

$\frac{\text{T-HMAP} \quad \frac{\overrightarrow{\vdash v_k : (\mathbf{Value} k)}}{\vdash \{(v_k v_v)\} : (\mathbf{HMap}^C p)} \quad \frac{\overrightarrow{\vdash v_v : \tau_v} \quad p = \{(k, \tau_v)\}}{\vdash \{(v_k v_v)\} : (\mathbf{HMap}^C p) ; \mathbb{tt} \mathbb{ff} ; \emptyset}}{\vdash \{(v_k v_v)\} : (\mathbf{HMap}^C p) ; \mathbb{tt} \mathbb{ff} ; \emptyset}$		$\text{T-KW} \quad \Gamma \vdash k : (\mathbf{Value} k) ; \mathbb{tt} \mathbb{ff} ; \emptyset$
$\text{T-GETHMAP} \quad \frac{\Gamma \vdash e_m : (\bigcup (\mathbf{HMap}^A p a)) ; \psi_+ \psi_- ; o \quad \Gamma \vdash e_k : (\mathbf{Value} k) \quad \overrightarrow{(k \tau) \in p}^i}{\Gamma \vdash (\text{get } e_m e_k) : (\bigcup \overrightarrow{\tau}^i) ; \mathbb{tt} \mathbb{tt} ; \mathbf{key}_k(x)[o/x]}$	$\text{T-GETHMAPABSENT} \quad \frac{\Gamma \vdash e_m : (\mathbf{HMap}^A p a) ; \psi_+ \psi_- ; o \quad \Gamma \vdash e_k : (\mathbf{Value} k) \quad k \in a}{\Gamma \vdash (\text{get } e_m e_k) : \mathbf{nil} ; \mathbb{tt} \mathbb{tt} ; \mathbf{key}_k(x)[o/x]}$	
$\text{T-GETHMAPPARTIALDEFAULT} \quad \frac{\Gamma \vdash e_m : (\mathbf{HMap}^P p a) ; \psi_+ \psi_- ; o \quad \Gamma \vdash e_k : (\mathbf{Value} k) \quad (k \tau) \notin p \quad k \notin a}{\Gamma \vdash (\text{get } e_m e_k) : \top ; \mathbb{tt} \mathbb{tt} ; \mathbf{key}_k(x)[o/x]}$	$\text{T-ASSOCHMAP} \quad \frac{\Gamma \vdash e_m : (\mathbf{HMap}^A p a) \quad \Gamma \vdash e_k : (\mathbf{Value} k) \quad \Gamma \vdash e_v : \tau \quad k \notin a}{\Gamma \vdash (\text{assoc } e_m e_k e_v) : (\mathbf{HMap}^A(p, (k \tau)) a) ; \mathbb{tt} \mathbb{ff} ; \emptyset}$	

Figure 5. Map Typing Rules

$\frac{\text{TA-LOCAL} \quad \Sigma(x) = \gamma}{\Sigma \vdash x : \gamma}$	$\text{TA-NIL} \quad \Sigma \vdash \text{nil} : ?$	$\text{TA-TRUE} \quad \Sigma \vdash \text{true} : \mathbf{Boolean}$	$\text{TA-FALSE} \quad \Sigma \vdash \text{false} : \mathbf{Boolean}$	$\text{TA-KW} \quad \Sigma \vdash k : \mathbf{Keyword}$	$\text{TA-CLASS} \quad \Sigma \vdash C : \mathbf{Class}$
$\text{TA-NEWSTATIC} \quad \Sigma \vdash (\text{new}_{[[\vec{C}_i], C_1]} C \vec{e}) : C_1$	$\text{TA-NEWREFL} \quad \Sigma \vdash (\text{new } C \vec{e}) : ?$	$\text{TA-ABS} \quad \Sigma \vdash \lambda x^\tau. e : ?$	$\frac{\text{TA-LETHINT} \quad \Sigma, x : C \vdash e : \gamma}{\Sigma \vdash (\text{let } [\wedge C x e_1] e) : \gamma}$	$\frac{\text{TA-LET} \quad \begin{array}{c} \Sigma \vdash e_1 : \gamma_1 \\ \Sigma, x : \gamma_1 \vdash e : \gamma \end{array}}{\Sigma \vdash (\text{let } [x e_1] e) : \gamma}$	
$\text{TA-FIELDREFL} \quad \Sigma \vdash (. e fld) : ?$	$\text{TA-FIELDSTATIC} \quad \Sigma \vdash (. e fld_C^{C_1}) : C$	$\text{TA-METHODREFL} \quad \Sigma \vdash (. e (mth \vec{e}_i)) : ?$	$\text{TA-METHODSTATIC} \quad \Sigma \vdash (. e (mth_{[[\vec{C}_i], C]}^{C_1} \vec{e}_i)) : C$	$\text{TA-APP} \quad \Sigma \vdash (e e') : ?$	
$\frac{\text{TA-APLOCAL} \quad \Sigma(x) = \gamma}{\Sigma \vdash (x e') : \gamma}$	$\frac{\text{TA-DO} \quad \Sigma \vdash e : \gamma}{\Sigma \vdash (\text{do } e_1 e) : \gamma}$	$\text{TA-DEFMULTI} \quad \Sigma \vdash (\text{defmulti } \tau e) : ?$	$\text{TA-DEFMETHOD} \quad \Sigma \vdash (\text{defmethod } e_1 e_2 e_3) : ?$	$\frac{\text{TA-IF} \quad \begin{array}{c} \Sigma \vdash e_2 : C \\ \Sigma \vdash e_3 : C \end{array}}{\Sigma \vdash (\text{if } e_1 e_2 e_3) : C}$	
	$\text{TA-IFUNKNOWN} \quad \Sigma \vdash (\text{if } e_1 e_2 e_3) : ?$	$\text{TA-ISA} \quad \Sigma \vdash (\text{isa? } e e) : \mathbf{Boolean}$	$\text{TA-CONST} \quad \Sigma \vdash c : ?$		

Figure 6. tools.analyzer Type Hints