

Logic Coursework 2024/25: Written Work

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Question 1

Answer the following questions about complete sets of logical connectives, in each case justifying your answer.

- (i) Show $\{\neg, \rightarrow\}$ is a complete set of connectives.
- (ii) Show $\{\rightarrow, 0\}$ is a complete set of connectives (where 0 is the constant false).
- (iii) Is $\{\text{NAND}, \wedge\}$ a complete set of connectives?
- (iv) Is $\{\wedge, \vee\}$ a complete set of connectives?

Answer. In order to determine whether the sets are complete, I will be showing whether \wedge, \vee and \neg can be expressed using the connectives in the set, in which case any logical expression can be written in CNF or DNF, meaning that it's a complete set.

Proof for part (i).

Can \vee be expressed using $\{\neg, \rightarrow\}$?

Yes: $\neg p \rightarrow q \equiv \neg(\neg p) \vee q \equiv p \vee q$

Can \wedge be expressed using $\{\neg, \rightarrow\}$?

Yes: $\neg(p \rightarrow \neg q) \equiv \neg(\neg p \vee \neg q) \equiv p \wedge q$

Since \neg is already in our set of logical connectives, we can then conclude that $\{\neg, \rightarrow\}$ is a complete set of logical connectives, as any logical expression can be expressed in CNF/DNF using the connectives within the set. \square

Proof for part (ii).

Can \neg be expressed using $\{\rightarrow, 0\}$?

Yes: $p \rightarrow 0 \equiv \neg p \vee 0 \equiv \neg p$

Can \vee be expressed using $\{\rightarrow, 0\}$?

Yes: $(p \rightarrow 0) \rightarrow q \equiv \neg(p \rightarrow 0) \vee q \equiv \neg(\neg p) \vee q \equiv p \vee q$

Can \wedge be expressed using $\{\rightarrow, 0\}$?

Yes: $(p \rightarrow (q \rightarrow 0)) \rightarrow 0 \equiv (p \rightarrow (\neg q \vee 0)) \rightarrow 0 \equiv (p \rightarrow \neg q) \rightarrow 0 \equiv (\neg p \vee \neg q) \rightarrow 0 \equiv \neg(\neg p \vee \neg q) \vee 0 \equiv p \wedge q$

Therefore $\{\rightarrow, 0\}$ is a complete set of logical connectives \square

Proof for part (iii).

To denote NAND, I will use the symbol: $\bar{\wedge}$

Can \neg be expressed using $\{\bar{\wedge}, \wedge\}$?

Yes: $p \bar{\wedge} p \equiv \neg(p \wedge p) \equiv \neg p$

Can \vee be expressed using $\{\bar{\wedge}, \wedge\}$?

Yes: $(p \bar{\wedge} p) \bar{\wedge} (q \bar{\wedge} q) \equiv \neg p \bar{\wedge} \neg q \equiv \neg(\neg p \wedge \neg q) \equiv p \vee q$

\wedge is already in the set of logical connectives, therefore the set of logical connectives $\{\text{NAND}, \wedge\}$ is complete. \square

Proof for part (iv).

The set of logical connectives $\{\wedge, \vee\}$ is not complete as there is no way to represent one of the propositional variables being equal to zero (i.e: negation) when writing an expression in CNF/DNF. That is to say there is no way to express a tautology or contradiction using these logical connectives due to their property of idempotence. \square

Question 2

Convert $((p \rightarrow q) \rightarrow r) \rightarrow (s \rightarrow t)$ to

- (i) Conjunctive Normal Form (CNF)
- (ii) Disjunctive Normal Form (DNF)

Answer.

This question is easier to approach by writing the expression $\varphi = ((p \rightarrow q) \rightarrow r) \rightarrow (s \rightarrow t)$ in DNF first:

$$\begin{aligned}
 \varphi &\equiv (((\neg p \vee q) \rightarrow r) \rightarrow (s \rightarrow t)) \\
 &\equiv ((\neg(\neg p \vee q) \vee r) \rightarrow (s \rightarrow t)) \\
 &\equiv (((p \wedge \neg q) \vee r) \rightarrow (s \rightarrow t)) \\
 &\equiv (\neg((p \wedge \neg q) \vee r) \vee (s \rightarrow t)) \\
 &\equiv ((\neg(p \wedge \neg q) \wedge \neg r) \vee (s \rightarrow t)) \\
 &\equiv (((\neg p \vee q) \wedge \neg r) \vee (s \rightarrow t)) \\
 &\equiv (((\neg p \wedge \neg r) \vee (q \wedge \neg r)) \vee (s \rightarrow t)) \\
 &\equiv ((\neg p \wedge \neg r) \vee (q \wedge \neg r) \vee (\neg s \vee t)) \\
 &\equiv (\neg p \wedge \neg r) \vee (q \wedge \neg r) \vee \neg s \vee t \\
 \therefore \varphi_{DNF} &= (\neg p \wedge \neg r) \vee (q \wedge \neg r) \vee \neg s \vee t
 \end{aligned}$$

Using this, we can repeatedly use the distributive property to convert this to conjunctive normal form:

$$\begin{aligned}
 \varphi_{DNF} &= (\neg p \wedge \neg r) \vee (q \wedge \neg r) \vee \neg s \vee t \\
 &\equiv ((\neg p \vee q) \wedge \neg r) \vee \neg s \vee t \\
 &\equiv (((\neg p \vee q) \vee \neg s) \wedge (\neg r \vee \neg s)) \vee t \\
 &\equiv (\neg p \vee q \vee \neg s \vee t) \wedge (\neg r \vee \neg s \vee t) \\
 \therefore \varphi_{CNF} &= (\neg p \vee q \vee \neg s \vee t) \wedge (\neg r \vee \neg s \vee t)
 \end{aligned}$$

$$(i) \varphi_{CNF} = (\neg p \vee q \vee \neg s \vee t) \wedge (\neg r \vee \neg s \vee t)$$

$$(ii) \varphi_{DNF} = (\neg p \wedge \neg r) \vee (q \wedge \neg r) \vee \neg s \vee t$$

Question 3

What is the purpose of Tseitin's Algorithm? Apply Tseitin's Algorithm to turn the propositional formula $((x_1 \wedge x_2 \wedge x_3) \rightarrow (y_1 \wedge y_2 \wedge y_3)) \vee z$ to CNF.

Answer.

The purpose of Tseitin's Algorithm is to take an arbitrary propositional formula φ , and transform it to a new propositional formula φ' which is equisatisfiable with φ , and in conjunctive normal form.

Let $\varphi = (((x_1 \wedge x_2 \wedge x_3) \rightarrow (y_1 \wedge y_2 \wedge y_3)) \vee z)$

Introduce new variables for each subformula:

$$\alpha_1 \leftrightarrow x_1 \wedge x_2 \wedge x_3$$

$$\alpha_2 \leftrightarrow y_1 \wedge y_2 \wedge y_3$$

$$\alpha_3 \leftrightarrow \alpha_1 \rightarrow \alpha_2$$

$$\alpha_4 \leftrightarrow \alpha_3 \vee z$$

Write each expression as conjunctions

From α_1 :

$$\begin{aligned} \alpha_1 \leftrightarrow (x_1 \wedge x_2 \wedge x_3) &\equiv (\alpha_1 \rightarrow (x_1 \wedge x_2 \wedge x_3)) \wedge (\alpha_1 \leftarrow (x_1 \wedge x_2 \wedge x_3)) \\ &\equiv (\neg \alpha_1 \vee (x_1 \wedge x_2 \wedge x_3)) \wedge (\alpha_1 \vee \neg(x_1 \wedge x_2 \wedge x_3)) \\ &\equiv (\neg \alpha_1 \vee x_1) \wedge (\neg \alpha_1 \vee x_2) \wedge (\neg \alpha_1 \vee x_3) \wedge (\alpha_1 \vee \neg x_1 \vee \neg x_2 \vee \neg x_3) \end{aligned}$$

Similarly for α_2 :

$$\alpha_2 \leftrightarrow (y_1 \wedge y_2 \wedge y_3) \equiv (\neg \alpha_2 \vee y_1) \wedge (\neg \alpha_2 \vee y_2) \wedge (\neg \alpha_2 \vee y_3) \wedge (\alpha_2 \vee \neg y_1 \vee \neg y_2 \vee \neg y_3)$$

For α_3 :

$$\begin{aligned} \alpha_3 \leftrightarrow (\alpha_1 \rightarrow \alpha_2) &\equiv (\alpha_3 \rightarrow (\alpha_1 \rightarrow \alpha_2)) \wedge (\alpha_3 \leftarrow (\alpha_1 \rightarrow \alpha_2)) \\ &\equiv (\neg \alpha_3 \vee (\neg \alpha_1 \vee \alpha_2)) \wedge (\neg(\neg \alpha_1 \vee \alpha_2) \vee \alpha_3) \equiv (\neg \alpha_3 \vee \neg \alpha_1 \vee \alpha_2) \wedge ((\alpha_1 \wedge \neg \alpha_2) \vee \alpha_3) \\ &\equiv (\neg \alpha_1 \vee \alpha_2 \vee \neg \alpha_3) \wedge (\alpha_1 \vee \alpha_3) \wedge (\neg \alpha_2 \vee \alpha_3) \end{aligned}$$

For α_4 :

$$\begin{aligned} \alpha_4 \leftrightarrow (\alpha_3 \vee z) &\equiv (\alpha_4 \rightarrow (\alpha_3 \vee z)) \wedge (\alpha_4 \leftarrow (\alpha_3 \vee z)) \\ &\equiv (\neg \alpha_4 \vee \alpha_3 \vee z) \wedge (\neg(\alpha_3 \vee z) \vee \alpha_4) \equiv (\neg \alpha_4 \vee \alpha_3 \vee z) \wedge ((\neg \alpha_3 \wedge \neg z) \vee \alpha_4) \\ &\equiv (\neg \alpha_4 \vee \alpha_3 \vee z) \wedge (\neg \alpha_3 \vee \alpha_4) \wedge (\neg z \vee \alpha_4) \end{aligned}$$

The conjunction of all these variables and the clause α_4

gives us the Tseitin Transformation of φ'

To save space, I will write this as a clause set

$$\begin{aligned} \therefore \varphi' = \{ &\{\alpha_4\}, \{\neg \alpha_1, x_1\}, \{\neg \alpha_1, x_2\}, \{\neg \alpha_1, x_3\}, \{\alpha_1, \neg x_1, \neg x_2, \neg x_3\}, \\ &\{\neg \alpha_2, y_1\}, \{\neg \alpha_2, y_2\}, \{\neg \alpha_2, y_3\}, \{\alpha_2, \neg y_1, \neg y_2, \neg y_3\}, \\ &\{\neg \alpha_1, \alpha_2, \neg \alpha_3\}, \{\alpha_1, \alpha_3\}, \{\neg \alpha_2, \alpha_3\}, \{\neg \alpha_4, \alpha_3, z\}, \{\neg \alpha_3, \alpha_4\}, \{\neg z, \alpha_4\} \} \end{aligned}$$

Question 4

State with justification if each of the following sentences of predicate logic is logically valid.

- (i) $(\forall x \exists y \forall z (E(x, y) \wedge E(y, z))) \rightarrow (\forall x \forall z \exists y (E(x, y) \wedge E(y, z)))$
- (ii) $(\forall x \exists y \exists u \forall v (E(x, y) \wedge E(u, v))) \rightarrow (\exists u \forall v \forall x \exists y (E(x, y) \wedge E(u, v)))$
- (iii) $(\forall x \exists y \forall z R(x, y, z)) \rightarrow (\exists x \forall y \exists z R(x, y, z))$
- (iv) $(\forall x \forall y \exists z (E(x, y) \wedge E(y, z))) \rightarrow (\forall x \forall y \forall z (E(x, y) \vee E(y, z)))$

Answer.

(i) Logically Valid.

Proof.

□

(ii) Logically Valid.

Proof.

□

(iii) Logically Invalid.

Proof. In order to demonstrate that this is a logically invalid sentence, I will provide a counter-model: Consider R a ternary relation over the domain $\{0, 1\}$, where

$$R := \{(0, 0, 0), (0, 0, 1), (1, 0, 0), (1, 0, 1)\}$$

From this we can see, that for all values of x , there in fact is a y , for all z . However there doesn't exist an x for all values of y , where there exists a z . This is because y only appears as a 0 in the relation, hence y cannot take each value in the domain in this relation.

□

(iv) Logically Valid.

Proof.

□

Question 5

Evaluate the given sentence on the respective relation E over domain $\{0, 1, 2\}$ with relation $E := \{(0, 1), (1, 0), (1, 2), (2, 1), (2, 0), (0, 2)\}$

- (i) $\forall x \forall y \forall z \exists w (E(x, w) \wedge E(y, w) \wedge E(z, w))$
- (ii) $\exists x \forall y \forall z \exists w (E(x, w) \wedge E(y, w) \wedge E(z, w))$
- (iii) $\forall y \exists x \forall z \exists w (E(x, w) \wedge E(y, w) \wedge E(z, w))$
- (iv) $\exists x \exists y \exists z \forall w (E(x, w) \wedge E(y, w) \wedge E(z, w))$
- (v) $\forall x_1 \exists x_2 \forall y_1 \exists y_2 \forall z_1 \exists z_2 \forall z \exists w (E(x_1, x_2) \wedge E(x_2, w) \wedge E(y_1, y_2) \wedge E(y_2, w) \wedge E(z_1, z_2) \wedge E(z_2, w) \wedge E(z, w))$
- (vi) $\forall x_1 \exists x_2 \forall y_1 \exists y_2 \forall z_1 \forall z \exists z_2 \exists w (E(x_1, x_2) \wedge E(x_2, w) \wedge E(y_1, y_2) \wedge E(y_2, w) \wedge E(z_1, z_2) \wedge E(z_2, w) \wedge E(z, w))$

Answer.

For the purposes of this question, it will be easier to express the relation E as such:

$$E := \{(u, v) \in \{0, 1, 2\} : u \neq v\}$$

(i) This sentence can be refuted by assigning the following variables:

$$x = 0, y = 1, z = 2$$

That is to say, the expression inside the brackets requires that $x \neq w, y \neq w, z \neq w$. However if we assign x, y, z uniquely, by the pigeonhole principle, it follows that there will exist no such w , that satisfies the sentence.

(ii) This sentence can be refuted by considering the fact that if x, y, z are unique, then there will exist no w , such that the conjunction of the atomic formulae $E(x, w), E(y, w), E(z, w)$ are satisfied. From the domain of the interpretation, we can deduce that there is no such value of x .

(iii) True. (iv) True. (v) (vi)