

Reasoning, Games, Action and Rationality

Lecture 2

ESSLLI'08, Hamburg

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Plan for Today

- ▶ Logic for Knowledge, beliefs and preferences in games.
- ▶ Expectation 1: Common knowledge of rationality

Logics, logic and logic

- ▶ Logic
- ▶ is
- ▶ great.

break

Expectation 1: Rationality and common belief of rationality

- ▶ What happens if all players are rational, believe that all players are rational, believe that all players believe that (...)?

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- ▶ What happens if all players are rational, believe that all players are rational, believe that all players believe that (...)?
- ▶ “Classical” assumption about game-theoretic analysis. See e.g. Myerson (1991).

Example

	A	B
a	1, 2	0, 1
b	0, 1	1, 0

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	A	B
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- B is a bad strategy for Bob.

Example

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- ▶ B is a bad strategy for Bob.
- ▶ It is *never* rational for him to choose B.

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$$EV_{t_B}(B) \geq EV_{t_B}(A)$$

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- A type t_B of Bob would be rational in choosing B iff:

$$v_{Bob}(aB)\lambda_{Bob}(t_{Bob})(a) + v_{Bob}(bB)\lambda_{Bob}(t_{Bob})(b) \geq \\ v_{Bob}(aA)\lambda_{Bob}(t_{Bob})(a) + v_{Bob}(bA)\lambda_{Bob}(t_{Bob})(b)$$

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- A type t_B of Bob would be rational in choosing B iff:

$$1\lambda_{Bob}(t_{Bob})(a)+0\lambda_{Bob}(t_{Bob})(b) \geq 2\lambda_{Bob}(t_{Bob})(a)+1\lambda_{Bob}(t_{Bob})(b)$$

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- A type t_B of Bob would be rational in choosing B iff:

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- A type t_B of Bob would be rational in choosing B iff:

$$0 \geq \lambda_{Bob}(t_{Bob})(a) + \lambda_{Bob}(t_{Bob})(b)$$

$$\text{But } \lambda_{Bob}(t_{Bob})(a) + \lambda_{Bob}(t_{Bob})(b) = 0!$$

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- ▶ Bob never plays **B** at state (σ, t) if he is rational at that state.
- ▶ But then if Ann's type at that state believes that Bob is rational, that type must assign probability 1 to Bob playing **A**.
- ▶ Given this belief, **a** is her only rational strategy.

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- ▶ If Ann and Bob are rational, and Ann believes that Bob is rational at state (σ, t) , then $\sigma = aA$.
- ▶ This strategy profile is the only one that survives *iterated elimination of strictly dominated strategies*.

Strictly dominated strategies

Definition

A strategy s_i is *strictly dominated* by another strategy s'_i iff for all combinations of choices of the other players σ_{-i} :

$$v_i(s_i, \sigma_{-i}) < v_i(s'_i, \sigma_{-i})$$

Iterated elimination of strictly dominated strategies

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	A
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Common knowledge of rational and elimination of strictly dominated strategies

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Common knowledge of rational and elimination of strictly dominated strategies

- ▶ If Ann and Bob are rational, and Ann believes that Bob is rational at state (σ, t) , then $\sigma = aA$.
- ▶ For this game we need rationality and only one level of higher-order information to conclude that aA will be played. But in the general case:

Theorem

For any state (σ, t) of a type structure for an arbitrary finite game \mathbb{G} , if all players are rational and it is common belief that all players are rational at (σ, t) , then σ is a iteratively non-dominated strategy profile.

A. Brandenburger and E. Denkel. *Rationalizability and correlated equilibria. Econometrica*, 55:13911402, 1987.

Comments on characterization results

- ▶ If [such and such expectations] at state (σ, t) , then [such and such *solution concept*] is played at that state.

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- ▶ *Solution concepts*: proposal as to what is *rational* in a game. (traditionally)
- ▶ What about the converse?
 - If [such and such *solution concept*] then one can build a state in a model such that [such and such expectations] hold.

Alternative (Dynamic) Characterization

J. van Benthem. Rational. *Rational dynamic and epistemic logic in games*. In S. Vannucci, editor, *Logic, Game Theory and Social Choice III*, 1923. University of Siena, department of political economy, 2003..

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- ▶ Recover the *dynamic* aspect of elimination of strictly dominated strategies.
- ▶ Understand how common knowledge of rationality *arises*.

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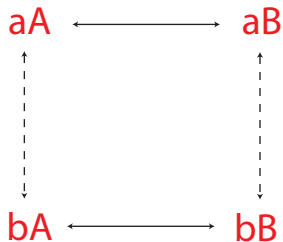
1. Consider only hard information (knowledge).
2. Work with *full models*:
 - ▶ 1 to 1 correspondence between the profiles and the states.
 - ▶ $\sigma_i(w) = \sigma_i(w')$ if and only if $w \sim_i w'$.

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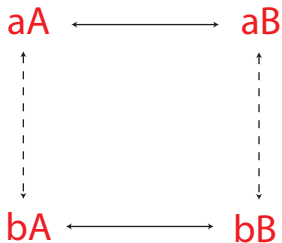


Announcing rationality

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aA



bA

After announcing that everybody is rational

Announcing rationality

- ▶ van Benthem's proposal: look at what happens when the players start to publicly announce that they are rational.

aA

After announcing once more that everybody is rational.

Announcing rationality

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Announcing rationality

- ▶ Again, here we only needed two announcements to reach the state where aA is played. In general:

Theorem

(van Benthem, 2003) *The following are equivalent for states w in full game models:*

1. $\sigma_i(w)$ is a iteratively non-dominated strategy for all i .
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Remarks and open questions

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- ▶ How to reach common knowledge.
 - ▶ This is an *if and only if*. Why?
 - Full models?
 - ▶ What happens in general?

Tomorrow

- ▶ Hard knowledge and Nash equilibrium.
- ▶ Prior beliefs, correlated beliefs and Nash equilibrium.