

# **RSA ALGORITHM TO ENCRYPT AND DECRYPT THE DATA**

# OBJECTIVE

To implement RSA  
Algorithm to Encrypt and  
Decrypt data.

# INTRODUCTION

- ❖ **Internet Security** is becoming more important everyday.
- ❖ **Cryptography** is the heart of security.
- ❖ In order to protect our data, **encryption** is done at sender side and **decryption** is carried out at receiver's side.

# INTRODUCTION

## ❖ SECURITY ASPECTS.

### ❖ AUTHENTICATION.

- ❖ Authentication provides the identification of the originator.
- ❖ It confirms to the receiver that the data received has been sent only by an identified and verified sender.

### ❖ CONFIDENTIALITY.

- ❖ Confidentiality is the fundamental security service provided by cryptography.
- ❖ It is a security service that keeps the information confidential from an unauthorized person.
- ❖ It is sometimes referred to as *privacy* or *secrecy*.

### ❖ NONREPUDIATION.

- ❖ It is a security service that ensures that an entity cannot refuse the ownership of a previous commitment or an action.
- ❖ It is an assurance that the original creator of the data cannot deny the creation or transmission of the said data to a recipient or third party.

### ❖ DATA INTEGRITY.

- ❖ It is security service that deals with identifying any alteration to the data.
- ❖ The data may get modified by an unauthorized entity intentionally or accidentally.
- ❖ Integrity service confirms that whether data is intact or not since it was last created, transmitted, or stored by an authorized user.

# INTRODUCTION

## ❖ HASHING.

- ❑ Creating a miniature version of a message instead of original message.

## ❖ AUTHENTICATION OF PEOPLE.

## ❖ KEY MANAGEMENT

# INTRODUCTION

## ❖ SECURITY AND INTERNET MODEL.

- ❖ OSI model provided encryption/ decryption services in the presentation layer which does not exist in the internet model.
- ❖ At which layer security needs to be implemented in internet model?
- ❖ Security breach can happen in any of the five layers.
- ❖ At physical layer, an intruder can wiretap into the transmission media and read or alter a sequence of bits.
- ❖ At data link layer, frames can be captured and read or altered.
  - ❖ Example: LAN, where transmission is broadcast and every station receives a copy of frame.
- ❖ At the network layer, an IP datagram can be removed, altered, or inserted into the network.
- ❖ At the transport layer, a user datagram or a segment can be captured or altered.
- ❖ Finally, at application layer, the whole message can be altered or read.

# INTRODUCTION

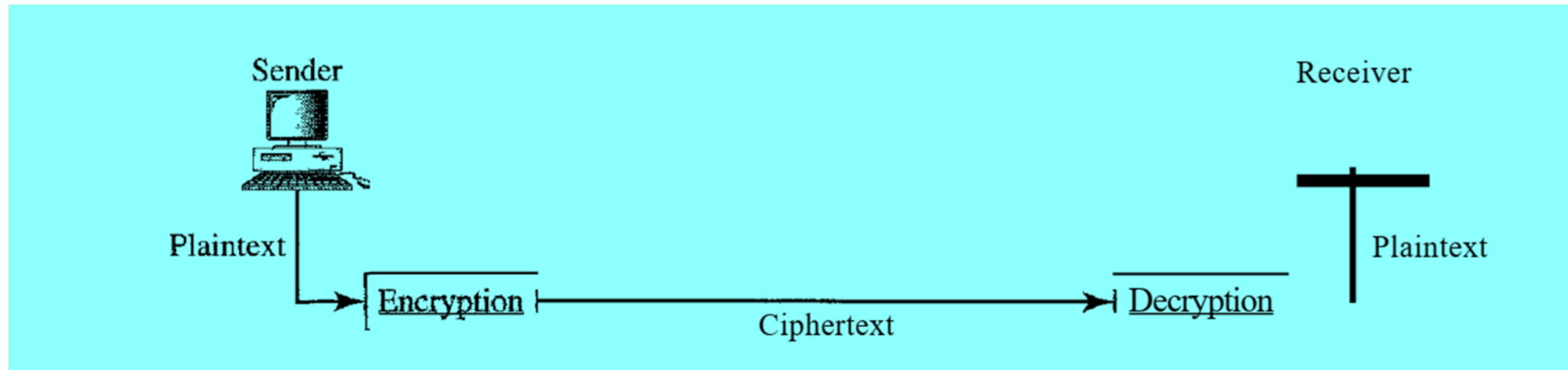
- ❖ **Concerns need be addressed before adding security to the Internet.**
  - ❖ **Security is more effective in which layer?**
  - ❖ **Security is easier to provide in which layer?**
  - ❖ **Security at one level, is that enough?**

# CRYPTOGRAPHY

- ❖ Cryptography in Greek means **“Secret Writing”**.
- ❖ It is an **art and science** of **transforming** messages to make them **secure** and **immune** to attacks.



# CRYPTOGRAPHY COMPONENTS BLOCK DIAGRAM



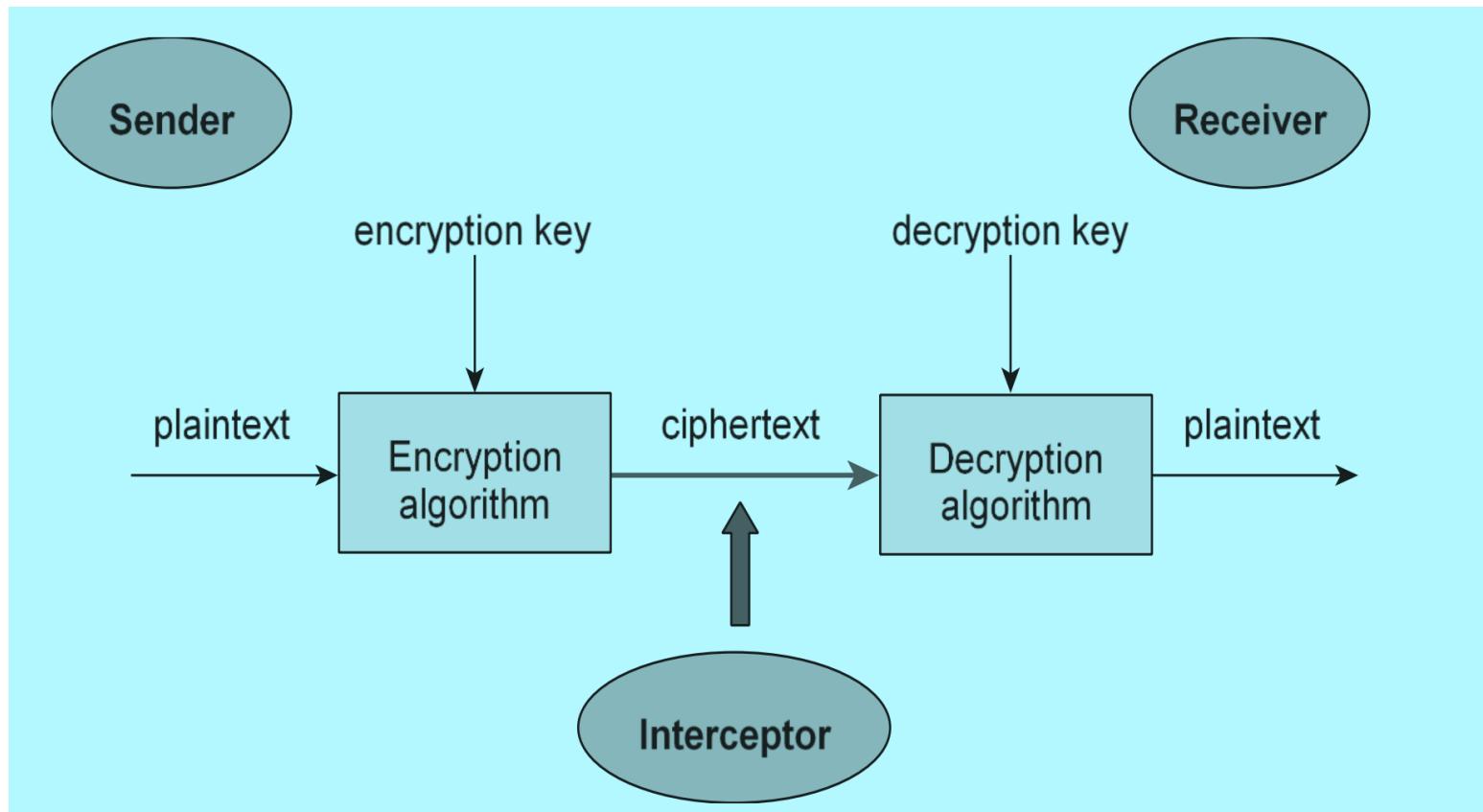
# CRYPTOGRAPHY COMPONENTS

- ❖ The original message, before being transformed, is called *plaintext*.
- ❖ After the message is transformed, it is called *ciphertext*.
- ❖ An *encryption algorithm* transforms the plaintext to ciphertext;
- ❖ A *decryption algorithm* transforms the ciphertext back to plaintext.
- ❖ The sender uses an encryption algorithm and the receiver uses a decryption algorithm.
- ❖ The term *cipher* refers to encryption/ decryption algorithms.
- ❖ A *key* is a number(value) that the cipher, as an algorithm, operates on.

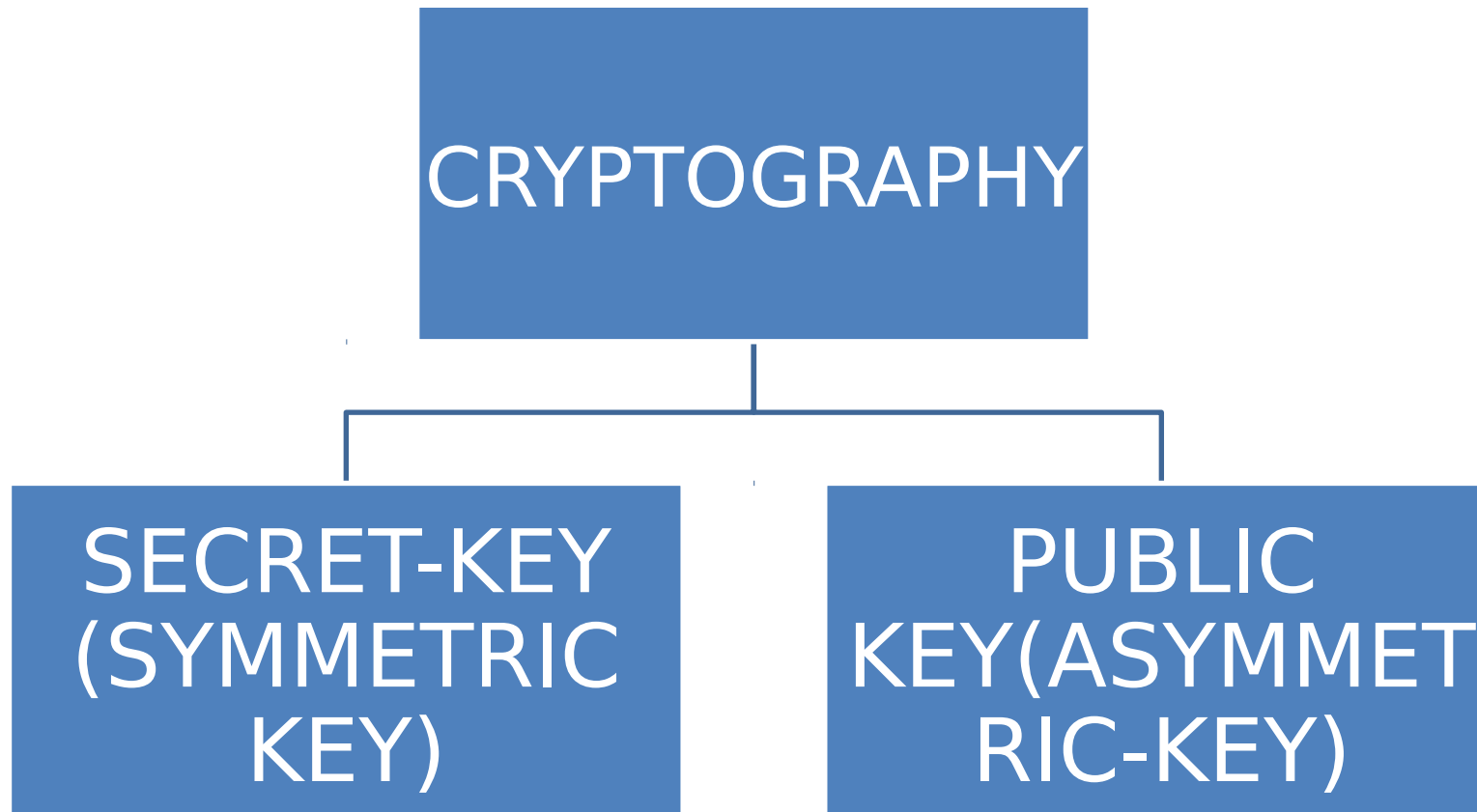
# ENCRYPTION AND DECRYPTION

- ❖ To encrypt a message, *an encryption algorithm, encryption key and plaintext* is required.
- ❖ To decrypt a message, a *decryption algorithm, a decryption key and ciphertext* is required.
- ❖ The encryption and decryption algorithms are public; anyone can access them.
- ❖ The **keys are secret**; they need to be protected.

# ENCRYPTION AND DECRYPTION BLOCK DIAGRAM



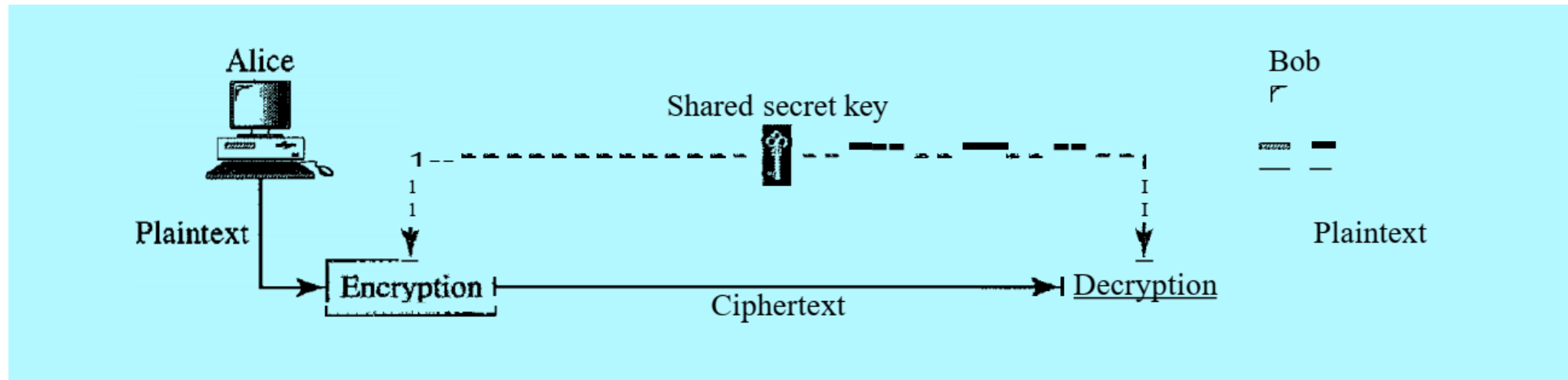
# TYPES OF CRYPTOGRAPHY



# **SYMMETRIC KEY(SECRET-KEY) CRYPTOGRAPHY**

- ❖ **In symmetric key cryptography, the same key is used by both sender and receiver.**
- ❖ **Sender uses this key and an encryption algorithm to encrypt data;**
- ❖ **Receiver uses the same key and a decryption algorithm to decrypt data.**
- ❖ **The key is shared.**

# SYMMETRIC KEY CRYPTOGRAPHY BLOCK DIAGRAM



# DIFFERENCES BETWEEN SYMMETRIC-KEY AND ASYMMETRIC-KEY CRYPTOGRAPHY

## Symmetric-key

- Same key is used for encryption and decryption.
- Used for long messages.
- Each pair of users must have a unique symmetric key: for 'N' users there need to be  $[N(N-1)/2]$  symmetric keys.
- Key distribution between sender and receiver is difficult.
- Simple algorithms.
- Ex: DES, Triple DES.

## Asymmetric-key

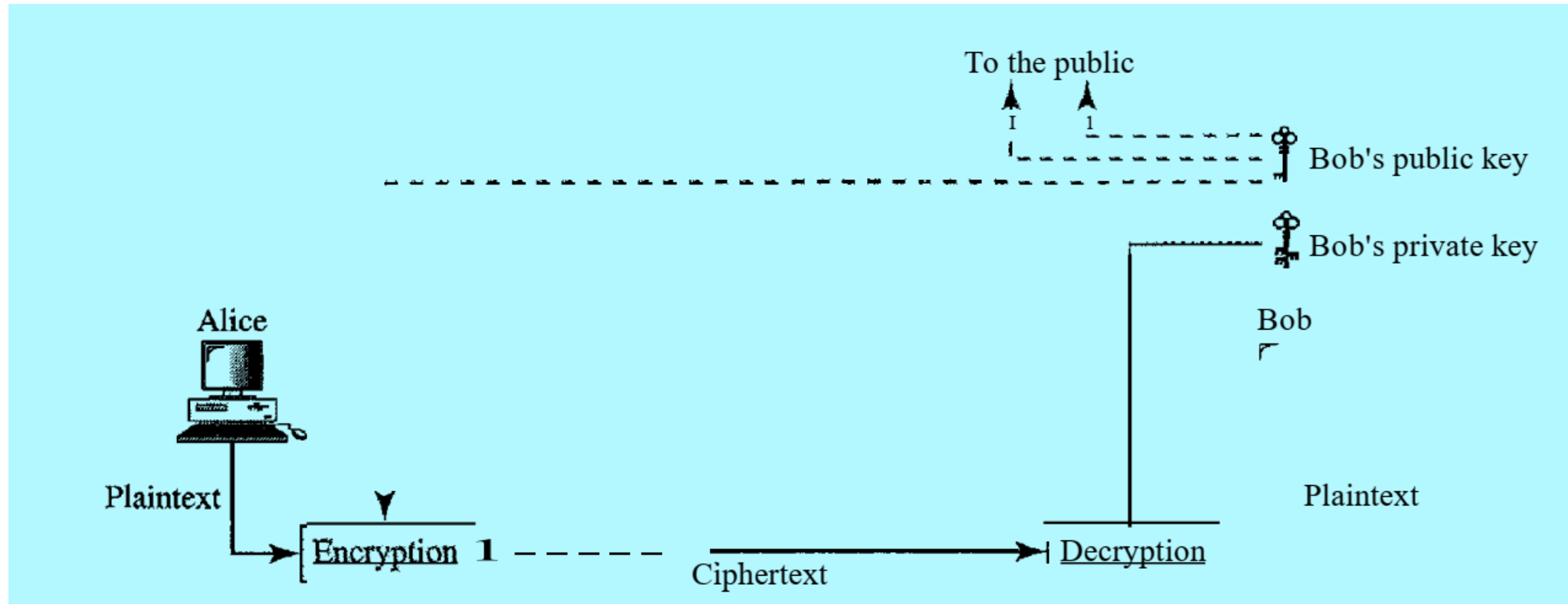
- Two different keys are used for encryption and decryption: Public key and Private key.
- More efficient for short messages.
- Private key is kept by receiver, public key is announced to the public.
- Overcomes the drawback of symmetric-key cryptography, by reducing the number of keys to  $2N$ .
- Complex algorithms.
- Ex: RSA algorithm.



# ASYMMETRIC KEY (PUBLIC-KEY) CRYPTOGRAPHY

- ❖ In public-key cryptography, there are two keys: a private key and a public key.
- ❖ The public key is used for encryption and private key for decryption.
- ❖ Advantages:
  - ❖ Removes the restriction of a shared symmetric key between two entities who need to communicate with each other.
  - ❖ Each entity creates a pair of keys: the private one is kept, and public one is distributed.
  - ❖ Each entity is independent, and a pair of keys created can be used to communicate with any other entity.
  - ❖ The number of keys needed is reduced tremendously to '2N' only for 'N' users.
- ❖ Disadvantages:
  - ❖ Complexity of algorithm: needs large numbers.
  - ❖ Calculating ciphertext from plaintext using the long keys takes a lot of time. Hence not recommended for large amounts of text.
  - ❖ The association between an entity and its public key must be verified. (can be overcome using certification authority)

# ASYMMETRIC KEY (PUBLIC-KEY) CRYPTOGRAPHY BLOCK DIAGRAM



# RSA ALGORITHM (Rivest, Shamir and Adleman)- EXAMPLE: GENERATING PUBLIC KEY

1. Choose two distinct prime numbers,  $p$ , and  $q$ .
2. Compute  $n = p * q$ .
3. Compute the totient of the product as  $\lambda(n) = (p - 1) * (q - 1)$ .  
 $\lambda(n)$  can also be,  $\lambda(n) = \text{lcm}(\lambda(p), \lambda(q)) = \text{lcm}(p - 1, q - 1)$ .
4. Choose any number  $1 < e < \lambda(n)$  that is coprime to  $\lambda(n)$ ,  $e$ .

**Coprime:** two integers  $a$  and  $b$  are said to be relatively, mutually or co prime if the only positive integer (factor) that divides both of them is 1  $a$ , and  $b$  themselves need not be prime.

**Example** 14, 15 are co prime, but they themselves are not prime, only common divisor is 1.

$(n, e)$  is the Public key

# RSA ALGORITHM- EXAMPLE: GENERATING PRIVATE KEY

5. Compute  $d$ , the modular multiplicative inverse of  $e$   
( $\text{mod } \lambda(n)$ )

Inverse  
Multiplicative  
Modular

$$d * e \text{ mod } \lambda(n) = 1$$

(  $n$  ,  $d$  ) is the Private key.

# RSA ALGORITHM- EXAMPLE: ENCRYPTION/ DECRYPTION

- ❖ To encrypt, message  $m$ , into cipher  $c$ ,  $c = (m^e) \bmod n$   
using public key  $(n, e)$  at sender
- ❖ To decrypt, cipher text  $c$  to message  $m$ ,  $m = (c^d) \bmod n$   
using private key,  $(n, d)$  at receiver

# **RSA ALGORITHM- EXAMPLE:**

**$p = 3, \quad q = 11, \quad n = 33$**

**$\lambda(n) = (p - 1) * (q - 1) = (3-1)*(11-1) = 2*10 = 20$**

**Choose any number  $1 < e < \lambda(n)$  that is coprime to  $\lambda(n)$ ,  $e = 3$**

**$(n, e)$ , is the Public key =  $(33, 3)$**

**Compute  $d$ , the modular multiplicative inverse of  $e \pmod{\lambda(n)}$**

**$d$  such that,  $d * e \pmod{\lambda(n)} = 1$**

**$d * e \pmod{\lambda(n)} = d * 3 \pmod{20} = 7 * 3 \pmod{20} = 21 \pmod{20} = 1$**

**$d = 7$**

**$(n, d)$ , is the Private key =  $(33, 7)$**

# **RSA ALGORITHM- EXAMPLE:**

**To encrypt, message m to Cipher text c,  $c = (m^e) \bmod n$**

**Consider m = 2 3 4 , say its b c d, b is 2, c is 3 and d is 4  
 $c = (m^3) \bmod 33$**

$$c = (2^3) \bmod 33 = 8$$

$$c = (3^3) \bmod 33 = 27$$

$$c = (4^3) \bmod 33 = 64 \bmod 33 = 31$$

**For message m = 2 3 4 , Cipher c = 8 27 31**

# RSA ALGORITHM- EXAMPLE:

To decrypt, cipher text  $c$  back to message  $m$ ,  $m = (c^d) \bmod n$

$$m = (c^7) \bmod 33$$

$$c = 8 \ 27 \ 31$$

$$m = (8^7) \bmod 33 = 2$$

$$m = (27^7) \bmod 33 = 3$$

$$m = (31^7) \bmod 33 = 4$$

For Cipher  $c = 8 \ 27 \ 31$  , message  $m = 2 \ 3 \ 4$



# SOURCE CODE

```
int main()
{
    int p, q, n, lambdaN, d, e, length, i;
    int message[10], cipher[10];

    printf("\n Enter two distinct prime numbers p and q: ");
    scanf("%d%d", &p, &q); // Choose two distinct prime numbers, p, and q

    n = p*q; // Compute n = p*q
    lambdaN = (p - 1) * (q - 1); // Compute totient of the product  $\lambda(n)=(p-1)*(q-1)$ 

    printf("\n Enter the Public and Private key, e and d, such that e and ");
    printf("(p-1)*(q-1) are co prime, and d * e mod (p-1)*(q-1) = 1 : ");
    scanf("%d%d", &e, &d);

    printf("\n Enter length of message: ");    scanf("%d", &length);

    printf("\n Enter the message, input 1 for a, 2 for b . . , ");
    printf("separated by space or line: ");
    for( i = 0; i < length; i++ )    scanf("%d", &message[i]);

    printf("\n At sender, encrypt message to cipher, cipher = ");
    for( i = 0; i < length; i++ )
    {
        cipher[i] = ( (long int) pow( message[i], e ) ) % n;
        printf("%d ", cipher[i]);
    }

    printf("\n At receiver, decrypt cipher to message, message = ");
    for( i = 0; i < length; i++ )
        printf("%ld ", ( (long int) pow( cipher[i], d ) ) % n );

    return 0;
}
```



```
#include <stdio.h>
#include <math.h>

int main()
{
    int p, q, n, lambdaN, d, e, length, i;
    int message[10], cipher[10];

    printf("\n Enter two distinct prime numbers p and q: ");
    scanf("%d%d", &p, &q); // Choose two distinct prime numbers, p, and q

    n = p*q; // Compute n = p*q
    lambdaN = (p - 1) * (q - 1); // Compute totient of the product  $\lambda(n)=(p-1)*(q-1)$ 

    printf("\n Enter the Public and Private key, e and d, such that e and ");
    printf("(p-1)*(q-1) are co prime, and d * e mod (p-1)*(q-1) = 1 : ");
    scanf("%d%d", &e, &d);

    printf("\n Enter length of message: ");    scanf("%d", &length);
```

```
printf("\n Enter the message, input 1 for a, 2 for b . . , ");
printf("separated by space or line: ");
for( i = 0; i < length; i++ )          scanf("%d", &message[i]);

printf("\n At sender, encrypt message to cipher, cipher = ");
for( i = 0; i < length; i++ )
{
    cipher[i] = ( (long int) pow( message[i], e ) ) % n;
    printf("%d ", cipher[i]);
}

printf("\n At receiver, decrypt cipher to message, message = ");
for( i = 0; i < length; i++ )
    printf("%ld ", ( (long int) pow( cipher[i], d ) ) % n );

return 0;
}
```



|

```
#include <stdio.h>
#include <math.h>
#include <string.h>
```

```
int gcd(int m , int n )// ALGORITHM Euclid(m, n), gcd: greatest common divisor
{
    // Computes gcd(m, n) by Euclid's algorithm
    int r = 0; // Input: Two nonnegative, not-both-zero integers m and n
    char temp; // Output: Greatest common divisor of m and n
    while( n!=0 ) // while n != 0
    {
        // do
        r = m % n; // r ← m mod n
        m = n; // m ← n
        n = r; // n ← r
    } // done
    return m; // return m
}
```

```

int main()
{
    int p, q, n, lambdaN, d, e, length, i;
    char string[20];
    int message[20], cipher[20];

    printf("\n Enter two distinct prime numbers p and q: ");
    scanf("%d%d",&p, &q); // Choose two distinct prime numbers, p, and q

    n = p*q;                // Compute n = p*q
    lambdaN = (p - 1) * (q - 1); // Compute totient of the product  $\lambda(n)=(p-1)*(q-1)$ 

    // choose number e, such that  $1 < e < \lambda(n)$  and is coprime to  $\lambda(n)$ 
    // Find e, such that  $\gcd(e, \lambda(n)) = 1$ , Greatest common divisor of e and  $\lambda(n)$ 
    // Two numbers e and  $\lambda(n)$  are co prime if  $\gcd(e, \lambda(n)) = 1$ ,
    // That is no common divisor other than 1

    e = 2; // e, such that,  $1 < e < \lambda(n)$  and coprime to  $\lambda(n)$ 
    while ( gcd( e, lambdaN ) != 1 && e < lambdaN )
    {
        e++;
    }
    printf("\n Public key = ( %d, %d )", n, e);
}

```

```

d = 1; // Private key, d, such that,  $d * e \bmod \lambda(n) = 1$ 
while ( ( ( d * e ) % lambdaN ) != 1 )
{
    d++;
}
printf("\n Private key = ( %d, %d )", n, d);

printf("\n Enter the message, lower case characters, no space in between: ");
scanf("%s", string );

length = strlen(string);

for( i = 0; i < length; i++ )
{
    message[i] = string[i] - 'a' ; // save a as 1, b as 2, c as 3 ...
}

printf("\n At sender, encrypt message to cipher, cipher = ");
for( i = 0; i < length; i++ )
{
    cipher[i] = ( (long int) pow( message[i], e ) ) % n;
    printf("\n %c as %d ", message[i] + 'a', cipher[i] );
}

```



```

printf("\n At receiver, decrypt cipher to message, message = ");
for( i = 0; i < length; i++ )
{
    printf("\n %d as %c ", cipher[i],
           (char)( ( (long int) pow( cipher[i], d ) % n ) + 'a' ) );
}

return 0;
}

```

```

/*Output:  ./a.out
Enter two distinct prime numbers p and q: 3 11
Public key = ( 33, 3 )
Private key = ( 33, 7 )
Enter the message, lower case characters, no space in between: dvg

At sender, encrypt message to cipher, cipher =
d as 27
v as 21
g as 18

At receiver, decrypt cipher to message, message =
27 as d
21 as v
18 as g
*/

```

# EXPECTED OUTPUT

## CASE:1

ENTER TWO DISTINCT PRIME NUMBERS P AND Q: **13 17**

**N=221;**

ENTER THE PUBLIC AND PRIVATE KEY, E AND D, SUCH THAT E AND  $(P-1)*(Q-1)$  ARE CO PRIME, AND  $D * E \text{ MOD } (P-1)*(Q-1) =$  **1: 35 11**

ENTER LENGTH OF MESSAGE: **3**

ENTER THE MESSAGE, INPUT 1 FOR A, 2 FOR B . . . , SEPARATED BY SPACE OR LINE: **1 2 3**

AT SENDER, ENCRYPT MESSAGE TO CIPHER, CIPHER = **1 59 61**

AT RECEIVER, DECRYPT CIPHER TO MESSAGE, MESSAGE = **1 2 3**

# EXPECTED OUTPUT

## CASE:2

ENTER TWO DISTINCT PRIME NUMBERS P AND Q: **3 11**

**N=33;**

ENTER THE PUBLIC AND PRIVATE KEY, E AND D, SUCH THAT E AND  $(P-1)*(Q-1)$  ARE CO PRIME, AND  $D * E \text{ MOD } (P-1)*(Q-1) =$  **1: 3 7**

ENTER LENGTH OF MESSAGE: **3**

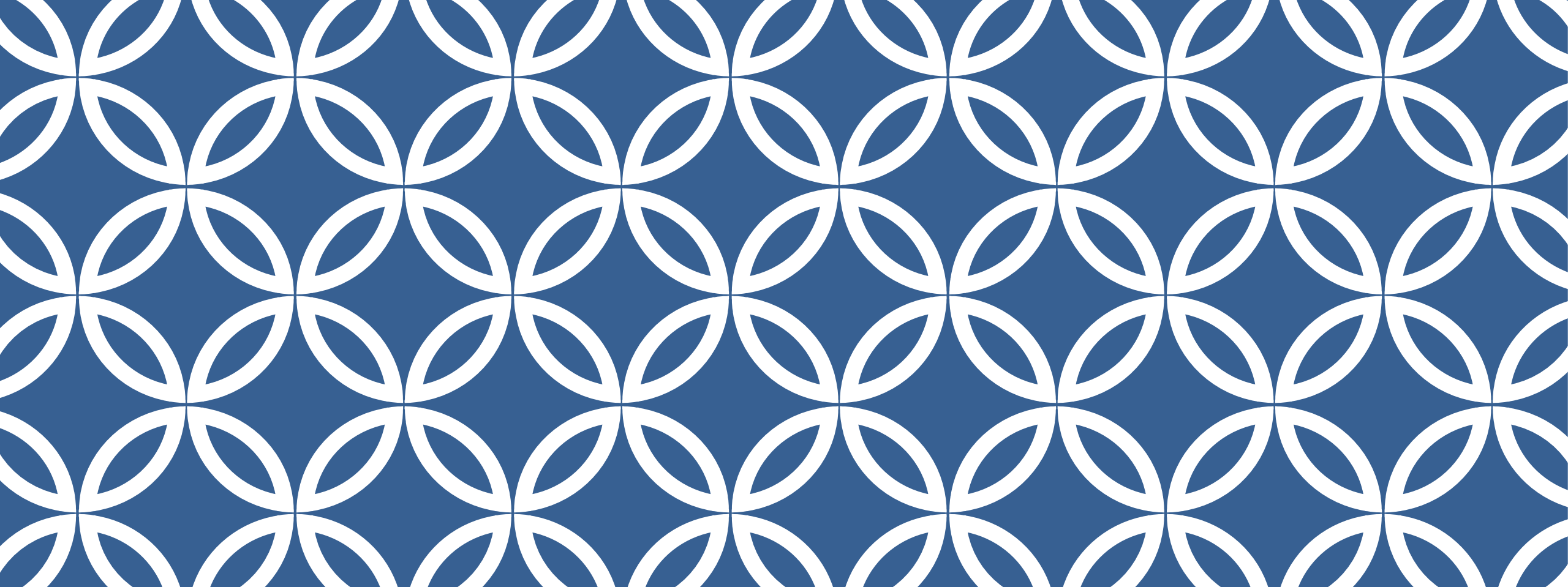
ENTER THE MESSAGE, INPUT 1 FOR A, 2 FOR B . . . , SEPARATED BY SPACE OR LINE: **2 3 4**

AT SENDER, ENCRYPT MESSAGE TO CIPHER, CIPHER = **8 27 31**

AT RECEIVER, DECRYPT CIPHER TO MESSAGE, MESSAGE = **2 3 4**

# REFERENCES

1. **Behrouz Forouzon** - **Data Communications and Networking, McGraw Hill Edition**
2. **Anany Levitin**, **Introduction to the design & analysis of algorithms**
3. **Thomas H. Cormen , Charles E. Leiserson , Ronald L. Rivest, Clifford Stein**  
**Introduction to Algorithms**



# EXPECTED OUTCOME

Students will be able  
to implement RSA  
Algorithm to Encrypt  
and Decrypt data.