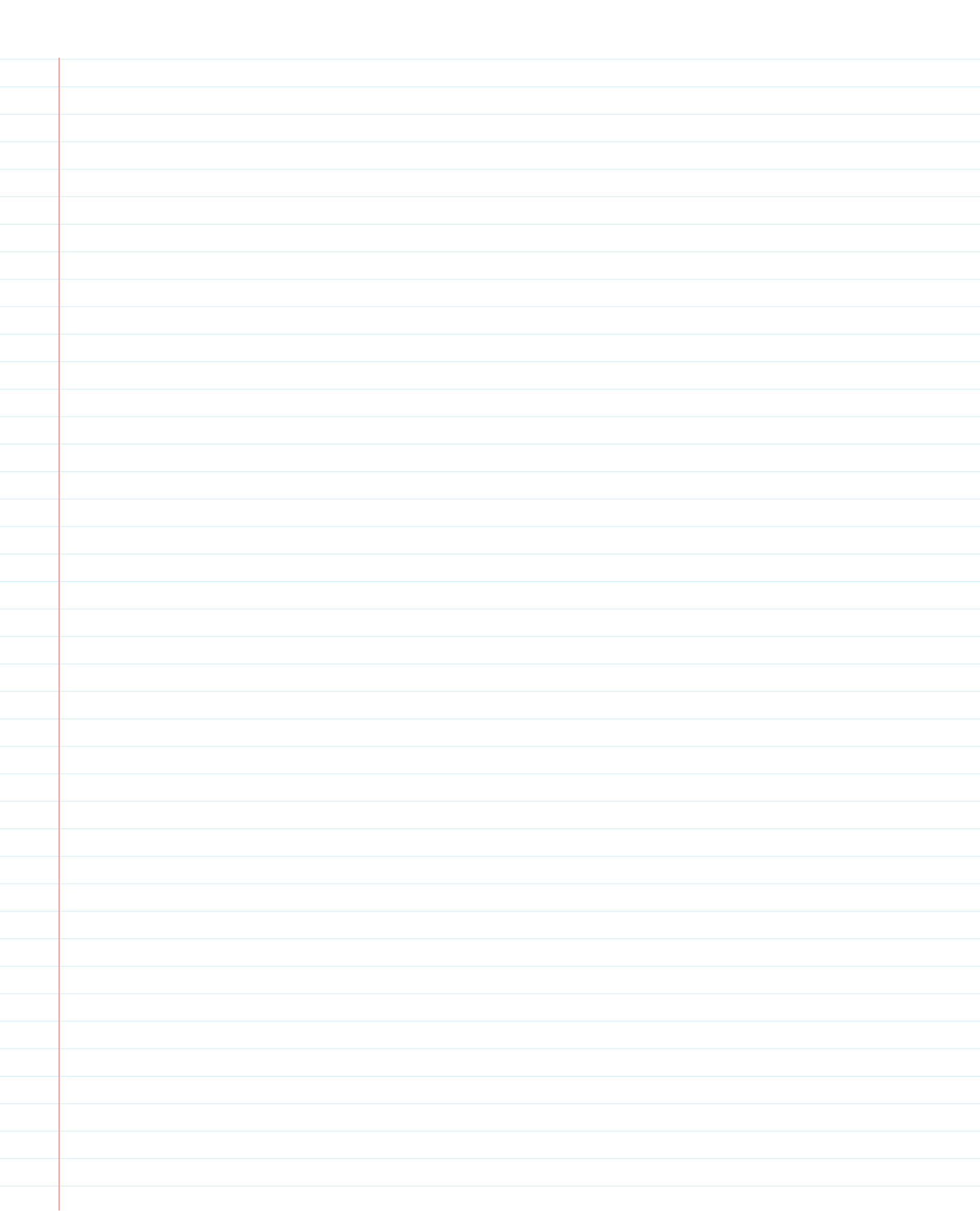
$$L = \{ a^n b^n : n = 1, 2, 3, \dots \}. \quad \Sigma = \{a, b\}.$$

$$= \{ ab, aabb, aaabbb, \dots \}.$$

$$L = \{ a^n b^n c^n : n = 1, 2, 3, \dots \}. \quad \Sigma = \{a, b, c\}.$$

$$= \{ abc, aabbcc, \dots \}.$$





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