



On the Foundations of Dynamic Games and Probability: Decision Making in Stochastic Extensive Form

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1. Game theory as a language

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2. Overview

Why we need an abstract graded language of dynamic games and probability

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3. Extensive form characteristics via stochastic decision forests

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Continuous-time stochastic timing

Section 1

Game theory as a language

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Why a theory? Language: Confusion and Comparison

Equilibrium, saddle point, Nash equilibrium, subgame-perfect Nash equilibrium, Bayesian Nash equilibrium, perfect equilibrium, Markov perfect equilibrium, perfect Bayesian equilibrium, correlated equilibrium, coarse-correlated equilibrium, nash equilibrium in stochastic differential games, equilibrium in mean-field games, ... strategy, pure strategy, mixed extended strategy, behaviour strategy, randomised strategy mixed extended strategy, ... zero-sum game, strategic / normal form game, extensive form game, stochastic game, dynamic game, differential game, stochastic differential game, Dynkin game, timing game, Bayesian game, Stackelberg game, principal-agent problem, interactive decision problem, stochastic optimal control, dynamic programming principle.

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- ▶ When using game-theoretic *words*, they must have a *meaning* in terms of first principles from decision theory.
- ▶ Not clear for complex **dynamics** and **uncertainty**: *Confusion*.
- ► → Abstract language for formally comparing usages and meanings.

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- a map from states to consequences.
- (3) Economically rational agents' decision making is in strategic **equilibrium**.



What about dynamics and probability?

(4) Dynamic decision making / choice under uncertainty = selecting a **strategy**, i.e. a complete contingent plan of action, alias

 $s: \{ Outside, \Pi, \bullet \} \rightarrow \{ enter \ \Pi \ , enter \ \bullet \ , eat pasta, eat fish \}.$

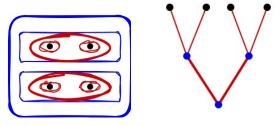


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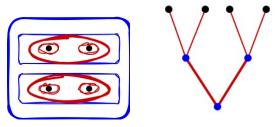


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(5) Representing preference via expected utility:

$$\mathrm{EU} = p_1 \cdot u(\blacksquare \, \mathsf{pasta}) + p_2 \cdot u(\blacksquare \, \mathsf{fish}) + p_3 \cdot u(\blacksquare \, \mathsf{pasta}) + p_4 \cdot u(\blacksquare \, \mathsf{fish}).$$

Core challenge: Complex dynamics and uncertainty



Figure: A more complex tree. By CMG Lee, CC BY-SA 4.0, https://commons.wikimedia.org/w/index.php?curid=153344713, 21st Jul 25.

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- ▶ When the horizontal structure ("local states") is complex, measurability becomes a technical and conceptual issue:
 - Precisely, even if one is able to distinguish all states, not all sets of states need be measurable.
 - Any a priori selection of σ -algebras on states is **not completely objective**. Hence, the game's rules cannot depend on it.

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- When the vertical structure ("dynamics") is complex, the notion "outcome generated by strategies" easily becomes ill-defined and specifying it gets subtle (e.g. reaction).

Section 2

Overview

Why we need an abstract graded language of dynamic games and probability

▶ Construct a — refined partitions / tree-based — extensive form theory incorporating the *measurability* of $g:[0,1] \to \mathbb{A}$ describing a choice g(v), based on a signal $v \sim \mathsf{Unif}[0,1]$.

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- ▶ Interpret the use of the words "strategy", "equilibrium", "subgame", etc. in terms of game-theoretic first principles for *continuous-time stochastic games*. Applies to important models used in finance, economics, engineering, e.g.:
 - Stochastic differential games:

$$\mathrm{d}\chi_t^s = V(\xi_t^s, \chi_t^s) \, \mathrm{d}\eta_t; \qquad \mathbb{E}_x u_i(\chi^s) \geq \mathbb{E}_x u_i(\chi^{(\tilde{s}^i, s^{-i})}), \quad \forall x, i, \, \tilde{s}^i;$$

► Timing games: $\chi_t^k = 1\{\tau^k > t\}$,

$$u_i(\chi^i,\chi^j) = L^i_{\tau^i} \, 1\{\tau^i < \tau^j\} + M^i_{\tau^i} \, 1\{\tau^i = \tau^j\} + F^i_{\tau^j} \, 1\{\tau^j < \tau^i\}.$$



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- ▶ Aim: *Intrinsic* language of **extensive form characteristics** that also arises *extrinsically* as a limit of strict stochastic extensive form theory.
 - EFC = the flow of information about past choices and exogenous events, along with a set of adapted choices locally available to decision makers

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Why stochastic decision forests and adapted choices?



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- Instead of a dynamically acting nature agent \to a static exogenous scenario space (Ω, \mathscr{E}) plus dynamic updates $\mathscr{F}_{\mathbb{Z}}$

What are stochastic decision forests and adapted choices?

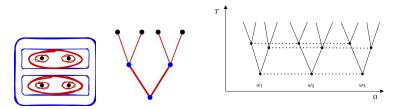


Figure: Decision trees via refined partitions and graphs; a quite simple stochastic decision forest

| Main notions | Main results | |
|----------------------------|---|--|
| Decision forest | Forest of decision trees = Decision forest | |
| Stochastic dec. for. (SDF) | Exogenous uncertainty for decision trees ✓, Induced decision tree | |
| | | |
| Exog. inform. structure | Filtration-like exogenous information ✓ | |
| Adapted choice | Exog. measurability + endog. partitions ✓ | |
| Action path SDF data | Consistent model based on outcomes | |
| | $(\omega,f)\in\Omega	imes\mathbb{A}^{\mathbb{T}}$, \mathbb{T} time, \mathbb{A} action space \checkmark | |

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| Main notions | Main results |
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| Information | = Endogenous Exogenous |
| Strategy | Adaptedness w.r.t. exog. info. ← choices |
| Well-posedness (WP) | ⇔ Well-posedness of scwise ext. forms |
| | \Leftrightarrow Simple order-theoretic properties of F |
| | \Rightarrow complex dynamic exogenous noise possible |
| (Random) histories | (Given WP:) \leftrightarrow (Random) moves, and |
| | provide a model of "subgames" (given WP) |
| Action path "SEF" | Covered by the theory of SEF, |
| | Well-posedness $\Leftrightarrow \mathbb{T}$ is (ess.ly) well-ordered |
| Equilibrium | Unifies a large class of equilibrium concepts |

Related to [Sel65, Har67,68,68a, AFR16].

Section 4

Extensive form characteristics beyond extensive forms

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Continuous-time stochastic timing

- (1) We are interested in games described via stochastic processes. By the preceding theory, we obtain:
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 - ▶ But otherwise: the model is based on outcomes, strategies induce outcome processes, and preferences are defined on outcome processes!

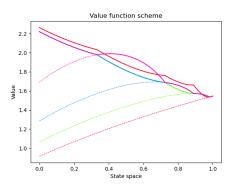


Figure: A symmetric four-player stochastic timing game in continuous time, with underlying diffusion $\mathrm{d}X_t=0.1(1-X_t^2)\,\mathrm{d}t+0.3\sqrt{X_t(1-X_t)}\,\mathrm{d}B_t$

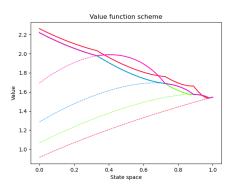


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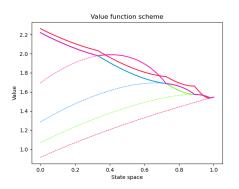


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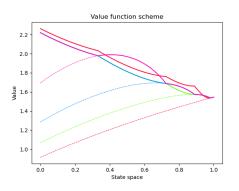


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- ▶ Approaches: additional sharing rules, extended mixed strategies for discrete-time "limit" payoffs (cf. [FT85, Ste16, RS17, Ste18])

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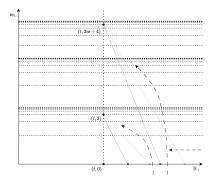


Figure: Vertically extended continuous time and tilting convergence

Stochastic analysis in vertically extended continuous time

| Main notions | Main results |
|---|--|
| $\overline{\mathbb{T}} = [\mathbb{R}_+ 	imes (\mathfrak{w}_1 + 1)] \cup \{\infty\}$ | Smallest complete lattice |
| | containing all tilted well-orders |
| | embeddable into \mathbb{R}_+ |
| $\overline{\mathbb{T}} \leftarrow \text{order topology; } \sigma\text{-algebra } \mathscr{P}_{\overline{\mathbb{T}}}$ | Continuous $f : \overline{\mathbb{T}} \to \mathbb{R}$ mb., |
| generated by $\overline{\mathbb{T}} \to \overline{\mathbb{R}_+} \times \{0,\dots,\alpha\}$, | Mb. Section and Projection Theorem, |
| and $\mathscr{B}_{\overline{\mathbb{R}_+} 	imes \{0,,lpha\}}$, $lpha \in \mathfrak{w}_1$ | Prog. measurability of $\llbracket 	au rbracket$ |
| Optional times | Characterisation of optional times, |
| | Début Theorem, |
| | Basic theory of optional / |
| | predictable processes |
| Tilting convergence | Optional processes arise from opt. |
| | processes defined on deterministic |
| | classical grids by closure under |
| | a) continuous binary operations, |
| | b) pointwise convergence, and |
| | c) tilting convergence. |

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- ▶ $W = \text{the set of } (\omega, h), \ \omega \in \Omega, \ h : \overline{\mathbb{T}} \to \mathbb{B} \text{ decreasing,}$ right-continuous, inert above $\mathfrak{w}, \ h(\infty) = 0 : \text{configuration space.}$

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- Let \mathscr{H}^i be a suitably augmented filtration on W such that \mathscr{H}^i_t contains $\mathscr{F}^i_t \otimes \mathscr{B}^{[0,t]_{\overline{\mathbb{T}}}}_{\mathbb{B}} \otimes \{\emptyset, \mathbb{B}^{[t,\infty]_{\overline{\mathbb{T}}}}\}|_W$, t>0, and $\mathscr{H}^i_0 = \mathscr{F}^i_0 \otimes \{\emptyset, \mathbb{B}^{\overline{\mathbb{T}}}\}|_W$: basic information structure.

- $(\Omega, \mathscr{E}, \mathscr{G}, \mathbb{P})$: a completely filtered probability space with time $\overline{\mathbb{T}}$.
- ▶ $I = \{1, 2\}$, $\mathbb{B} = \{0, 1\}^{I}$: a two-player set, action / state space.
- ▶ $W = \text{the set of } (\omega, h), \ \omega \in \Omega, \ h \colon \overline{\mathbb{T}} \to \mathbb{B} \text{ decreasing,}$ right-continuous, inert above $\mathfrak{w}, \ h(\infty) = 0 \colon \text{configuration space.}$
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- ▶ Let $\mathscr{M}^i = \operatorname{Prd}(\mathscr{H}^i) \vee \operatorname{Opt}(\mathscr{F}^i \otimes \{\emptyset, \mathbb{B}^{\overline{\mathbb{I}}}\})$: Meyer information structure.

For any $i \in I$, let \mathcal{S}_0^i be the set of \mathscr{H}^i -progressively measurable, processes $s^i \colon \overline{\mathbb{T}} \times W \to \{0,1\}$, lower semicontinuous from the right, inert above \mathfrak{w} , and satisfying both $s_t^i(\omega,h) \leq h^i(t-)$, for all $(t,\omega,h) \in \overline{\mathbb{T}} \times W$, and $s_\infty^i = 0$; then, let \mathcal{S}^i be the set of $s^i \in \mathcal{S}_0^i$ such that any optional time τ^i for i admits \mathscr{M}^i -measurable $\widetilde{s}^i \colon \overline{\mathbb{T}} \times W \to \{0,1\}$ with $\widetilde{s}^i \in \mathcal{S}_0^i$ and $s_{\tau^i}^i = \widetilde{s}_{\tau^i}^i$: strategies for i. Let $\mathcal{S} = \times_{i \in I} \mathcal{S}^i$: strategy profiles.

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Theorem (Well-posedness). For any player $i \in I$, any optional time τ^i for i, any state process $\tilde{\chi}$, there is a unique state process χ within the information set of i given by $(\tau^i, \tilde{\chi})$ such that, for all $k \in I$, all $(t, \omega) \in [\![\tau^i(., \chi(.)), \infty]\!]$:

$$s_t^k(\omega, \chi(\omega)) = \chi_t^k(\omega).$$

(Related to [Sti92])



Application to a stylised preemption game

Notation: p and π denote the projections from $\overline{\mathbb{T}}$ onto $\overline{\mathbb{R}_+}$ and \mathfrak{w}_1+1 .

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$$s_t^i(\omega,h) = \begin{cases} 1, & \text{if } \Big(\pi(t) < \mathfrak{w}, \ \upsilon^{i,\pi(t)}(\omega) \geq \frac{\eta_{\rho(t)}^i}{1 + \eta_{\rho(t)}^i}(\omega)\Big), \\ 0, & \text{else}. \end{cases}$$

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Theorem (simplified). s defines an equilibrium. The corresponding action process is a tilting limit of sequences of processes ξ^n along sequences of classical, deterministic grids G_n , which arise from the respective discrete-time equilibrium along G_n , $n \to \infty$.

Comments on the continuous-time preemption example

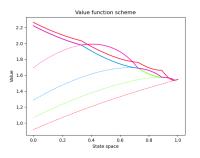


Figure: A symmetric four-player stochastic timing game in continuous time, with underlying diffusion $\mathrm{d}X_t=0.1(1-X_t^2)\,\mathrm{d}t+0.3\sqrt{X_t(1-X_t)}\,\mathrm{d}B_t$

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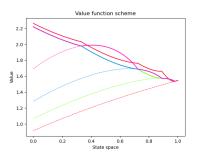


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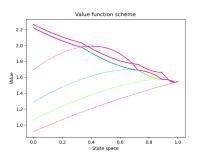


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- One strategy process, derived counterfactuals, intrinsic model of timing outcomes: A vertically-extended-continuous-time model for a vertically-extended-continuous-time problem
- Equilibrium at preemption boundary can be explained as well (though in different form); no tilting limit because both stopping times "evaporate" at infinity; no complete equilibrium convergence (payoffs are discontinuous)

Section 5

Conclusion

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- ▶ A new perspective on preemption and, more generally, stochastic differential games

Non-exhaustive list of references (for convenience)

For convenience, a non-exhaustive list of the most important references is given.

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A full bibliography can be found in the doctoral dissertation:

E. Emanuel Rapsch. On the Foundations of Dynamic Games and Probability: Decision Making in Stochastic Extensive Form. Doctoral Dissertation, TU Berlin, Berlin 2025.

It is accessible on arXiv under the DOI

https://doi.org/10.48550/arXiv.2508.04752.

- ▶ The "Tower of Babel" painting by Pieter Bruegel the Elder, in the collection of KHM Vienna, is in the public domain and has last been downloaded from https://en.wikipedia.org/wiki/The Tower of Babel (Bruegel)#/media/File:Pieter Bruegel the Elder - The Tower of Babel (Vienna) - Google Art Project - edited.jpg, On 26th August 2025.
- ▶ The figure showing the "complex tree" is licensed under CC BY-SA 4.0. The reference is: CMG Lee. (no caption). https://commons.wikimedia.org/w/index.php?curid=153344713, 21st July 2025.
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- The remaining figures have been generated by the author using the graphics software draw.io, version v27.0.9, the Matplotlib package in Python, and the tikz package in LATEX. 4□ > 4同 > 4 □ > 4 □ > □
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Thank you for your attention!



Figure: CMG Lee, ibid.; TU Berlin

Section 6

${\sf Appendix}$

Stochastic decision forests

Definition (simplified). A stochastic decision forest on (Ω, \mathscr{E}) is given by a decision forest F, a surjection $\pi \colon F \to \Omega$ whose fibres $\pi^{-1}(\{\omega\}) = T_{\omega}$ are the connected components of (F, \supseteq) , and a set $\mathbb X$ such that, if X denotes the set of F's moves:

- 1. any element $x \in X$ is a section of moves, that is, it is a map $x \colon \Omega \to X$ satisfying $\pi \circ x = \mathrm{id}_{\Omega}$;
- 2. \mathbb{X} induces a covering of F's moves, that is, $\{\mathbb{X}(\omega) \mid \mathbb{X} \in \mathbb{X}, \ \omega \in \Omega\} = X$.

The elements of X are called random moves.

Let (F, π, \mathbb{X}) be a stochastic decision forest and $W = \bigcup F$ be the set of outcomes.

Definition (simplified). An exogenous information structure is a family $\mathscr{F} = (\mathscr{F}_{\mathbb{x}})_{\mathbb{x} \in \mathbb{X}}$ of sub- σ -algebras of \mathscr{E} .

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Definition (simplified; cf. [AFR05]). A *choice* is a non-empty union of nodes. For any subset $c \subseteq W$, let

$$\downarrow c = \{x \in F \mid c \supseteq x\}, \qquad P(c) = \{x \in F \mid \exists y \in \downarrow c \colon \uparrow x = \uparrow y \setminus \downarrow c\}.$$

A choice c is

- ▶ non-redundant iff for any ω ∈ Ω with $P(c) ∩ T_ω = \emptyset$, we have $c ∩ W_ω = \emptyset$;
- ▶ complete iff for all random moves $x \in X$, $x^{-1}(P(c)) \in \{\emptyset, \Omega\}$.

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A reference choice structure is a family $\mathscr{C} = (\mathscr{C}_{x})_{x \in \mathbb{X}}$ of sets of \mathscr{C}_{x} non-redundant, complete choices c with $x^{-1}(P(c)) = \Omega$, $x \in \mathbb{X}$.

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$$\mathbf{x}^{-1}(P(c \cap c')) = \{\omega \in D_{\mathbf{x}} \mid \mathbf{x}(\omega) \in P(c \cap c')\} \in \mathscr{F}_{\mathbf{x}}.$$

Simple examples for stochastic decision forests and adapted choices



Figure: Left: A simple stochastic decision forest represented as a directed graph. Moves are indicated by circles. — Right: The induced decision tree of random nodes.

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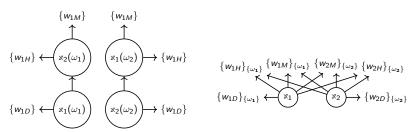


Figure: The absent-minded driver sdf, following [Gil97], and the induced poset of random nodes. Non-minimal elements are indicated by circles, respectively.

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$$\begin{aligned} x_t(\omega,f) &= \{(\omega,f') \in W \mid f'|_{[0,t)_{\mathbb{T}}} = f|_{[0,t)_{\mathbb{T}}}\}, \quad t \in \mathbb{T}, \ (\omega,f) \in W; \\ \mathsf{nodes} &= F = \{x_t(w) \mid t \in \mathbb{T}, \ w \in W\} \cup \{\{w\} \mid w \in W\}; \\ \mathsf{random moves:} \qquad & = \mathsf{x}_t(f) \colon \Omega \mapsto \mathsf{x}_t(\omega,f), \quad t \in \mathbb{T}, \ f \in \mathbb{A}^{\mathbb{T}}. \end{aligned}$$

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- ▶ "Extensive form characteristics" via **paths of action** $f: \mathbb{T} \to \mathbb{A}$ (\mathbb{T} : total order with minimum 0, "time"; \mathbb{A} : Polish action space).
- ▶ Outcomes = $W = \Omega \times \mathbb{A}^{\mathbb{T}}$:

$$\begin{aligned} x_t(\omega,f) &= \{(\omega,f') \in W \mid f'|_{[0,t)_{\mathbb{T}}} = f|_{[0,t)_{\mathbb{T}}}\}, \quad t \in \mathbb{T}, \ (\omega,f) \in W; \\ \text{nodes} &= F = \{x_t(w) \mid t \in \mathbb{T}, \ w \in W\} \cup \{\{w\} \mid w \in W\}; \\ \text{random moves:} \qquad & \texttt{x} = \texttt{x}_t(f) \colon \Omega \mapsto \texttt{x}_t(\omega,f), \quad t \in \mathbb{T}, \ f \in \mathbb{A}^{\mathbb{T}}. \end{aligned}$$

- **E**xogenous information structure: $\mathscr{F}_{\mathbb{X}_t(f)} = \mathscr{G}_t$, for filtration $(\mathscr{G}_t)_{t \in \mathbb{T}}$
- ▶ Reference choice structure: $\mathscr{C}_{\mathbb{x}_t(f)}$, for $t \in \mathbb{T}$, $f \in \mathbb{A}^{\mathbb{T}}$, is the set of all

$$\{(\omega, f') \in W \mid f'(t) \in A\}, \qquad A \in \mathscr{B}_{\mathbb{A}} \setminus \{\emptyset, \mathbb{A}\}.$$

▶ Choices at $x_t(f)$ are given by

$$c_t(f;g)=\{(\omega,f')\in\mathbb{A}^{\mathbb{T}}\mid f'|_{[0,t)_{\mathbb{T}}}=f|_{[0,t)_{\mathbb{T}}},\,f'(t)=g(\omega)\},\quad g\colon\Omega\to\mathbb{A}.$$

▶ **Theorems** (much simplified). The data induces a stochastic decision forest. $c_t(g)$ is a "suitable" choice. It is \mathscr{F} - \mathscr{C} -adapted iff g is \mathscr{G}_t -measurable.

Application to a stylised preemption game with "diffusion"

Notation: p and π denote the projections from $\overline{\mathbb{T}}$ onto $\overline{\mathbb{R}_+}$ and \mathfrak{w}_1+1 .

(1) The game: Alice and Bob can try to grab a bill worth p(t) US dollar at time t. The first who grabs gets it. However, in case of simultaneous first grabbing both are fined one unit of local currency instead (below level \mathfrak{w}).

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- (2) Exchange rates are denoted by $\eta^i > 0$, $i \in I$. Define the equilibrium candidate $s = (s^i)_{i \in I}$ as follows. Let $i \in I$, j = 3 i, and $(t, \omega, h) \in \overline{\mathbb{T}} \times W$. If $h^i(t-) = 0$, or if $t = \infty$, let $s^i_t(\omega, h) = 0$. Else, let

$$s_t^i(\omega,h) = \begin{cases} 1, & \text{if } \left(\pi(t) < \mathfrak{w}, \ \upsilon^{i,\pi(t)}(\omega) \ge \frac{\rho(t) \, \eta_{\rho(t)}^i}{1 + \rho(t) \, \eta_{\rho(t)}^i}(\omega) > 0 \right) \\ & \text{or } \left(\pi(t) < \mathfrak{w}, \ \rho(t) = 0, \ i = 2 \right), \\ 0, & \text{else.} \end{cases}$$

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Theorem (simplified). s defines an equilibrium. On p(t) > 0, the corresponding action process is a tilting limit of sequences of processes ξ^n along sequences of classical, deterministic grids G_n , which arise from the respective discrete-time equilibrium along G_n , $n \to \infty$.