

CSIE 2344, Spring 2019 — Midterm Solution Sketch

Name: _____ SID: _____ Email: _____

1 True or False (12pts)

Determine whether the following statements are true or false by circling the correct choice. No explanation is required.

- T F 1. (1pt) Use 2's complement and 5-bit word, $-10_{10} = 10100_2$.
- T F 2. (1pt) Use 2's complement, 4-bit words, and a 4-bit parallel adder, assume A_3, B_3 are the first bits (*i.e.*, highest bits or most significant bits) of the two input words and S_3 is the first bit of the output sum, there is an overflow if and only if $(A_3 \equiv B_3) \oplus S_3 = 1$.
- T F 3. (1pt) Assume X, Y, Z are Boolean variables, $X \oplus YZ = (X \oplus Y)(X \oplus Z)$ is always true.
- T F 4. (1pt) Assume X, Y, Z are Boolean variables, $(X + Y)(X + Z) = X + YZ$ is always true.
- T F 5. (1pt) Follow the numbering method for minterms and maxterms in the lecture, $(m_i)' = M_i$.
- T F 6. (1pt) It is possible to reorder indexes (00-01-11-10 of both sides) in a Karnaugh map so that m_{15} is at the bottom-right corner and the Karnaugh map can still work.
- T F 7. (1pt) With three Boolean variables, $F = \sum m(0, 1, 2, 3)$ if and only if $F' = \prod M(4, 5, 6, 7)$.
- T F 8. (1pt) With three Boolean variables, $F_1 = \sum m(0, 1, 2, 3)$ and $F_2 = \sum m(0, 1, 4, 5)$ if and only if $F_1 \cdot F_2 = \sum m(0, 1)$.
- T F 9. (1pt) {XOR} is a functional complete set.
- T F 10. (1pt) {AND, XNOR} is a functional complete set.
- T F 11. (1pt) Given a Boolean function $F = A'B'C' + C$, $A'B'$ is a prime implicant.
- T F 12. (1pt) Given a Boolean function $F = A'B'C' + C$, $A'B'$ is an essential prime implicant.

Answer: F, F, F, T, T, T, F, F, F, T, T, T.

2 Conversion to Base 7 (8pts)

Convert 71.75_{10} to base 7.

Answer: $131.\underline{51}$ from

$$71 = 7 \cdot 10 + 1;$$

$$10 = 7 \cdot 1 + 3;$$

$$1 = 7 \cdot 0 + 1,$$

and

$$0.75 \cdot 7 = 5.25;$$

$$0.25 \cdot 7 = 1.75,$$

where there is a repeat of 0.75.

3 7-ELEVEN (8pts)

Design a combinational logic circuit which has one output F and a 4-bit input $ABCD$ representing a binary number ($ABCD = 0001$ represents 1_{10}). F should be 1 if and only if the input is at least 7 but no greater than 11 (*i.e.*, $7 \leq \text{input} \leq 11$). Use one OR gate and two AND gates.

Answer: Mark $m_7, m_8, m_9, m_{10}, m_{11}$ in a Karnaugh map, use the Karnaugh map to get the minimum sum-of-products expression $F = AB' + A'BCD$, and draw it.

4 Karnaugh Maps (16pts)

Given $F(A, B, C, D) = \sum m(3, 4, 5, 6, 7, 9, 11, 12, 15) + \sum d(2, 13)$.

1. (8pts) Find a minimum sum-of-products expression for F . Only the K-map and the final expression are required.
2. (8pts) Find a minimum product-of-sums expression for F . Only the K-map and the final expression are required.

Answer:

1. Use a Karnaugh map to get the minimum sum-of-products expression $F = A'C + AD + BC'$.
2. Use a Karnaugh map to get the minimum product-of-sums expression $F' = B'D' + A'B'C' + ACD'$ and $F = (B + D)(A + B + C)(A' + C' + D)$.

5 Static Hazards (16pts)

Given $F(A, B, C, D) = A'C' + AD + BCD'$.

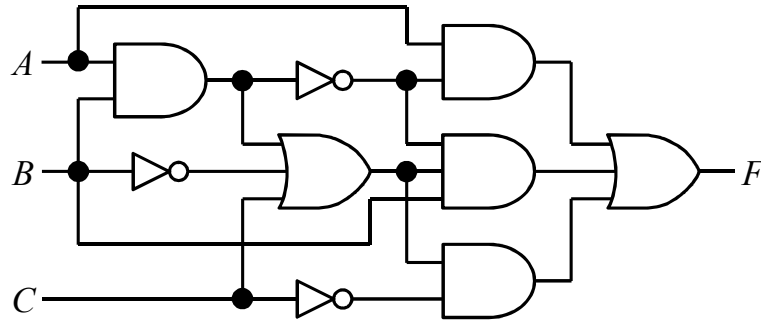
1. (8pts) Draw a minimum two-level AND-OR gate circuit (AND at the second level, closer to the inputs, and OR at the first level, closer to the output) for F without static-1 hazards.
2. (8pts) Draw a minimum two-level AND-OR gate circuit for F without static-0 hazards.

Answer:

1. Use a Karnaugh map to get $F = A'C' + AD + BCD' + C'D + ABC + A'BD'$ so that any pair of neighboring 1's is covered by a term and draw it.
2. Use a Karnaugh map to get the minimum sum-of-products expression $F = A'C' + AD + BCD'$ and draw it. Note that you do not need to add any additional gate.

6 Circuit Conversion (16pts)

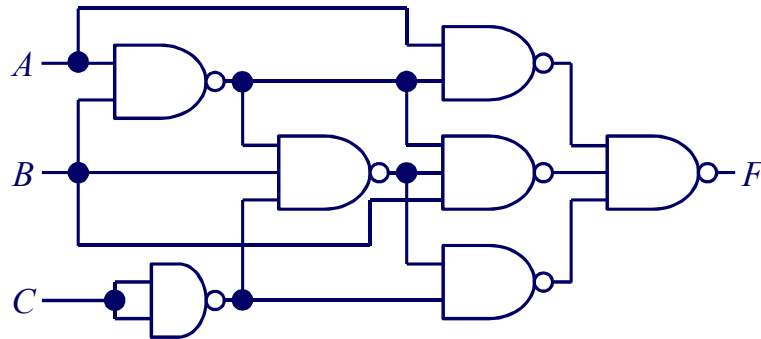
- (8pts) Use gate equivalences to convert the following circuit into a four-level circuit containing only NAND gates (NOT gates are not allowed) and circuit inputs A, B, C (A', B', C' are not allowed as circuit inputs). To get the all points, the number of NAND gates should be 7. No explanation is required.



- (8pts) Convert the circuit to a minimum three-level 2-input or 3-input NAND gate circuit where A', B', C' are ALLOWED as circuit inputs, but A, A' are not allowed as the inputs of third-level gates. (**Warning: this one may take some time!**)

Answer:

- The circuit is shown below:



- Simply F and use a Karnaugh map to get $F = AB' + AC' + B'C' + A'BC = A(B' + C') + B'C' + A'BC$, so the circuit is the corresponding one of $F = \{[A(BC)']'[B'C']'[A'BC]'\}'$.

7 Proving DeMorgan's Law (8pts)

Assume X_1, X_2, \dots, X_n are Boolean variables, prove $(X_1 + X_2 + \dots + X_n)' = X_1' X_2' \dots X_n'$ for $n \geq 1$. You should not assume that DeMorgan's Laws are true before proving them. (Hint: how to prove it for any positive integer n ?)

Answer: Proof by induction:

When $n = 1$, LHS = X_1' and RHS = X_1' , which are equivalent.

When $n = 2$, using truth tables, LHS = $(X_1 + X_2)'$, which is 1 if and only if $X_1 = X_2 = 0$, and RHS = $X_1' X_2'$, which is 1 if and only if $X_1 = X_2 = 0$, so they are equivalent.

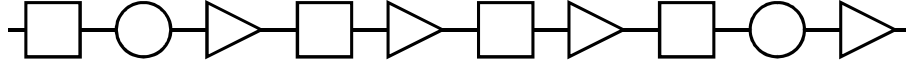
Assume that, when $n = k$, $(X_1 + X_2 + \dots + X_k)' = X_1' X_2' \dots X_k'$.

When $n = k+1$, $(X_1 + X_2 + \dots + X_k + X_{k+1})' = ((X_1 + X_2 + \dots + X_k) + X_{k+1})' = (X_1 + X_2 + \dots + X_k)' X_{k+1}' = (X_1' X_2' \dots X_k') X_{k+1}' = X_1' X_2' \dots X_k' X_{k+1}'$.




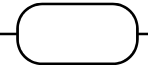



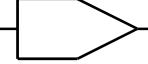



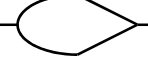


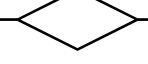







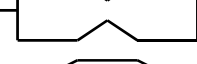



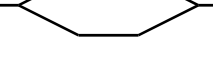
By induction, the claim is proved.

8 Logic Minimization (16pts)

Assume that each shape is one type of gates, and the input and output lines of each gate are 2019 bits. The goal of this problem is to convert the following sequence of gates into another sequence of gates so that the total cost is minimized.



The rules of conversion are shown below. For example, if you do not convert anything, the total cost is $6 + 3 + 2 + 6 + 2 + 6 + 2 + 6 + 3 + 2 = 38$. If you convert the last two gates (circle and triangle) by the rule #3, the total cost is $6 + 3 + 2 + 6 + 2 + 6 + 2 + 6 + 4 = 37$. Note that you can choose not to convert a gate if the total cost can be minimized. Also, a rule of conversion can be applied more than once.

Gate	Cost	Cost Before Conversion	Rule of Conversion	Cost after Conversion
	6	9 (6+3)	  \rightarrow #1 	7
	3	8 (6+2)	  \rightarrow #2 	7
	2	5 (3+2)	  \rightarrow #3 	4
		8 (2+6)	  \rightarrow #4 	6
		11 (6+3+2)	   \rightarrow #5 	8
		14 (6+2+6)	   \rightarrow #6 	12
		10 (2+6+2)	   \rightarrow #7 	7

- (4pts) Write down the minimum total cost after converting the given sequence. No explanation is required. No partial credit will be given.
- (4pts) Write down the applied rules of conversion from left of the sequence to right of the sequence. No explanation is required. No partial credit will be given.
- (8pts) Explain(*e.g.*, write down Pseudocode) how to systematically and efficiently convert any sequence of gates (square, circle, and triangle) into another sequence of gates by the rules of conversion above and minimize the cost.

Answer:

- (4pts) 28.
- (4pts) #1,#4,#7,#5.

3. (8pts) Omitted, but it can be solved by dynamic programming. For each gate from left to right, compute possible conversions and pick the solution with the minimum total cost until that gate. The following is the process of the example question here.

Gate	1	2	3	4	5	6	7	8	9	10
Rule	Keep	Keep	Keep	Keep	Keep	Keep	Keep	Keep	Keep	Keep
Cost	6	6 + 3	7 + 2	8 + 6	13 + 2	14 + 6	19 + 2	20 + 6	25 + 6	27 + 2
Rule		#1	#3	#4	#2	#4	#2	#4	#1	#3
Cost		7	6 + 3	7 + 6	8 + 7	13 + 6	14 + 7	19 + 6	20 + 7	25 + 4
Rule			#5		#5	#6	#7	#6		#5
Cost			8		7 + 7	8 + 12	13 + 7	14 + 12		20 + 8
Minimum	6	7	8	13	14	19	20	25	27	28