

# Discrete Math Notes

B.Tech. CSE

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# 1 UNIT 1

## 1.1 Set Theory

Schaum series- Lipschitz

### 1.1.1 Sets

Sets are well defined collection of mathematical objects.

*Example:*

The collection of best mathematicians in the world is not a set as there is no fixed criteria for being the best mathematicians.

Notation: Sets are denoted by capital letters such as  $A, B, X, Y$ .  
the elements are denoted by small letters such as  $a, b, x, y$ .

Def<sup>n</sup> :

A set  $A$  is called to be a subset of  $B$  iff

$$a \in A \implies a \in B$$

It is denoted by  $A \subseteq B$ .

### 1.1.2 Empty and Universal set

Def<sup>n</sup> :

An empty set is a set which contains no elements. It is either denoted by empty braces or the greek letter  $\phi$ .

Def<sup>n</sup> :

A Universal set is a set which contains all the elements (in the context).

Def<sup>n</sup> :

A set which contains only one element is called a singleton set.  
for example:  $\{5\}$ .

Note:  
for any set  $A$ ,  $\phi$  and  $A$  are always subsets called improper subsets.

### 1.1.3 Power Set

Def<sup>n</sup> :

A power set of a set is the collection of all the subsets of  $A$ . It is denoted by  $2^A$ .

## 1.2 Representation of Sets

There are 2 ways to represent sets:

1. Set builder form
2. Roaster form

### 1.2.1 Set builder form

Def<sup>n</sup> :

It is based on the unique property of the collection. The iterator is set and a property is defined in curly braces

*Example:*

$$A = \{x : x = 2y, y \in \mathbb{Z}\}$$

OR

$$A = \{2x : x \in \mathbb{Z}\}$$

### 1.2.2 Roaster form

Def<sup>n</sup> :

In this representation we list the elements in curly braces seperated by commas.

*Example:*

$$A = \{\dots, -4, -2, 0, 2, 4, \dots\}$$

## 1.3 Operations on sets

We have defined the following functions on sets

1. Union
2. Intersection
3. Difference
4. Symmetric Difference

### 1.3.1 Union of Sets( $\cup$ )

Def<sup>n</sup> :

Collection of all the elements of the sets

*Example:*

$$A \cup B = \{x : x \in A \text{ or } x \in B\}$$

### 1.3.2 Intersection of Sets( $\cap$ )

Def<sup>n</sup> :

Collection of all the elements in both the sets

*Example:*

$$A \cap B = \{x : x \in A \text{ and } x \in B\}$$

### 1.3.3 Difference of Sets(-)

Def<sup>n</sup> :  
Collection of all the elements one set but not the other

*Example:*

$$A - B = \{x : x \in A \text{ and } x \notin B\}$$

### 1.3.4 Symmetric difference of Sets( $\Delta$ )

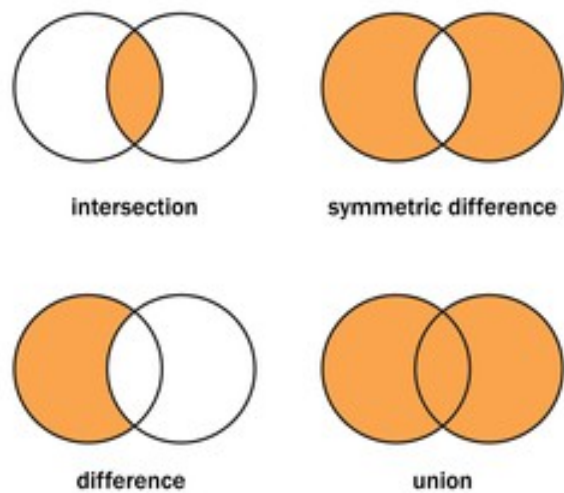
Def<sup>n</sup> :  
Collection of all the elements which exist in exactly one of the sets

*Example:*

$$A \Delta B = (A - B) \cup (B - A) = (A \cup B) - (A \cap B)$$

## 1.4 Venn diagram

A pictorial representation of sets is called a venn diagram



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## 1.5 De-morgan's Law

Let  $A$  and  $B$  be two sets then

1.  $(A \cup B)^c = A^c \cap B^c$
2.  $(A \cap B)^c = A^c \cup B^c$

*Proof:* Let  $x \in (A \cup B)^c$

$$\begin{aligned}\implies x &\notin A \cup B \\ \implies x &\notin A, x \notin B \\ \implies x &\in A^c, x \in B^c \\ \implies x &\in A^c \cap B^c\end{aligned}$$

Thus we can say that  $(A \cup B)^c \subseteq A^c \cap B^c$   
Similarly Let  $x \in A^c \cap B^c$ .

$$\begin{aligned}\implies x &\in A^c, x \in B^c \\ \implies x &\notin A, x \notin B \\ \implies x &\notin A \cup B \\ \implies x &\in (A \cup B)^c\end{aligned}$$

Thus we can say that  $A^c \cap B^c \subseteq (A \cup B)^c$   
This is possible iff  $(A \cup B)^c = A^c \cap B^c$

Q.E.D.

## 1.6 Partition of sets

Let  $S$  be a non-empty set. Then  $S$  has the partition if it has a collection of subsets  $A_i$  such that:

1.  $\forall a \in S, \exists$  unique  $i$  such that  $a \in A_i$
2.  $A_i \cup A_j = \phi, i \neq j$

*Example:*

Consider the set  $S = \{1, 2, \dots, 9\}$

1.  $\{\{1, 3, 5\}, \{2, 6\}, \{4, 9\}\}$
2.  $\{\{1, 3, 5\}, \{2, 4, 6, 8\}, \{7, 9\}\}$

then 1 is not a partition of  $S$  as the element 7 is missing. However 2 is a partition of the set  $S$ .

## 1.7 Relations

A relation  $R$  from set  $A$  to set  $B$  is subset of  $A \times B$  where :

$$A \times B = \{(a, b) : a \in A, b \in B\}$$

$$R \subseteq A \times B = \{(a, b) : a \in A, b \in B\}$$

- Domain  $\rightarrow$  All the elements of set  $A$
- Codomain  $\rightarrow$  All the elements of set  $B$
- Range  $\rightarrow$  All the second elements of  $R$

*Example:*

$$\begin{aligned} R &= \{(a, b) : b = 2a + 1, a \in [1, 10], b \in [1, 10]\} \\ A &= [1, 10] \\ R &\subseteq A \times A \\ &= \{(1, 3), (2, 5), (3, 7), (4, 9)\} \end{aligned}$$

- Domain =  $\{1, 2, 3, 4\}$
- Range =  $\{3, 5, 7, 9\}$

### 1.7.1 Composition of Relations

Let  $R$  be a relation from  $A$  to  $B$

Let  $S$  be a relation from  $B$  to  $C$

then

$$R \circ S = \{(a, c) : \exists b \in B \text{ s.t. } (a, b) \in R, (b, c) \in S\}$$

*Example:*

$$\begin{aligned} \text{Let } A &= \{1, 2, 3, 4\} \\ B &= \{a, b, c, d\} \\ C &= \{x, y, z\} \\ R &= \{(1, a), (2, a), (3, a), (4, d)\} \\ S &= \{(a, y), (b, x), (c, z)\} \\ \implies R \circ S &= \{(1, y), (2, y), (3, z)\} \end{aligned}$$

### 1.7.2 Equivalence Relations

A relation  $R$  from  $A$  to  $A$  is said to be equivalence if it satisfies the following conditions:

1. Reflexivity
2. Transitivity
3. Symmetricity

Def<sup>n</sup> :

A relation is said to be reflexive iff:

$$(a, a) \in R \quad \forall a \in A$$

Def<sup>n</sup> :

A relation is said to be Transitive iff:

$$(a, b) \in R, (b, c) \in R \implies (a, c) \in R$$

Def<sup>n</sup> :

A relation is said to be Symmetric iff:

$$(a, b) \in R \implies (b, a) \in R$$

*Example:*

Consider the relation  $R$  on  $A = \{1, 2, 3, 4\}$

$$R_1 = \{(1, 1), (1, 2), (2, 3), (1, 3), (4, 4)\}$$

$$R_2 = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (4, 4)\}$$

$$R_3 = \{(1, 3), (2, 1)\}$$

$$R_4 = A \times A$$

Then

- $R_1$  is Transitive only
- $R_2$  is Equivalence
- $R_3$  is neither Symmetric, Reflexive or Transitive only
- $R_4$  is Equivalence

Def<sup>n</sup> :

A relation on a set  $A$  is said to be Anti-symmetric if and only if:

$$(a, b) \in A, (b, a) \in A \implies a = b$$

## 1.8 Functions

A relation from set  $A$  to  $B$  such that each element of  $A$  has a unique mapping in  $B$  then it is called a function and is denoted as

$$f : A \rightarrow B$$

All functions are relations but not vice-versa.

### 1.8.1 Domin, Range and Codomain

for  $f : A \rightarrow B$

- Domain:  $A$
- Codomain:  $B$
- Range:  $\{b \in B : \exists a \in A \text{ such that } f(a) = b\}$

Note: Domain( $f$ )  $\subseteq A$

Range( $f$ )  $\subseteq B$



One-one    Each horizontal line cuts the graph at atmost one point  
 Onto        Each horizontal line cuts the graph at one or more points

### 1.8.2 One-one and Onto function

Let  $f : A \rightarrow B$  be a function then  $f$  is called one-one if distinct elements of  $A$  have different image. Mathematically:

$$f(x_1) = f(x_2) \implies x_1 = x_2$$

$f$  is said to be onto if each element of  $B$  has a preimage in  $A$ . Mathematically

$$\forall b \in B, \exists a \in A \text{ such that } f(a) = b$$

Def<sup>n</sup> :  
 $f$  is bijective iff it is both one-one and onto

### 1.8.3 Geometrical characterisation of one-one and onto functions

*Example:*

Let  $f : \mathbb{R} \rightarrow \mathbb{R}$  such that  $f(x) = x^2$   
 One-one:

$$\begin{aligned} \text{Let } f(x_1) &= f(x_2) \\ \implies x_1^2 &= x_2^2 \\ \implies x_1^2 - x_2^2 &= 0 \\ \implies (x_1 - x_2)(x_1 + x_2) &= 0 \\ \implies x_1 = x_2 \text{ or } x_1 &= -x_2 \end{aligned}$$

Onto: We have  $f(x) = x^2 \geq 0$   
 $\nexists x < 0 \in \mathbb{R}$  such that  $f(x) = x^2$

## 1.9 Composition of Functions

Let  $f : A \rightarrow B$  and  $g : B \rightarrow C$  be two functions then  $g \circ f : A \rightarrow C$  is called composition of  $f$  and  $g$ .

$$g \circ f(x) = g(f(x))$$

Result:

Let  $f : A \rightarrow B$  and  $g : B \rightarrow C$  be two functions then

1.  $g \circ f$  is one-one if both  $f$  and  $g$  are one-one
2.  $g \circ f$  is onto if both  $f$  and  $g$  are onto

*Proof:*

1. Let  $(g \circ f)(x_1) = (g \circ f)(x_2)$

$$\implies g(f(x_1)) = g(f(x_2))$$

$$\implies f(x_1) = f(x_2) \text{ as } g \text{ is one-one}$$

$$\implies x_1 = x_2 \text{ as } f \text{ is one-one}$$

2.  $g \circ f : A \rightarrow C$

Let  $z \in C$ , we will show that  $\exists x \in A$  such that  $g \circ f(x) = z$

Since  $g$  is onto  $\implies \exists y \in B$  such that  $g(y) = z$

Since  $f$  is onto and  $y \in B \implies \exists x \in A$  such that  $f(x) = y$

Q.E.D.

## 1.10 Induction

Used to prove that a statement is true for all integers or natural numbers.

$$P(n) \text{ holds } \forall n$$

### 1.10.1 Principle of Mathematical Induction

Let  $P(n)$  be the given statement.

1. Basic Step:  $P(1)$  is true

2. Induction Step:  $P(k) \implies P(k+1)$  is true

This causes a domino effect and effectively proves that  $P(n)$  is true  $\forall n$  *Example*:

Using Induction, Prove that:

$$1 + 2 + 3 + 4 + 5 + \dots + n = \frac{n(n+1)}{2}, n \in \mathbb{Z}^+$$

1. Base Case:

$$1 = \frac{1(1+1)}{2} = 1$$

Q.E.D.

2. Induction Step: Let  $P(k)$  is true, then we have to show that  $P(k+1)$  is true.

$$\begin{aligned} 1 + 2 + \dots + k &= \frac{k(k+1)}{2} \\ \implies 1 + 2 + \dots + k + (k+1) &= \frac{k(k+1)}{2} + (k+1) \\ \implies 1 + 2 + \dots + k + (k+1) &= \frac{k(k+1)}{2} + \frac{2(k+1)}{2} \\ \implies 1 + 2 + \dots + k + (k+1) &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

Q.E.D.

## 1.11 Recursion

This is used when we cannot explicitly represent any object or mathematical term. so we use recursion.

Def<sup>n</sup> :

Recursive function

Let  $\{a\}_n, n \in \mathbb{Z}_0^+$  such that  $a_n \in \mathbb{Z} \forall n$  then:

1. Basic Step:  $a_n$  is given at  $n = 0$
2. Recursive step:  $a_n = f(a_{n-1}, a_{n-2} \dots)$

Well known examples of recursive expressions are:

1. Arithmetic progression

$$a_n = a_{n-1} + d$$

2. Geometric progression

$$a_n = r a_{n-1}$$

3. Factorial function

$$f(n) = n \times f(n-1)$$

$$f(0) = 1$$

4. Fibonacci function

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 0$$

$$f(1) = 1$$

### 1.11.1 Linear recurrence relations with constant coefficients

Let  $a_n = \phi(a_{n-1}, a_{n-2}, \dots, a_0, m)$

It can be written as:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_{n-1} a_1 + c_n a_0 + f(m)$$

If  $f(m) = 0$  then the function is called homogeneous

If  $f(m) \neq 0$  then the function is called non-homogeneous

### HOMOGENEOUS SOLUTION

Note:  $k$ th order recurrence relation is one such that

$$a_n = \phi(a_{n-1}, a_{n-2}, \dots, a_{n-k}, m)$$

or

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + f(m)$$

so a second order recurrence relation will look like:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + c_3 \quad (1)$$

Then characteristic equation of (1) is given by:

$$x^2 - c_1 x - c_2 = 0, c_1, c_2 \in \mathbb{R}$$

Come forth 3 cases:

1. When roots are real and distinct:

$$x^2 - c_1 x - c_2 = 0 \implies \begin{cases} x = r_1 \\ x = r_2 \end{cases} \quad r_1 \neq r_2$$

Then the solution of (1) is given by:

$$a_n = p_1 r_1^n + p_2 r_2^n$$

2. Roots are real and equal:

$$r_1 = r_2 = r$$

Then the solution of (1) is given by:

$$a_n = (p_1 + n p_2) r^n$$

*Example:*

Solve  $a_n = 2a_{n-1} + 3a_{n-2}$

$$\begin{aligned} a_n - 2a_{n-1} - 3a_{n-2} &= 0 \\ \implies x^2 - 2x - 3 &= 0 \\ \implies (x - 3)(x + 1) &= 0 \\ \implies x &= -1, 3 \end{aligned}$$

Solution:

$$a_n = p_1(-1)^n + p_2(3)^n$$

Suppose:  $a_0 = 1, a_1 = 2$

$$\begin{aligned} n = 0 &\implies a_0 = p_1(-1)^0 + p_2(3)^0 \\ &\implies 1 = p_1 + p_2 \\ n = 1 &\implies a_1 = p_1(-1)^1 + p_2(3)^1 \\ &\implies 2 = -p_1 + 3p_2 \end{aligned}$$

From the above 2 equations we get:

$$\begin{aligned} p_1 &= \frac{1}{4} \\ p_2 &= \frac{3}{4} \end{aligned}$$

$$\implies a_n = \frac{1}{4}(-1)^n + \frac{3}{4}(3)^n$$

*Example:*

Solve:  $a_n = 6a_{n-1} - 9a_{n-2}$

with  $a_1 = 3, a_2 = 27$

## NON-HOMOGENEOUS SOLUTION

We have

$$a_n = c_1a_{n-1} + c_2a_{n-2} + \cdots + c_ka_{n-k} + f(n) \quad (2)$$

where  $f(n) \neq 0$

From (2) the associated homogeneous linear recurrence relation is

$$a_n = c_1a_{n-1} + \cdots + c_ka_{n-k} \quad (3)$$

We can get the solution of (3), Let say  $(a_n^c)$

Let  $a_n^p$  is the particular solution of (2). Then the general solution of (2) can be written as

$$a_n = a_n^c + a_n^p$$

### **Theorem:**

Suppose that  $a_n$  satisfies the relation

$$a_n = c_1a_{n-1} + c_2a_{n-2} + \cdots + c_ka_{n-k} + f(n)$$

where  $c_i \in \mathbb{R} \forall i$  and

$$f(n) = (b_t n^t + b_{t-1} n^{t-1} + \cdots + b_1 n + b_0) s^n$$

where  $b_i, s \in \mathbb{R} \forall i$

Then the following cases.

1. If  $s$  is not the root of the characteristic equation then the particular solution of (2) is of the type

$$(p_t n^t + p_{t-1} n^{t-1} + \cdots + p_1 n + p_0) s^n$$

2. if  $s$  is root of the characteristic equation then the particular solution of (2) is of the type

$$n^m (p_t n^t + p_{t-1} n^{t-1} + \cdots + p_1 n + p_0) s^n$$

where  $m$  is the multiplicity of the root

*Example:*

Write the type of particular solution for the recurrence relation,  $a_n = 6a_{n-1} - 9a_{n-2} + f(n)$

1.  $f(n) = 3^n$
2.  $f(n) = n3^n$
3.  $f(n) = n^2 2^n$

*Solution:*

The associated homogeneous linear recurrence relation:

$$a_n - 6a_{n-1} + 9a_{n-2} = 0$$

So characteristic equation :

$$x^2 - 6x + 9 = 0$$

$$\implies (x - 3)^2 = 0$$

$$\implies x = 3, 3$$

1.  $f(n) = 3^n$

comparing to  $f(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n$  we get  
so

$$b_0 = 1, b_i = 0 \forall i > 0$$

$s = 3$  is root of the characteristic equation with multiplicity 2

So particular solution would be:

$$n^2 p_0 3^n$$

2.  $f(n) = n 3^n$

comparing to  $f(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n$  we get  
so

$$b_0 = 1, b_1 = 1, b_i = 0 \forall i > 1$$

$s = 3$  is root of the characteristic equation with multiplicity 2

So particular solution would be:

$$n^2 (p_1 n + p_0) 3^n$$

3.  $f(n) = n^2 2^n$

comparing to  $f(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n$  we get  
so

$$b_0 = 1, b_1 = 1, b_2 = 1, b_i = 0 \forall i > 2$$

$s = 2$  is root of the characteristic equation with multiplicity 2

So particular solution would be:

$$a_n^p = (p_2 n^2 + p_1 n + p_0) 2^n$$

$$a_n^c = (c_1 + c_2 n) 3^n$$

From the above 2 equations ;

$$(p_2 n^2 c p_1 n + p_0) 2^n = 6(p_2 (n-1)^2 + p_2 (n-1) + p_0) 2^{n-1} - 9(p_2 (n-2)^2 + p_1 (n-2) + p_0) 2^{n-2} + n^2 2^n$$

*Example:*

Solve:

$$a_n = 6a_{n-1} - 9a_{n-2} + n^2 2^n, a_0 = 1, a_1 = 2$$

*Solution:*

The associated homogeneous linear recurrence relation is  $a_n = 6a_{n-1} - 9a_{n-2}$  So the characteristic equation is

$$\begin{aligned}x^2 - 6x + 9 &= 0 \\ \implies x &= 3, 3\end{aligned}$$

So

$$a_n^c = (c_1 + c_2n)3^n$$

for particular solution:

$$a_n^p = (b_2n^2 + b_1n + b_0)2^n$$

so the general solution :

$$a_n = (c_1 + c_2n)3^n + (b_2n^2 + b_1n + b_0)2^n$$

putting the values of  $a_0$  and  $a_1$  we get

$$\begin{aligned}1 &= c_13^0 + b_02^0 \\ \implies 1 &= c_1 + b_0\end{aligned}$$

$$\begin{aligned}2 &= (c_1 + c_2)3 + (b_2 + b_1 + b_0)2 \\ \implies 2 &= 3c_1 + 3c_2 + 2b_2 + 2b_1 + 2b_0\end{aligned}$$

Comparing the coefficients of  $2^n$  and  $3^n$  in both sides:

$$\begin{aligned}b_0 + b_1 + b_2 &= 1 \\ c_1 + c_2 &= 0 \\ c_1 &= 1 \\ b_0 &= 0 \\ c_2 &= -c_1 = -1\end{aligned}$$

## 2 Logic and Proof Techniques

Def<sup>n</sup> :

Proposition(statement): is a declarative sentence that is either true or false but not both.

*Example:*

1. Ice floats on water.
2.  $2 - x = 3$

### 2.1 Proof Techniques

Def<sup>n</sup>:

**Conjecture:** Any statement given is a conjecture.

**Theorem:** A statement which is to be proved true.

$$\text{Conjecture} + \text{proof} \rightarrow \text{Theorem}$$

**Proof:** The technique by which we check whether the given statement is true or not

There are the following proof techniques:

1. Direct Proof
2. Proof by Contrapositive
3. Proof by Contradiction
4. Proof by Counterexample
5. Proof by Cases

### 2.1.1 Direct Proof

If we are given  $p \rightarrow q$ , then take  $p$  and by implications of  $p$ , we arrive at  $q$ .

*Example:*

If  $n$  is even then prove that  $n^2$  is even.

Given:  $n$  is even.

To prove:  $n^2$  is even

Proof: let us say that

$$\begin{aligned} n &= 2k, k \in \mathbb{Z} \\ \implies n^2 &= 4k^2 \\ \implies n^2 &= 2(2k^2) \end{aligned}$$

*Example:*

For all integers  $a, b$  and  $c$  if  $a|b$ , and  $b|c$  prove that  $a|c$

### 2.1.2 Proof by Contrapositive

Suppose we have to prove that  $p \implies q$ , then it is equivalent to proving  $\neg q \implies \neg p$

*Example:*

For all integers  $a$  and  $b$  if  $a + b$  is odd then either  $a$  is odd or  $b$  is odd

$p$  :  $a + b$  is odd

$q$  :  $a$  is odd or  $b$  is odd

$\neg p$  :  $a + b$  is even

$\neg q$  :  $a$  is even and  $b$  is even



since  $p \implies q \equiv \neg q \implies \neg p$   
 We have  $a$  is even and  $b$  is even

$$a = 2m, b = 2n, m, n \in \mathbb{Z}$$

$$\begin{aligned} \implies a + b &= 2m + 2n \\ &= 2(m + n) \\ &= 2k, k \in \mathbb{Z} \end{aligned}$$

*Example:*

For every prime number  $r$ , if  $r \neq 2$  then  $r$  is odd

### 2.1.3 Proof by Contradiction

We have to prove  $p \implies q$

Then it is equivalent to  $\neg p \implies \neg q$

for that we take  $p$  is not true and then by implications, we arrive at some Contradiction.

*Example:*

$\sqrt{2}$  is irrational.

suppose  $\sqrt{2}$  is a rational number.

Then we can represent it in form of  $\frac{p}{q}$  such that  $p$  and  $q$  are coprimes

$$\begin{aligned} \implies \sqrt{2} &= \frac{p}{q} \\ \implies 2 &= \frac{p^2}{q^2} \\ \implies 2q^2 &= p^2 \end{aligned}$$

hence  $p$  is an even number. therefore we can write  $p$  as

$$\begin{aligned} p &= 2k \\ p^2 &= 4k^2 \\ 2q^2 &= 4k^2 \\ q^2 &= 2k^2 \end{aligned}$$

thus  $q$  is also an even number.

this is a contradiction to the fact that  $p$  and  $q$  are coprime.

### 3 Lattice

#### 3.1 Complement of an element

Let  $L$  be a Lattice then an element  $b \in L$  is said to be complement of  $a \in L$  if

$$a * b = 0$$

$$a \oplus b = 1$$

An element may have no complement, unique complement or more than one complements

**Theorem:**

The following conditions always hold true:

$$0' = 1$$

$$1' = 0$$

and these complements are unique

Def<sup>n</sup> :

A lattice  $L$  is complemented if each element has atleast one complement.

Def<sup>n</sup> :

A Lattice is said to be distributive if

$$a * (b \oplus c) = (a * b) \oplus (a * c)$$

$$a \oplus (b * c) = (a \oplus b) * (a \oplus c)$$

This is true if each element has atleast one complement.

#### 3.2 Boolean Algebra

Def<sup>n</sup> :

A Lattice is Boolean if it is both complementary and distributive i.e. each element of lattice has exactly one complement. Alternatively, A Lattice is Boolean if it is isomorphic to a  $D_n$ (divisors of  $n$ ).

*Example:*

**Theorem:**

Here are some results to keep in mind.

1. If a lattice does not contain  $2^n$  elements then it cannot be a boolean lattice.
2. if  $|L| = 2^n$  Then it may or may not be boolean.
3. if  $n = p_1 \times p_2 \times \dots \times p_k$  where each  $p_i$  is a distinct prime then  $L$  will be a boolean algebra.

4. Let  $n \in \mathbb{Z}$  and  $p^k | n$  for some  $k > 1$  then  $D_n$  is not a boolean algebra

### 3.3 Logic Gates and Circuits

Basic type of logic gates:

1. OR Gate
2. AND Gate
3. NOT Gate

#### 3.3.1 OR Gate

$$Y = A * B \equiv A + B$$

$Y$  will be 0 if and only if  $A$  and  $B$  both are zero.

#### 3.3.2 AND Gate

$$Y = A \oplus B \equiv A \cdot B$$

$Y$  will be 1 if and only if  $A$  and  $B$  both are 1.

#### 3.3.3 NOT Gate

$$Y = A' = \overline{A}$$

$Y$  will just be the opposite of  $A$ .

#### 3.3.4 Logic Circuit

Def<sup>n</sup>:

Any combination of logic gates to get the desired output is called logic circuit.

Note that logic circuit follows the rules of boolean algebra.

#### 3.3.5 NAND and NOR Gate

NAND Gate is the combination of AND and NOT Gate.

NOR Gate is the combination of OR and NOT Gate.

## 4 GROUP THEORY yeah imma kms

### 4.1 Binary Operations

Def<sup>n</sup>:

Let  $A$  be a set then any operation  $*$  on  $A$  is binary operation if  $\forall a, b \in A, a * b \in A$ . For example,  $(\mathbb{Z}, +)$  is a group, but  $(\mathbb{Z}, /)$  is not.

### 4.2 Group

Def<sup>n</sup> :

Let  $G$  be a nonempty set with a defined binary operation  $*$  then  $G$  is called group if it satisfies a certain condition.

1. Closure:  $a * b \in G \forall a, b \in G$ .
2. Associativity:  $(a * b) * c = a * (b * c)$ .
3. Existence of Identity:  $\exists e \in G, a * e = a \forall a \in G$ .
4. Existence of Inverse:  $\forall a \in G \exists b \in G$  such that  $a * b = e$

*Example:*

$(\mathbb{Z}, +)(\mathbb{Q}, +)(\mathbb{R}, +)$  are groups

$(\mathbb{Z}, -)$  is not a group because  $-$  is not associative

$(\mathbb{N}, +)$  is not a group because there does not exist an identity in  $\mathbb{N}$

$(\mathbb{Z}, \cdot)$  is not a group because there is no inverse for  $a \in \mathbb{N}$

$(\mathbb{R}, \cdot)$  is not a group because there is no inverse for  $0 \in \mathbb{R}$

$(\mathbb{R} \setminus \{0\}, \cdot)$  is a group.

*Example:*

let  $X = \mathbb{M}(\mathbb{R})_{2 \times 2}$ , then  $G = (X, +)$  is a group.

However  $G' = (X, \cdot)$  is not a group as there may or may not exist an inverse for a real matrix.

### 4.3 Abelian Group

Let  $G$  be a group with binary operation  $*$  then  $G$  is abelian if

$$a * b = b * a \forall a, b \in G$$

*Example:*

$(\mathbb{Z}, +)$  is a group where  $a + b = b + a \forall a, b \in \mathbb{Z}$

Thus  $(\mathbb{Z}, +)$  is an abelian group.

*Example:*

$GL(2, \mathbb{R})$  = group of non singular matrices of order 2 with real entities

note that it is not necessary that the product of two matrices be commutative.

*Example:*

$SL(2, \mathbb{R})$  = group of matrices of order 2 with real entities and determinant 1

note that it is not necessary that the product of two matrices be commutative.

### 4.4 Unitary group

Def<sup>n</sup> :

$$U(n) = \{a : a \in \mathbb{Z}, \gcd(a, n) = 1, a < n\}$$

*Example:*

$$U(10) = \{1, 3, 7, 9\}$$

Note that  $(U(10), \odot_n)$   
 $\odot_n$  is multiplication modulo

Note:

$$a \oplus_n b = a + b \pmod{n}$$

$$a \odot_n b = a \cdot b \pmod{n}$$

#### 4.4.1 Cayley Table

It is a visual table representation of a group. we just multiply the group with itself to check for inverses

*Example:*

$$\mathbb{Z}_n = \{0, 1, \dots, n-1\}$$

then  $(\mathbb{Z}_n, \oplus_n)$  is a group with identity 0.

and  $\langle \mathbb{Z}_n, \oplus_n \rangle$  is an abelian group.

#### 4.5 Properties of a group

1. Uniqueness of identity

The identity of the group must be unique. To prove take two distinct identities and show that they are equal.

2. Cancellation Laws:-

Let  $G$  be a group and  $a, b, c \in G$ . Then

(a)  $ab = ac \implies b = c$  (left cancellation law)

(b)  $ab = cb \implies a = c$  (right cancellation law)

3. Uniqueness of inverse

4. Let  $G$  be a group then  $(ab)^{-1} = b^{-1}a^{-1}$

5. Let  $G$  be an abelian group then  $(ab)^{-1} = b^{-1}a^{-1} = a^{-1}b^{-1}$

*Proof:*

- 1.

2. (a) To show:  $ab = ac \implies b = c$

Proof:

$$ab = ac$$

since  $G$  is a group  $\exists a^{-1} \in G$ . So premultiply by  $a^{-1}$

$$a^{-1}ab = a^{-1}ac$$

$$eb = ec$$

$$b = c$$

3. Let  $G$  be a group and  $a \in G$ . Let  $b, c$  be the inverses of  $a$ . since  $b$  is inverse of  $a$ :

$$ab = ba = e$$

similarly

$$ac = ca = e$$

From the two above statements we can infer that:

$$ab = ac = e$$

From the left cancellation law

$$b = c$$

Q.E.D.

## 4.6 Subgroups

Def<sup>n</sup>:

The order of a group is the number of elements in a group.

*Example:*

$$G = U(10) = \{1, 3, 7, 9\}$$

$$O(G) = 4$$

Def<sup>n</sup>:

**The order of an element in a group**

Let  $G$  be a group and  $a \in G$  then order of  $a$  is defined as:

$$O(a) = \{n : a^n = e, a \in G\}(\cdot)$$

$$O(a) = \{n : na = e, a \in G\}(+)$$

then  $n$  is called order of  $a$

Def<sup>n</sup>:

Let  $H$  be subset of  $G$  then  $H$  is called subgroup of  $G$  if  $H$  itself is group under the same operation as that of  $G$ .

Notation:  $H \leq G$

*Example:*

$$(\mathbb{Z}, +) \leq (\mathbb{R}, +)$$

*Example:*

$$(\mathbb{Z}_{10}, \oplus_{10}) \not\leq (\mathbb{R}, +) \text{ as the operation is not the same}$$

### 4.6.1 One step subgroup test

Let  $H$  be subset of the group  $G$ . Then  $H$  is subgroup of  $G$  if

$$ab^{-1} \in H \forall a, b \in H (\text{wrt multiplication})$$

$$a - b \in H \forall a, b \in H (\text{wrt addition})$$

*Example:*

Let  $G$  be a Abelian group with identity  $e$  and  $H = \{x \in G : x^2 = e\}$  then check for  $H \leq G$  by one step test:

$$ab^{-1} \in H \forall a, b \in H (\text{wrt multiplication})$$

$$\text{since } a, b \in H, a^2 = e, b^2 = e$$

$$\text{so } (ab^{-1})^2 = ab^{-1}ab^{-1} = aab^{-1}b^{-1} = a^2b^{-2} = ee^{-1} = e$$

Therefore  $H \leq G$  ■

### 4.6.2 Two step subgroup test

Let  $G$  be a group and  $H \subseteq G$  then iff:

1.  $ab \in H \forall a, b \in H$

2.  $a \in H \forall a \in H$

*Example:*

Let  $G = GL(2, \mathbb{Z})$  with addition.

$$\text{Let } H = \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} : a + b + c + d = 0 \right\}$$

*Example:*

$$G = GL(2, \mathbb{R})$$

$$H = \left\{ \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix} : a, b \in \mathbb{Z} \setminus \{0\} \right\}$$

Then  $H$  will not be a subgroup because there may or may not exist an inverse of  $a \in H$  in  $H$

## 4.7 Cyclic groups

a group  $G$  is called cyclic if  $G = \langle a \rangle$  where

$$\langle a \rangle = \{a^n : n \in \mathbb{Z}\} \text{ With respect to multiplication}$$

$$\langle a \rangle = \{na : n \in \mathbb{Z}\} \text{ With respect to addition}$$

*Example:*

$U(10) = \{1, 3, 7, 9\}$  under multiplication  $\odot_{10}$

$$3^1 = 3$$

$$3^2 = 9$$

$$3^3 = 7$$

$$3^4 = 1$$

$$3^5 = 3$$

So  $U(10)$  is cyclic

Note: in  $\mathbb{Z}_n$ , the generators are the numbers which are coprime to  $n$ . so  $\mathbb{Z}_{10}$  has 1,3,7,9 as it's generators under addition modulo  $n$

#### 4.7.1 Euler's phi (totient) function

$$\phi(1) = 1$$

for  $n > 1$

$$\phi(n) = \begin{cases} n-1, & n \text{ is prime} \\ p^m - p^{m-1}, & n = p^m \\ (p-1)(q-1), & n = pq \end{cases}$$

Note that  $\mathbb{Z}$  has only 1 and  $-1$  as it's generators. It is also the only infinite cyclic group

#### 4.8 Subgroups of $\mathbb{Z}_n$

For any divisor  $k$  of  $n$ , the set  $\langle \frac{n}{k} \rangle$  is a unique subgroup of order  $k$  and thees are the only subgroup of  $\mathbb{Z}_n$ .

*Example:*

Let us take  $\mathbb{Z}_{30}$

$$\begin{aligned} \langle \frac{30}{1} \rangle &= \langle 30 \rangle = \{0\} \\ \langle \frac{30}{2} \rangle &= \langle 15 \rangle = \{0, 15\} \\ \langle \frac{30}{3} \rangle &= \langle 10 \rangle = \{0, 10, 20\} \\ \langle \frac{30}{5} \rangle &= \langle 6 \rangle = \{0, 6, 12, 18, 24\} \\ \langle \frac{30}{6} \rangle &= \langle 5 \rangle = \{0, 5, 10, 15, 20, 25\} \\ \langle \frac{30}{10} \rangle &= \langle 3 \rangle = \{0, 3, 6, 9, 12, 15, 18, 21, 24, 27\} \\ \langle \frac{30}{15} \rangle &= \langle 2 \rangle = \{0, 2, 4, 6, 8, 10, 12, \dots, 28\} \\ \langle \frac{30}{30} \rangle &= \langle 1 \rangle = \{0, 1, 2, 3, 4, 5, 6, \dots, 29\} \end{aligned}$$

#### 4.9 Center of a group

$$Z(G) = \{x \in G; xa = ax \forall a \in G\}$$

in leyman terms, collection of all elements which commute with each other

Note: if  $G$  is abelian then  $Z(G) = G$



*Example:*

$$\text{if } G = GL(2, \mathbb{R}) = \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} : a, b, c, d \in \mathbb{R} \right\}$$

$$Z(G) = \left\{ \begin{pmatrix} e & 0 \\ 0 & e \end{pmatrix} : e \in \mathbb{R} \right\}$$

Note:  $Z(G) \leq G$

#### 4.10 Centralizer of an element

for an element  $a \in G$

$$C(a) = \{x \in G : xa = ax\}$$

Note: if  $G$  is abelian then  $C(a) = G$

*Example:*

$$G = GL(2, \mathbb{R})$$

$$\text{find } C \left( \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \right)$$

$$\text{Suppose } \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in G$$

$$\begin{pmatrix} a & b \\ c & b \end{pmatrix} \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} a & b \\ c & b \end{pmatrix}$$

$$\begin{pmatrix} a+b & a \\ c+d & c \end{pmatrix} = \begin{pmatrix} a+c & b+d \\ a & b \end{pmatrix}$$

#### 4.11 Partition of numbers

Let  $n$  be a positive integer then the ways in which the number can be written form partition of  $n$ .

*Example:*

$$n = 1, 1$$

$$n = 2, 2$$

$$, 1 + 1$$

$$n = 3, 3$$

$$, 2 + 1$$

$$, 1 + 1 + 1$$

$$n = 4, 4$$

$$, 3 + 1$$

$$, 2 + 2$$

$$, 2 + 1 + 1$$

$$, 1 + 1 + 1 + 1$$

Formula:

$$P(n+1) = \sum_{r=0}^n \binom{n}{r} P(r)$$

given  $P(0) = 1$

So

$$\begin{aligned}P(4) &= P(3+1) \\&= \sum_{r=0}^3 \binom{3}{r} P(r) \\&= \binom{3}{0} P(0) + \binom{3}{1} P(1) + \binom{3}{2} P(2) + \binom{3}{3} P(3) \\&= 1 + 3 \times 1 + 3 \times 2 + 1 \times 3 \\&= 13\end{aligned}$$

#### 4.12 Order of elements in $S_n$

Take  $n$  and make the partitions of  $n$ .

*Example:*

$$\begin{array}{ll}S_4 & 4 \rightarrow LCM = 4 \\& 1 + 3 \rightarrow LCM = 3 \\& 2 + 2 \rightarrow LCM = 2 \\& 1 + 1 + 2 \rightarrow LCM = 1 \\& 1 + 1 + 1 + 1 \rightarrow LCM = 1\end{array}$$

It means that  $S_4$  has elements of order 1, 2, 3, 4.

To find the number of elements of order  $m$  in  $S_n$ .

Let  $n = p_1 + p_2 + \dots$

$$\frac{n!}{p_1^{m_1} \cdot m_1! \times p_2^{m_2} \cdot m_2!}$$

## 5 Cosets of $H$ in $G$

Let  $G$  be a group and  $H \leq G, a \in G$

then

$$aH = \{ah : h \in H\}$$

and

$$Ha = \{ha : h \in H\}$$

are called left and right cosets of  $H$  in  $G$

*Example:*

$G = S_3, H = \{I, (13)\}$ . Find all the left cosets of  $H$  in  $G$

Left Coset:

$$aH = \{ah : h \in H\}$$

$$G = S_3 = \{I, (12), (23), (13), (123), (132)\}$$

$$IH = \{I, (13)\} = H$$

$$(12)H = \{(12), (12)(13)\} = \{(12), (132)\}$$

$$IH = \{I, (13)\} = H$$

$$(23)H = \{(23), (23)(13)\} = \{(23), (123)\}$$

$$IH = \{I, (13)\} = H$$

$$(13)H = \{(13), (13)(13)\} = \{(13), I\}$$

$$IH = \{I, (13)\} = H$$

$$(123)H = \{(123), (123)(13)\} = \{(123), (23)\}$$

$$IH = \{I, (13)\} = H$$

$$(132)H = \{(132), (132)(13)\} = \{(132), (12)\}$$

So there are 3 distinct left cosets of  $H$  in  $G$

*Example:*

Let  $G = \mathbb{Z}_9$  under addition

and  $H = \{0, 3, 6\}$

Then find all the right cosets of  $H$  in  $G$

$$H + a = \{h + a, h \in H\}$$

$$H + \{0\} = \{0, 3, 6\}$$

$$H + \{1\} = \{1, 4, 7\}$$

$$H + \{2\} = \{2, 5, 8\}$$

$$H + \{3\} = \{3, 6, 9\} = \{0, 3, 6\}$$

$$H + \{4\} = \{4, 7, 10\} = \{1, 4, 7\}$$

$$H + \{5\} = \{5, 8, 11\} = \{2, 5, 8\}$$

$$H + \{6\} = \{6, 9, 12\} = \{0, 3, 6\}$$

$$H + \{7\} = \{7, 10, 13\} = \{1, 4, 7\}$$

$$H + \{8\} = \{8, 11, 14\} = \{2, 5, 8\}$$

Thus we get 3 distinct cosets

*Example:*

Find all the left cosets of  $\{1, 11\}$  in  $U(30)$

$$U(30) = \{1, 7, 11, 13, 17, 19, 23, 29\}$$

$$\begin{aligned}
\{1\}H &= \{1, 11\} = H \\
\{7\}H &= \{7, 77\} = \{7, 17\} \\
\{11\}H &= \{11, 121\} = \{1, 11\} \\
\{13\}H &= \{13, 143\} = \{13, 23\} \\
\{17\}H &= \{17, 187\} = \{7, 17\} \\
\{19\}H &= \{19, 209\} = \{19, 29\} \\
\{23\}H &= \{23, 253\} = \{13, 23\} \\
\{29\}H &= \{29, 309\} = \{19, 29\}
\end{aligned}$$

Thus we get 4 distinct cosets

## 5.1 Normal Subgroup

Let  $G$  be a group and  $H \leq G$ . Then  $H$  is called normal in  $G$  if  $aH = Ha \forall a \in G$ .  
 It means if  $a \in G$  and  $h \in H$  then there exists  $h' \in H$  such that  $ah' = ha$

Notation:  $H \triangleleft G$

*Example:*

$G = S_3$  and  $H = \{I, (12)\}$

Is  $H$  normal to  $G$ ?

Let  $a = (13) \in S_3$   
 and  $h = (12) \in H$

Note: We have  $aH = Ha \forall a \in G$

$$\implies a^{-1}Ha \subseteq H$$

By counterexample we can show that  $H$  is not normal in  $G = S_3$ .

$$H \not\triangleleft G$$

*Example:*

$$G = GL(2, \mathbb{R})$$

and

$$H = \left\{ \begin{bmatrix} a & b \\ 0 & d \end{bmatrix} \mid a, b, d \in \mathbb{R}, ad \neq 0 \right\}$$

$$\text{Let } \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in GL(2, \mathbb{R})$$

$$\text{Then } \begin{pmatrix} a & b \\ c & d \end{pmatrix}^{-1} = \frac{1}{ad - bc} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$$

$$\begin{pmatrix} a_1 & b_1 \\ 0 & c_1 \end{pmatrix} \in H$$

$$\text{so } \frac{1}{ad-bc} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix} \begin{pmatrix} a_1 & b_1 \\ 0 & c_1 \end{pmatrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix} =$$

In the end,  $H$  is not normal to  $G$

## 6 Graph Theory

A graph is a pair of set  $(V, E)$  where  $V$  is the set of vertices and  $E$  is the set of edges. Here vertices are represented by dots and edges are represented by lines/curves joining the vertices.

### 6.1 Types of graph

1. Directed graph
2. Undirected graph

#### 6.1.1 Directed graph

The graph in which the edge set is an ordered pair of vertices.  
Here  $(v_i, v_j)$  represents the edge from  $v_i$  to  $v_j$

#### 6.1.2 Undirected graph

The graph in which the edge set is an unordered pair of vertices.  
Here  $\{v_i, v_j\}$  represents the edge from  $v_i$  to  $v_j$  and vice-versa.

### 6.2 Null graph

A graph having zero edges.

### 6.3 Self loop

An edge joining a vertex to itself is called self loop.

### 6.4 Degree of a vertex in an undirected graph

Let  $G$  be a graph and  $v$  be any vertex. Then degree of  $v$  is number of edges incident with it and denoted by  $\deg_G(v)$

Note:

1. if  $\deg_G(v) = 0 \implies$  isolated vertex
2. if  $\deg_G(v) = 1 \implies$  pendent vertex
3. any self loop is counted twice.

## 6.5 Degree of a vertex in a directed graph

1. Indegree ( $\deg_G^+(v)$ ):  
The number of edges incident to it (towards vertex)
2. Outdegree ( $\deg_G^-(v)$ ):  
The number of edges incident from it (away from vertex)

Note:

Self loop is counted each time in both cases.

## 6.6 Path

In an undirected graph a path is a sequence of zero or more edges of type:

$$\{v_0, v_1\}, \{v_1, v_2\}, \dots, \{v_{n-1}, v_n\}$$

OR

$$v_0, v_1, \dots, v_n$$

where  $v_0 \rightarrow$  initial vertex

$v_n \rightarrow$  final vertex

Note:

1. In a path edges and vertices may be seen multiple times.
2. The number of edges in a path is called length of path.
3. If the length of path is zero then it is called trivial path.

### 6.6.1 Types of path

1. Open path
2. Closed path

If the initial and final vertices are same then it is closed path otherwise open.

1. Simple Path

A path is simple if all the vertices are different except initial or final vertices.

## 6.7 Representation of graph

A matrix is a convenient way to describe a graph.

We can use the properties of matrix algebra to find out many useful characteristics of a graph.

There are 3 types of matrix:

1. adjacency matrix
2. path matrix
3. incidence matrix

## 6.8 Matrix representation by Adjacency matrix

In case of undirected graph the adjacency matrix  $A$  will be defined as:

$$a_{ij} = \begin{cases} 1 & , \text{ If } v_i \text{ and } v_j \text{ are adjacent to each other} \\ 0 & , \text{ Otherwise} \end{cases}$$

In case of directed graph the adjacency matrix  $A$  will be defined as:

$$a_{ij} = \begin{cases} 1 & , (v_i, v_j) \in G \\ 0 & , \text{ Otherwise} \end{cases}$$

## 6.9 Bipartite graph $(G(V_1, V_2, E))$

Let  $G$  be a graph with vertex set  $V$ . Then if we can write  $V$  as disjoint union of two vertex sets  $V_1$  and  $V_2$  such that each edge of  $E$  has one vertex in  $V_1$  and other vertex in  $V_2$  then the graph is called bipartite graph.

A bipartite graph is called complete bipartite graph  $K(m, n)$  if each element of  $V_1$  is connected to each element of  $V_2$ , here  $m$  and  $n$  are the number of elements in set one and two respectively.

$K(3, 3)$