

$$x := x + 1;$$

# Expressions and assignments

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## Arithmetic expressions

Operator precedence or strength: order of evaluation of expressions with more than one operator. (Parentheses can always be used to specify order explicitly.)

Operators are usually grouped:

- exponentiation ( $**$ )
- unary operators (abs, unary + and - etc.)  
(not is a unary operator, too)
- multiplicative ( $*$ ,  $/$ ,  $\text{div}$ ,  $\text{mod}$  etc.)  
(and is multiplicative, too)
- additive (binary +, binary - etc.)  
(or is additive, too)

Rules are confusing and widely different.

Pascal:            multiplicative > additive

C: self-increment, unary > multiplicative > additive

Ada:     $**$  > multiplicative > unary +, - > additive

Fortran:     $**$  > multiplicative > additive

## Associativity

Let  $\mathbb{Y}$  be any binary operator.

Left to right:  $x \mathbb{Y} y \mathbb{Y} z = (x \mathbb{Y} y) \mathbb{Y} z$

Pascal, Ada, C (all normal operators)

Right to left:  $x \mathbb{Y} y \mathbb{Y} z = x \mathbb{Y} (y \mathbb{Y} z)$

C (self-increment operators ++ and --)

Nonassociative (exponentiation in Ada):

$x**y**z$  is syntactically incorrect, though  
( $x**y$ ) $**z$  and  $x**(y**z)$  are OK.

## No precedence, one associativity rule

This is what we have in APL: always right to left.

$x + y * z$  means  $x + (y * z)$

$x * y + z$  means  $x * (y + z)$  (!?)

In Smalltalk: always left to right.

$x + y * z$  means  $(x + y) * z$  (!?)

$x * y + z$  means  $(x * y) + z$

## Evaluation of arguments, side-effects

A function that appears in an expression may have a side-effect (change to some non-local object, not mentioned in the expression). Example:

```
function twice(var x: real): real;
begin x := x + x; twice := x end;
```

In the statement

```
z := twice(y);
```

the value of  $y$  is changed "secretly".

Such effects are to be avoided, if possible. Built-in functions seldom have side-effects, arithmetic functions—never!

## Conditional expressions

In Algol 60:

```
if x > 0 then 1 else
if x = 0 then 0 else -1
```

The same in C:

```
(x > 0 ? 1 : (x == 0 ? 0 : -1))
```

## Overloading

One name or symbol—more than one distinct use.

Examples in Pascal:

```
+
integer addition, floating-point addition,
string concatenation, set union

*
integer and floating-point multiplication,
set intersection

abs
integer → integer, real → real

mod, div
no overloading
```

Overloading can be always resolved by context (all operands have known types):

```
2 + 3: integers, ['a'] + ['c', 'd']: sets
```

In Ada, overloading is an important element of design. Ada is extendible: a new meaning can be given to an operator in addition to what it already means. Overloading is also possible in C++.

```
type date is
record
  day: 1..31;
  month: 1..12;
  year: 1000..9999;
end record;

DT: date;

function "+"(D: date; I: integer)
  return date is
NewD: date;
begin
  -- code that assigns D + I days
  -- to NewD;
  return NewD;
end;
```

Now, we can use this operator as follows:

```
DT := (18, 10, 1994) + 2;
```

Ada also has only two standard I/O procedure names: get, put—both heavily overloaded.

Overloading may be quite confusing. In C:

& means "bitwise conjunction" and "address of",

\* means "multiplication" and "dereference".

In PL/I, = means "equality" and "assignment".

In Prolog, the comma is overloaded in a rather careless manner.

This comma can be read "and":

```
a :- b, c, d.
```

This comma separates arguments:

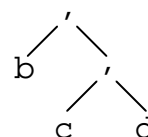
```
a(b, c, d)
```

This comma separates list elements:

```
[b, c, d]
```

This comma is a functor:

```
(b, c, d) means
```



## Coercion

If objects of two numeric types appear as operands, we "upgrade" the "lower" type.

Hierarchy in Fortran:

```
integer < real < double < complex
```

In Pascal:

```
integer < real
```

## Logical expressions

### Six comparison operators

equal	=	==	.EQ.	
not equal	<>	/=	!=	.NE.
less	<			.LT.
less or equal	<=	=<		.LE.
greater	>			.GT.
greater or equal	>=	=>		.GE.

Equality is well-defined for all types, but no natural ordering usually exists for non-scalar types.

```
function "<"(L, R: date) return boolean is
earlier: boolean := true;
begin
  if L.year > R.year then
    earlier := false;
  elsif L.year = R.year then
    if L.month > R.month then
      earlier := false;
    elsif L.month = R.month then
      if L.day >= R.day then
        earlier := false;
      end if;
    end if;
  endif;
  return earlier;
end;
```

### Testing that involves sets

equality: =, <>  
 membership: in, not in  
 inclusion (only in Pascal): <=, >=

### Logical operators

Pascal: not, and, or

Ada: not, and, and then, or, or else, xor

Short-circuit operations are based on these facts:

true or else ANYTHING   ≡ true  
 false and then ANYTHING   ≡ false

Evaluating P and Q and R may mean evaluating all of P, Q, R, or stopping after the first false.

P and then Q and then R must stop after computing the first false.

Similarly, P or else Q or else R must stop after computing the first true.

A: array(1..10) of integer;  
 Unsafe: if n>10 or (A(n)=0) then --  
 Safe: if n>10 or else (A(n)=0) then --

## Assignment

A single assignment is obvious:

```
target := expression
target = expression
target ← expression
```

Multiple assignment is more interesting.

PL/I: A, B := EXPR;

Algol 60: A := B := EXPR;

- (1) Find the value of EXPR.
- (2') Assign this value to A, then B.
- (2'') Assign this value to B, then A.

This is not quite unimportant. Consider

```
I := 5; A[I] := I := 10;
```

The order in which target addresses are found also matters!

- (1) Find all target addresses.
- (2) Find the value of EXPR.
- (3) Assign this value to A and B.

With this method, A[5] := 10.

- (1) Find the value of EXPR.
- (2) Find target addresses left-to-right, assign the value to every address.

With this method, still A[5] := 10.

- (1) Find the value of EXPR.
- (2) Find target addresses right-to-left, assign the value to every address.

With this method, A[10] := 10.

This statement in C is not a multiple assignment:

```
A = B = EXPR;
```

Here, B = EXPR has a value (the value of EXPR) that is next assigned to A: the assignment operator in C associates right-to-left.

Another syntactic addition in C: mixing assignment with arithmetic.

```
A += B; means A = A + B;
A *= B; means A = A * B; etc.
```

Finally, we can have conditional targets (in C++):

```
(x != 0 ? y : z) = 17;
or (even less readable):
x ? y : z = 17;
```

Summary

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