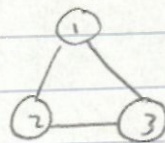


#5 Algorithms HW 1

3.1

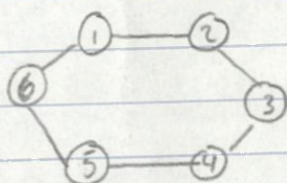
$n \times n$



Elizabeth Port

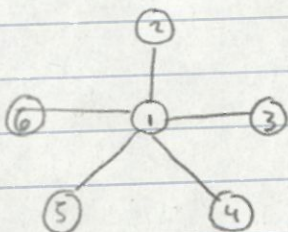
$$\begin{matrix} n=1 \\ n=2 \\ n=3 \end{matrix} \Rightarrow \begin{matrix} n=1 & n=2 & n=3 \\ \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

For Ring, boolean matrix might look like



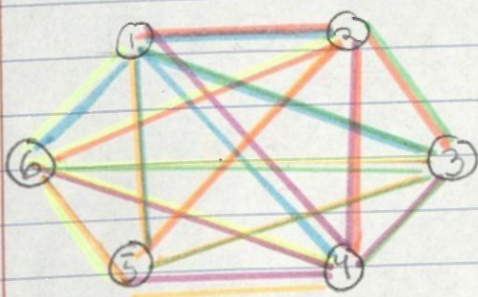
$n =$	1	2	3	4	5	6
$n=1$	0	1	0	0	0	1
$n=2$	1	0	1	0	0	0
$n=3$	0	1	0	1	0	0
$n=4$	0	0	1	0	1	0
$n=5$	0	0	0	1	0	1
$n=6$	1	0	0	0	1	0

For Star, matrix might look like



	1	2	3	4	5	6
1	0	1	1	1	1	1
2	1	0	0	0	0	0
3	1	0	0	0	0	0
4	1	0	0	0	0	0
5	1	0	0	0	0	0
6	1	0	0	0	0	0

For fully connected, matrix might look like



	1	2	3	4	5	6
1	0	1	1	1	1	1
2	1	0	1	1	1	1
3	1	1	0	1	1	1
4	1	1	1	0	1	1
5	1	1	1	1	0	1
6	1	1	1	1	1	0

Identifying a Ring:

```
// Input: Boolean adjacency matrix
// Output: True / False if Matrix is
// a representation of a ring network topology
For y = 0 → n-1 {
    • Check that exactly two indices
      are non zero
      For x = 0 → n-1 {
          Array[x] != 0 { n++ }
          if n++ > 2 return Not Ring }
      }
    repeat for all y-indices.
    return is Ring
```

Identifying a Star:

```
// Input: Boolean adjacency matrix
// Output: True / False if Matrix is Star network topology

For y = 0 → n-1 {
    • Check if all indices are non-zero except 1
      For x = 0 → n-1 {
          Array[x] = 0 { n++ } }
      if n ≤ 1: fully Connected ++ }
    • if less or more than 2 fully connected lines exist,
      adjacency matrix does not represent a star
      if (fullyConnected == 2): return is Star
      else (return NotStar)
```


Identifying a fully connected mesh.

// input: Boolean adjacency matrix

// output: T/F depending if matrix represents fully connected mesh

• General approach: ensure each x row & y column does not contain more than 1 non-zero

```
for  $y = 0 \rightarrow n-1$  {  
    for  $x = 0 \rightarrow n-1$  {  
        if (Array[x] == 0) {  $n++$  } // check all  $x$  indices  
        if  $n > 1$  { return NotMesh } // only contain 1  $\phi$   
    }  
    if (Array[y] == 0) {  $z++$  } // check all  $y$  indices  
    if  $z > 1$  { return NotMesh } // only contain 1  $\phi$   
}
```

Run the above three programs sequentially to identify which of the 3 topologies the matrix represents.
If none return true, matrix is not a known topology.

Because the time this algorithm takes is directly related to the # of nodes (and is the size of the matrix) it is a linear time complexity.

3.1 #7

a) // Assume weight^{of each stack} is stored in $A + B$
if $(A > B)$ {
 A is stack of fake coins
else { A is stack of real coins }

This algorithm is $O(1)$, because no matter how many coins are in the stacks, only 1 comparison is needed.

b) 2 weighings are required each time new coins are added. This is because we only need the current weights of each stack to determine which is heavier.

3.2 #5

How many comparisons will be made in searching for the following patterns in a text of 1,000 zero's

a) ... 0 0 0 0 0 0 0 0 0 0 ... looking for: 00001
 0 0 0 0 1 • Five comparisons made
 0 0 0 0 1 • Five comparisons made.

For each digit of the 1,000 0's 5 comparisons will be made
 $1,000 \times 5 = 5,000$ comparisons will be made

b) ... 0 0 0 0 0 0 0 0 0 0 ... looking for: 10000
 | • 1 comparison
 | • 1 comparison.

For each digit of the 1,000 0's, 1 comparison is made.

So 1,000 total comparisons will be made

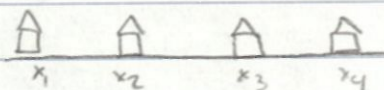
c) ... 0 0 0 0 0 0 0 0 0 0 ... looking for: 01010
 0 1 • 2 comparisons

For each digit of the 1,000 0's, 2 comparisons will be made. So 2,000 total comparisons

3.2 #7

Depending on the string & pattern be searched / searched for, there may be a very great advantage. For example for text = 0000000000 & pattern = 0000000001 searching left to right would require 90 comparisons, but searching right to left would require 10. The opposite could also be true for a different string, or it may make no difference: it depends on the string & pattern used.

3.3 #3



a) Find average distance: $\frac{1}{n} \sum |x_i - x_j|$

//input: Array of n ordered values

//output: location of post office w/ min average distance
if ($n > 1$)

index $\leftarrow \lceil n/2 \rceil$

else

index $\leftarrow 0$

return index

b) //input: Array of n ordered values

//output: location x so that x to post office is minimized

define: $m = \frac{x_1 + x_n}{2}$

while $x_i < m$: $i = i + 1$

if $x_i - x_1 < x_n - x_{i-1}$

location = i

else

location = $i - 1$

return location