

Digital Systems Design and Laboratory

[8. Combinational Circuit Design]

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Outline

- ❑ **Review of Combinational Circuit Design**

- ❑ Design of Circuits with Limited Gate Fan-In

- ❑ Gate Delays and Timing Diagrams

- ❑ Hazards in Combinational Logic

- ❑ Simulation and Testing of Logic Circuits

Combinational Circuit Design

□ Generic combinational circuit design steps

- Translate the word description into a switching function (Unit 4)
- Simplify the function
 - Boolean algebra (Units 2 and 3)
 - Karnaugh map (Unit 5)
 - Quine-McCluskey (Unit 6)
 - Other methods
- Realize it using available logic gates (Unit 7)
 - Hardware cost = (#terms, #literals)
 - Start from minimum SOP \equiv AND-OR \rightarrow NAND-NAND/OR-NAND/NOR-OR
 - Start from minimum POS \equiv OR-AND \rightarrow NOR-NOR/AND-NOR/NAND-AND
 - # of levels = maximum # of cascaded gates between I/Os
 - Change level by factoring or multiplying out

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Example 1: Single Output

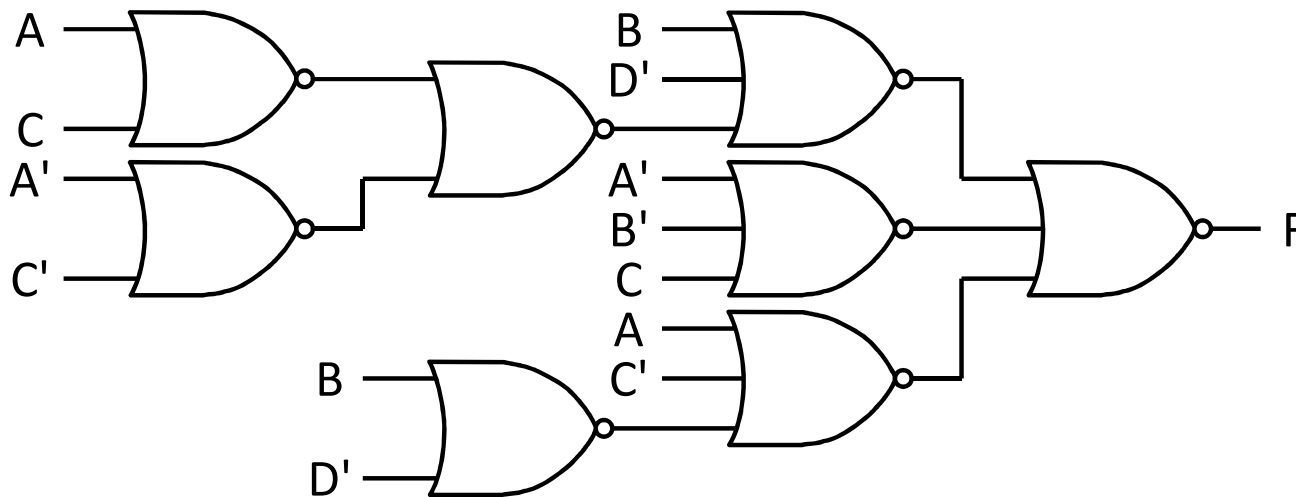
□ Realize $f(a, b, c, d) = \sum m(0, 3, 4, 5, 8, 9, 10, 14, 15)$ using 3-input NOR gates

➤ $F' = A'B'C'D + AB'CD + ABC' + A'BC + A'CD'$ (from K-map)

➤ $F = (A + B + C + D')(A' + B + C' + D')(A' + B' + C)(A + B' + C')(A + C' + D)$

- Two 4-input OR gates (X) → factoring
- Three 3-input OR gates (O)
- One 5-input AND gate (X)

➤ $F = [B + D' + (A + C)(A' + C')] [A' + B' + C] [A + C' + B'D]$



Example 2: Multiple Outputs (1/2)

- Realize the functions using only 2-input NAND gates and inverters

BC \ A	0 1	
	0	1
00	1	1
01		1
11	1	
10	1	

$$F_1 = B'C' + AB' + A'B$$

BC \ A	0 1	
	0	1
00	1	1
01		
11	1	1
10	1	

$$F_2 = B'C' + BC + A'B$$

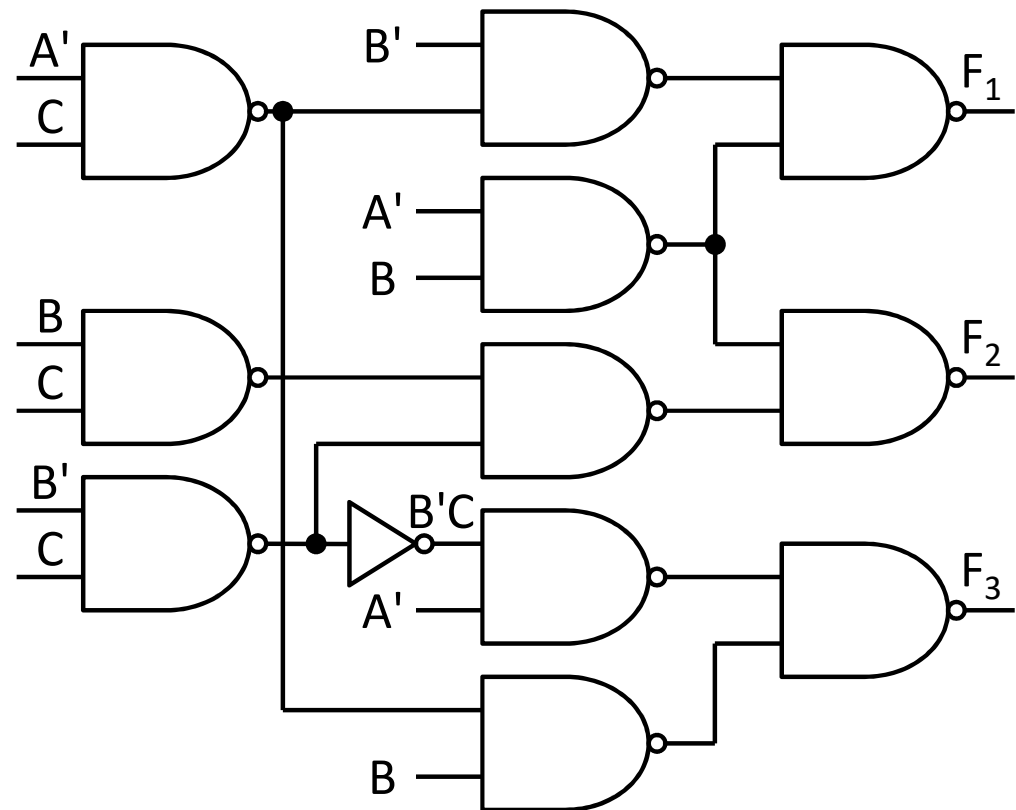
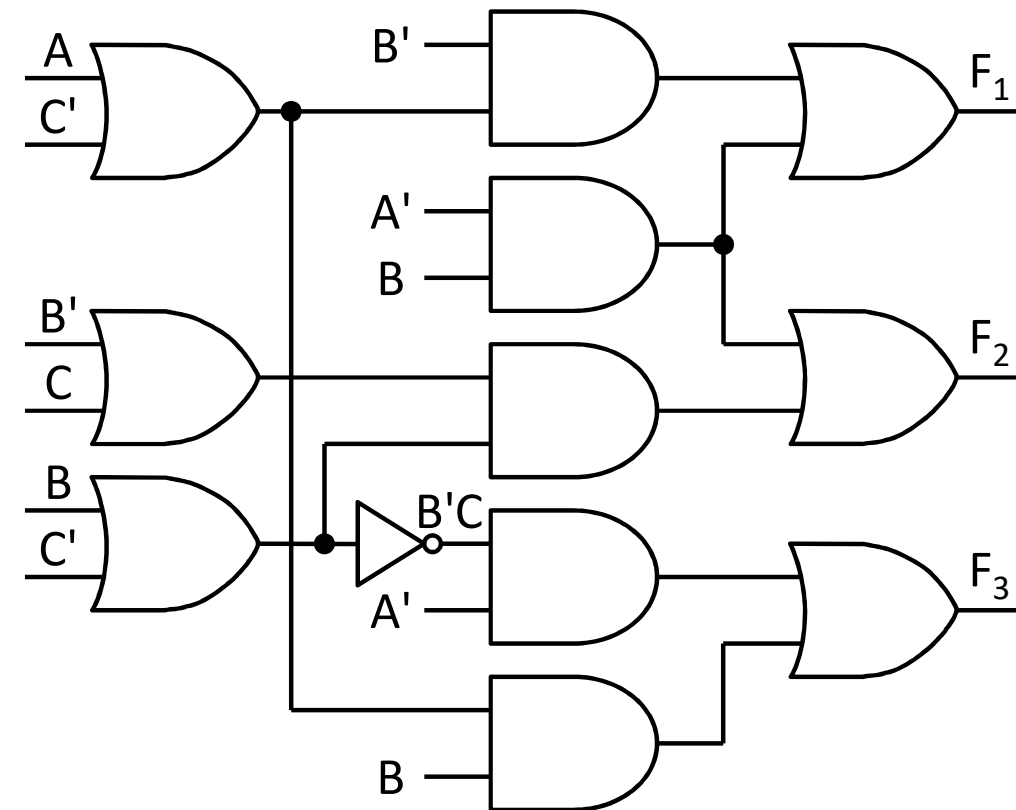
BC \ A	0 1	
	0	1
00		
01	1	
11		1
10	1	1

$$F_3 = \underline{A'B'C} + AB + BC'$$

Example 2: Multiple Outputs (2/2)

□ Two-level

- $F_1 = B'C' + AB' + A'B = B'(A + C') + A'B$
- $F_2 = B'C' + BC + A'B = B'C' + B(A' + C)$
- $F_3 = \underline{A'B'C} + AB + BC' = \underline{A'(B + C)'} + B(A + C')$



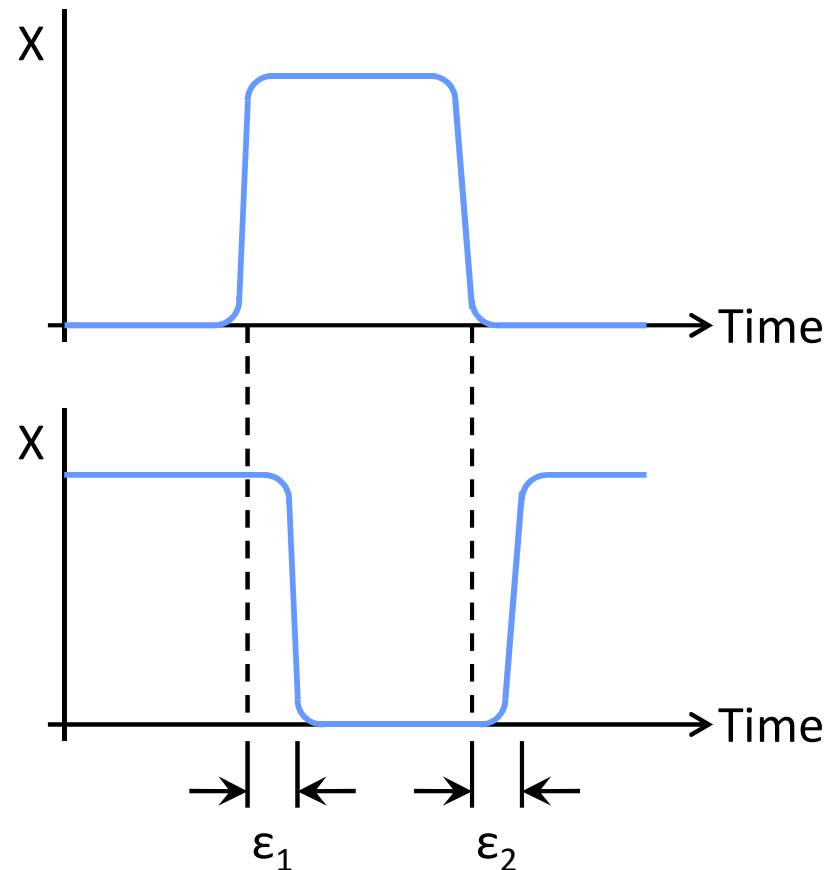
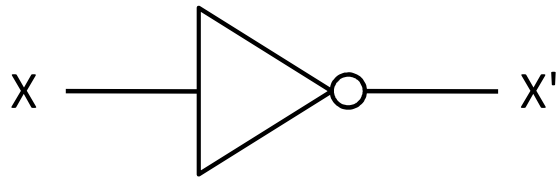
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Gate Delay

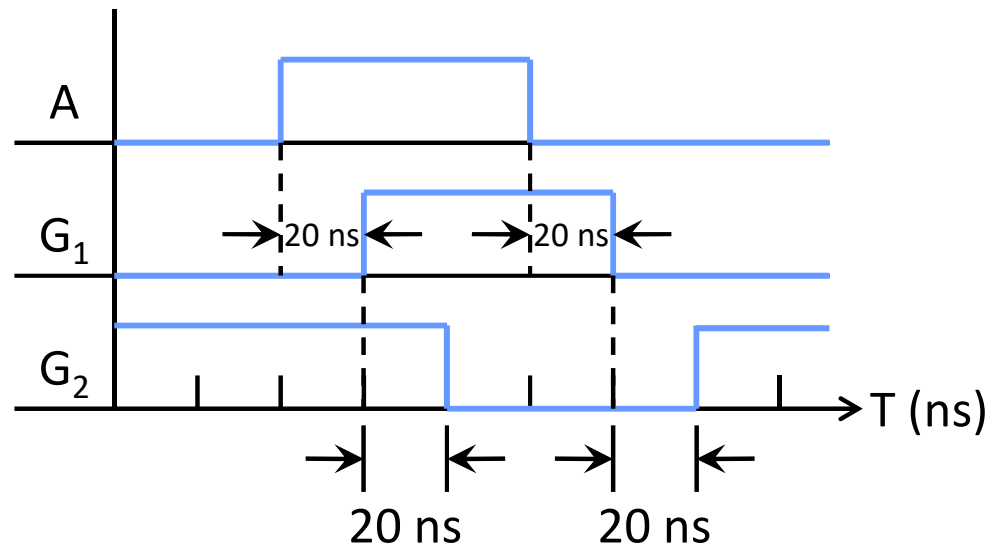
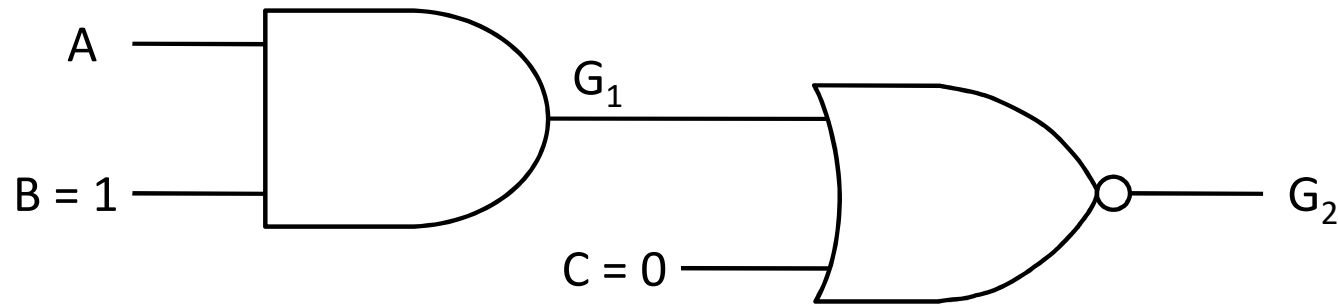
□ Propagation delay

- When the input to a logic gate is changed, the output takes some time to change value
- Usually different for input rising $0 \rightarrow 1$ and falling $1 \rightarrow 0$



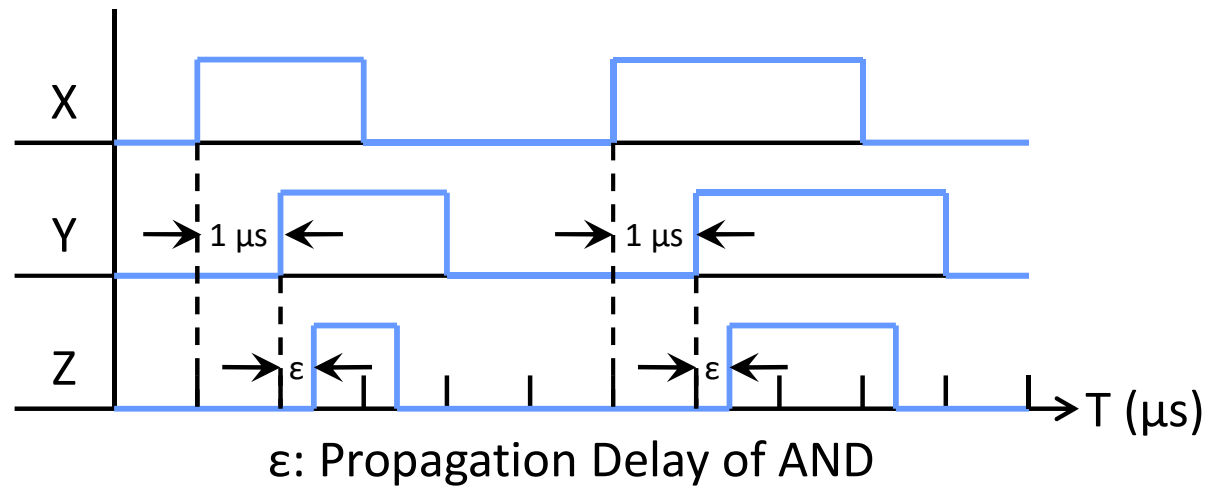
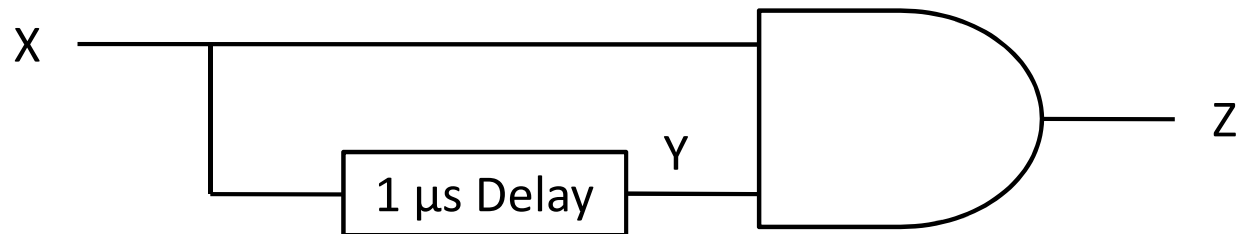
Sample Timing Diagram (1/2)

- Assume each gate has a propagation delay of 20 ns



Sample Timing Diagram (2/2)

- A circuit with delay element



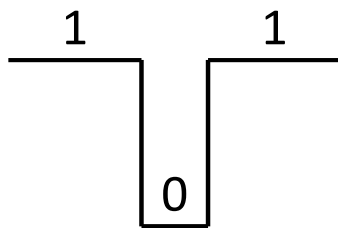
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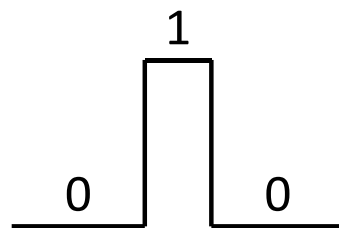
Hazards

□ Hazard

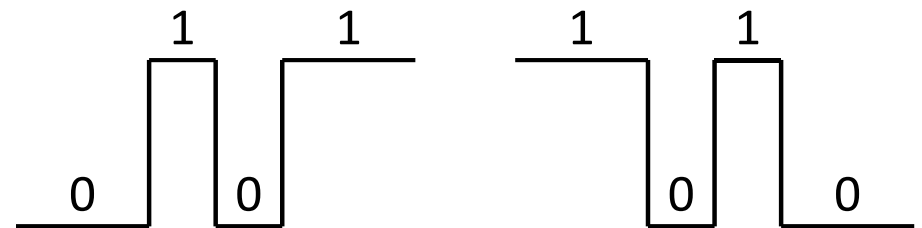
- Unwanted switching transients appearing at the output when the input to a combinational circuit changes
 - Unequal propagation delays for different paths from inputs to outputs
- Static 1-/0-hazard
 - Output momentarily goes to 0/1 when it should remain a constant 1/0
- Dynamic hazard
 - Output change 3 or more times when the output changes from 0 to 1 (1 to 0)



Static 1-Hazard



Static 0-Hazard

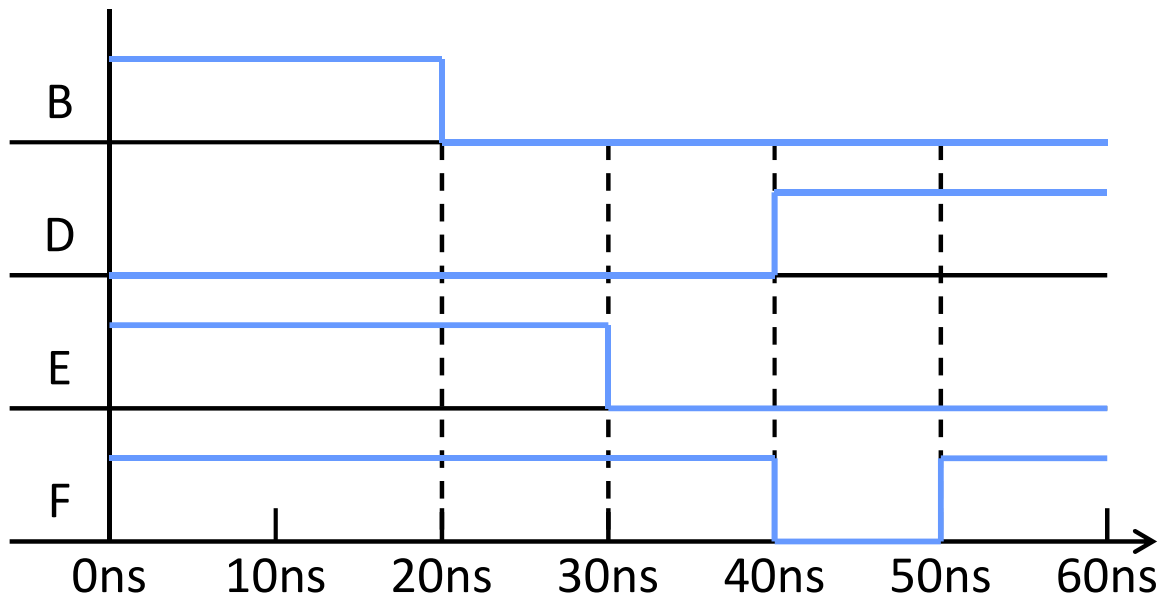
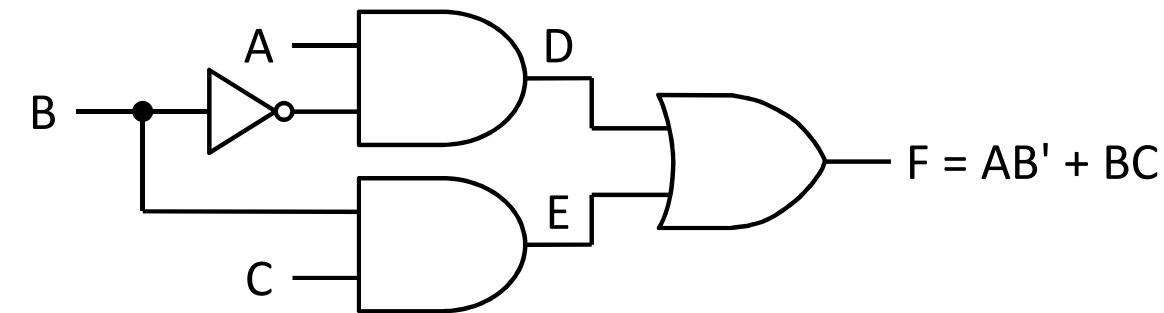


Dynamic Hazard

Static 1-Hazard Example

□ Example: assume each gate has a propagation delay of 10 ns

➤ If $A = C = 1$, static-1 hazard occurs when B changes from 1 to 0



A	0	1
BC		
00	0	1
01	0	1
11	1	1
10	0	0

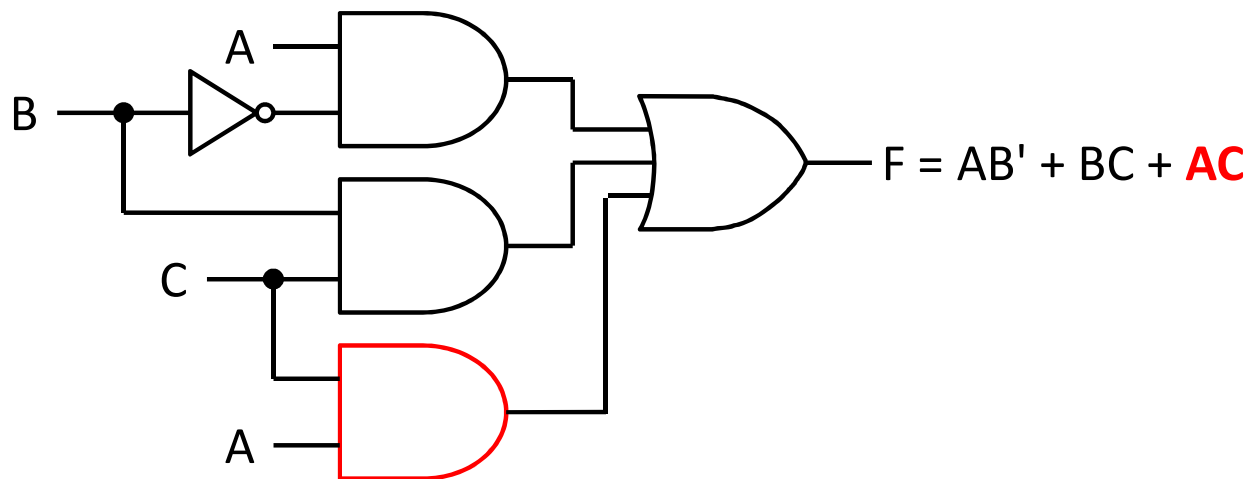
Static 1-Hazard Removal

□ When?

- In Karnaugh maps, if any two adjacent 1's are not covered by the same loop, a 1-hazard exists for the transition between the two 1's

□ How to remove hazards?

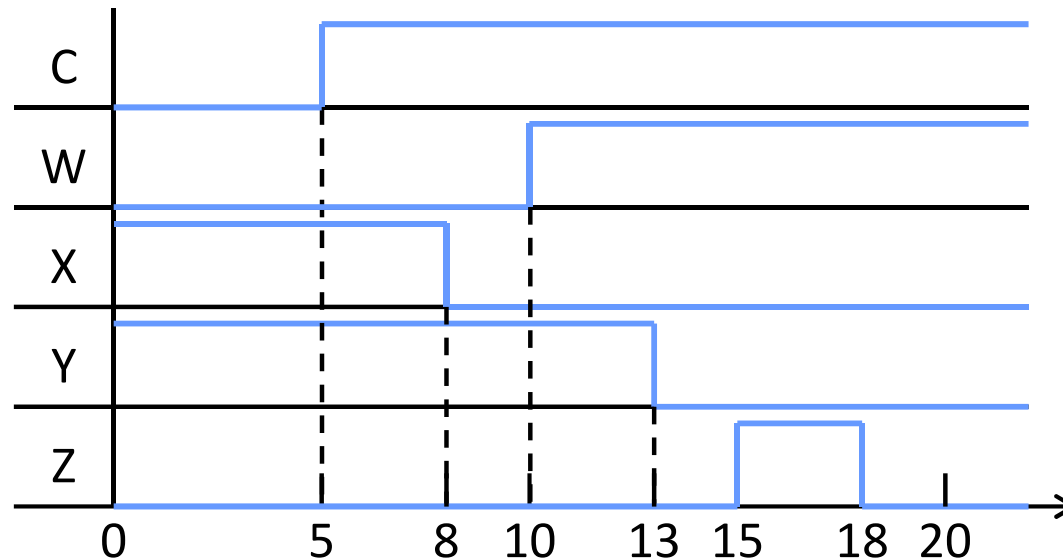
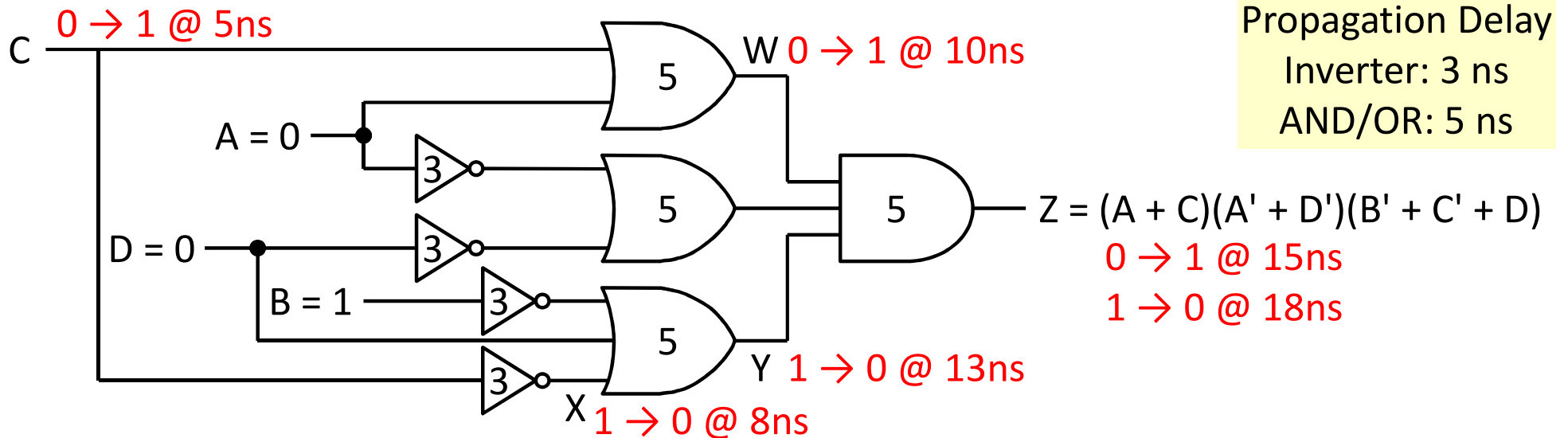
- Add additional loops such that all adjacent 1's are covered by some loop



BC \ A	0	1
00	0	1
01	0	1
11	1	1
10	0	0

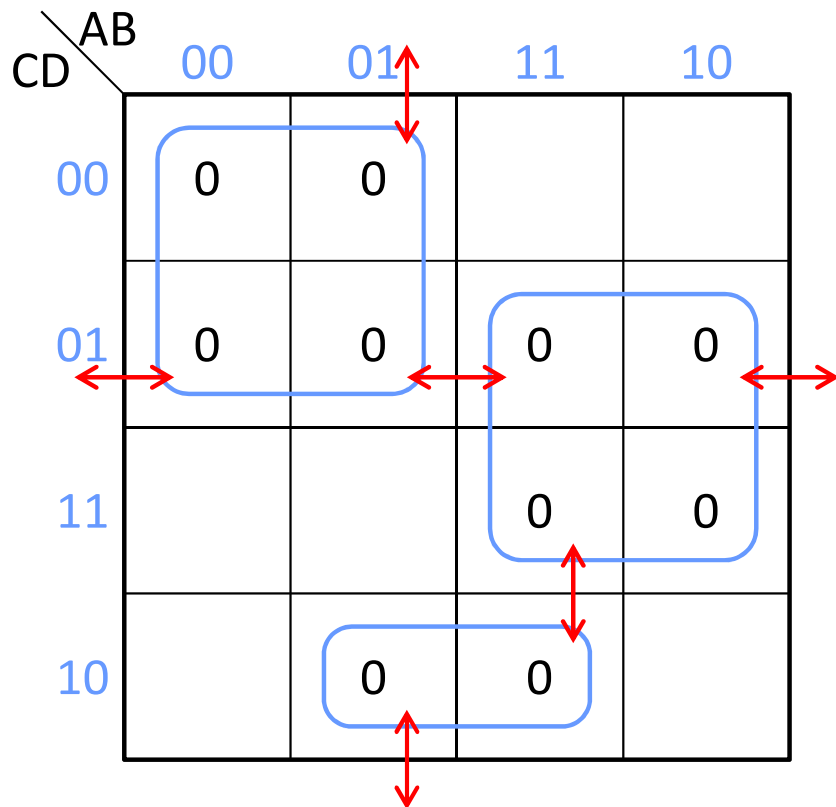
Static 0-Hazard Example

□ Example: $A = 0$, $B = 1$, $D = 0$, and C changes from 0 to 1

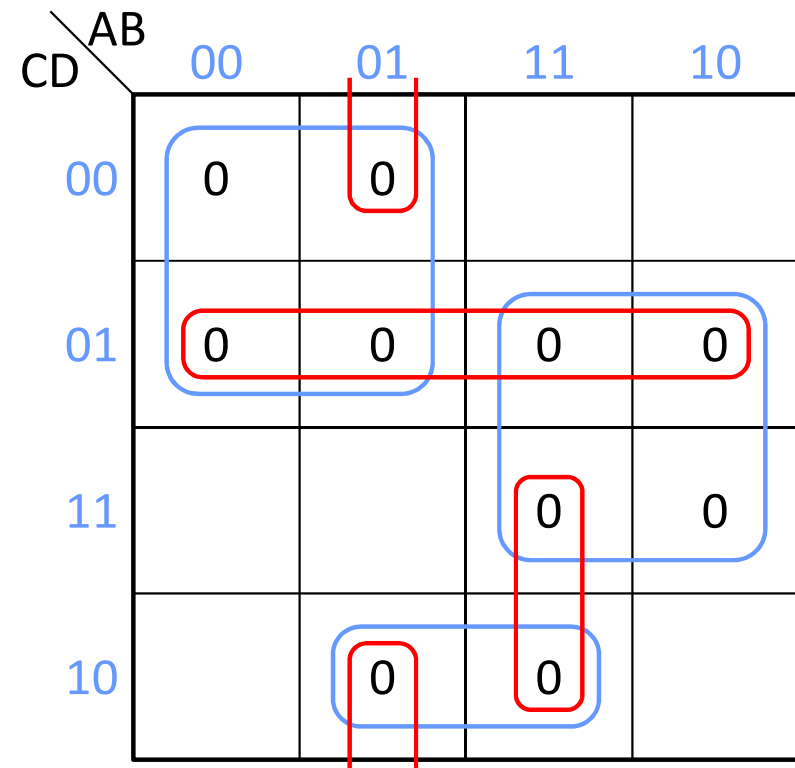


Static 0-Hazard Removal

- Similarly, add additional loops for adjacent 0's



$$Z = (A + C)(A' + D')(B' + C' + D)$$



$$Z = (A + C)(A' + D')(B' + C' + D) \\ (C + D')(A + B' + D)(A' + B' + C')$$

Stop Here (20190325)

Hazards in Multilevel Logic (1/2)

- ❑ Derive a SOP/POS for a multilevel circuit, but the complementation laws ($XX' = 0$, $X + X' = 1$) are NOT used
 - Some redundant terms exist in the SOP (POS)
 - $XX'\alpha$
 - α is a product of literals or it may be null
 - $XX'\alpha$ represents a pseudo gate that may temporarily have the output value 1 as X changes and if $\alpha = 1$
 - $X + X' + \beta$
 - β is a sum of literals or it may be empty
 - $X + X' + \beta$ represents a pseudo gate that may temporarily have the output value 0 as X changes and if $\beta = 0$
- ❑ For SOP, ensure every pair of adjacent 1's is covered together
 - The sum of all prime implicants is safe!
- ❑ Treat each X and X' as independent variables to prevent introduction of hazards

Hazards in Multilevel Logic (2/2)

❑ Static 1-hazards: analyze the same as for 2-level

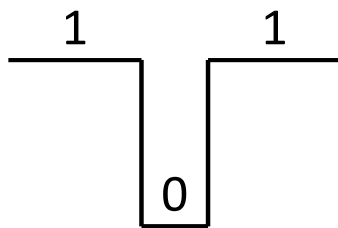
❑ Static 0-hazards: $XX'\alpha$ in SOP

➤ Adjacent 0's differ in the value of X with $\alpha = 1$

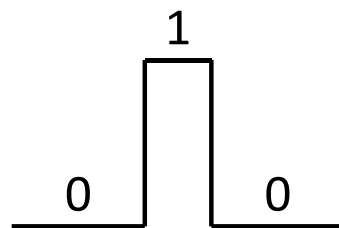
❑ Dynamic hazards: $XX'\alpha$ in SOP

➤ Adjacent input combinations on K-map differ in the value of X with $\alpha = 1$ and with opposite function values, and

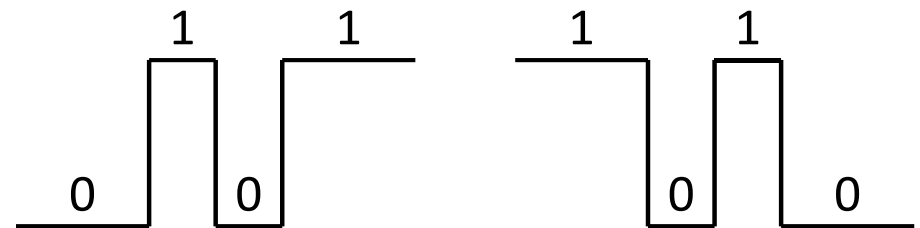
➤ For these combinations, the change in X propagates over ≥ 3 different paths



Static 1-Hazard



Static 0-Hazard



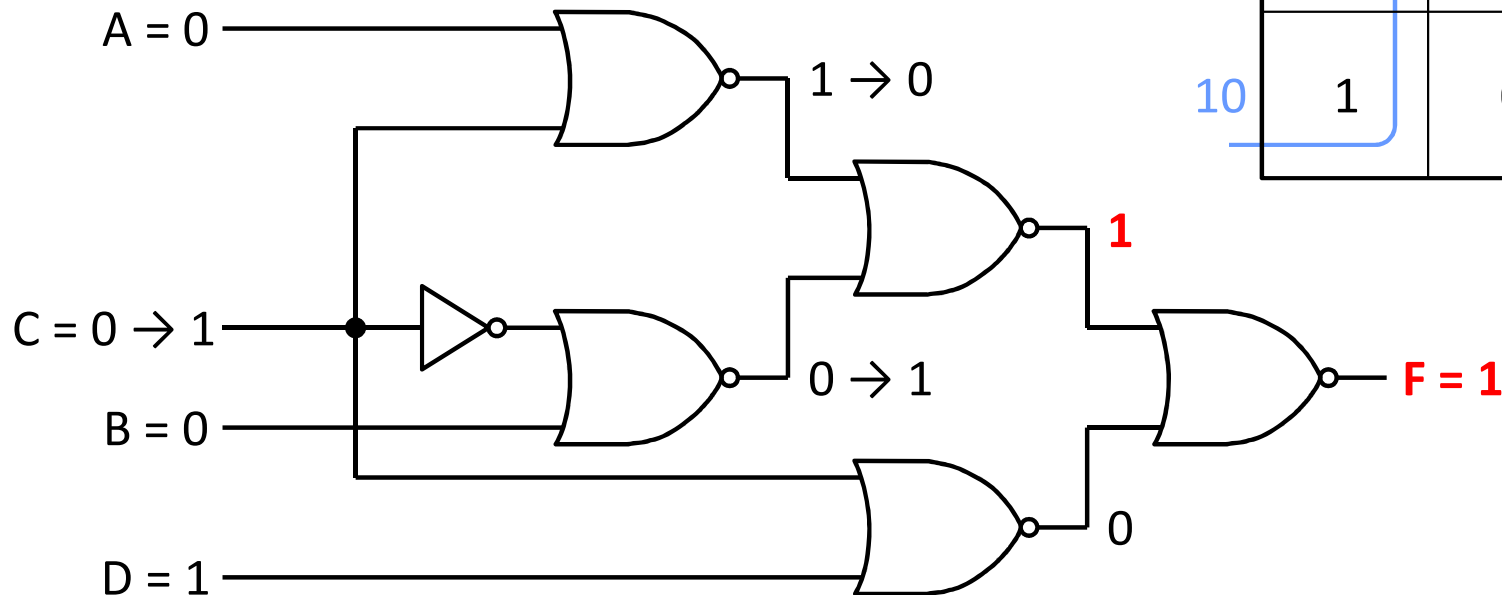
Dynamic Hazard

Example (1/3)

$$\begin{aligned} F &= (A'C' + B'C)(C + D) \\ &= A'C'C + A'C'D + B'C + B'CD \\ &= A'C'C + A'C'D + B'C \end{aligned}$$

Static 1-hazard

➤ 0001 \leftrightarrow 0011



AB \ CD	00	01	11	10
00	0	0	0	0
01	1	1	0	0
11	1	0	0	1
10	1	0	0	1

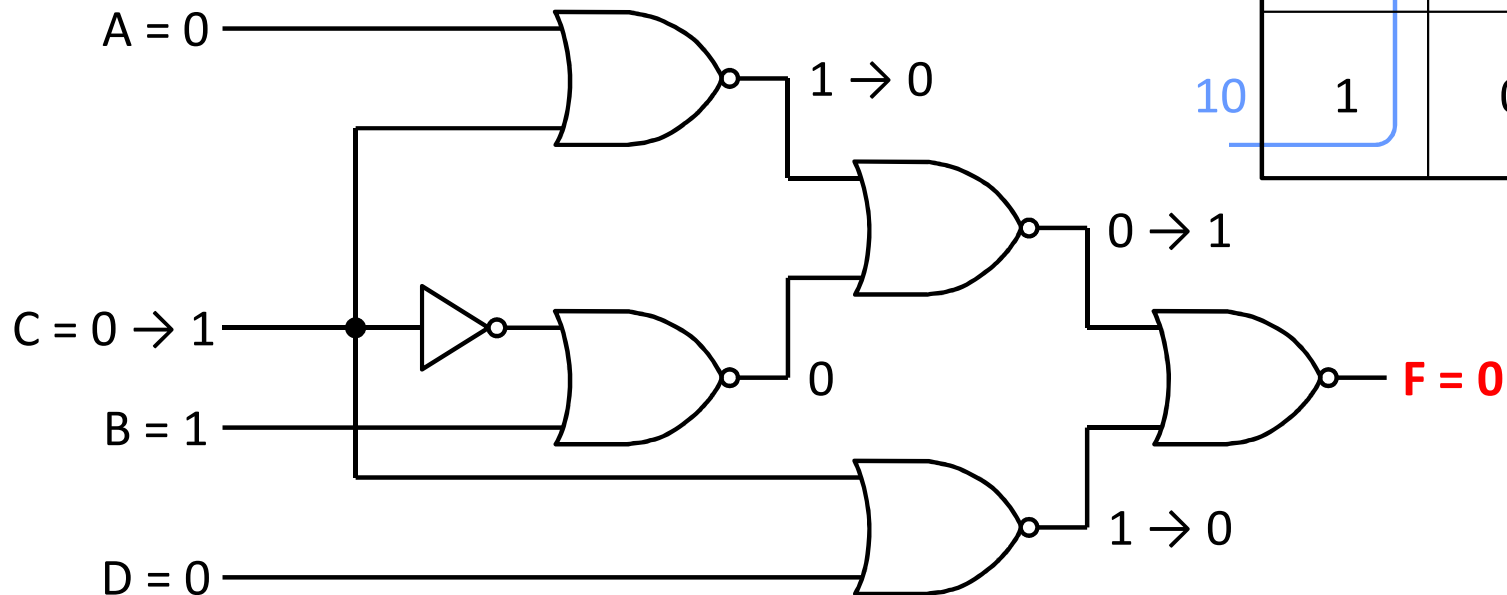
Example (2/3)

❑ $F = A'C'C + A'C'D + B'C$

❑ Static 0-hazard

- Adjacent 0's differ in the value of C with $A' = 1$
- $0100 \leftrightarrow 0110$

AB		00	01	11	10
CD	00	0	0	0	0
	01	1	1	0	0
	11	1	0	0	1
	10	1	0	0	1

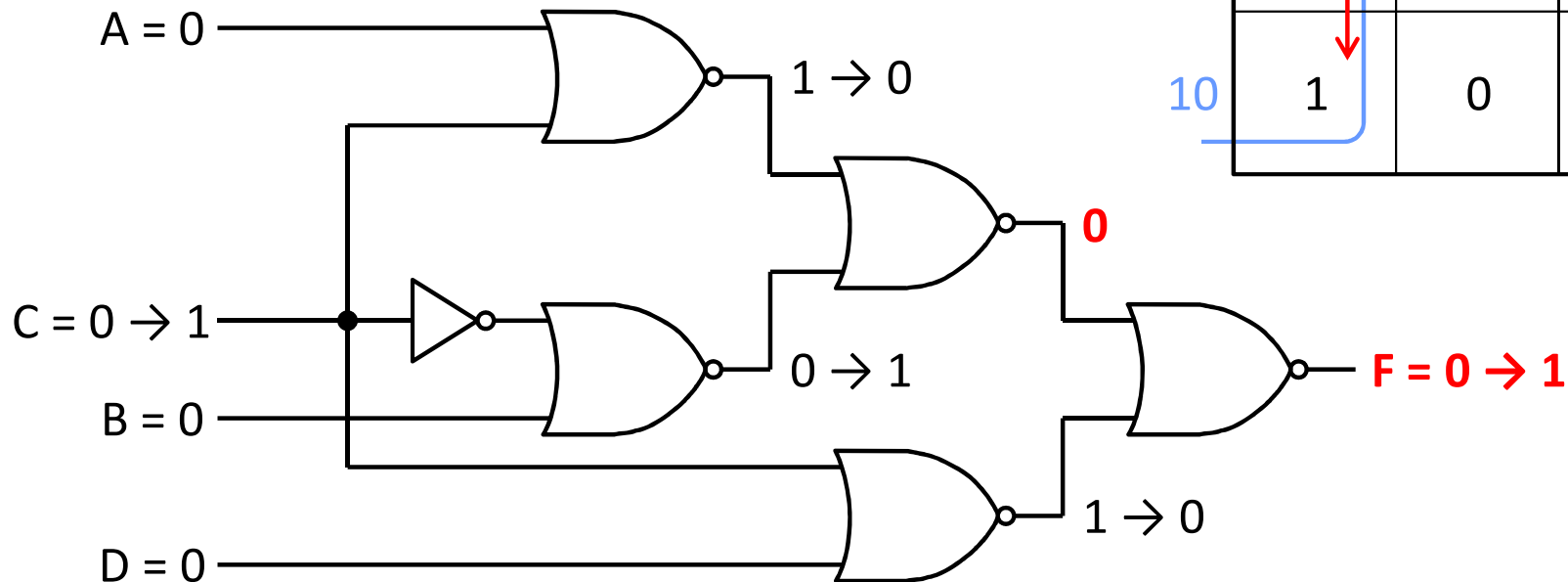


Example (3/3)

□ $F = A'C'C + A'C'D + B'C$

- Adjacent input combinations on K-map differ in the value of C with $A' = 1$ and with opposite function values, and
- For these combinations, the change in C propagates over ≥ 3 different paths
- $0000 \leftrightarrow 0010$ (how about $0101 \leftrightarrow 0111$?)

AB \ CD	00	01	11	10
00	0	0	0	0
01	1	1	0	0
11	1	0	0	1
10	1	0	0	1



Designing a Hazard-Free Circuit

□ Hazard-free AND-OR circuit

- Find a SOP expression for the output where each pair of adjacent 1's is covered by a 1-term
 - The sum of all prime implicants always satisfies this condition
 - A two-level AND-OR circuit based on this will be free of 1-, 0-, and dynamic hazards
- Transform the SOP expression into the desired form by simple factoring, DeMorgan's laws, etc.
 - Treat each variable X and X' as independent variables

□ Hazard-free OR-AND circuit

- Alternatively, start with a POS form where every pair of adjacent 0's is covered by a 0-term
- Dual procedure

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Simulation

❑ Verify logic circuits by

- Actually building them
- Simulating them on a computer

❑ 4 logic values

- 0 (low), 1 (high), X (unknown), Z (high-impedance/open circuit)

❑ AND and OR functions for 4-valued simulation

AND	0	1	X	Z
0	0	0	0	0
1	0	1	X	X
X	0	X	X	X
Z	0	X	X	X

OR	0	1	X	Z
0	0	1	X	X
1	1	1	1	1
X	X	1	X	X
Z	X	1	X	X

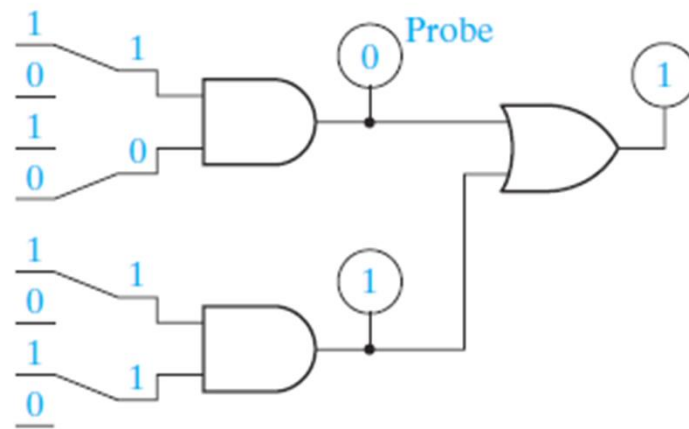
4-Valued Logic Simulator (1/2)

□ Simulate a circuit

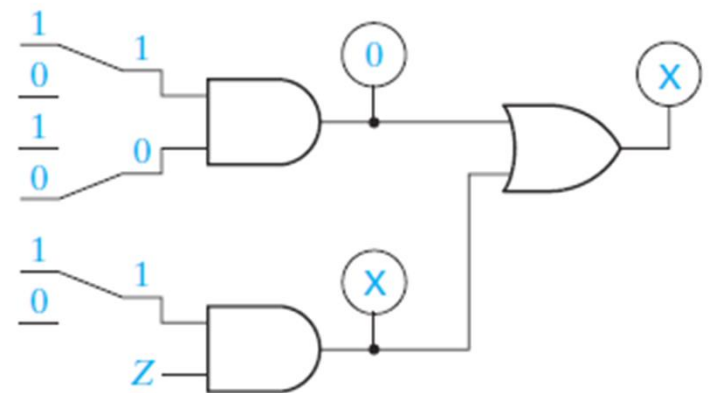
- Computations: level-by-level from inputs to outputs
- Evaluations: once inputs change

FIGURE 8-13

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(a) Simulation screen showing switches



(b) Simulation screen with missing gate input

4-Valued Logic Simulator (2/2)

❑ Possible errors in designs

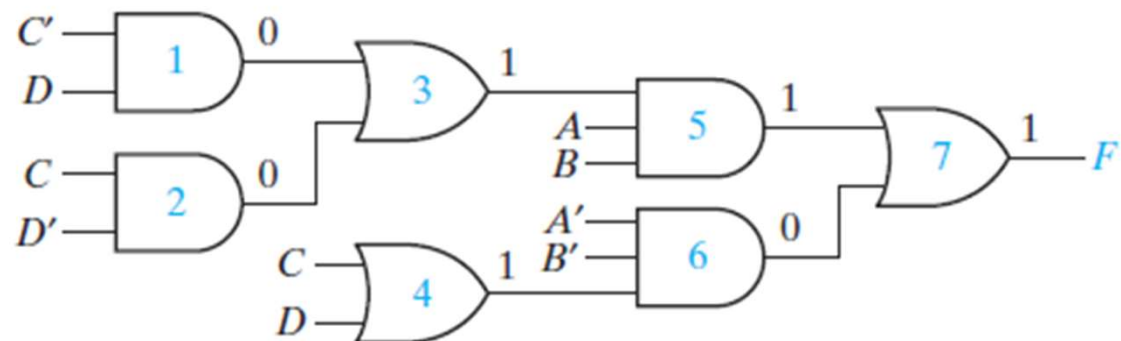
- In simulation
 - Incorrect design
 - Gates with wrong connection
 - Wrong input signals to the circuits
- In real circuits (testing)
 - Defective gates
 - Defective connecting wires

❑ Debug

- Start from output, work back until the trouble is located

FIGURE 8-14
Logic Circuit with
Incorrect Output

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Q&A