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Assignment 3:

Height-weighted Quick Union with Path Compression

Tasks

- Implemented Height-weight Quick Union with Path Compression.
- Created a new class called UF_HWQUPC_Client that takes input, (number of sites) from the command line, connects the pairs if they are not connected and returns the number of connections/pairs.
- Determined the relationship between number of sites (N) and number of connections (M)

Relationship/Conclusion:

According to the data shown in the screenshot and spreadsheet, it can be clearly concluded that

M = NLogN/2,

where M is the number of connections/pairs and N is the number of sites.

Worst case time complexity for weighted quick union is O(logN).

This relationship holds true for increasing values of the number of sites.

In case of weighted Quick Union with Path Compression, initial operations might be expensive because the tree would be of log height, but then it goes cheaper as the tree becomes flatter and flatter, with every element becoming the direct descendent of the root.

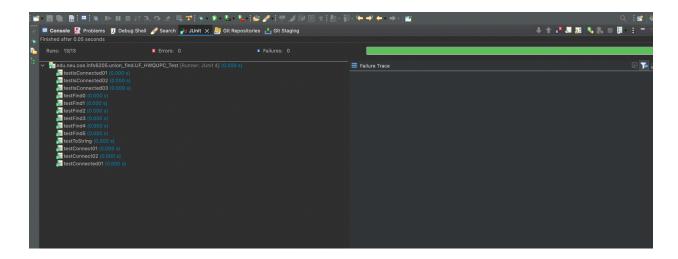
Operations in case of weighted Quick Union with Path Compression are very close to amortized constant time. Hence this algorithm is optimal.

M operations on N nodes is O (N + M Log N). This is practical for any realistic input.

Evidence:

1. Screenshot of the test cases:

The following is the snapshot of the unit test cases



2. Console Output:

On executing the main method in the UF_HWQUPC_Client class, the following output is retrieved. This output is further analyzed and plotted in the graph to reach the conclusion derived in the above section.

```
Console X Problems Debug Shell Search Julia Git Repositories Git Staging

determinated JUL-HWOUPC_Client (I) Jurya Application) Julbrary/Internet Plug-Ins/JavaApplictPlugin-plugin/Contents/Home/bin/java (Mar 4, 2022, 8-58:08 PM - 9:00:27 PM)

Enter initial number of sites:
The number you enter will be doubled in each iteration.

100

Number of sites: 200,
Number of sites: 200,
Number of sites: 400,
Number of sites: 800,
Number of sites: 800,
Number of sites: 800,
Number of sites: 12800,
Number of sites: 2200,
Number of sites: 2500,
Number of sites: 3200,
Number of connections: 29105
Number of sites: 3200,
Number of connections: 29105
Number of sites: 3200,
Number of connections: 2508326
Number of sites: 3276800,
Number of connections: 24516669
Number of sites: 3276800,
Number of connections: 24516669
Number of sites: 32020,
Number of connections: 24516669
Number of sites: 52428800,
Number of connections: 234728377
Number of connections: 234728377
Number of connections: 2445123109
```

3. Spreadsheet Data:

Here I have used sufficiently large different values of N

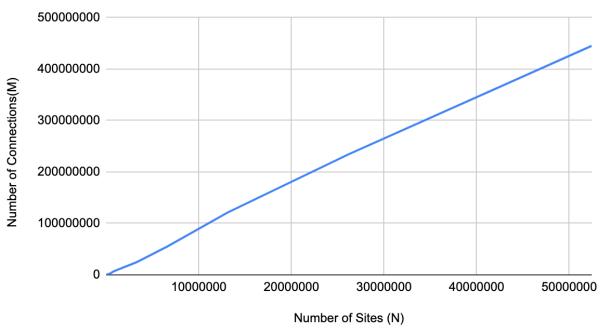
Number of Sites (N)	Number of Connections(M)
100	214
200	560
400	1230
800	2712
1600	6785
3200	13164
6400	29561
12800	69268
25600	181201
51200	279105
102400	750067
204800	1197436
409600	2508326
819200	6814379
1638400	12901713

3276800	24516669
6553600	54678886
13107200	121064916
26214400	234728377
52428800	445123109

4. Graphical Representation:

Observations from the experiments are analyzed and plotted on a graph. On close observations it can be clearly derived that the following graph is similar to M = NLogN/2.

Number of Connections(M) vs. Number of Sites (N)



Code Snippet:

```
| Tamindonthembrains | DULINGONCOLOUR | Demonstrate | Demonstrate | DULINGONCON |
| Tamindonthembrains | DULINGONCOLOUR | DEMONSTRATE | DEMONSTRATE |
| Tamindonthembrains | DULINGONCOLOUR | DEMONSTRATE | DEMONSTRATE |
| Tamindonthembrains | DULINGONCOLOUR | DEMONSTRATE |
| Tamindonthembrains | DULINGONCOLOUR | DEMONSTRATE |
| Tamindonthembrains | DULINGONCOLOUR |
```