# UFSC Livro do time (2024-2025)

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## 1 Data Structures

## 1.1 Centroid Decomposition

```
#include <algorithm>
#include <iostream>
#include <utility>
#include <utility>
#include <vector>

using namespace std;
using pii = pair<int, int>;

// 'closest_red', guery and update were used for solving xenia and the tree.
struct CentroidDecomposition {
    vector<utor<int>> tree;
    vector<int>> subtrees_sz, closest_red;
    vector<vector<oli>yetpii>> parents;
```

```
vector<bool> removed;
  CentroidDecomposition(vector<vector<int>> adj)
    :tree{adj} {
    int n = tree.size();
    subtrees_sz.resize(n);
    removed.assign(n, false);
    closest_red.assign(n, 1e9);
    parents.resize(n);
   centroid_decomposition(0, -1);
  void calculate_subtree_sizes(int u, int p = -1) {
   subtrees_sz[u] = 1;
    for (auto v : tree[u]) {
     if (v == p || removed[v])
       continue;
      calculate_subtree_sizes(v, u);
      subtrees_sz[u] += subtrees_sz[v];
  int find_centroid(int u, int p, int n) {
    for (auto v : tree[u]) {
      if (v == p || removed[v])
       continue;
      if (subtrees_sz[v] > n / 2)
       return find_centroid(v, u, n);
    return u;
  void calculate_distance_to_centroid(int u, int p, int centroid, int d) {
    for (auto v : tree[u]) {
      if (v == p || removed[v])
        continue;
      calculate_distance_to_centroid(v, u, centroid, d + 1);
    parents[u].push_back({centroid, d});
  void centroid_decomposition(int u, int p = -1) {
    calculate_subtree_sizes(u);
    int centroid = find_centroid(u, p, subtrees_sz[u]);
    for (auto v : tree[centroid]) {
     if (removed[v])
        continue:
      calculate_distance_to_centroid(v, centroid, centroid, 1);
    removed[centroid] = true;
    for (auto v : tree[centroid]) {
     if (removed[v])
       continue;
      centroid_decomposition(v, u);
  int query(int u) {
    int ret = closest_red[u];
    for (auto&[p, pd] : parents[u])
     ret = min(ret, pd + closest_red[p]);
    return ret;
  void update(int u) {
   closest_red[u] = 0;
    for (auto &[p, pdist] : parents[u])
     closest_red[p] = min(closest_red[p], pdist);
};
```

## 1.2 Ordered Set

```
// C++ program to demonstrate the
// ordered set in GNU C++
#include <iostream>
using namespace std;
// Header files, namespaces,
// macros as defined above
#include <ext/pb_ds/assoc_container.hpp>
```

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace gnu pbds;
#define ordered_set tree<int, null_type,less<int>, rb_tree_tag,tree_order_statistics_node_update>
// Driver program to test above functions
int main()
    // Ordered set declared with name o_set
    ordered_set o_set;
    // insert function to insert in 
// ordered set same as SET STL
    o_set.insert(5);
    o set insert (1):
    o_set.insert(2);
    // Finding the second smallest element
    // in the set using * because
    // find_by_order returns an iterator
    cout << *(o_set.find_by_order(1))</pre>
         << endl;
    // Finding the number of elements
    // strictly less than k=4
    cout << o_set.order_of_key(4)</pre>
         << endl:
    // Finding the count of elements less
    // than or equal to 4 i.e. strictly less
    // than 5 if integers are present
    cout << o_set.order_of_key(5)</pre>
         << endl;
    // Deleting 2 from the set if it exists
    if (o_set.find(2) != o_set.end())
        o_set.erase(o_set.find(2));
    // Now after deleting 2 from the set
    // Finding the second smallest element in the set
    cout << *(o_set.find_by_order(1))</pre>
         << endl;
    // Finding the number of
    // elements strictly less than k=4
    cout << o_set.order_of_key(4)</pre>
         << endl;
    return 0;
```

## 1.3 Sparse Table

```
#include<vector>
#include<utility>
// preprocessing: O(n log n)
// range minimum query (minimum element in [L, R] interval): O(1)
struct SparseTable {
  vector<vector<int>> st;
  int k = 25;
  int n;
  SparseTable(const vector<int>& vec) {
    n = vec.size();
    st.assign(k+1, vector<int>(n));
    st[0] = vec;
    for (int i = 1; i <= k; ++i)
  for (int j = 0; j + (1 << i) <= n; ++j)
   st[i][j] = min(st[i-1][j], st[i-1][j + (1 << (i-1))]);</pre>
  int query(int 1, int r) {
    int i = bit_width((unsigned long) (r - 1 + 1)) - 1; // change to log2 and memoization if <math>c++20 is
    return min(st[i][1], st[i][r - (1 << i) + 1]);</pre>
};
```

## 1.4 Sparse Table

```
from math import log2
from typing import Callable
class SparseTable:
    - Consegue responder range minimum queries em O(1) [P.s: range minimum queries --> menor número
          dentro de um intervalo]
    - Não pode ser atualizada
    - construida em O(n log n)
    - como funciona: The main idea behind Sparse Tables is to precompute all answers for range queries
           with power of two length.
                     Afterwards a different range query can be answered by splitting the range into
                           ranges with power of two lengths,
                     looking up the precomputed answers, and combining them to receive a complete
    def __init__(self, arr: list, func,max_n: int):
        func: função que depende do tipo de query.
        max_n: maior quantidade possível de elementos
        k = int(log2(max_n)) + 1
        self.st = [[0 for i in range(k)] for j in range(max_n)]
        self.func = func
        for i, v in enumerate(arr):
            self.st[i][0] = v
        for j in range(1, k):
            for i in range(len(arr)):
                if i + (1 << j) - 1 >= len(arr):
                   break
                self.st[i][j] = func(
                    self.st[i][j-1],
                    self.st[i + (1 << (j-1))][j-1]
    def query(self, L: int, R: int):
        # Assumes 1-indexed range
        R-=1; L-=1
        k = int(log2(max(R-L, 1)))
        return self.func( self.st[L][k], self.st[R - (1 \ll k) + 1][k])
```

## 1.5 Fenwick Tree

```
class FenwickTree:
    - Consegue responder range queries em O(log n)
    - Consegue atualizar um elemento por vez em O(log n)
    - O(n log n) pra construir
    - indexada por 1
    - queries tambem devem ser indexadas por 1
    Dado um nodo n
         - posição do filho: n - n&-n
         - posição do pai: n + n&-n
    def __init__(self, n):
         self.ft = [0 for _ in range(n+1)]
    def update(self, i: int, v: int):
          """incremento o item na posição 'i' com o valor 'v'"""
         while i < len(self.ft):</pre>
             self.ft[i] += v
             i += i&-i
    def query(self, L: int, R: int):
    """retorna soma no intervalo [L, R]"""
         return self.__sum_from_beginning(R) - self.__sum_from_beginning(L-1)
    def __sum_from_beginning(self, i: int):
    """Retorna soma no intervalo [1, i]"""
         sum = 0
```

```
while i:
    sum += self.ft[i]
    i -= i&-i
return sum
```

## 1.6 Suffix Array

```
#include <bits/stdc++.h>
#include <vector>
using namespace std:
using vi = vector<int>;
using ii = pair<int, int>;
class SuffixArray {
private:
  vi RA;
                                   // rank array
  void countingSort(int k) {
                                   // up to 255 ASCII chars
   int \max i = \max(300, n);
                                   // clear frequency table
   vi c(maxi, 0);
   int t = c[i];
     c[i] = sum;
     sum += t;
   vi tempSA(n);
   for (int i = 0; i < n; ++i) // sort SA
     tempSA[c[SA[i] + k < n ? RA[SA[i] + k] : 0]++] = SA[i];
   swap(SA, tempSA); // update SA
  void constructSA() { // can go up to 400K chars
   SA.resize(n);
   iota(SA.begin(), SA.end(), 0); // the initial SA
   RA.resize(n):
   for (int i = 0; i < n; ++i)
     RA[i] = T[i];
                                   // initial rankings
   for (int k = 1; k < n; k <<= 1) { // repeat log_2 n times</pre>
     // this is actually radix sort
     countingSort(k); // sort by 2nd item
     countingSort(0); // stable-sort by 1st item
     vi tempRA(n);
     tempRA[SA[0]] = r;
     for (int i = 1; i < n; ++i) // compare adj suffixes</pre>
       (RA[SA[i] + k] == RA[SA[i - 1] + k]))
              : ++r:
     swap(RA, tempRA); // update RA
if (RA[SA[n - 1]] == n - 1)
       break; // nice optimization
  const char *T; // the input string
  const int n; // the length of T
               // Suffix Array
  SuffixArray(const char *initialT, const int _n) : T(initialT), n(_n) {
   constructSA(); // O(n log n)
  };
```

# 2 Dynamic Programming

### 2.1 Coin Change

```
dp[0] = 0

for amount in range(1, val + 1):
    for coin in coins:
        if amount - coin >= 0:
            temp = 1 + dp[amount - coin]
            if temp < dp[amount]:
                 dp[amount] = temp
return dp[val] if dp[val] != float("inf") else -1</pre>
```

### 2.2 Knapsack

### 2.3 Longest Common Subsequence

```
Returns the length of the longest common subsequence
between strings "p" and "q". The sequence doesn't need
to be contiguous.
Example:
p -> "ABCXYZ"
q -> "ABCYZ"
longest_common_subsequence(p, q) -> "ABXZ", of length 4.

'''
def longest_common_subsequence(p, q):
    dp = [[0] * (len(q) + 1) for _ in range(len(p) + 1)]
    for i in range(len(p)):
        if p[i] == q[j]:
             dp[i+1][j+1] = 1 + dp[i][j]
        else:
             dp[i+1][j+1] = max(dp[i+1][j], dp[i][j+1])
    return dp[len(p)][len(q)]
```

# 3 Graphs

# 3.1 Class Graph

```
class Graph:

def __init__(self):
    self.graph = {}

def add_unidirectional_edge(self, u, v):
    if u in self.graph:
        self.graph[u].add(v)
    else:
        self.graph[u] = {v}

def add_bidirectional_edge(self, u, v):
    self.add_unidirectional_edge(u, v)
    self.add_unidirectional_edge(v, u)

def get_next_vertices(self, u):
    if u in self.graph:
        return self.graph[u]
    else:
        return set()
```

```
,,,
Returns best path from "start" to "end" in a weighted graph.
If such path doesn't exist, returns -1.
Uses the "Graph" class.
"weights" list is 0 indexed (i.e. index 0 is not empty).
class Graph:
    def __init__(self):
        self.graph = {}
    def add_unidirectional_edge(self, u, v):
        if u in self.graph:
            self.graph[u].add(v)
            self.graph[u] = \{v\}
    \begin{tabular}{ll} \bf def & add\_bidirectional\_edge(self, u, v): \\ \end{tabular}
        self.add_unidirectional_edge(u, v)
        self.add_unidirectional_edge(v, u)
    def get_next_vertices(self, u):
        if u in self.graph:
            return self.graph[u]
            return set()
def get_closest_node(results, unvisited):
    val = float("inf")
    for node in unvisited:
        if results[node] <= val:</pre>
            val = results[node]
            res = node
    return res
def dijkstra(graph, start, end, vertice_count, weights):
    unvisited = set()
    for i in range(1, vertice_count+1):
        unvisited.add(i)
    results = [float("inf") for _ in range(vertice_count + 1)]
    results[start] = 0
    while unvisited:
        node = get_closest_node(results, unvisited)
        unvisited.remove(node)
        if node == end:
            return results[node] if results[node] != float("inf") else -1
            break
        for next in graph.get_next_vertices(node):
            if next in unvisited:
                 d = results[node] + weights[node][next]
                if d < results[next]:</pre>
                    results[next] = d
```

#### 3.3 Max Flow

```
#include <hits/stdc++ h
using namespace std;
int n:
vector<vector<int>> capacity;
vector<vector<int>> adi:
int bfs(int s, int t, vector<int> &parent) {
  fill(parent.begin(), parent.end(), -1);
  parent[s] = -2;
  queue<pair<int, int>> q;
  q.push({s, 1e9});
  while (!q.empty()) {
    int cur = q.front().first;
    int flow = q.front().second;
    q.pop();
    for (int next : adj[cur]) {
  if (parent[next] == -1 && capacity[cur][next]) {
        parent[next] = cur;
        int new_flow = min(flow, capacity[cur][next]);
        if (next == t)
```

```
return new flow;
        q.push({next, new_flow});
  return 0;
int maxflow(int s, int t) {
 int flow = 0;
  vector<int> parent(n);
  int new flow;
  while (new_flow = bfs(s, t, parent)) {
    flow += new_flow;
    int cur = t;
    while (cur != s) {
     int prev = parent[cur];
     capacity[prev][cur] -= new_flow;
     capacity[cur][prev] += new_flow;
      cur = prev;
  return flow:
```

#### 3.4 Floyd Warshall

# 4 Linear Sorting

### 4.1 Radix Sort

```
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
using ull = unsigned long long;
// If numbers are too large, MAYBE increase base.
// Once I had to use 2^15 to get AC // Some numbers are using 11. In practice, this is only needed for very large
// bases.
// int base = 32768;
int base = 512; // IDK a good value
int get_digit(int a, int divisor) { return a / divisor % base; }
bool cmp(int a, int b, int divisor) {
 return get_digit(a, divisor) < get_digit(b, divisor);</pre>
void counting_sort(vector<int> &vec, vector<int> &output, int divisor) {
  int 1 = INT32\_MAX, r = 0;
  vector<int> aux(vec.begin(), vec.end());
  for (auto &v : aux) {
    v = get_digit(v, divisor);
    1 = \min(1, v);
    r = max(r, v);
```

```
vector < int > f(r - 1 + 1);
  for (auto &v : aux)
    ++f[v-1];
  for (int i = 1; i < f.size(); ++i)</pre>
    f[i] = f[i - 1] + f[i];
  for (11 i = vec.size() - 1; i >= 0; --i) {
    int d = aux[i];
    \operatorname{output}[f[d-1]-1] = \operatorname{vec}[i];
    f[d - 1]--;
void radix_sort(vector<int> &vec) {
  auto [il, ir] = minmax_element(vec.begin(), vec.end());
  int 1 = *i1, r = *ir;
  int num_digits = 1;
  int tmp = r;
  while (tmp >= base)
    num_digits++;
    tmp = tmp / base;
  vector<int> a(vec.begin(), vec.end()), b(vec.begin(), vec.end());
  auto *output = &a;
  auto *aux = &b;
  int divisor = 1;
  for (int i = 0; i < num_digits; ++i) {</pre>
    swap(aux, output);
    counting_sort(*aux, *output, divisor);
    divisor *= base;
  for (int i = 0; i < vec.size(); ++i)
    vec[i] = (*output)[i];
```

### 5 Math

#### 5.1 Eratostenes

#### 5.2 Factorize

#### 5.3 Is Prime

```
from math import sqrt

def is_prime(n):
    sq = int(sqrt(n))

if n % 2 == 0:
    return False
    for i in range(3, sq + 1, 2):
        if n % i == 0:
            return False
    return True
```

#### 5.4 Mdc E Mmc

```
def mdc(a, b):
    dividendo = a
    divisor = b
    if b > a:
        divisor = a
    while 1:
        resto = dividendo % divisor
        if resto == 0:
            break
        dividendo, divisor = divisor, resto
    return divisor
#Dont use without defining mdc().
def mmc(a, b):
    return a * b // mdc(a, b)
```

# 6 String Algorithms

#### 6.1 Knuth Morris Pratt

```
Returns list with the "string" indexes on which the pattern "substr" starts.
def build_lps_list(sub):
    lps = [0] * len(sub)
    for i in range(1, len(sub)):
         j = lps[i-1]
        while j > 0 and sub[i] != sub[j]:
        j = lps[j-1]
if sub[i] == sub[j]:
            j += 1
        lps[i] = j
    return lps
def knuth_morris_pratt(substr, string):
    lps = build_lps_list(substr)
    i = 0
    j = 0
    res = []
    while i < len(string):</pre>
        if string[i] == substr[j]:
            i += 1
             j += 1
             if j == len(substr):
                res append(i-j)
                 j = lps[j-1]
        elif j == 0:
             j = lps[j-1]
    return res
```

## 3.2 Longest Palindromic Substring

```
Manacher's algorithm. There are two implementations of it below.
First implementation:
 - Returns the longest palindromic substring of "s" with smallest starting index.
Second implementation:
 - Returns the length of the longest palindromic substring of "s" (a little bit
   faster because doesn't need to slice the substring out of "s").
All of this in O(n) time complexity.
First time writing this algo. Do not ask me how it works.
\#Returns the actual substring.
def longest_palindromic_substring(s):
    str = "#" + "#".join(s) + "#"
    c = 0
    r = 0
    lps = [0] * len(str)
    best_length = 0
    best_idx = 0
    for i in range(1, len(str) - 1):
        if i < r:
            lps[i] = min(r-i, lps[2*c - i])
        lps[i] += 1
        if lps[i] > best_length:
            best_length = lps[i]
            best_idx = i
        if i + lps[i] > r:
           c = i
            r = i + lps[i]
    return s[(best_idx - best_length)//2 : (best_idx + best_length)//2]
#Returns the length of the substring.
def longest_palindromic_substring(s):
    str = "#" + "#".join(s) + "#"
    c = 0
    \mathbf{r} = 0
    lps = [0] * len(str)
    best_length = 0
    for i in range(1, len(str) - 1):
       if i < r:
    lps[i] = min(r-i, lps[2*c - i])</pre>
        while len(str) - 1 - lps[i] > i and str[i + 1 + lps[i]] == str[i - 1 - lps[i]]:
            lps[i] += 1
        if lps[i] > best_length:
           best_length = lps[i]
        if i + lps[i] > r:
            r = i + lps[i]
    return best_length
```

# 7 Utils

### 7.1 Binary Search

```
///
Returns index of target, in ascending iterables.
For descending iterables, swap "<" with ">".
If number doesn't exist, returns -1.
def binary_search(iterable, target, down_idx=0, top_idx=None):
    if top_idx is None:
        top\_idx = len(iterable) - 1
    while down_idx <= top_idx:</pre>
        cur = (down_idx + top_idx) // 2
        if iterable[cur] == target:
            return cur
        elif iterable[cur] < target: #swap here
            down_idx = cur + 1
            top_idx = cur - 1
    return -1
Returns index of smallest number in iterable bigger than or equal
to target, in ascending iterables.

Swapping ">=" with "<=" returns index of biggest number in iterable
smaller than or equal to target, in descending iterables.
If number doesn't exist, returns -1.
def binary_search(iterable, target, down_idx=0, top_idx=None):
```

```
if top_idx is None:
       top_idx = len(iterable) - 1
    while down_idx <= top_idx:</pre>
        cur = (down_idx + top_idx) // 2
        if iterable[cur] >= target: #swap here
            res = cur
            top_idx = cur - 1
        else:
           down_idx = cur + 1
    return res
Returns index of smallest number in iterable bigger than or equal
to target, in descending iterables.
Swapping ">=" with "<=" returns index of biggest number in iterable
smaller than or equal to target, in ascending iterables.
If number doesn't exist, returns -1.
def binary_search(iterable, target, down_idx=0, top_idx=None):
    if top_idx is None:
       top_idx = len(iterable) - 1
    while down_idx <= top_idx:</pre>
        cur = (down_idx + top_idx) // 2
        if iterable[cur] >= target: #swap here
            res = cur
            down_idx = cur + 1
        else:
            top_idx = cur - 1
    return res
```

#### 7.2 Fast Io

```
import sys
input = lambda: sys.stdin.readline().removesuffix('\n')
print = lambda s="", end="\n": sys.stdout.write(str(s)+end)
```

# 7.3 Inversion Counting

```
a realização de queries para saber quantos valores já apareceram em um determinado intervalo de nú meros.

por exemplo:
    suponha que um loop itere sobre os valores 5 4 3 2 1.

    na terceira iteração do for loop (quando o valor for 3),
    a árvore indicará que foram encontrados dois valores no intervalo [3:5]

isso permite encontrar inversões de forma rápida, uma vez que tudo que precisamos fazer para encontrar todas
as inversões de um número n é descobrir quantos números maiores que ele aparecem antes dele.
    i.e: bast fazer uma query a fenwick tree no intervalo [n:len(iteratable)].

### If = FenwickTree(len(iteratable))
mp = normalize(iteratable)
inv = 0

for val in iteratable:
    inv += ft.query(mp[val], len(iteratable))
    return inv
```

#### 7.4 Binary Search For Smallest Possible Value

```
using namespace std;
bool valid(ull time, ull goal, vull& machines) {
  ull sum = 0;
  for (auto@ m : machines)
    sum += time/m;
  return sum >= goal;
int main() {
  fast_io();
  ull n, t;
 cin >> n >> t;
  vull machines;
  while (n--) {
   ull tmp;
    cin >> tmp;
    machines.push_back(tmp);
  ull boundary = t*(*max_element(machines.begin(), machines.end())) + 1;
  DEBUG (boundary);
  for (ull b = boundary/2; b >= 1; b /= 2) {
   DEBUG(valid(k+b, t,machines));
while (!valid(k+b, t,machines)) k+=b;
  cout << k+1 << '\n';
```

# 8 Vscode Debug Config

#### 8.1 Launch

```
1 {
2     // Use IntelliSense to learn about possible
     attributes.
3     // Hover to view descriptions of existing attributes
4     // For more information, visit: https://go.microsoft
     .com/fwlink/?linkid=830387
```

```
"name": "(qdb) Launch",
                "type": "cppdbg",
                "request": "launch",
10
                "program": "${workspaceFolder}/main",
11
                "args": [],
12
                "stopAtEntry": false,
13
                "cwd": "${fileDirname}",
14
                "environment": [],
15
                "externalConsole": false,
16
                "MIMode": "qdb",
17
                "setupCommands": [
18
19
                         "description": "Enable pretty-
20
                            printing for gdb",
                         "text": "-enable-pretty-printing",
21
                         "ignoreFailures": true
22
                    },
23
24
                         "description": "Set Disassembly
25
                            Flavor to Intel",
                         "text": "-qdb-set disassembly-flavor
26
                             intel",
                         "ignoreFailures": true
27
28
29
30
31
32
```

# 9 Geometry

#### 9.1 Circle Intersection

"version": "0.2.0",

"configurations": [

6

```
struct Circle {
   int x, y, r;
};

inline bool is_inside(Circle &a, Circle &b) {
   double d = sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));

   return d <= a.r + b.r || d <= a.r - b.r || d <= b.r - a.r;
}</pre>
```